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typeclasses allows us to expand the functionality of **existing types**. In Java, if String doesn't implement the interface you want, you

can't do anything. With typeclasses you can do it;)

with OO polymorphism we have 2 steps, interface and declaration+implementation

scala> trait Serializable
defined trait Serializable

scala> case class Car(brand: String) extends Serializable defined class Car

typeclasses add another step to polymorphism:

serializableCar: Serializable[Car] = \$anon\$10412db8f0

Abstraction

Let's try to identify a pattern in those functions:

```
scala> def sum(a: Int, b: Int): Int = a + b
sum: (a: Int, b: Int)Int

scala> def concat(a: String, b: String): String = a + b
concat: (a: String, b: String)String

scala> def and(a: Boolean, b: Boolean): Boolean = a && b
and: (a: Boolean, b: Boolean)Boolean
```

declaration

Laws

Another important feature of typeclasses are laws. Laws are properties that ensure that typeclass instances are correct.

For example, we can derive from our Sumable that it's associative.
sum(a, sum(b, c)) == sum(sum(a, b), c)

Typeclasses, together with laws, provide

- Recognizability: when we see the use of Sumable we'll know that it's an associative binary operation, regardless of the type.
 - ► Generality: If we create a datatype, and we see it's Sumable, we'll be able to use all functions that operate on Sumables.

Using typeclasses

```
def needsTypeclassContextBound[A: Sumable] = ???
```

def needsTypeclassImplicit[A](implicit x: Sumable[A])



There are lots of typeclasses libraries for Scala, but we'll use cats

in our examples.

We've already seen a very common typeclass in our examples, Sumable. It's normally called Semigroup in Functional programming.



defined trait Semigroup

Monoid

Monoids are semigroups that have an identity element. What's an identity element? one that used in combine has no effect.

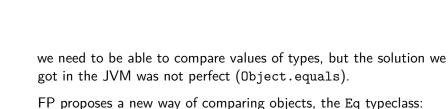
defined trait Monoid

What could created?	be the iden	tity element	t for the thre	e semigroups we

Now that we have identity we can add a couple of more laws to Monoid:

```
sum(a, sum(b, c)) == sum(sum(a, b), c) // associativity
sum(a, identity) == a // right identity
sum(identity, a) == a // left identity
```

Eq



```
trait Eq[A] {
  def eqv(a: A, b: A): Boolean
}
```

This approach to object equality has another benefit: trying to compare values of different types will result in a compiler error.

Show

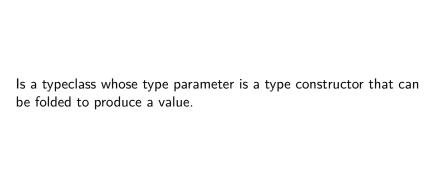
The same happens with the string representation of values. Java
tries to solve it with Object.toString, but that's not perfect

either. We might not want to be able to print passwords, for

example.

```
trait Show[A] {
  def show(a: A): String
}
```





```
trait Foldable[F[_]] {
```

def foldLeft[A, B](fa: F[A], b: B)(f: (B, A) => B): B
def foldRight[A, B](fa: F[A],lb: Eval[B])(f: (A, Eval[B])

Functor

Is a typeclass for type see how it's declared.	constructors ·	that can b	e mapped o	over. L	₋et's

```
trait Functor[F[_]] {
  def map[A, B](fn: A => B)(fa: F[A]): F[B]
}
```

Applicative

Applicatives are Functors that have two features:

- can put pure values into its context
- > can map a function lifted to its context over all its elements

```
trait Applicative[F[_]] extends Functor[F] {
  def pure[A](a: A): F[A]
```

def ap[A, B](fn: F[A => B])(fa: F[A]): F[B]

Monad

Monads are Applicatives that can sequence operations with a

flatMap method.

```
trait Monad[F[_]] extends Applicative[F] {
  def flatMap[A, B](fn: A => F[B])(fa: F[A]): F[B]
}
```

Since monads have a flatMap method, we can use any Monad $[F]$ a for comprehension!

in

Exercise 5