

## 填空题

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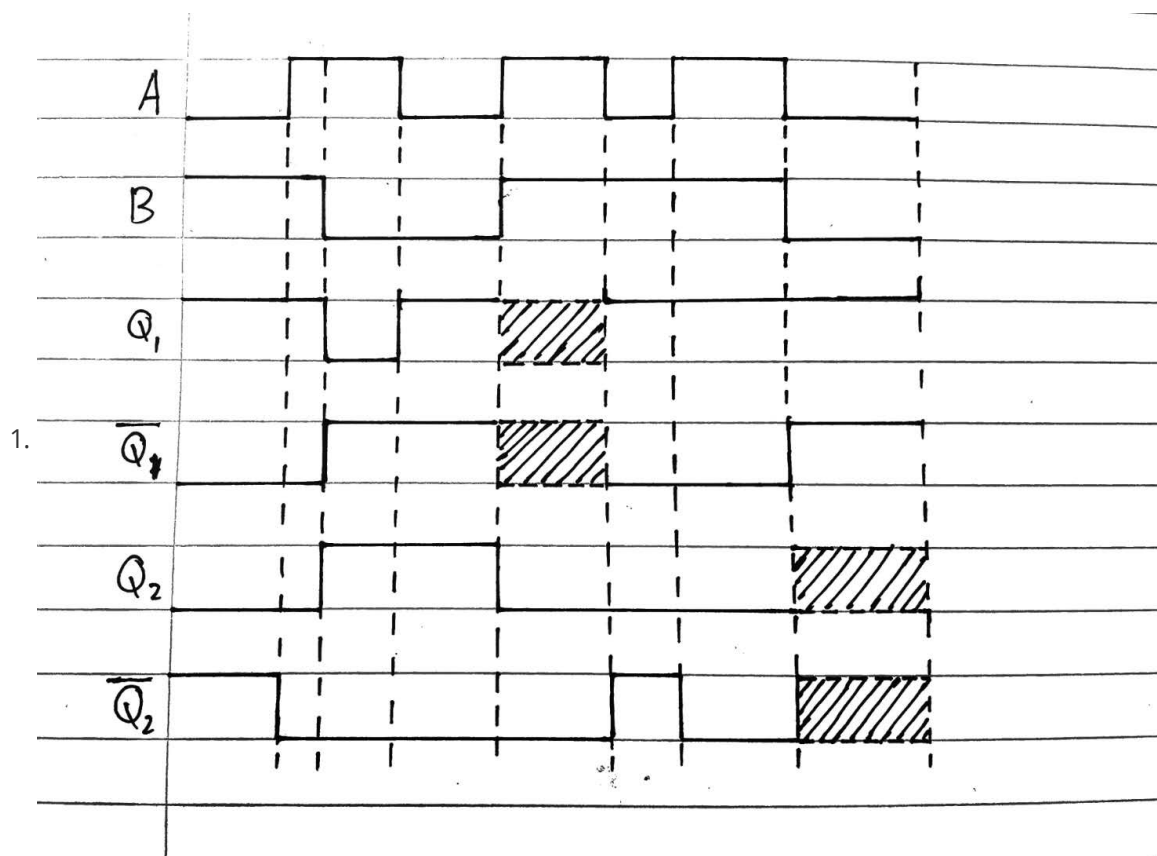
1.  $D, JK, RS$
2.  $\overline{R}_D + \overline{S}_D = 1$
3.  $J\overline{Q}^n + \overline{K}Q^n$
4. 同步时序逻辑，异步时序逻辑
5.  $n$
6. 反馈复位法，预置法
7. 100
8. 256

## 选择题

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1.  $C$
2.  $D$
3.  $B$
4.  $D$
5. ②
6. ④
7. ②
8. ②
9. ④
10. ②
11. ④
12. ①

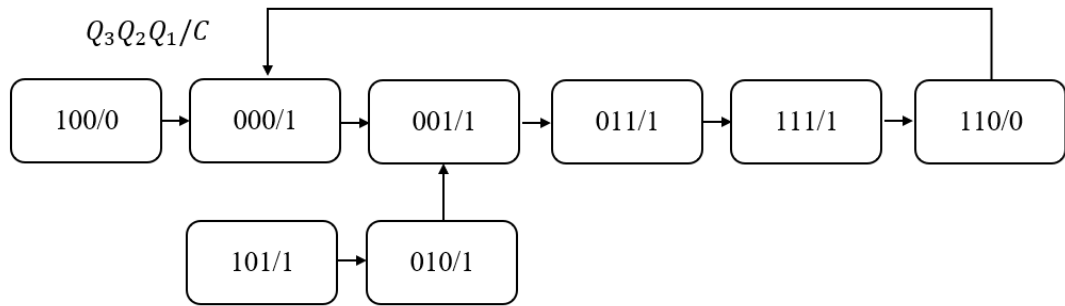
# 分析与设计题



2. (1) 驱动方程:  $D_1 = \overline{Q_3}^n, D_2 = Q_1^n, D_3 = Q_1^n Q_2^n$

状态方程:  $Q_1^{n+1} = \overline{Q_3}^n, Q_2^{n+1} = Q_1^n, Q_3^{n+1} = Q_1^n Q_2^n, C = \overline{Q_1}^n Q_3^n$

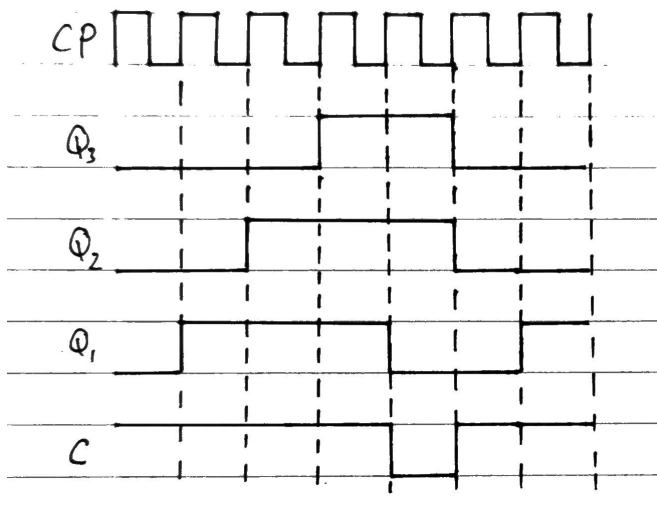
(2) 状态转换图



状态转换表

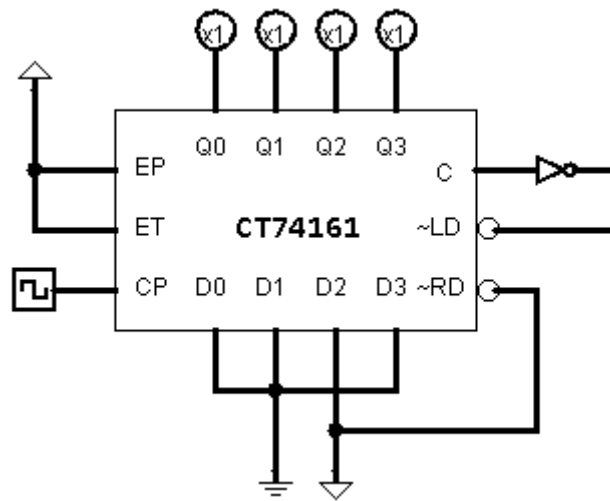
$Q_3^n$	$Q_2^n$	$Q_1^n$	$Q_3^{n+1}$	$Q_2^{n+1}$	$Q_1^{n+1}$	$C$
0	0	0	0	0	1	1
0	0	1	0	1	1	1
0	1	0	0	0	1	1
0	1	1	1	1	1	1
1	0	0	0	0	0	0
1	0	1	0	1	0	1
1	1	0	0	0	0	0
1	1	1	1	1	0	1

时序图



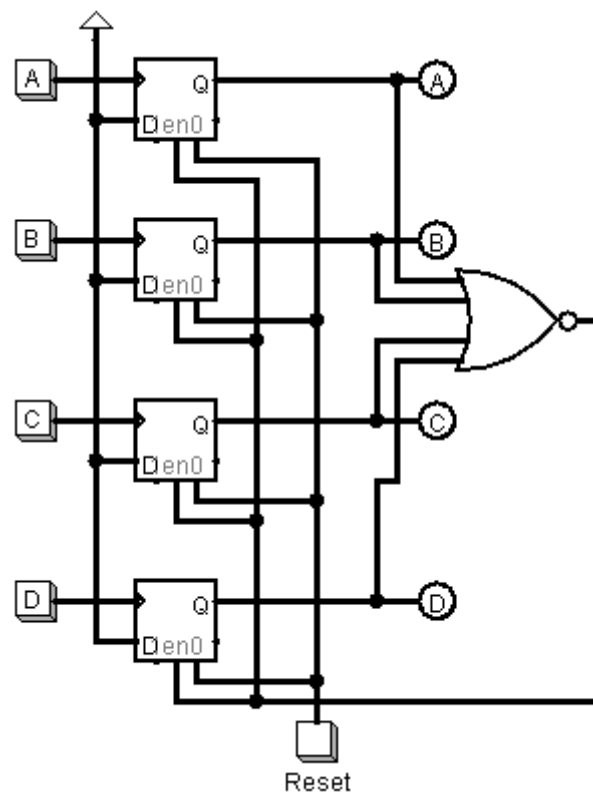
(3) 具有自启动特性的同步五进制加法计数器

3.



```
4. module cnt6 (clk, q, cout);
    input clk;
    output reg [2:0] q = 0; // 初始化q
    output reg cout = 0; // 初始化cout
    parameter [2:0] st0 = 3'b010,
                   st1 = 3'b011,
                   st2 = 3'b111,
                   st3 = 3'b110,
                   st4 = 3'b100,
                   st5 = 3'b000; // parameter 声明位宽
    always @ (posedge clk) begin
        case (q) // 状态的转移
            st0 : q = st1;
            st1 : q = st2;
            st2 : q = st3;
            st3 : q = st4;
            st4 : q = st5;
            st5 : q = st0;
            default : q = st0;
        endcase
        cout = q == st5; // cout的输出
    end
endmodule
```

5.



6. (1)

状态机定义：

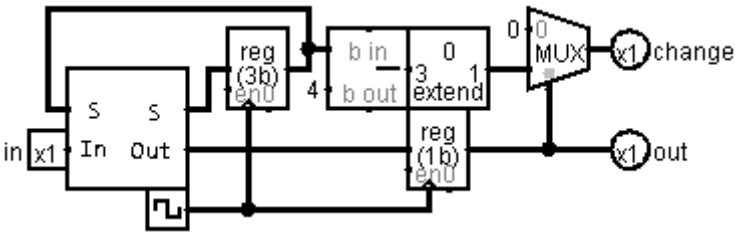
- $s = 0$ (未投入硬币)
- $s = 1$ (已投入5分)
- $s = 2$ (已投入1角)
- $s = 3$ (已投入1角5分)
- $s = 4$ (已投入2角)
- $s = 5$ (已投入2角5分)

输入定义：

- $in = 0$ (投入5分)
- $in = 1$ (投入1角)

输出定义：

- $out = 0$ (不吐出货物)
- $out = 1$ (吐出货物)
- $change = 0$ (不找钱)
- $change = 1$ (找5分)



(2)

状态转换表

$S_{20}$	$S_{10}$	$S_{00}$	in	$S_{21}$	$S_{11}$	$S_{01}$
0	0	0	0	0	0	1
0	0	0	1	0	1	0
0	0	1	0	0	1	0
0	0	1	1	0	1	1
0	1	0	0	0	1	1
0	1	0	1	1	0	0
0	1	1	0	1	0	0
0	1	1	1	1	0	1
1	0	0	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	0	0	0
1	0	1	1	0	0	0

输出逻辑真值表

$S_{20}$	$S_{10}$	$S_{00}$	out	change
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	1	0
1	0	1	1	1

次态逻辑表达式

$$\begin{aligned} S_{01} &= m_0 + m_3 + m_4 \\ &= \overline{S_{20}}\overline{S_{10}}\overline{S_{00}}in + \overline{S_{20}}\overline{S_{10}}S_{00}in + \overline{S_{20}}S_{10}\overline{S_{00}}in \\ &= \overline{S_{20}}\overline{S_{00}}in + \overline{S_{20}}S_{10}S_{00}in \\ S_{11} &= m_1 + m_2 + m_3 + m_4 \\ &= \overline{S_{20}}S_{10}\overline{S_{00}}in + \overline{S_{20}}S_{10}S_{00}in + \overline{S_{20}}\overline{S_{10}}\overline{S_{00}}in + \overline{S_{20}}\overline{S_{10}}S_{00}in \\ &= \overline{S_{20}}\overline{S_{10}}in + \overline{S_{20}}\overline{S_{10}}S_{00}in + \overline{S_{20}}S_{10}\overline{S_{00}}in \\ &= \overline{S_{20}}\overline{S_{10}}in + \overline{S_{20}}\overline{S_{10}}S_{00} + \overline{S_{20}}S_{10}\overline{S_{00}}in \\ S_{21} &= m_5 + m_6 + m_7 \\ &= \overline{S_{20}}S_{10}\overline{S_{00}}in + \overline{S_{20}}S_{10}S_{00}in + \overline{S_{20}}S_{10}S_{00}in \\ &= \overline{S_{20}}S_{10}in + \overline{S_{20}}S_{10}S_{00}in \\ &= \overline{S_{20}}S_{10}in + \overline{S_{20}}S_{10}S_{00} \end{aligned}$$

输出逻辑表达式

$$\begin{aligned} out &= S_{20} \\ change &= S_{20}\overline{S_{10}}S_{00} \end{aligned}$$