# **Assignment 1**

# 1. Simple Recursive Solution

#### Brief introduction of the solution

I used the maxmin algorithm to decide which move is the best. A tree can be used to represent all the possible boards after some legal moves. The height of the tree is decided by the search level. Every node has a score which is decided by its children, except for the leaf nodes whose score is decided by the count of bits of the according color. So this can be easily solved by a recursive algorithm as following:

To make this algorithm parallel, we can simply use cilk\_for to let the program search different branches of the tree at the same time. We also need to use a reducer to remember both the best score during the searching. Actually what we need is the best move, so the reducer I use is op\_max\_index and op\_min\_index. Because the board is 8\*8, every move can be represent by a 64 bit integer. And I consider this integer to be the index of the max/min score.

## Cilkview profiles analysis

Since the table is too long, I only show the Cilk Parallel Region part:

Level Parallelism	ı Profile	Speed Estimate					
Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs Average / strand Strands along span Average / strand on span Total number of atomic Frame count Entries to parallel region	297 480	6 - 5 - 4 - 3 - 2 - 1 - 0 - 0	) 50	100	150	200	250

	Work	21,386,301		
	Span	800,100	25 -	
	Burdened span	24,114,852		
	Parallelism	26.73	20 -	
	Burdened parallelism	0.89	15 -	
_	Number of spawns/sync			
2	Average / strand	271	10 -	
	Strands along span	956		
	Average / strand on spa		5 -	
	Total number of atomic	26517		
	Frame count	55649	0 -	
	Entries to parallel region	60	ó 50 100 150 200 250	
	Work	154,739,121		
	Span	1,204,837	120 -	
	Burdened span	36,025,094	100 -	
	Parallelism	128.43		
	Burdened parallelism	4.3	80 -	
	Number of spawns/syncs		60 -	
3	Average / strand	265	40 -	
	Strands along span Average / strand on spar	1428		
	Total number of atomic	194933	20 -	
	Frame count	413533	0 -	
	Entries to parallel region	60	0 50 100 150 200 250	
	Work	645,375,394	250 -	
	Span	1,522,881		
	Burdened span	48,119,202	200 -	
	Parallelism Burdened parallelism	423.79 13.41	150	
	Number of spawns/syncs			
4	Average / strand	242	100 -	
-	Strands along span	1896		
	Average / strand on spa		50 -	
	Total number of atomic	887581	0.1	
	Frame count	1885410	0 50 100 150 200 250	
	Entries to parallel region	60		
	Work	9,992,509,676	250 -	
	Span	2,016,643	250	
	Burdened span	59,745,042	200 -	
	Parallelism	4955.02		
	Burdened parallelism	167.25	150 -	
_	Number of spawns/syncs	11955960		
5	Average / strand	278	100 -	
	Strands along span Average / strand on span	2364 853	50 -	
	Total number of atomic	11956261		
	Frame count	25406355	0	
	Entries to parallel region	60	0 50 100 150 200 250	
	Work	79,496,532,634	250 -	
	Span	2,310,359		
	Burdened span	71,855,804	200 -	
	Parallelism Burdened parallelism	34408.74		
	Number of spawns/syncs	1106.33	150 -	
C	Average / strand	266	100 -	
6	Strands along span	2820		
	Average / strand on spa		50 -	
	Average / straing on spa			
	Total number of atomic	99422501	0-1	
	Total number of atomic Frame count	211272115	0 50 100 150 200 250	
	Total number of atomic			
	Total number of atomic Frame count Entries to parallel region	211272115 60		
	Total number of atomic Frame count Entries to parallel region	211272115 60 2,234,513,357,696	0     50     100     150     200     250	
7	Total number of atomic Frame count Entries to parallel region Work Span	211272115 60	0 50 100 150 200 250	
7	Total number of atomic Frame count Entries to parallel region Work Span Burdened span Parallelism	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56	250 - 200 -	
7	Total number of atomic Frame count Entries to parallel region Work Span Burdened span Parallelism Burdened parallelism	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66	0     50     100     150     200     250	
7	Total number of atomic Frame count Entries to parallel region Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66 2579520768	250 - 200 -	
7	Total number of atomic Frame count Entries to parallel region  Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs Average / strand Strands along span	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66 2579520768 288 3276	250 - 200 - 150 -	_
7	Total number of atomic Frame count Entries to parallel region  Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs Average / strand Strands along span Average / strand on span	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66 2579520768 288 3276	250 - 200 - 150 -	
7	Total number of atomic Frame count Entries to parallel region  Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs Average / strand Strands along span Average / strand on span Total number of atomic	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66 2579520768 288 3276 797 2579521069	250 - 200 - 150 - 100 - 50 -	
7	Total number of atomic Frame count Entries to parallel region  Work Span Burdened span Parallelism Burdened parallelism Number of spawns/syncs Average / strand Strands along span Average / strand on span Total number of atomic Frame count	211272115 60 2,234,513,357,696 2,611,369 83,662,516 855686.56 26708.66 2579520768 288 3276	250 - 200 - 150 -	



We can see that when the search level is low, there are less parallelism in our program (the burden parallelism is less than 8 which means at most we only need 8 cores), so the heavy overheads makes the parallel program slower than serial version (the speed up < 1). When the search level is high, we have enough parallelism and get a linear speed up. The results above suggest me that it is better to do serial calculation when the search level is less than 4.

# 2. Upgrade the Solution with Truncating

#### The brief introduction of truncating

I defined a serial version of s\_FindBestMove and a parallel version of p\_FindBestMove. The pseudocode is:

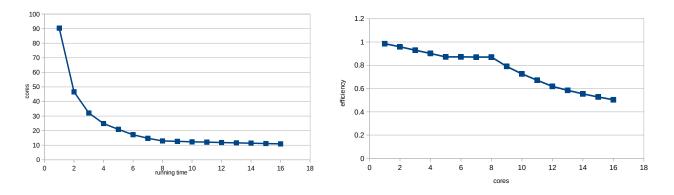
```
p_FindBestMove(node) {
    score = Nan
    IF level <= 3
        update score according to s_FindBestMove(node)

ELSE
    CILK LOOP for every children node:
        IF level <= 4
             update score according to s_FindBestMove(children node)
        ELSE
             update score according to p_FindBestMove(children node)
             END
             END LOOP
    RETURN score
}</pre>
```

s\_FindBestMove(node) is as same as the FindBestMove function is section 1.

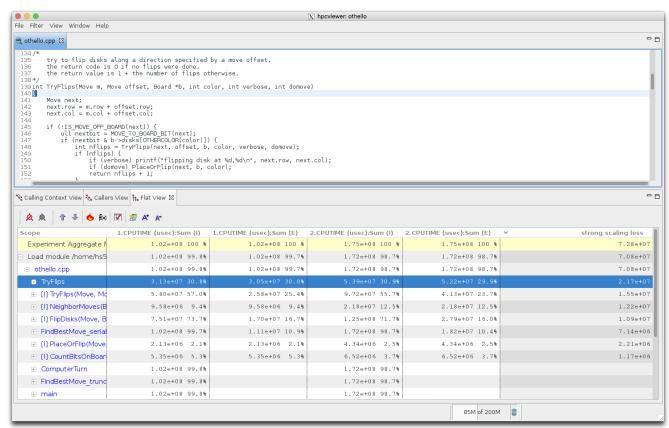
### The parallel efficiency of solution with truncating

The real running time of serial version program is 1m40.419s. So the parallel efficiency figure is as following:



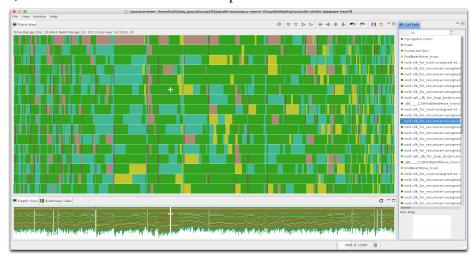
Since the efficiency begin to drop down quickly when the number of cores is larger than 8, I use hpcview and hpctraceview to see what happens there.

# 3. HpcToolkits Usage



I compared the performance under 4 cores and 16 cores. As we can see, the FindBestMove\_serial and TryFlips are two main scaling loss. I do not make any change of TryFlips because the granularity will be too fine when this function is paralleled. Actually, I have tried to do so, but the heavy overhead slowed down my program. So I recovered the change.

As for FindBestMove\_serial, it is related to the truncate level, so I tuned the parameter of truncate level to the best, which is 3. And I move to use Hpctraceviewer to see what else to do.

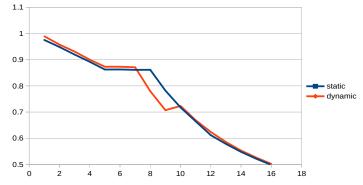


The figure above is the trace figure of program with 16 cores. We can see multiple vertical lines in the graph, which means some processors are waiting results from FindBestMove\_trunc(). Besides waiting time is longer in middle than the beginning and the end of the program. By look into the call stack, I think one reason must be that at the middle possible legal moves are more than legal moves at the beginning and the end. As a result, with the same truncate parameter, the job of FindBestMove\_serial is also increasing in the middle. So the waiting time is longer.

### Improvement by dynamic truncating

Because the jobs of FindBestMove\_serial become more when it is in the middle of the program, it maybe become worthy to make FindBestMove\_serial parallel in the middle of the program. One easiest way to do this is to low down the truncating parallel to decrease FindBestMove\_serial's jobs. I used a global variable total\_bit\_count to record the total bits on the board, the truncate parameter will be decreased to 2 when 33 < total\_bit\_count < 42.

The comparison of static truncating and dynamic truncating is as below:



The dynamic truncating worked, but only little progress is made. I think I can get a better performance after tuning the parameters. (Sorry for time is up).