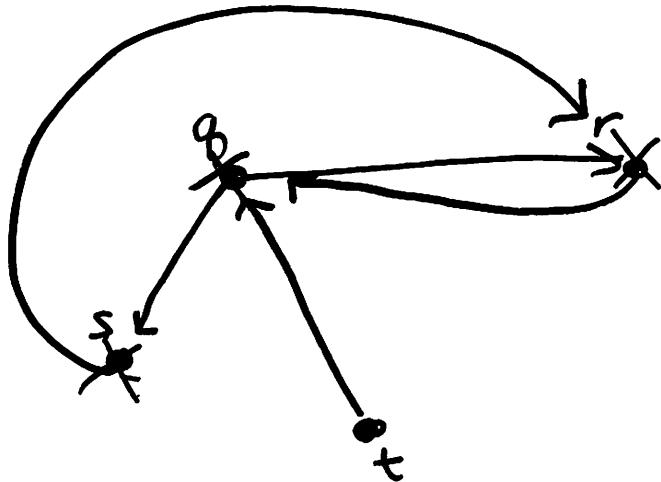


$G :$



Problem : Is there a path from s to t ?

1. $A = \{ \langle G, s, t \rangle \mid G \text{ is a D.G. that has a } \}$
Path from s to t

2. $M =$ " 1. Place mark on node s
2. Repeat until there are no nodes that can
be marked
3. → Scan all edges of G . If (a, b) exists
where a is marked and b is not, mark
 b .
4. If t is marked, accept.
If t is not marked, reject.

3. Argue that M computes in polynomial time.

$m = \text{Number of nodes}$ $e = \text{number of edges}$

- Step 1 is one operation
- Step 4 is one operation
- Steps 2+3 $m \times e$

$$1 + 1 + m \times e$$

$$n = |<G, s, t>|$$

m : number of nodes $e = \text{num edges}$

n is impacted by m and e

~~Algorithm~~

$$\frac{1 + 1 + m \cdot e}{\text{Size of input}} = O(2 + n^2) = O(n^2)$$