1. Brenda says: $A \rightarrow C \equiv \neg A \lor C$

Anthony says: $B \land \neg C$

Charles says: $\neg C \land (A \lor (A \land B) \lor B \lor (A \land B)) \equiv \neg C \land (A \lor B)$

- $\textbf{(a)} \ (\neg A \lor C) \land (B \land \neg C) \land (\neg C \land (A \lor B))$
- (b) No, because the formula is satisfiable due to $(\neg A \land B \land \neg C)$ is a model.
- (c) If all were guilty Anthony and Charles would be lying. This is because if they weren't guilty they would be telling the truth.
- (d) $\neg A \lor C = T$, A is False & C is False

$$B \land \neg C = T$$
, B is True & C is False

$$\neg C \land (A \lor B) = T$$
, C is False & A is False so B is True

Brenda is Guilty and Anthony and Charles are innocent.

2.
$$F = (P \lor Q) \rightarrow ((P \lor Q \lor \neg R) \land (R \lor P \lor Q))$$

$$\equiv (P \lor Q) \rightarrow ((P \lor Q) \lor (\neg R \land R))$$

$$\equiv$$
 (P \vee Q) \rightarrow ((P \vee Q) \vee False)

$$\equiv (P \lor Q) \to (P \lor Q)$$

$$\equiv \neg (P \lor Q) \lor (P \lor Q)$$

 $\equiv T$

F is a Tautology thus,¬F is a contradiction.

3.

$$(3.1) (P \rightarrow Q) \rightarrow ((Q \rightarrow R) \rightarrow (P \rightarrow R))$$

$$\equiv \neg (P \rightarrow Q) \lor ((Q \rightarrow R) \rightarrow (P \rightarrow R))$$

$$\equiv \neg (P \to Q) \lor (\neg (Q \to R) \lor (P \to R))$$

$$\equiv \neg(\neg P \lor Q) \lor (\neg(\neg Q \lor R) \lor (\neg P \lor R))$$

$$\equiv (P \land \neg Q) \lor (Q \land \neg R) \lor (\neg P \lor R)$$

$$\equiv P \wedge (\neg Q \vee Q) \wedge \neg R \vee (\neg P \vee R)$$

$$\equiv (P \land \neg R) \lor (\neg P \lor R)$$

$$\equiv (P \lor \neg P \lor R) \land (\neg R \lor \neg P \lor R)$$

$$\equiv (\neg P \lor R \lor \neg R) \lor (P \land \neg P)$$

$$\equiv (\neg P \lor R \lor \neg R) \lor False$$

$$\equiv \neg P \lor (R \lor \neg R)$$

$$\equiv \neg P \lor T$$

$$\equiv T$$

CNF, Tautology

$$(3.2) (P \rightarrow Q) \leftrightarrow (P \rightarrow R)$$

$$\equiv ((P \to Q) \to (P \to R)) \land ((P \to R) \to (P \to Q))$$

$$\equiv (\neg (P \to Q) \lor (P \to R)) \land (\neg (P \to R) \lor (P \to Q))$$

$$\equiv (\neg(\neg P \lor Q) \lor (\neg P \lor R)) \land (\neg(\neg P \lor R) \lor (\neg P \lor Q))$$

$$\equiv ((P \land \neg Q) \lor (\neg P \lor R)) \land ((P \land \neg R) \lor (\neg P \lor Q))$$

$$\equiv ((P \vee \neg P \vee R) \wedge (\neg Q \vee \neg P \vee R)) \wedge ((P \vee \neg P \vee Q) \wedge (\neg R \vee \neg P \vee Q))$$

$$\equiv ((\neg P \lor R \lor \neg Q) \lor (P \land \neg P)) \land ((\neg P \lor Q \lor \neg R) \lor (P \land \neg P))$$

$$\equiv ((\neg P \lor R \lor \neg Q) \lor False) \land ((\neg P \lor Q \lor \neg R) \lor False)$$

$$\equiv (\neg P \lor R \lor \neg Q) \land (\neg P \lor Q \lor \neg R)$$

CNF

$$(3.3) (P \land Q) \rightarrow (\neg P \leftrightarrow Q)$$

$$\equiv \neg (P \land O) \lor (\neg P \leftrightarrow O)$$

$$\equiv \neg (P \land Q) \lor ((\neg P \rightarrow Q) \land (Q \rightarrow \neg P))$$

$$\equiv \neg (P \land Q) \lor ((\neg (\neg P) \lor Q) \land (\neg Q \lor \neg P))$$

$$\equiv (\neg P \lor \neg Q) \lor (P \lor Q) \land (\neg Q \lor \neg P)$$

$$\equiv (\neg P \lor P) \lor (\neg Q \lor Q) \land (\neg Q \lor \neg P)$$

$$\equiv T \vee T \wedge (\neg Q \vee \neg P)$$

$$\equiv (\neg Q \vee \neg P)$$

CNF

$$(3.4) ((P \rightarrow Q) \land \neg Q) \rightarrow \neg P$$

$$\equiv \neg((P \rightarrow Q) \land \neg Q) \lor \neg P$$

$$\equiv \neg((\neg P \lor Q) \land \neg Q) \lor \neg P$$

$$\equiv ((P \land \neg Q) \lor Q) \lor \neg P$$

$$\equiv ((P \lor Q) \land (\neg Q \lor Q)) \lor \neg P$$

$$\equiv (P \lor Q) \lor \neg P$$

$$\equiv (P \lor \neg P) \lor Q$$

$$\equiv T \lor Q$$

 $\equiv T$

CNF, Tautology

4.

$$\begin{aligned} & (\textbf{4.1}) \ Q \wedge S \wedge \neg W \wedge (\neg P \vee \neg Q \vee V \vee \neg S) \wedge (S \vee \neg V) \wedge R \\ & \equiv (1 \to Q) \wedge (1 \to S) \wedge (W \to 0) \wedge (P \wedge Q \wedge S \to V) \wedge (V \to S) \wedge (1 \to R) \\ & \mathcal{A}(Q) = 1 \\ & \mathcal{A}(S) = 1 \\ & \mathcal{A}(W) = 0 \\ & \mathcal{A}(P) = 0 \\ & \mathcal{A}(R) = 1 \end{aligned}$$
 This formula is Satisfiable, $\mathcal{A}(Q) = 1$, $\mathcal{A}(S) = 1$, $\mathcal{A}(W) = 0$, $\mathcal{A}(P) = 0$, $\mathcal{A}(V) = 0$, $\mathcal{A}(R) = 1 \vDash F$

$$\textbf{4.2} \ (\neg A \lor E) \land (\neg B \land (\neg C \lor (A \to B)) \land A \land (\neg E \lor C \lor \neg D) \land (D \land (D \lor F))$$

$$\equiv (\neg A \lor E) \land (\neg B \land (\neg C \lor (A \rightarrow B)) \land A \land (\neg E \lor C \lor \neg D) \land D$$

$$\equiv (A \to E) \land (B \to 0) \land (C \land A \to B) \land (1 \to A) \land (D \land E \to C) \land (1 \to D)$$

$$\mathcal{A}(\mathbf{A}) = 1$$

$$\mathcal{A}(E) = 1$$

$$\mathcal{A}(\mathbf{B}) = 0$$

$$\mathcal{A}(\mathbf{C}) = 1/0? \leftarrow$$

$$\mathcal{A}(\mathbf{D}) = 1$$

$$A(F) = 0$$

This formula is unsatisfiable because there is a contradiction regarding the proposition C where it can be neither 1 nor 0 because either answer would give an unsatisfiable formula. Thus, there exists no model for this formula.

4.3
$$(P \land Q) \rightarrow (\neg P \leftrightarrow Q)$$

$$\equiv \neg (P \land Q) \lor (\neg P \leftrightarrow Q)$$

$$\equiv \neg (P \land Q) \lor ((\neg P \to Q) \land (Q \to \neg P))$$

$$\equiv \neg (P \land Q) \lor ((\neg (\neg P) \lor Q) \land (\neg Q \lor \neg P))$$

$$\equiv (\neg P \vee \neg Q) \vee (P \vee Q) \wedge (\neg Q \vee \neg P)$$

$$\equiv (\neg P \vee P) \vee (\neg Q \vee Q) \wedge (\neg Q \vee \neg P)$$

$$\equiv T \vee T \wedge (\neg Q \vee \neg P)$$

$$\equiv (\neg Q \lor \neg P)$$

CNF

$$\equiv (Q \land P \rightarrow 0)$$

$$\mathcal{A}(\mathbf{Q}) = 1$$

$$\mathcal{A}(\mathbf{P}) = 0$$

This formula is Satisfiable, $\mathcal{A}(Q) = 1 \& \mathcal{A}(P) = 0 \models F$