

Collision Handling

$\left\{ \begin{array}{l} \text{march 6} \rightarrow 204 \\ \text{march 17} \rightarrow 206 \\ \text{march 15} \rightarrow 407 \end{array} \right.$

ASCII

march 6 — Hash function

$\% m$



7

march 17

7

0	
1	
2	
3	
4	
5	
6	
7	<u>march 6 204</u>

Hashtable

$\text{arr}(7)$

Loose the data

collision

march 17 —

Hash function



7

1) chaining → Linked List / List

extra space

Live Session ✓

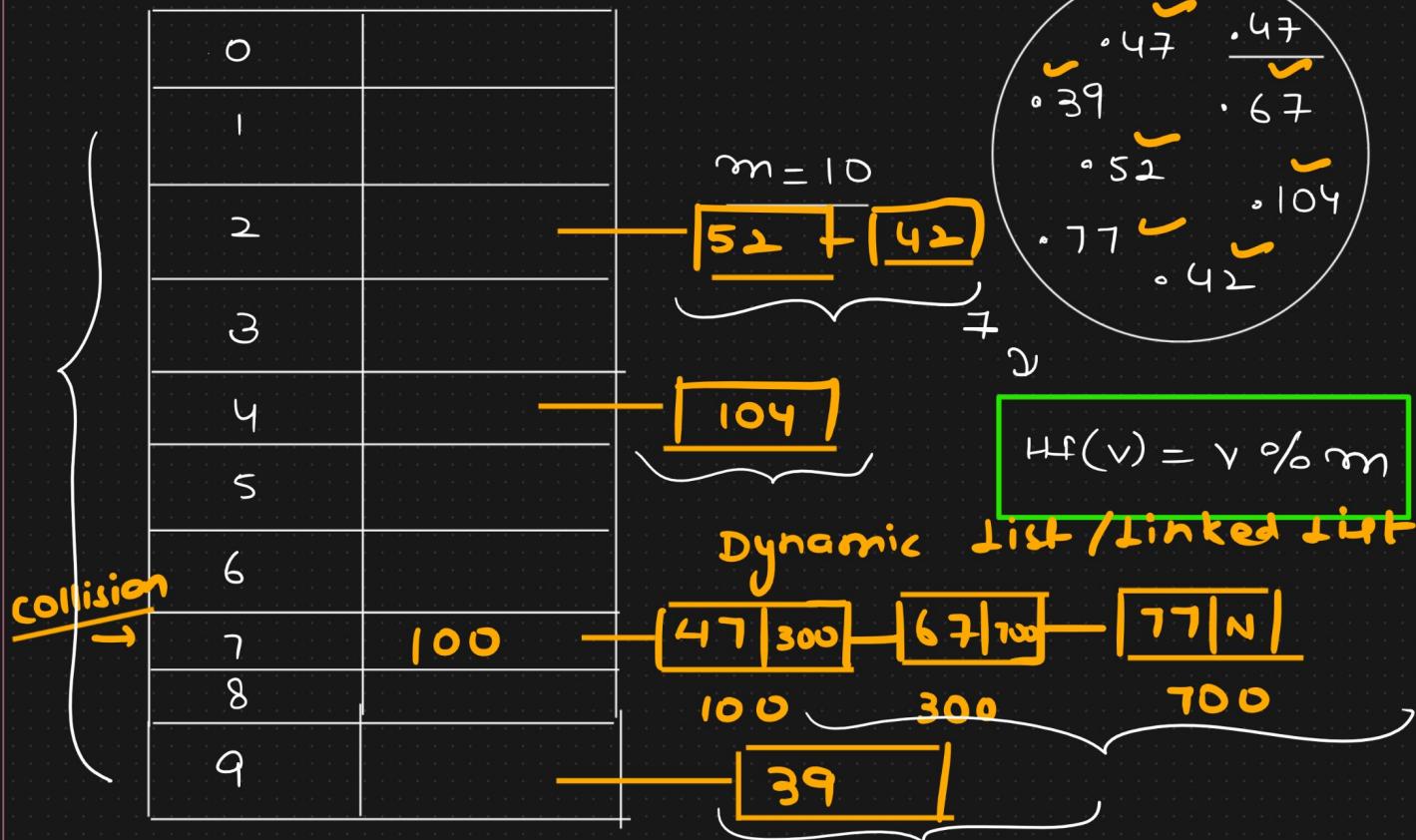
2) Open addressing → not using any extra space

- a) Linear Probing ✓
 - b) Quadratic Probing ✓
 - c) Double Hashing ✓
- change Hash Function

3) Load Factor ✓

Hash Table

chaining



Drawback

Worst case scenario

↳ 1) extra time & extra space

search $\rightarrow \Theta(n)$ TC

space complexity $\rightarrow \Theta(n)$

Open Addressing

1) Linear Probing

$$hf(v) = v \% m \quad \text{size of hash table}$$

$$\Rightarrow LP(v, i) = (hf(v) + i) \% m \quad \# \text{ collisions}$$

collision happened

↓

move the value to
next empty slot

$$50, 75, 99, 20, 35, 88, 45, 23, 55, 67 \quad hf(v) \% m$$

$$50 \% 10 \\ = 0$$

$$m = 10$$

$$LP(20, 0) = 0$$

$$LP(20, 1) = (0+1) \% 10 = 1$$

$$hf(20) = 0$$

size of Hash table
Only value = m

0	50	\leftarrow
1	20	\leftarrow
2	55	\leftarrow Empty slot
3	23	\leftarrow
4	67	\leftarrow
5	75	\leftarrow
6	35	\leftarrow
7	45	\leftarrow
8	88	\leftarrow
9	99	\leftarrow

$$hf(v) = (v \% m) \% m$$

$$= v \% m$$

value = 45

$$hf(45) = 45 \% 10 = 5$$

$$\lfloor P(45, 0) = 5 \% 10 = 5 \rightarrow \text{collision}$$

$$\lfloor P(45, 1) = (5+1) \% 10 = 6 \rightarrow \text{collision}$$

$$\lfloor P(45, 2) = (5+2) \% 10 = 7$$

value = 55

(55, 65)



$$hf(55) = 55 \% 10 = 5$$

55, 0

$$\lfloor P(55, 0) = (5+0) \% 10 = 5$$

$$\lfloor P(55, 1) = (5+1) \% 10 = 6$$

$$\lfloor P(55, 2) = (5+2) \% 10 = 7$$

$$\lfloor P(55, 3) = (5+3) \% 10 = 8$$

$$\lfloor P(55, 4) = (5+4) \% 10 = 9$$

$$\lfloor P(55, 5) = (5+5) \% 10$$

$$10 \% 10 = 0$$

$$\lfloor P(55, 6) = (5+6) \% 10$$

$$= 1$$

55, 7

$$\lfloor P(55, 7) = (5+7) \% 10$$

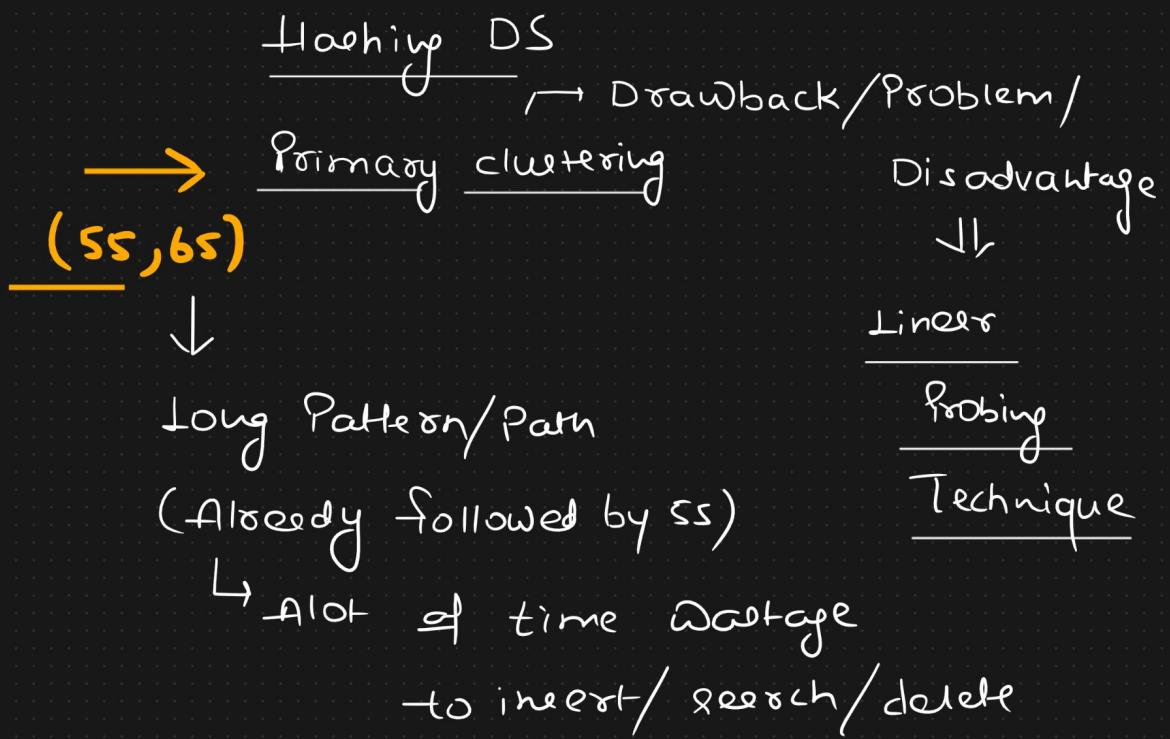
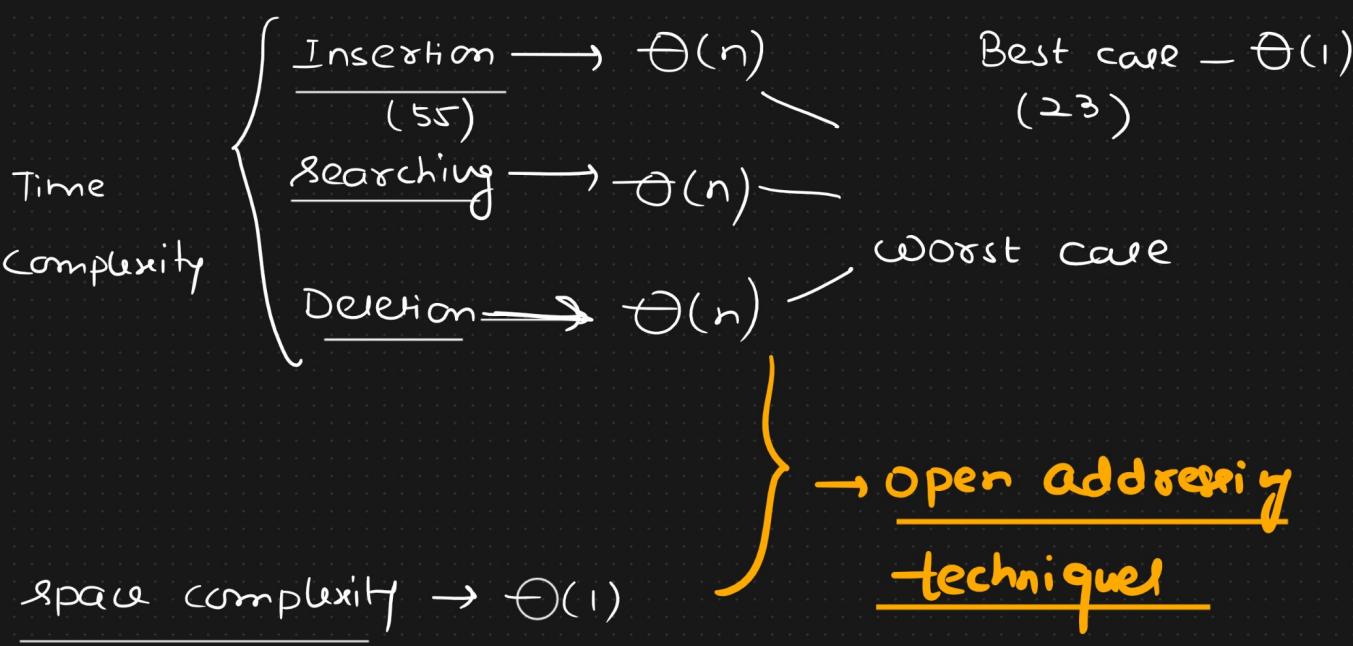
$$= 2$$

$$\lfloor P(65, 8) = (5+8) \% 10 = 3$$

lot of
time
wastage



→ traversed hash table > $m/2$



Quadratic Probing

$$\rightarrow QP(v,i) = \left(hf(v) + c_1 i + c_2 i^2 \right) \bmod m$$

$$hf(v) = v \bmod m \quad \rightarrow \text{collision num}$$

$C_1, C_2 \rightarrow \text{constant}$

$$m = 10$$

1

✓ 50, ✓ 75, ✓ 99, ✓ 20, ✓ 35, ✓ 88, ✓ 45, ✓ 23, (55), ✓ 67 | P not able to get an

0	50
1	45
2	20
3	23
4	
5	75
6	
7	35
8	88
9	99

not able to

get an
index

Data Loss

$$hf(20) = 0 \quad hf'(20)$$

$$Q_P(20, 0) = \left(\frac{\downarrow}{\circ} + \circ + \circ\right) / 10 = 0 \rightarrow \text{collision}$$

$$Q.P(20,1) = (0+1+1) \div 10 = 2 \div 10 = 2$$

$$\begin{array}{l} \text{Value} = 35 \\ h_f(35) = 35 \% 10 = 5 \end{array} \quad \boxed{C_1 = C_2 = 1}$$

$$QP(v, i) = (h_f(v) + c_1 \cdot i + c_2 \cdot i^2) \% m$$

$$\begin{aligned} QP(35, 0) &= (5 + 0 + 0) \% 10 \\ &= 5 \end{aligned}$$

$$QP(35, 1) = (5 + 1 + 1) \% 10$$

$$7 \% 10 = 7$$

major Limitation of Quadratic Probing

→ $\Rightarrow \text{Value} = 55$

already
↓ Occupied

(5, 1, 7)

cycle

$$QP(55, 0) = 5 \checkmark$$

$$QP(55, 1) = (5 + 1 + 1) \% 10 = 7 \checkmark$$

$$QP(55, 2) = (5 + 2 + 4) \% 10 = 1$$

$$QP(55, 3) = (5 + 3 + 9) \% 10 = 7 \checkmark$$

$$QP(55, 4) = (5 + 4 + 16) \% 10 = 5 \checkmark$$

$$QP(55, 5) = (5 + 5 + 25) \% 10 = 5$$

→ Secondary clustering

Double Hashing

↳ Two Hash functions

$$hf_1(v) = v \% m$$

$$\star hf_2(v) = 1 + v \bmod (m - 2)$$

$$DH(v, i) = \left(hf_1(v) + i \cdot hf_2(v) \right) \% m$$
$$(0 + 2 * 5) \% 10$$

collisions

50, 75, 99, 20, 35, 88, 45, 23, 55, 67

$$20 \rightarrow 20, 50$$

↳ cycle

0	50
1	
2	
3	
4	
5	75
6	
7	
8	
9	99

$$hf_1(20) = 20 \% 10$$
$$= 0$$

$$hf_2(20) = 1 + 20 \% 8$$
$$= 5$$

$$DH(20, 3) =$$
$$(0 + 3 * 5) \% 10$$
$$= 5$$

$$DH(20, 1) = (0 +$$
$$1 * 5) \% 10$$
$$= 5$$
$$DH(20, 2) = (0 +$$
$$2 * 5) \% 10$$
$$= 0$$

how much $\leftarrow \frac{\text{Load factor } (\alpha)}{m \leq m > 0.5}$

is the
load of
the hash table

$\alpha = \frac{m}{n}$ $n = \# \text{ data}$ $0 \leq 1$

$m = \text{size of hash table}$

$$m = 10$$

$$n = 8$$

$$\alpha = 0.8 \left(\frac{8}{10} \right)$$

$$m = 100 \quad \alpha = \frac{8}{100} = 0.08$$