

COL 100

Oct 8, 2018

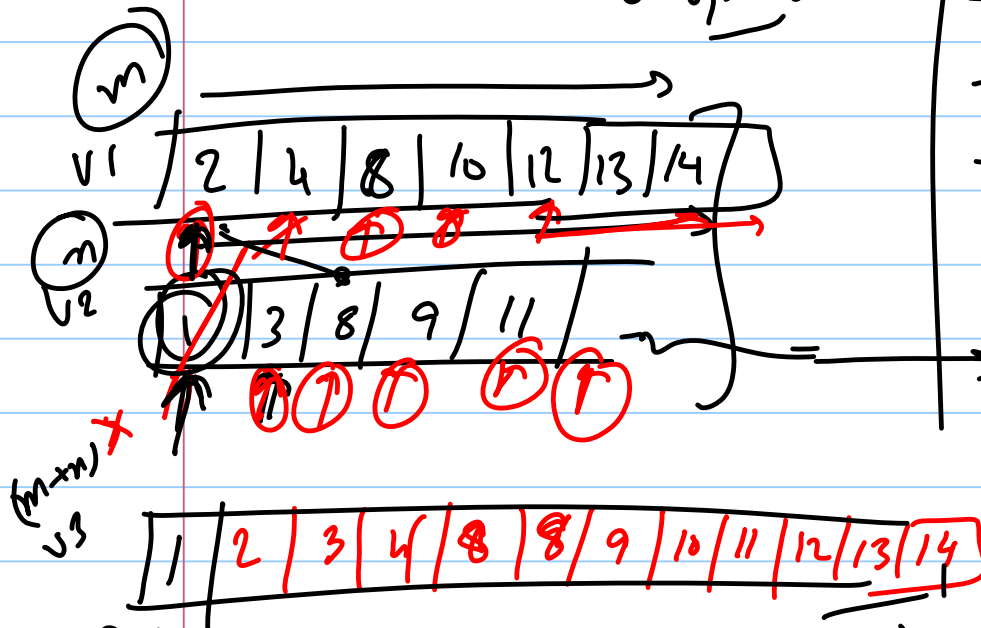
(Tentative)
Plan for Rnt

08-10-2018

→ Complexity
→ Stacks / Queues / Maps
→ Recursion
Problem Solving

①

```
int index1;
int index2;
int index3;
int size1 = v1.size();
int size2 = v2.size();
while (index1 < size1
       || index2 < size2
```



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$(n \times m) + (m \times n) = O(mn)$ where $(m+n)$ — Order Notation

$$\tan^2(bn+c) \sim \ln^2$$

$$for (int i = 1; i <= n; i++) \{$$
$$\text{für } 1 \leq j \leq n-1 \text{ gilt } \partial \leq i \leq n-1$$

~~part~~ Cont c "HELLO"; $O(n^2)$

$$\sum_{i=1}^n i = \frac{(n)(n+1)}{2} = O(n^2)$$

$$f_n(x) \quad k=1, \quad 1 \leq n, \quad k=1, \dots, n$$

lowly " " "

$$O(n^2 + n)$$

$$= O(n^2)$$

169-05

for (int x = N; x > 0; x = x - 1) {
 v.add(x);
 }
 Vector<int> v

N = 10

for (int x = 1; x <= N; x = x + 2) {

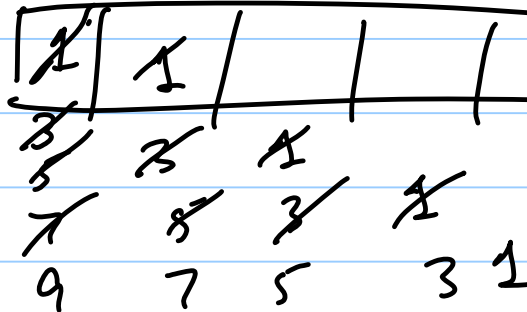
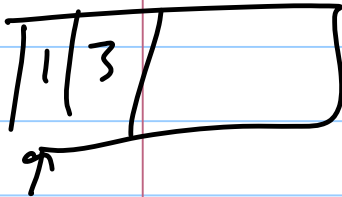
v.insert(0, x);

Cost < "N/2";

N/2
 O(N)

↓

3



O(N)

$$\sum_{i=1}^{n/2} i$$

$$= \frac{n/2 (n/2 + 1)}{2}$$

$$= O(n^2)$$

Input: n

$$N! \geq (n/2)^{n/2}$$

If you double n ,
Time $N \rightarrow 2N$

Complexity

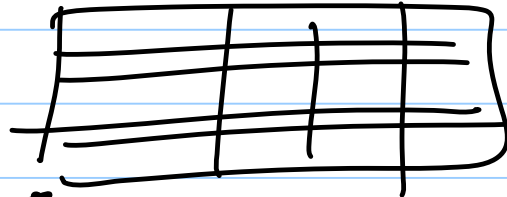
1.	Constant	$O(1)$	remain same
2.	Logarithmic	$O(\log_2 n)$	Add constant (2)
3.	Linear	$O(n)$	Double
4.	log x linear	$O(n \log n)$	slightly more than double
5.	quadratic	$O(n^2)$	four times
6.	log x quadratic	$O(n^2 \log n)$	slightly more than 4 times
7.	Cubic	$O(n^3)$	8 times the complexity
8.	Exponential	$O(2^n)$	grows extremely fast
9.	Factorial	$O(n!)$	grows extreme fast

Stacks / Queues: -

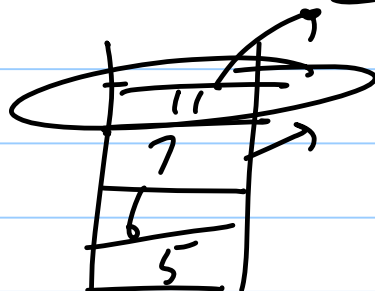
ADT

Abstract Data Type

- (1) Vectors
- (2) Arrays



(3)



First in Last out
→ Queue