

CO2100

Nov 1, 2018

Recursion

→ Solving problems which
can be broken into
self-similar (but smaller)
problems

→ Base Case(s)

Imp:-

All graded
assignments
count towards
final score

Remaining
count → $\rightarrow \square \square \square \square \square \square \square \square$
count(line) = count(line-1)

Base Case: - count(line-1) = 1 \rightarrow no one behind
you.

Pr. condition: $n \geq 0$

```
int fact(int n) {
    if (n == 0)
        return 1;
    return n * fact(n-1);
}
```

if (n < 0) {
 cout << "Error";
 return -1;
}

error handling

Base Case

fact(5)

5 * fact(4)
4 * fact(3)

3 * fact(2)

2 * fact(1)

1 * fact(0)

1

return n * fact(n-1);

work

smaller 'self-similar' problem

fact
5 * fact(4)
4 * fact(3)
fact(5)

1 n > 0

int mystery (int n) {

if (n < 10) {
return n;
}

else {

int a = n/10;
int b = n%10;

return mystery(a+b);

}

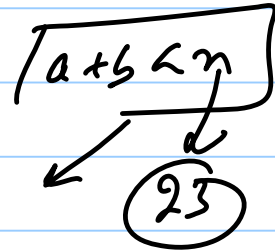
}

mystery(648) = mystery(64 + 8)

↪ mystery(7 + 2)

↪ 9

mystery(648)



1234
=

int mystery(int n) {

if (n < 10)

return n;

else {

int a = n / 10;

int b = n % 10;

return

b + mystery(a);

mystery(148)

~~148~~

← 18

= 8 + mystery(14)

← 4 + mystery(1)

← 1

648

Base Case

Palindrom:- "aka" "madam"

"step on no pats"

"able was I ex I saw elba"

↑ madam ↓
↑ ↑

```
bool checkPalindrom ( string str ) {  
    int len = str.size();  
    for ( int i = 0; i < n/2; i++ ) {  
        char ch1 = str.at(i);  
        char ch2 = str.at(n-1-i);  
        if ( ch1 != ch2 )  
            return false;  
    }  
    return true;  
}
```

palindrome recursive program

Base Case:-

1 madam 1

$$\left[\begin{array}{l} \text{str.length}() == 0 \\ \text{str.length}() == 1 \end{array} \right] \text{ true}$$

```
int bool checkPalindrome(string str) {
    int len = str.length();
    if (len == 0 || len == 1) {
        return true;
    }
    char ch1 = str.at(0);
    char ch2 = str.at(len-1);
    ?
}
```

`if (ch1 != ch2) {`
`return false;`
`}` → work

`3 else {` → `2`

`return checkPalindrome(str.substr(1, len-1));`

smaller sub problem

`}`
`3`

gkcdcfedc
↙ recursive call

$O(n)$

Be Big-O (Time Complexity) for complexity of code.

for cnt func(int n) $\sum_{\text{count} = \sum_{i=0; i \leq n; i++}} \left[\begin{array}{l} \text{count} \leq \text{"endh"} \\ \text{count} \leq \text{"endh"} \end{array} \right] \times n$

finding O (code snippet) {
 if (code snippet is single line) // Row Cms
 → O(1); // Per line
 }

findBigO(codeSnippet) {

if (code snippet is a single line) // best case
return O(1),

if (code snippet is a for-loop) // for loop
return #times-loop-executed * findBigO(code ^{(n), the} loop);

for codeBlock in code snippet
return sum_of - findBigO(codeBlock);

}
[]
[[for(—) {
[{
}]
}]

```

for (int i = 0; i < N * N; i += 3) {
    for (int j = 3; j < 219; j += 1) {
        cout << "sum: " << i + j
              << endl;
    }
}
}
cout << "sum: " << 100 << endl;
}

```

$O(N^3)$