

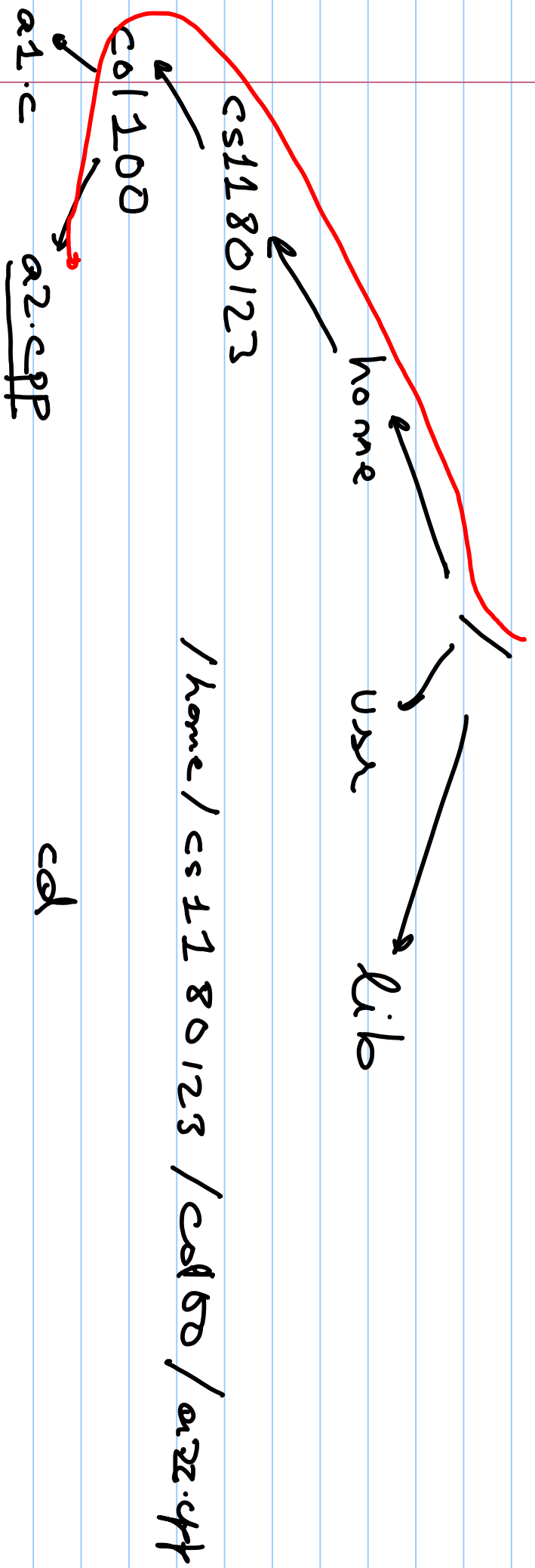
COL 100 Lecture 18

Reviews:

File 1/0

```
ifstream infile;  
if (infile.open("filename.txt")) {  
    :  
    infile.get(ch);  
    infile.getline(word);  
    infile.close();  
}
```

File names



oper car fail.

what if it fails?

if

(infile.open("filename"))

exp with side effects



Output file streams

ofstream ofile ;

ofile.open("filename.txt");

string word = "output";

int x = 3;

ofile << word << " "
<< x << endl;

ofile.close();

filename.txt:

Output 3 \n

Efficiency

- Lots of ways to solve a problem:
- some are more efficient

- How to measure

time

efficiency

memory

~~both power~~

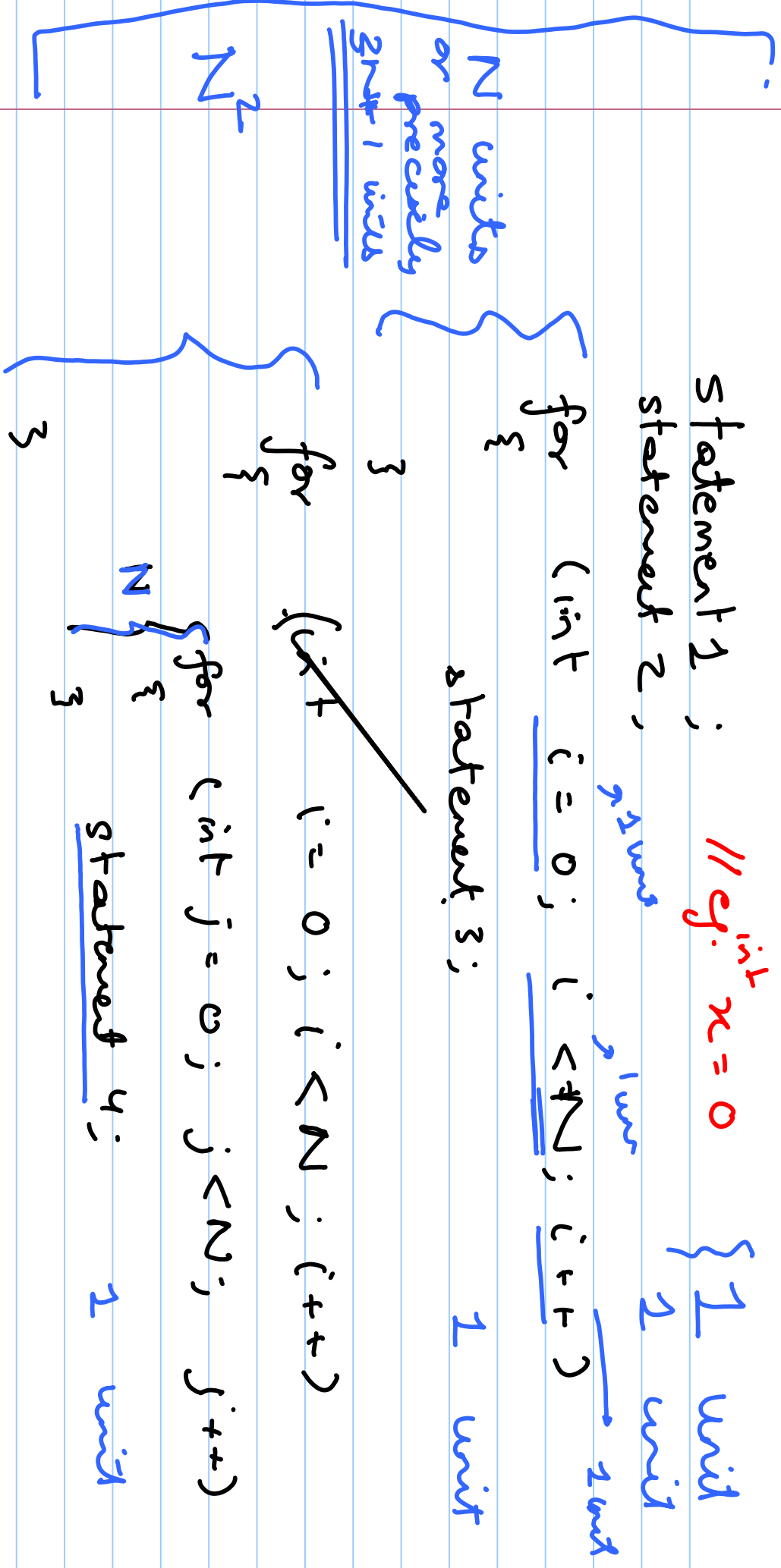
Programs

- Algorithms are better if they take less time

Unit of time

→ use count the number of units of time and get the routine

if input size = \sqrt{N} , then program takes \downarrow time



Total runtime (roughly):

$$N^2 + \underbrace{(3N+1)} + 2$$

$N^2 + N + 1$: ignore all constants

↓ only the highest powers of N matter

$$\underbrace{N^2}$$

this is the "dominating" term

Another way to think about this:

$$N^2 + N + 1 < \sim 2N^2$$

We will say that the code snippet has

$\underline{O}(N^2) \rightarrow$ Big-O notation

Runtime:

order $- N^2$
Big O $- N^2$ ($O - N^2$)

Finding Big - O ~~notable~~

- Work from the inner-most nesting level code
- Nested code multiplies
- Code at same nesting level adds

int sum = 0;

for (int i = 1; i < ~~100,000~~^N; i++)

for (int j = 1; j <= i; j++)

for (int k = 1; k <= N; k++)

sum++;

}

}

}

~~100,000 * (100,000 + 1) / 2~~

$\rightarrow \sum_{i=1}^{100,000} i$

$$\sum_{i=1}^N i$$

$$\frac{N(N+1)}{2} \times N$$

$$O(N^3)$$

1 [Vector < int > v;

$\sum_{i=1}^{N/2} i$ for (int x=1; x<=N; x+=2) {
 v.insert(0, x); v.length()
 } }
cout << v << endl;

$O(N^2)$

Complexity class

a category of algorithmic efficiency based on the of algorithm's relationship to the input size N

~~$\Theta(1)$~~

constant



$O(1)$

