Threading Support in Xv6

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Introduction

1.1 What are threads?

Within a program, a Thread is a separate execution path. It is a lightweight process that the operating system can schedule and run concurrently with other threads. The operating system creates and manages threads, and they share the same memory and resources as the program that created them. This enables multiple threads to collaborate and work efficiently within a single program. Each such thread has its own CPU state and stack, but they share the address space of the process and the environment.

1.2 Project Overview

The aim of this project was to add threading support in xv6 via adding kernel threads. The core idea is to add 2 main systemcalls Clone and Join. After that we had to implement Locks to ensure the safety of execution of our threads. Finally we programmed a file called Test-Thread to test the functionality of our program.

System Calls

2.1 Clone

This call creates a new kernel thread which shares the calling process's address space. File descriptors are copied as in fork(). The new process uses stack as its user stack, which is passed two arguments (arg1 and arg2) and uses a fake return PC (0xfffffff); a proper thread will simply call exit() when it is done (and not return). The stack should be one page in size and page-aligned. In our implementation of this system call the parent would allocate 4096 bytes (1 page) of dynamic memory using malloc and set the pointer pointing at the child's stack to this block of memory. There are a few changes made to struct proc to make this happen:

New Member	Info
int tid	New Thread Id
char * tstack	Thread stack
int Is-Thread	Set to 1 if it is a thread
int Thread-Num	The number of threads of a pro-
	cess

2.2 Join

The other new system call is int join(int tid). This call waits for a child thread that shares the address space with the calling process to exit. It returns 1 upon success and 0 if failed. After join system call the user stack dynamically allocated by the parent will be freed. It is really good to mention that the kernel stack allocated for this thread will freed when exit() is called.

2.3 Wrappers and a few details

In our implementation there are 2 user functions called :

- 1. int thread-create(void (*worker)(int*,int*),int* arg1,int* arg2)
- 2. **int** thread-join(**int** thread-id)

These 2 functions serve as wrappers to the clone and join system calls. Now i would like to show the changes before and after creating a thread in the virtual address space of the parent process:

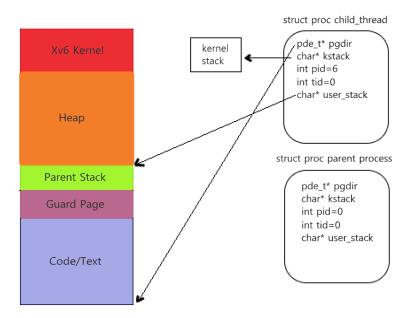


Figure 2.1: Virtual address space before clone

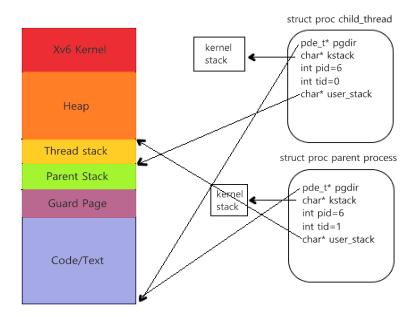


Figure 2.2: Virtual address space after clone

Lock system

3.1 Lock-init

for the safety of our program and conserving the logic of our code we need to implement a lock system(spin lock). There should be a type lock that one uses to declare a lock, and two routines(lock-acquire and lock-release). One last routine, void lock-init(lock *), is used to initialize the lock as need be (it should only be called by one thread).

3.2 Lock-acquire

The lock-acquire(lock*) is the function responsible for taking the lock if availale if not it will loop over and over until the lock is released.

3.3 Lock-release

The lock-release(lock*) is the function responsible for releasing the lock.

Test-Program

4.1 Test-Thread.c

In Our implementation we added a new c file that uses the user-space functions thread-create and thread-join to test the multithreaded functionality.