



The Frogs Game

- Given
 - n pairwise disjoint sets of size δ
- Task
 - Merging consecutive pairs of sets until there is only one set left
 - Every time cost is the sum of the sizes of k consecutive sets, left or right, whatever turns out cheaper
- Prove
 - The total cost an adversary can generate here is $O(\delta k n \log n)$



The Frogs Game

- Algorithm
 - Iteratively merge adjacent sets, choosing pairs where the minimum sum of k left/right neighbors exists
 - Calculate merge cost as the minimum between left k elements before or right k elements after the merge position
 - Repeat greedily until one set remains, accumulating minimal total merging cost
 - Pre-Compute left sum and right sum for each index to reduce the time complexity