An introduction to the Aa language

Madhav P. Desai

March 8, 2018

The Aa language

- Serves as an intermediate representation in the AHIR-V2 flow.
- Control-flow (imperative) language with support for parallelism and branching.

A simple program in Aa

Program-structure in Aa

A module

Aa Types

Aa provides a comprehensive set of types.

- Unsigned integers
 \$uint<1>, \$uint<32> etc.
- Signed integers
 \$int<1>, \$int<32> etc.
- Sized floats
 \$float<8,32>, \$float<11,52>
- Pointers:
 \$pointer<\$uint<32>> etc..
- Arrays:
 \$array[32][4] \$of \$uint<32> etc.
- Records: \$record \$uint<32> \$uint<1>
- Named Records:
 \$record [MyRec] \$pointer<MyRec>) \$uint<32>

Aa Objects

Storage objects:

\$storage A: \$array [1024] \$of \$uint<32>

Constant objects:

```
$constant A: $uint<4> := _b0011
$constant B: $uint<32> := _hf0f0f0f0
$constant C: $float<8,23> := _f2.3465e+0
```

Pipe objects:

```
$pipe A: $uint<32> $depth 4
$lifo $pipe B: $uint<8> $depth 8
$noblock $pipe C: $uint<39> $depth 2
```

Implicit objects: these are defined by statements:

$$a := (A + B)$$

They are also called *static-single-assignment* or SSA variables.

► Interface objects: Inputs and outputs of modules are treated like SSA variables. Inputs cannot be written into.



Aa Storage Objects

- ▶ Will be implemented in memory spaces.
- ▶ Access to a storage object takes 3+ cycles.
- Should use sparingly.

Aa Pipe Objects

- ▶ Blocking in nature: provide synchronization mechanism.
- ▶ FIFO or LIFO data transfer within and across modules.
- ▶ Non-blocking pipes: return a 0 if read is attempted on an empty pipe.

Aa SSA Objects

- Implemented as registers.
- ► Fast access (immediate read, single cycle write).
- Should use as often as possible.
- Note that the type of these objects is inferred from the context in which they appear.

$$a := (A + B)$$

A and B must have the same type, and the type of a is inferred to be the type of A.

Aa Expressions

Constants:

```
_b00011
_habf1
```

Simple references:

a

Array references:

```
a[0][1]
a[(I+1)][J][K]
```

Unary expressions:

Binary expressions:

```
(a <op> B)

<op> can be +,-,*,/,<<,>>,|,&,~|,~&,^,~^

==,!=,>,>=,<,<=
```

Aa More Expressions

Ternary expressions:

```
($mux <test-expr> <if-expr> <else-expr> )
e.g. ($mux (a > 0) (b+1) (b-1))
```

► Concatenation expression:

```
(a && b)
```

Bit-select expression:

Address-of expression:

```
@(a)
@(a[I])
```

Pointer-dereference expression:

```
->(ptr)
```

If it appears on the left-hand-side, it is a store, else it is a load.

Aa Statements

- Atomic statements.
- non-Atomic statements.

Atomic Aa Statements

- Simple statements.
- ► *Block* statements.

Aa Atomic Simple Statements

Assignment statements:

```
target-expression := source_expression
e.g.
a := (b + c)
```

Call statements:

```
$call fpadd32 (A B) (C)
```

Aa Atomic Block Statements

- Series-block statements.
- Parallel-block statements.
- Branch-block statements.
- Fork-block statements.

Aa Series Block Statements

```
$seriesblock [SB] {
    $storage a: $uint<32>
    a := (b + c)
    d := (a + e)
} (d => D)
```

Control-flow is sequential: statements are executed in order, token leaves statement after last statement finishes. A module body is

also a series-block.

Aa Parallel Block Statements

```
$parallelblock [PB] {
    a := (b + c)
    p := (q + r)
}
```

Control-flow: both statements started in parallel, token leaves statement after both statements have finished.

Aa Branch Block Statements

```
$branchblock [BB] {
   $merge $entry loopback
        \pi := (\text{suint}(\text{uint}(32))) \ on \pi := (\text{uint}(32)) \
                       NI $on loopback
   $endmerge
   a[I] := (b[I] + c[I])
   NI := (I+1)
   $if (NI < 16) $then
       $place [loopback]
   $endif
}
```

Control-flow: sequential, but control flow is altered by merge, place, if and switch statements.

Aa Phi Statements

- Can appear only in branch blocks, and within merge statements.
- ▶ Indicate the action to be taken depending on how control got to this statement.
- Example:

Aa Do-Pipeline-While Statements

These are not atomic, and can occur only inside a branch-block.

Control-flow: sequential, controlled by the condition check. The places \$entry and \$loopback are implicitly defined. When control enters the do-while, the token gets placed in \$entry and when control loops-back from the condition check, the token gets placed in \$loopback. The compiler will pipeline the loop by keeping \$depth iterations alive.

Aa Fork Block Statements

```
$forkblock [FB] {
   $seriesblock [S1] { ... }
   $seriesblock [S2] { ... }
   $join S1 S2 $fork
      $seriesblock [S3] { ... }
      $seriesblock [S4] { ... }
   $endjoin
   $join S3 S4 $endjoin
```

Control-flow: all statements will start in parallel, join-forks will trigger new statements etc. Token exits block when all statements finish.

Aa Volatile Statements

▶ If you do not want to use a register:

```
volatile p := (q + r)
```

- ▶ p will not necessarily be implemented as a register. You can only assume that p is equivalent to (q + r).
- Useful in describing combinational logic.

Aa Summary

- An Aa program consists of a collection of module descriptions and object declarations.
- Each module is a sequence of statements.
- Statements can be simple or block-structured.
- Specific blocks allow branching, while others describe parallelism.
- Objects in an Aa program can be FIFOs (pipes), storage or implicitly inferred.
- Almost everything you can express in C can be expressed, but we are closer to hardware.

Look at the FIR again