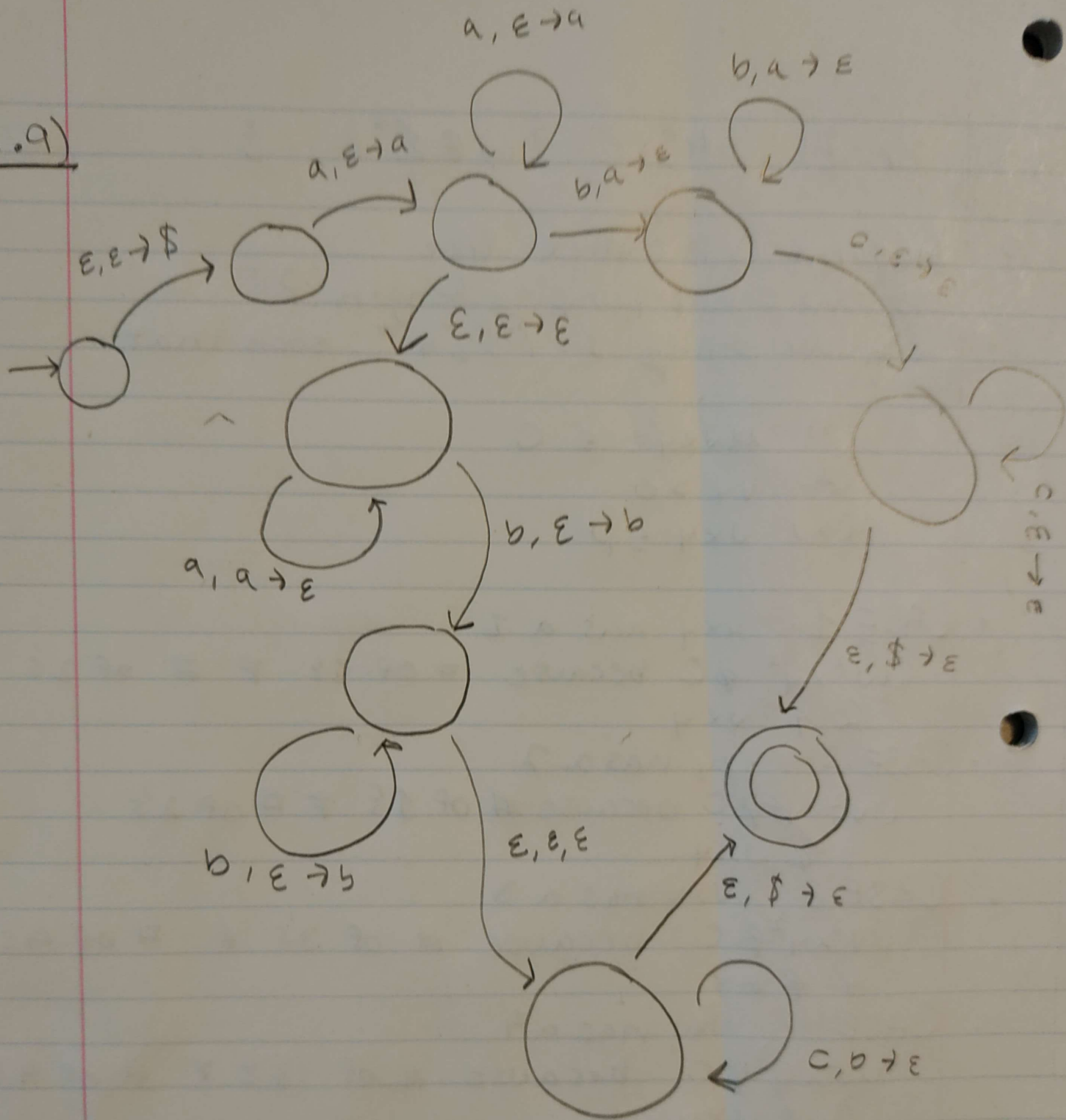


2.9)



2.11)

$E \rightarrow E + T \mid T$

$T \rightarrow T \times F \mid F$

$F \rightarrow (E) \mid a$

$\epsilon, E \rightarrow T$

$\epsilon, T \rightarrow F$

$\epsilon, F \rightarrow a$

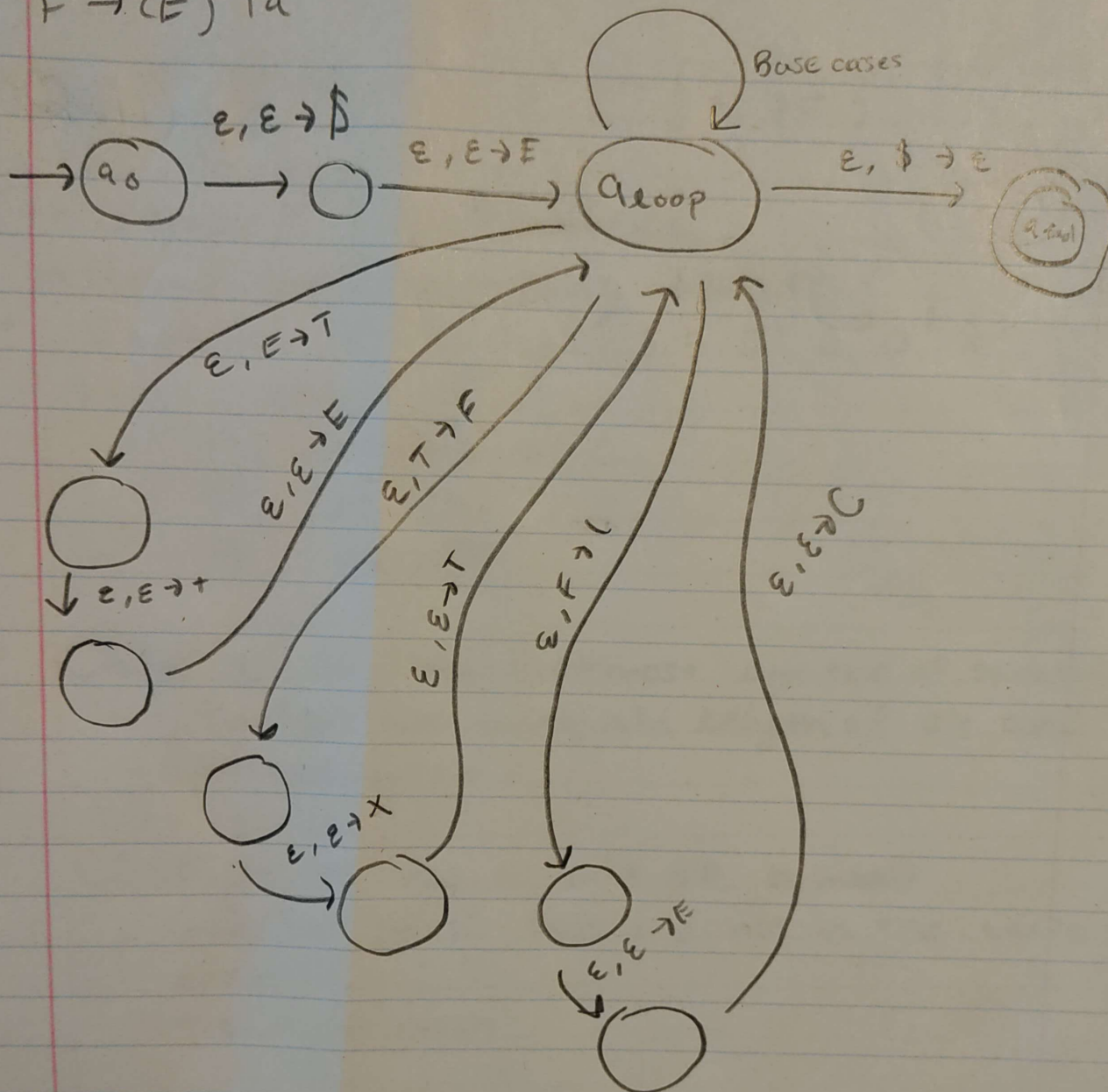
$\epsilon, + \rightarrow \epsilon$

$\epsilon, \times \rightarrow \epsilon$

$\epsilon, (\rightarrow \epsilon$

$\epsilon,) \rightarrow \epsilon$

$\epsilon, a \rightarrow \epsilon$



2.18)

- a) let D be a DFA that recognizes R
let P be a PDA that recognizes C

* Prove C recognizes $C \cap R$

Q set of states of P

Q set of states of D

F_P is set of states accepted by P

F_D is set of states accepted by D

PDA that recognizes $C \cap R$ stops if
 $q \in F_P \times F_D$

So $C \cap R$ is recognized by \bar{P}

- Therefore by proof of construction $C \cap R$ is context free

- b)
- R is the regular language $a^*b^*c^*$
 - If A were a CFL then $A \cap R$ would be a CFL
 - But since $A \cap R = \{a^n b^n c^n \mid n \geq 0\}$ proves $A \cap R$ is not context free via pumping lemma contradiction, then A is not context free.

2.30a) $C = \{0^n 1^n 0^n 1^n \mid n \geq 0\}$

- Assume C is context free
- Assume C has pumping length P
- Consider the string $s = 0^P 1^P 0^P 1^P = UVXYZ$ such that
 - 1) $UV^iXY^iZ \in C$
 - 2) $|VY| > 0$
 - 3) $|VXY| \leq P$

CASE 1: VY have at most one type of symbol
 UV^2XY^2Z has unequal length of 0's and 1's
 - contradiction

CASE 2: VY has a mix of symbols
 UV^2XY^2Z has symbols not in the correct order
 - contradiction

By proof of contradiction C is not a context free language

2.32 $\Sigma = \{1, 2, 3, 4\}$ $C = \{w \in \Sigma^+ \mid \dots\}$

- Assume C is context free
- Assume C has pumping length ' p '
- consider string $s = 1^p 2^p 4^p$ such that:

- (1) $uv^i xy^i z \in C$
- (2) $v \neq \epsilon$
- (3) $|vxy| \leq p$

- CASE 1: vxy has a 1
 $uv^2xy^2 \notin C$ because # of 1's \neq # of 2's
 $2 \notin vxy$

- CASE 2: vxy has a 2
 $uv^2xy^2 \notin C$ because # of 1's \neq # of 2's
 $1 \notin vxy$

- CASE 3: vxy has a 3
 $uv^2xy^2 \notin C$ because # of 3's \neq # of 4's
 $4 \notin vxy$

- CASE 4: vxy has a 4
 $uv^2xy^2 \notin C$ because # of 3's \neq # of 4's
 $3 \notin vxy$

By proof of contradiction, C is not context free.