SRM

SRM INSTITUTE OF SCIENCE AND TECHNOLOGY

Ramapuram Campus

Department Of Computer Science and Engineering



18CSC261T / FORMAL LANGUAGE AND AUTOMATA CONTINUOUS LEARNING ASSESMENT-II

SET-I

Branch: BTech CSBS

Year/Sem: II/III

Marks: 50

Part – A (10 X 1 = 10 marks) Answer all the questions

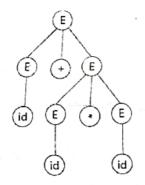
- 1. Which one of the following languages over $\Sigma = \{a,b\}$ is NOT context-free?
 - A. $\{wwR|w\in\{a,b\}*\}$
 - B. $\{\text{wanbnwR}|\text{w}\in\{\text{a,b}\}*,\text{n}\geq 0\}$
 - C. $\{\text{wanwRbn}|\mathbf{w}\in\{a,b\}*,n\geq0\}$
 - D. $\{anbi|i \in \{n,3n,5n\}, n \ge 0\}$
- 2. Automaton accepting the regular expression of any number of a's is
 - A. a*
 - B. ab*
 - C. (a/b)*
 - D. a*b*c
- 3.Consider the following context-free grammar over the alphabet $\Sigma = \{a,b,c\}$ with S as the start symbol:S \rightarrow abScT|abcTT \rightarrow bT|bWhich one of the following represents the language genrated by the above grammar?
 - A. $\{(ab)n(cb)n|n\geq 1\}$
 - B. {(ab)ncbm1cbm2
 - C. cbmn $[n,m1,m2,mn\geq 1]$
 - D. {(ab)n(cbm)n|m,n≥1;
- 4. Assume the statements S1 and S2 given as:S1: Given a context free grammar, there exists an algorithm for determining whether L (G) is infinite.S2: There exists an algorithm to determine whether two context free grammars generate the same language. Which of the following is true?
 - A. SI is correct and S2 is not correct
 - B. Both S1 and S2 are correct
 - C. Both S1 and S2 are not correct
 - D. S1 is not correct and S2 is correct
- 5. A student wrote two context-free grammars G1 and G2 for generating a single C-like array declaration. The dimension of the array is at least one. For example, int a[10][3]; The grammars use D as the start symbol, and use six terminal symbols int; id[] num.Grammar G1 Grammar

- A. Both G1 and G2
- B. Only G1
- C. Only G2
- D. Neither G1 nor G2

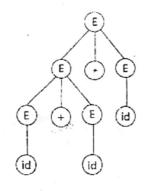
6. For any two languages L1 and L2 such that L1 is context-free and L2 is recursively enumerable but not recursive, which of the following is/are necessarily true?

- L1 (complement of L1) is recursive
- II. L2 (complement of L2) is recursive
- III. L1 is context free
- IV. L1UL2 is recursively enumerable *
- A. I
- B. III
- C. III and IV
- D. I and IV

7. ConsiderE \rightarrow E + E E \rightarrow E * E E \rightarrow id *



Parse tree 1



Parse tree 2

- A. $E \rightarrow E + T, T \rightarrow T * F, F \rightarrow id$
- B. $E \rightarrow E + T$, $T \rightarrow F$
- $C,\ E \to E + T \ , E \to T \ T \to T * F, T \to F$
- D. $E \rightarrow E + T$, $E \rightarrow T$

8.Push down automata accepts _____languages.

- A. Type 3
- · B. Type 2
 - C. Type 1
 - D. Type 0

- 9.A push down automaton with only symbol allowed on the stack along with fixed symbol.
 - A. Embedded PDA
 - B. Nested Stack automata
 - · C. DPDA
 - D. Counter Automaton
- 10. Which of the following option resembles the given PDA?
 - A. $\{0^n1^n|n>=0\}$
 - B. $\{0^n1^2n|n>=0\}$
 - C. $\{0^2n1^n|n>=0\}$
 - D. $\{0^21^n|n>=0\}$

PART-B $(4 \times 4 = 16 \text{ marks})$

- M. Define CFG.
- 12, Construct a CFG over {a,b} generating a language consisting of equal number if a's and b's.
- 13, Eliminate all unit production S->Aa|B;B->A|bb;A->a|bc|B.
- 14. Define Push Down Automata.
- 15. Is it true that non-deterministic PDA is more Powerful than that of deterministic PDA? Justify your answer.
- Construct a PDA to accept a language {(ab)ⁿ|n>=1} by empty stack.

Part - C
$$(2*12 = 24 \text{ marks})$$

- V1. a. Convert to Chomsky Normal Form S \rightarrow ASA | aB, A \rightarrow B | S, B \rightarrow b | ϵ OR
- 17. b. Consider the grammar, $S \rightarrow iCts$, $S \rightarrow iCtSeS$, $S \rightarrow a$, $C \rightarrow b$
 - 1) Construct leftmost derivation for the sentence W=libtibtaea.
 - 2) Show the corresponding Parse tree for the above sentence.
 - 3) Is the above grammar ambiguous? If so prove it.
- 18 .a. Design a PDA for $L = \{WCW^R \mid W \text{ in } (0+1)^*\}$ OR
- 18 .b. Obtain CFG for the PDA as given below

$$\partial(q_0,0,Z)=(q_0,AZ)$$

$$\partial(q_0,0,A)=(q_0,AA)$$

$$\partial(q_0,1,A)=(q_0,\varepsilon)$$

$$\hat{c}(q_1, 1, A) = (q_1, \varepsilon)$$

$$\partial(q_0, \varepsilon, A) = (q_1.\varepsilon)$$

$$\partial(q_0, \varepsilon, Z) = (q_1, \varepsilon)$$

HODESTE

7.21/1/21