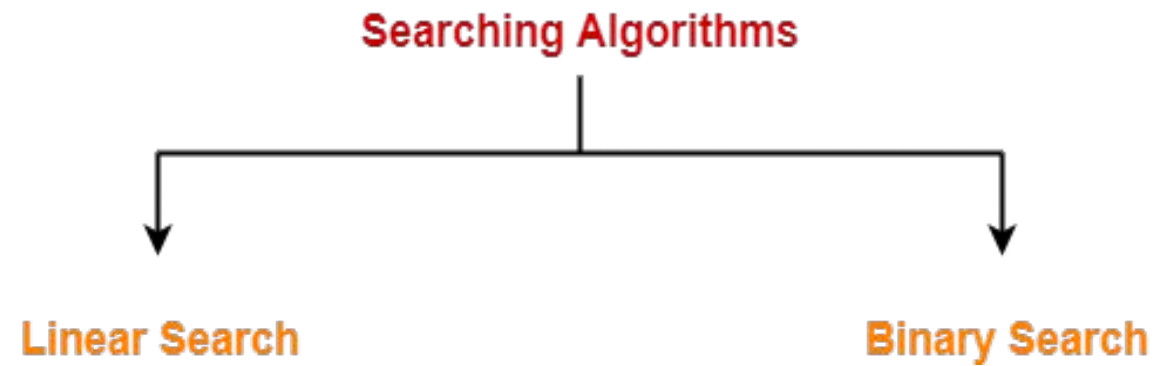


# **BINARY SEARCH**

BinarySearch

Complexity of binary search

# **BINARY SEARCH**



# Binary Search

**Session Learning Outcome-SLO-**Solve problems using divide and conquer approaches

- **Motivation of the topic**

Binary Search is one of the fastest searching algorithms.

- It is used for finding the location of an element in a linear array.
- It works on the principle of divide and conquer technique.

*Binary Search Algorithm can be applied only on **Sorted arrays**.*

- So, the elements must be arranged in-
  - Either ascending order if the elements are numbers.
  - Or dictionary order if the elements are strings.
- To apply binary search on an unsorted array,
  - First, sort the array using some sorting technique.
  - Then, use binary search algorithm.

- Binary search is an efficient searching method. While searching the elements using this method the most essential thing is that the elements in the array should be a sorted one.
  - An element which is to be searched from the list of elements stored in the array  $A[0 \dots n-1]$  is called as Key element.
  - Let  $A[m]$  be the mid element of array  $A$ . Then there are three conditions that need to be tested while searching the array using this method.
  - They are given as follows
    - ❖ If  $key == A[m]$  then desired element is present in the list.
    - ❖ Otherwise if  $key < A[m]$  then search the left sub list
    - ❖ Otherwise if  $key > A[m]$  then search the right sub list
- The following algorithm explains about binary search.

# Recursive Binary search algorithm

## Algorithm

Input: A array  $A[0...n-1]$  sorted in ascending order and search key  $K$ .

Output: An index of array element which is equal to  $k$

Low  $\leftarrow$  0; high  $\leftarrow$  n-1

while low  $\leq$  high do

    mid  $\leftarrow$  (low+high)/2

    if  $K=A[\text{mid}]$  return mid

    else if  $K < A[\text{mid}]$  high  $\leftarrow$  mid-1

    else low  $\leftarrow$  mid+1

return

# EXAMPLE FOR BINARY SEARCH

**Step:1** Consider a list of elements sorted in array A as

0	1	2	3	4	5	6	
10	20	30	40	50	60	70	
Low			high				

The Search key element is key=60

Now to obtain middle element we will apply formula

$$\text{mid} = (\text{low} + \text{high}) / 2$$

$$\text{mid} = (0 + 6) / 2$$

$$\text{mid} = 3$$

Then check  $A[\text{mid}] == \text{key}$ ,  $A[3] = 40$

$A[3] \neq 60$  Hence condition failed

Then check  $\text{key} > A[\text{mid}]$ ,  $A[3] = 40$

$60 > A[3]$  Hence condition satisfied so search the Right Sublist

## Step 2:

- The Right Sublist is

50	60	70
----	----	----

- Now we will again divide this list to check the mid element



- $\text{mid} = (\text{low} + \text{high}) / 2$
- $\text{mid} = (4 + 6) / 2$
- $\text{mid} = 5$
- Check if  $A[\text{mid}] == \text{key}$
- (i.e)  $A[5] == 60$ . Hence condition is satisfied. The key element is present in position 5.
- **The number is present in the Array A[ ] at index position 5.**

Thus we can search the desired number from the list of the elements.

# ANALYSIS

- The basic operation in binary search is comparison of search key with array elements.
- To analyze efficiency of binary search we must count the number of times the search gets compared with the array elements.
- The comparison is also called three way comparisons because the algorithm makes the comparison to determine whether key is smaller, equal to or greater than  $A[m]$ .
- In this algorithm after one comparison the list of  $n$  elements is divided into  $n/2$  sub list.
- The worst case efficiency is that the algorithm compares all the array elements for searching the desired element.
- In this method one comparison is made and based on the comparison array is divided each time in  $n/2$  sub list.



- Hence worst case time complexity is given by

$$C_{\text{worst}}(n) = C_{\text{worst}}(n/2) + 1 \quad \text{for } n > 1$$

Time required to      one comparison made

Compare left or      with middle element

Right sub list

Also  $C_{\text{worst}}(1) = 1$

- But as we consider the rounded down values when array gets divided the above equation can be written as

$$C_{\text{worst}}(n) = C_{\text{worst}}(n / 2) + 1 \quad \text{for } n > 1$$

$$C_{\text{worst}}(1) = 1$$

- We can analyse the best case , Worst case and Average case . The time complexity of binary search is given as follows

Best case	Average case	Worst Case
$\theta(1)$	$\theta(\log n)$	$\theta(\log n)$

- In conclusion we are now able to completely describe the computing time of binary search by giving formulas that describe best, average and worst cases
- Successful searches      unsuccessful searches  
 $\theta(1)$      $\theta(\log n)$      $\theta(\log n)$        $\theta(\log n)$
- best   average   worst      best, average, worst

- **Advantages of Binary search:**

- Binary search is an optimal searching algorithm using which we can search the desired element very efficiently

- **Disadvantages of binary Search :**

- This Algorithm requires the list to be sorted . Then only this method is applicable

- **Applications of binary search:**

- The binary search is an efficient searching method and is used to search desired record from database
- For solving with one un known this method is used

## Summary:

- Binary Search time complexity analysis is done below-
  - In each iteration or in each recursive call, the search gets reduced to half of the array.
  - So for  $n$  elements in the array, there are  $\log_2 n$  iterations or recursive calls.
- **Time Complexity of Binary Search Algorithm is  $O(\log_2 n)$ .**
  - Here,  $n$  is the number of elements in the sorted linear array.

## Home assignment:

- Search the Element 15 from the given array using Binary Search Algorithm.

3	10	15	20	35	40	60
a[0]	a[1]	a[2]	a[3]	a[4]	a[5]	a[6]

**Binary Search Example**