

# Merge Sort with Time complexity analysis

Design and Analysis of Algorithms

# Session Learning Outcome-SLO

- At the end of this session, you should be able to perform merge sort and also evaluate its time complexity

# Introduction

- Attributed to a Hungarian mathematician John von Neumann
- Used for sorting unordered arrays
- Uses divide-and-conquer strategy
- **1<sup>st</sup> phase** is to **divide** the array of numbers into 2 equal parts
  - If necessary, these subarrays are divided further
- **2<sup>nd</sup> phase** is the **conquer** part
  - Involves sorting of subarrays recursively and combine the sorted arrays to give the final sorted list

# Merge sort – informal procedure

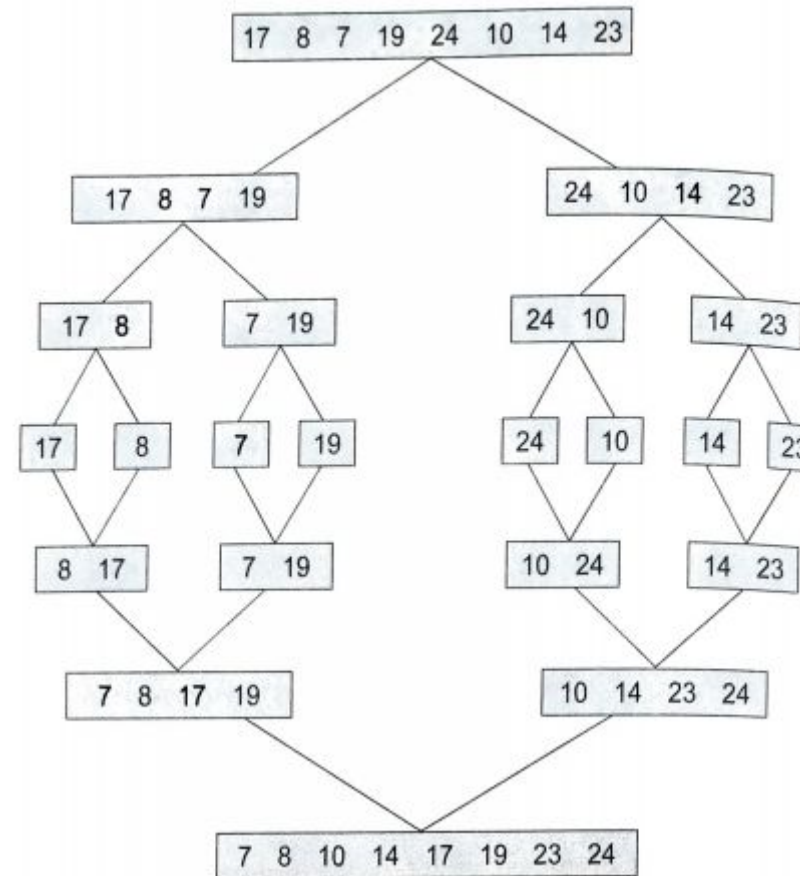
Step1: Divide an array into subarrays B and C

Step 2: Sort subarray B recursively; this yields B sorted subarray

Step 3: Sort subarray C recursively; this yields C sorted subarray

Step 4: Combine B and C sorted subarrays to obtain the final sorted array A

# MergeSort- An Example



# Algorithm mergesort(A[first .. Last])

**Input:** Unsorted array A with first =1 and last=n

**Output:** Sorted array A

Begin

    if (first == last) then

        return A[first]

    else

        mid = (first+last)/2

    for  $i \in \{1, 2, \dots, \text{mid}\}$

        B[i] = A[i]

    End for

    for  $i \in \{\text{mid} + 1, \dots, n\}$

        C[i] = A[i]

    End for

    mergesort(B[1 .. mid])

    mergesort(C[mid + 1 .. n])

    merge(B,C,A)

End

# Algorithm merge(B,C,A)

**Input:** Two sorted arrays B and C

**Output:** Sorted array A

Begin

i=1

j=1

k=1

m = length(B)

n = length(C)

while ((i<=m) and (j <= n)) do

if(B[i] < C[j]) then

A[k] = B[i]

i=i+1

else

A[k] = C[i]

j=j+1

End if

k=k+1

end while

if (i > m) then

while k <= m + n do

A[k] = C[j]

j,k = j+1, k+1

end while

else if (j < n) then

while (k <= m+n) do

A[k] = B[j]

i,k = i+1,k+1

end while

end if

return (A)

end

# Complexity Analysis

- $T(n) = 2T\left(\frac{n}{2}\right) + n - 1, \text{ for } n \geq 2$   
     $= 1, \text{ for } n=2$   
     $= 0, \text{ when } n < 2$



# Summary

- Uses divide and conquer strategy
- a comparison based sort
- out of place merge sort
- Stable Algorithm
- Merging method is used
- Variants of merge sort
  - in place
    - bottom up merge sort
    - top down merge sort

# Questions

- Why is merge sort an out-of-place sorting technique?
- Does it use recursive procedure?
- What are the best, worst and average case time complexities?

# Reference

- S. Sridhar, Design and Analysis of Algorithms, Oxford University Press, 2015

# Thank You