

Answer key - 18CSC301T - FLA

① CFG - parse Tree always binary - convert CFA - CNF

$\text{stmt} \rightarrow \text{If stmt} \mid \text{print stmt}$

→ 7 marks

~~stmt~~ $\rightarrow \text{if (Expr) stmt}$

$\text{print stmt} \rightarrow \text{print (Expr)}$

$\text{Expr} \rightarrow \text{ID} \mid \text{num} \mid \text{Expr} + \text{Expr} \mid \text{Expr} * \text{Expr} \mid (\text{Expr})$

→ no useless symbols and no ϵ transition

Elimination of unit production

$\text{stmt} \rightarrow \text{if (Expr) stmt} \mid \text{print (Expr)}$

CNF $X_1 \rightarrow \text{if}$

$X_2 \rightarrow \text{print}$

$X_3 \rightarrow ($

$X_4 \rightarrow)$

$X_5 \rightarrow +$

$X_6 \rightarrow *$

$X_7 \rightarrow X_1 X_3$

$X_8 \rightarrow X_4 \text{ stmt}$

$X_9 \rightarrow X_7 * \text{Expr}$

$X_{10} \rightarrow X_2 X_3$

$X_{11} \rightarrow \text{Expr } X_4$

$X_{12} \rightarrow \text{Expr } X_5$

$X_{13} \rightarrow \text{Expr } X_6$

$X_{14} \rightarrow X_3 \text{ Expr}$

$\text{If stmt} \rightarrow \text{if (Expr) stmt}$

$\rightarrow X_1 X_3 \text{ Expr } X_4 \text{ stmt}$

$\rightarrow X_7 \text{ Expr } X_4 \text{ stmt}$

$\rightarrow X_7 \text{ Expr } X_8$

$\text{If stmt} \rightarrow X_9 X_8$

$\text{print stmt} \rightarrow \text{print (Expr)}$

$\rightarrow X_2 X_3 \text{ Expr } X_4$

$\text{print stmt} \rightarrow X_{10} X_{11}$

$\text{stmt} \rightarrow X_9 X_8 \mid X_{10} X_{11}$

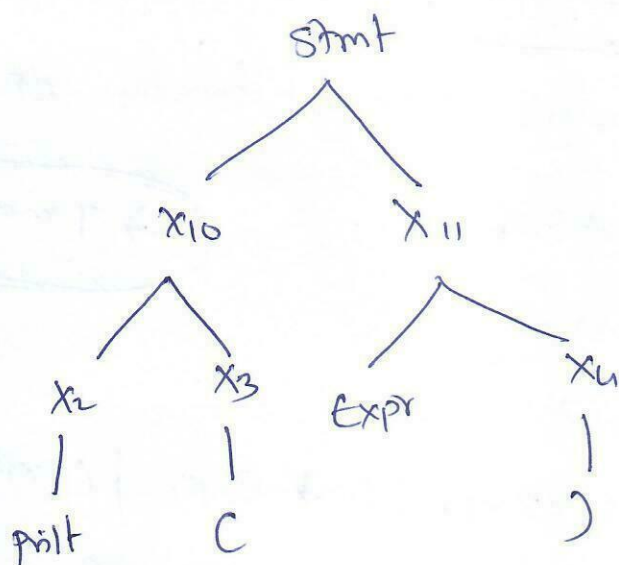
$\text{Expr} \rightarrow \text{ID} / \text{num} / \text{Expr } X_5 \text{ Expr} / \text{Expr } X_6 \text{ Expr} / X_3 \text{ Expr } X_4$

$\text{Expr} \rightarrow \text{ID} / \text{num} / X_{12} \text{ Expr} / X_{14} X_4$

Parse Tree

— 3 marks

②



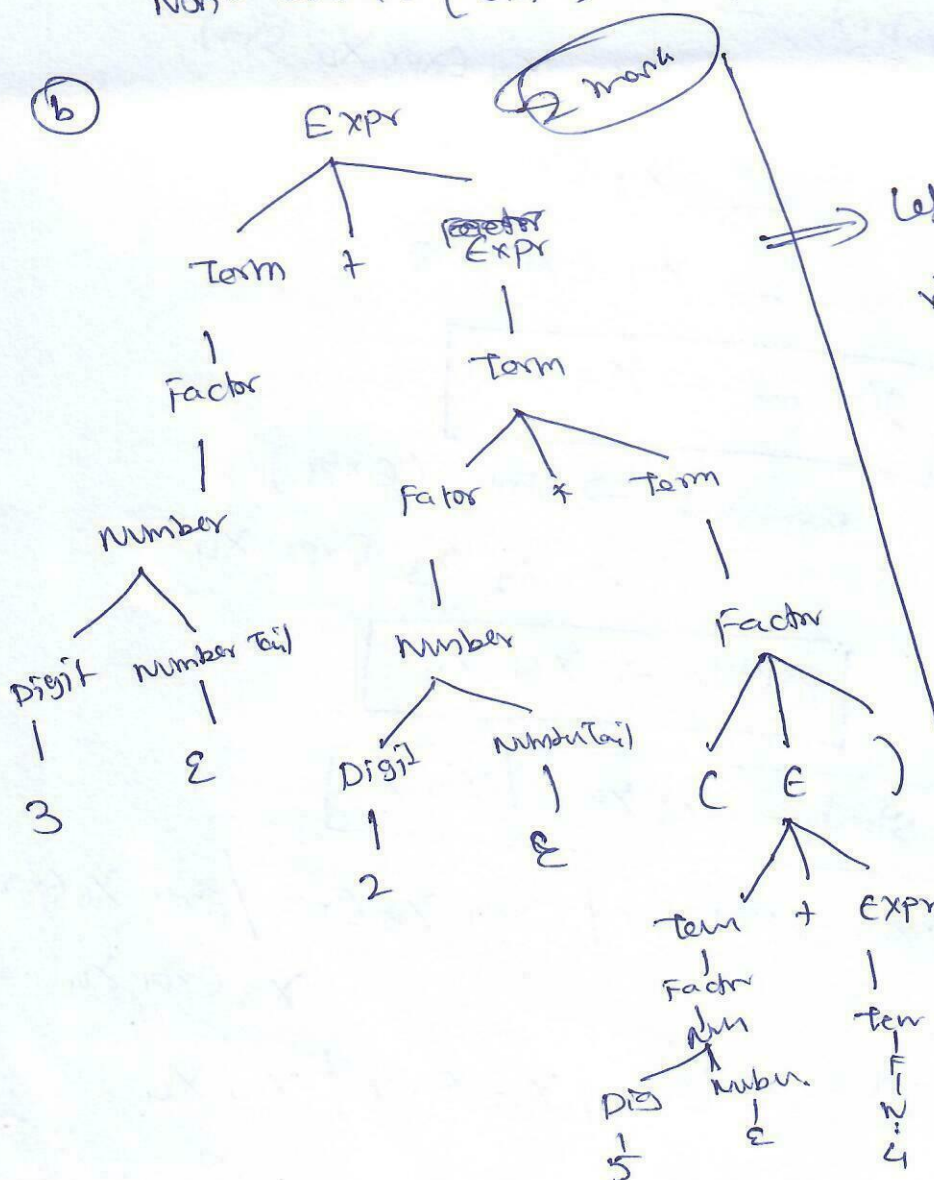
② CFG

① Terminals = $\{ +, *, (,), 0, 1, 2, 3, 4, 5, \dots, 9 \}$

Non-Terminals = $\{ \text{Expr, Term, Factor, Number, Number Tail, Digit} \}$

⇒ 4 marks

②



Left most & right most derivation.

2 + 2 marks

c) simplification

→ 5 marks

→ no useless symbols
 Σ - production \Rightarrow 'Number Tail' \rightarrow Digit number Tail | Σ

elimination

Number Tail \rightarrow Digit number Tail | Digit

Number \rightarrow Digit number Tail | Digit

unit production elimination

~~Expr~~ number \rightarrow Digit number Tail | 0 | 1 | 2 ... | 9

number \rightarrow Digit number Tail | 0 | 1 | 2 ... | 9

Factor \rightarrow Digit number Tail | 0 | 1 | 2 ... | 9 | (E)

term \rightarrow Factor * term | Digit number Tail | 0 | 1 ... | 9 | (E)

Expr \rightarrow Term + Expr | Factor * term | Digit number Tail | 0 | 1 | 2 ... | 9 | (E)

d)

conversion to CNF

- $X_1 \rightarrow ($
- $X_2 \rightarrow)$
- $X_3 \rightarrow X_1 E$
- $X_4 \rightarrow +$
- $X_5 \rightarrow *$
- $X_6 \rightarrow \text{Factor } X_5$

Factor $\rightarrow (E)$
 Factor $\rightarrow X_1 E X_2$
 Factor $\rightarrow X_3 X_2$

term \rightarrow Factor * term.
 term \rightarrow Factor X_5 term
 term $\rightarrow X_6$ term

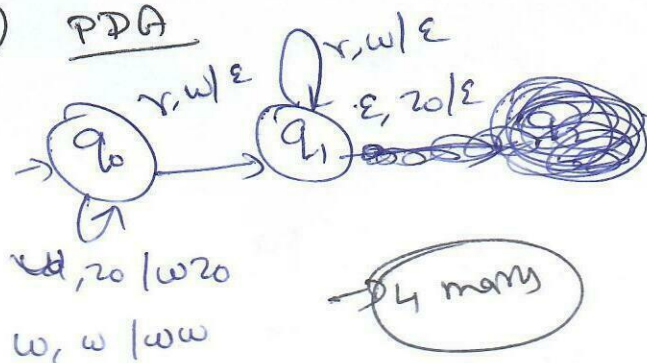
Expr \rightarrow term + Expr
 Expr \rightarrow term X_4 Expr
 $X_7 \rightarrow$ term X_4
 Expr $\rightarrow X_7$ Expr

③ $\{w^n r^n \mid n \geq 0\}$

(a) $n = \text{finite} \Rightarrow \text{Finite Automata}$ — 1 mark

(b) $n = \text{infinite} \Rightarrow \text{push down automata}$ — 1 mark

(c) PDA tuple notation



$PDA = \{q_0, q_1, q_2\}, \{w, r\}, \{z_0, w\}$

$\delta(q_0, w, z_0) = (q_0, w z_0)$

all transition

an example + ID for it \Rightarrow 2 marks

(d) CNF PDA to cfa $[V, T, P, S]$ $S = S$ $T = \{w, r\}$

$V = \{S, [q_0 z_0 q_0], [q_0 z_0 q_1], [q_1 z_0 q_0], [q_1 z_0 q_1], [q_0 w q_0], [q_0 w q_1], [q_1 w q_0], [q_1 w q_1]\}$

$\delta(q_0, w, z_0) = (q_0, w z_0)$

$\delta(q_0, w, w) = (q_0, w w)$

$[q_0 z_0 q_0] \rightarrow w [q_0 w q_0] [q_0 z_0 q_0]$

$[q_0 w q_0] \rightarrow w [q_0 w q_0] [q_0 w q_0]$

$[q_0 z_0 q_0] \rightarrow w [q_0 w q_1] [q_1 z_0 q_0]$

$[q_0 w q_0] \rightarrow w [q_0 w q_1] [q_1 w q_0]$

$[q_0 z_0 q_1] \rightarrow w [q_0 w q_0] [q_0 z_0 q_1]$

$[q_0 w q_1] \rightarrow w [q_0 w q_0] [q_0 w q_1]$

$[q_0 z_0 q_1] \rightarrow w [q_0 w q_1] [q_1 z_0 q_1]$

$[q_0 w q_1] \rightarrow w [q_0 w q_1] [q_1 w q_1]$

$\delta(q_0, r, w) = (q_1, \epsilon)$

$\delta(q_1, r, w) = (q_1, \epsilon)$

$[q_0 w q_1] \rightarrow r$

$[q_1 w q_1] \rightarrow r$

$\delta(q_1, \epsilon, z_0) = (q_1, \epsilon)$

$[q_1 z_0 q_1] \rightarrow \epsilon$

(e) No, NPDA is more powerful than deterministic PDA.

few problems can be solved only by NPDA not

by DPDA.

\Rightarrow 2 marks

$$g(a_0, r, r) = (a_0, rr)$$

$$[a_0 r a_0] \rightarrow r [a_0 r a_0] [a_0 r a_0]$$

$$[a_0 r a_0] \rightarrow r [a_0 r a_1] [a_1 r a_0]$$

$$[a_0 r a_1] \rightarrow r [a_0 r a_0] [a_0 r a_1]$$

$$[a_0 r a_1] \rightarrow r [a_0 r a_1] [a_1 r a_1]$$

$$g(a_0, r, z_0) = (a_0, rz_0) \quad \textcircled{2}$$

$$[a_0 z_0 a_0] \rightarrow r [a_0 r a_1] [a_1 z_0 a_0]$$

$$[a_0 z_0 a_0] \rightarrow r [a_0 r a_0] [a_0 z_0 a_0]$$

$$[a_0 z_0 a_1] \rightarrow r [a_0 r a_0] [a_0 z_0 a_1]$$

$$[a_0 z_0 a_1] \rightarrow r [a_0 r a_1] [a_1 z_0 a_1]$$

→ 6 marks