

1) Time Complexity : $O(n^2)$
Better approach

```
int sum = 0
```

```
int n = 4
```

```
for (int i = 1; i <= n; i++) {  
    sum += i  
}
```

}

sum = $n * (n+1) / 2$;

T.C = $O(1)$

S.C = $O(1)$

2) $T(n) = 3T(n-1) + 12n$
 $T(0) = 5$, given
 $T(2) = ?$

$$T(1) = 3T(0) + 12n$$

$$T(1) = 15 + 12n$$

$$T(2) = 3T(1) + 12n$$

$$= 3[15 + 12n] + 12n$$

$$= 45 + 36n + 12n$$

$$T(2) = 45 + 48n$$

Q4 $T(n) = 16T(\frac{n}{4}) + n^2 \log n$

$$T(n) = aT(\frac{n}{b}) + \theta(n^k \log^p n)$$

$a = 16$, $k = 1$
 $b = 4$, $p = 1$

compare a & b^k
 $16 > 4^1$

$a > b^k$

$$T(n) = \theta(n^{\log_b a})$$

$$= \theta(n^{\log_4 16})$$

$$= \theta(n^{\log_4 4^2})$$

$$\boxed{T(n) = \theta(n^2)}$$

$$T(n) = T(n-1) + C \quad \text{--- (1)}$$

$$n \rightarrow n-1$$

$$T(n-1) = T(n-2) + C$$

$$T(n) = T(n-2) + C + C \quad \text{--- (2)}$$

$$n \rightarrow n-2$$

$$T(n-2) = T(n-3) + C$$

$$T(n) = T(n-3) + C + C + C \quad \text{--- (3)}$$

⋮
k times

$$T(n) = T(n-k) + C + C + C + \dots \text{--- k times}$$

$$T(n) = T(n-k) + kC$$

$$T(1) = 1 \quad \text{Base case}$$

$$T(n-k) = 1$$

$$n-k = 1$$

$$n = k+1$$

$$\boxed{k = n-1}$$

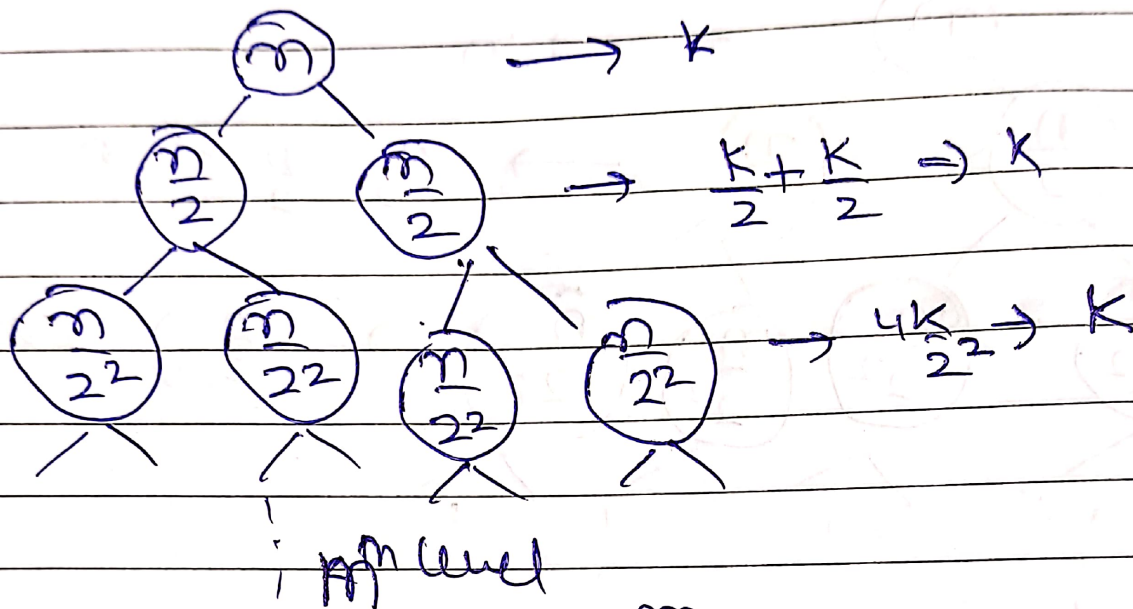
$$T(n) = 1 + (n-1)C$$

$$T(n) = 1 + nC - C$$

$$\boxed{T(n) = O(n)} \quad \text{ans}$$

26 // $T(n) = 2T\left(\frac{n}{2}\right) + K$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + K$$



$$\frac{n}{2^m} = 1 \quad n = 2^m$$

$$m = \log_2 n$$

level separation

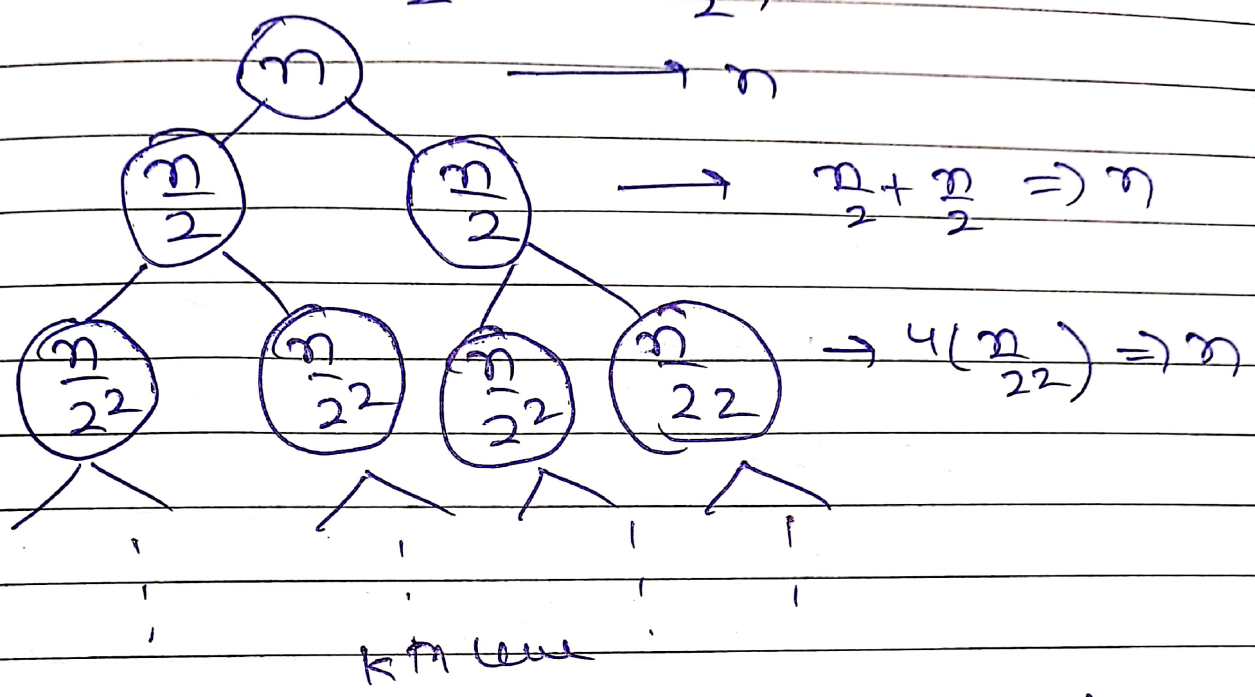
$$K + K + K \dots m^{\text{th}} \text{ level} = \log_2 m$$

$$\Rightarrow \Theta(K \log_2 m) \text{ ans}$$

Q5

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + n$$



$$\frac{n}{2^k} = 1 \quad (\text{Base case condition})$$

$$n = 2^k$$

$$k = \log_2 n$$

level repetition

$$n + n + n + \dots \log_2 n \text{ times}$$

$$O(n \log_2 n) \text{ ans}$$