

Tutorial 5

1

Name - Amandeep Singh
Section - G
Roll No. - 53
Uni Roll No. - 2016623

Course - B.Tech CSE
Sem - 4

<u>Ans 1</u>	BFS	DFS
	<ul style="list-style-type: none">(i) uses queue data structure.(ii) Stands for Breadth first search(iii) Can be used to find the single source shortest path in an unweighted graph.(iv) Siblings are visited before the children	<ul style="list-style-type: none">→ uses stack data structure→ stands for Depth first search.→ we might traverse through more edges to reach a destination vertex from a source.→ children are visited before the siblings.
	<p>Applications</p> <ul style="list-style-type: none">→ Minimum spanning tree for an unweighted graph.→ peer to peer Networks→ Social Networking websites→ GPS Navigation systems	<p>Applications</p> <ul style="list-style-type: none">→ Detecting cycle in a graph.→ Path finding.→ Topological Sorting→ Solving puzzles with only one solution.

Ans 2 In BFS we use queue data structure as queue is used when things don't have to be processed immediately but have to be processed in FIFO order like a BFS.

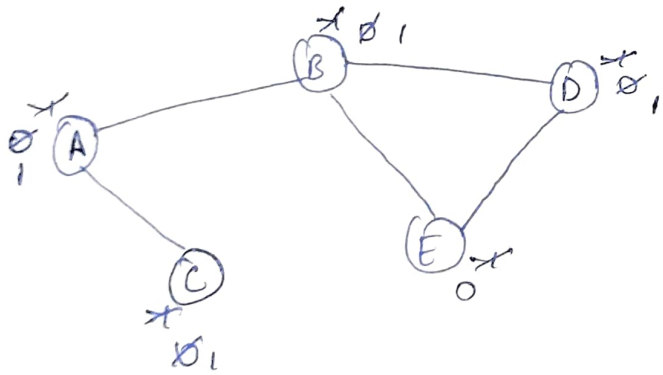
→ It is used in social networking sites and GPS Navigation Systems.

- In DFS stack is used as DFS uses backtracking. (2)
- For DFS, we retrieve it from root to the farthest node as much as possible, this is the same idea as LIFO (Last In First Out).
- It is used in topological sortings.

Ans 3 → Dense graph is a graph in which the no. of edges is close to the maximal no. of edges

- Sparse graph is a graph in which the no. of edges is close to the minimal no. of edges. It can be disconnected graph.
- * Adjacency lists are preferred for sparse graph & adjacency matrix for dense graph.

Ans 4 Cycle detection in Undirected graph (BFS)



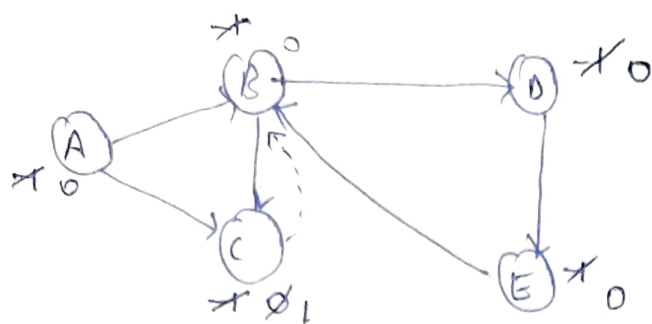
-1 → unvisited
 0 → into the queue
 1 → traversed

Queue → A B C D E

Visited set → A B C D

- When D checks its adjacent vertices it finds E with 0
- If any vertex finds the adjacent vertex with flag 0, then it contains cycle.

→ Cycle detection in Directed Graph (DFS)



Stack : E D B A

visited set : A B C D E

Parent Map

vertices	Parent
A	-
B	A
C	B
D	B
E	D

⇒ B → D → E → B

Here E finds B (adjacent vertex of E) with 0.

⇒ It contains a cycle.

Ans 5 The disjoint set data structure is also known as union-find data structure & merge-find set. It is a data structure that contains a collection of disjoint or non-overlapping sets.

→ The disjoint set means that when the set is partitioned into the disjoint subsets, various operations can be performed on it.

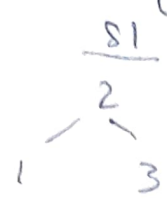
→ In this case, we can add new sets, we can merge the sets, & we can also find the representatives of a set.

* Operations on disjoint Set :-

(1) Union → If S_1 & S_2 are two disjoint set, their union $S_1 \cup S_2$ is a set of all elements x such that x is in either S_1 or S_2 .

(2) Find \rightarrow Given an element to find the set containing it.

eg \rightarrow



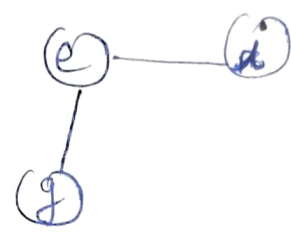
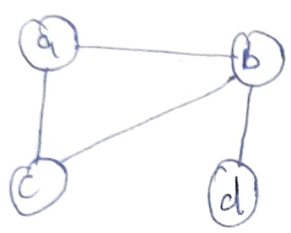
find(3) = s1

find(5) \Rightarrow s2

return in which set x belongs.

(3) Make-Set(x): Creates a set containing x .

Ans 7



(j)

$V = \{a, b, c, d, e, g, h, i, j, l\}$

$E = \{(a, b), (a, c), (b, c), (b, d), (e, i), (e, g), (h, l), (j)\}$

	{a}	{b}	{c}	{d}	{e}	{g}	{h}	{i}	{j}	{l}
(a,b)	{a, b}	{c}	{d}	{e}	{g}	{h}	{i}	{j}	{l}	
(a,c)	{a, b, c}	{d}	{e}	{g}	{h}	{i}	{j}	{l}		
(b,c)	{a, b, c}	{d}	{e}	{g}	{h}	{i}	{j}	{l}		
(b,d)	{a, b, c, d}	{e}	{g}	{h}	{i}	{j}	{l}			
(e,i)	{a, b, c, d}	{e, i}	{g}	{h}	{j}	{l}				
(e,g)	{a, b, c, d}	{e, i, g}	{h}	{j}	{l}					
(h,l)	{a, b, c, d}	{e, i, g}	{h, l}	{j}						
(j)	{a, b, c, d}	{e, i, g}	{h, l}	{j}						

we have,

{a, b, c, d}

{e, i, g}

{h, l}

{j}

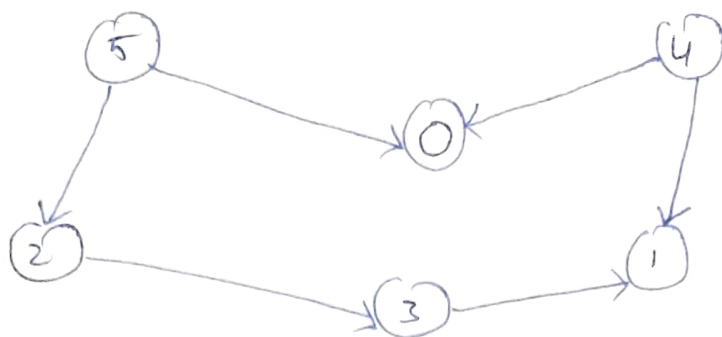
}

Ans

we have,

- {a, b, c, d}
- {e, i, g}
- {h, l}
- {j}

Ans



0

1

2 → 3

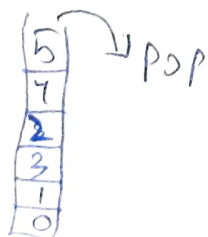
3 → 1

4 → 0 → 1

5 → 2 → 0

Algorithm:-

- (1) Go to node 0, it has no outgoing edges so push node 0 into the stack & mark it visited.
- (2) Go to node 1, again it has no outgoing edges, push node 1 into stack.
- (3) Go to node 2, process all the adjacent nodes & mark 2 visited.
- (4) Node 3 is already visited to continue with next node.
- (5) Go to node 4, all its adjacent nodes are already visited so, push node 4 into the stack & mark it visited.
- (6) Go to node 5, all its adjacent nodes are already visited so, push node 5 into the stack & mark visited.



5 4 2 3 1 0

Output

Ans 9 Heap is generally preferred for priority queue implementations because it helps provide better performance compared to array or linked list. (6)

→ Algorithm where priority queue is used :-

- (1) Dijkstra's shortest path Algo. :- when the graph is stored in the form of an adjacency list or matrix, priority queue can be used to extract minimum efficiently. when implementing Dijkstra's algorithm.
- (2) Prim's algo. :- To store keys of nodes & on another side extract the minimum key nodes at every step.

Ans 10

Min Heap

Max Heap

(i) For every pair of the parent & descendant child node, the parent node always has lower value than its child node.

→ The parent node has greater value than descendant child node.

(ii) The value of nodes inc. as we transverse from root to leaf node.

→ The value of nodes decreases as we transverse from root to leaf node.

(iii) Root node has the lowest value.

→ The root node has the greatest value.