



# **INDIAN INSTITUTE OF INFORMATION TECHNOLOGY DESIGN AND MANUFACTURING KANCHEEPURAM**

LAB ASSIGNMENT 5 - REPORT  
ON  
MATRIX ADDITION  
AND  
MATRIX MULTIPLICATION  
IN CUDA

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## Strategy

In my program for matrix addition, the instruction which is running in parallel is in the “for” loop i.e  $C[i] = A[i] + B[i];$

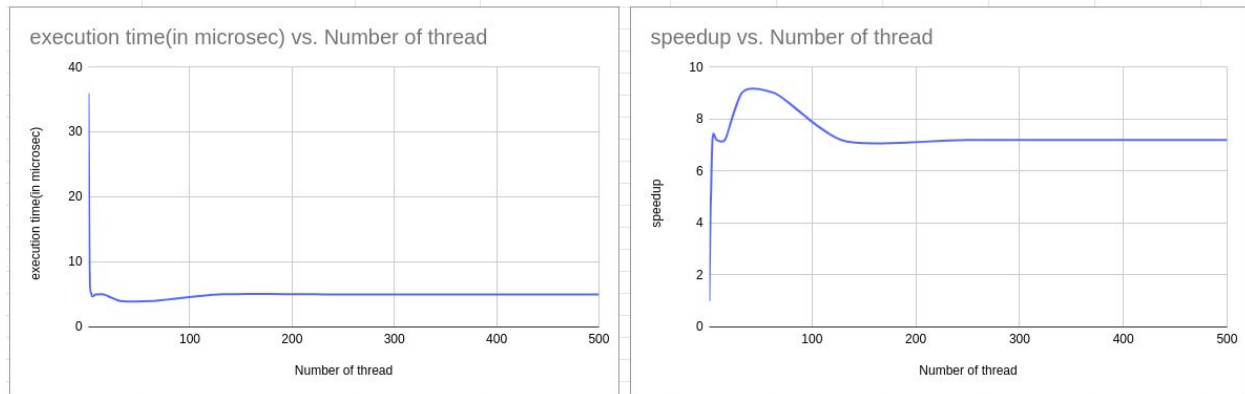
Instead of running the program serially, we can distribute the task of adding matrices between threads so that my program could run parallelly and in turn save the execution time.

Therefore , firstly, i converted the matrix in an array and divided the task of adding. For example, if my matrix size is  $4 \times 4$  then i have converted it in array of size 16. Suppose i have 4 thread then thread1 will execute 1,5,9,13, thread2 will execute 2,6,10,14, thread3 will execute 3,7,11,15 and thread4 will execute 4,8,12,16. In this way matrix addition will happen with the help of threads in CUDA

## Graph and tables

<https://docs.google.com/spreadsheets/d/1DM3hmeEPAiy6AFi0OVakTaHoM5Bg47a1P4gIVVsIYhU/edit#gid=0>

| Question1        |                             |         |                             |
|------------------|-----------------------------|---------|-----------------------------|
| Number of thread | execution time(in microsec) | speedup | parallelization fraction(f) |
| 1                | 36                          | 1       | 0                           |
| 2                | 9                           | 4       | 1.5                         |
| 4                | 5                           | 7.2     | 1.148148148                 |
| 8                | 5                           | 7.2     | 0.9841269841                |
| 16               | 5                           | 7.2     | 0.9185185185                |
| 32               | 4                           | 9       | 0.917562724                 |
| 64               | 4                           | 9       | 0.9029982363                |
| 128              | 5                           | 7.2     | 0.8678915136                |
| 256              | 5                           | 7.2     | 0.8644880174                |
| 500              | 5                           | 7.2     | 0.8628367847                |



## Calculation of parallelization fraction

$T(1) = 36$  microsecond

Here, for  $P = 32$  the execution time is minimum

$T(P) = 4$  microsecond

$$\text{Speedup} = \frac{T(1)}{T(P)} = \frac{36}{4} = 9$$

From Amdahl's Law,

$$\text{Speedup} = \frac{1}{(f/P) + (1-f)} \quad \text{Where, } f = \text{Parallelization factor } P = \text{Thread Number}$$

$$\text{So, } f = \frac{(1 - T(P)/T(1))}{(1 - (1/P))}$$

Therefore,  $f = 0.917562724$  which means that approx 91% of the program is parallelizable.

# MATRIX MULTIPLICATION

## Strategy

In my program for vector addition, the instruction which is running in parallel is in the "for" loop i.e  $C[x][y] += A[x*N+k] * B[k*N+y];$

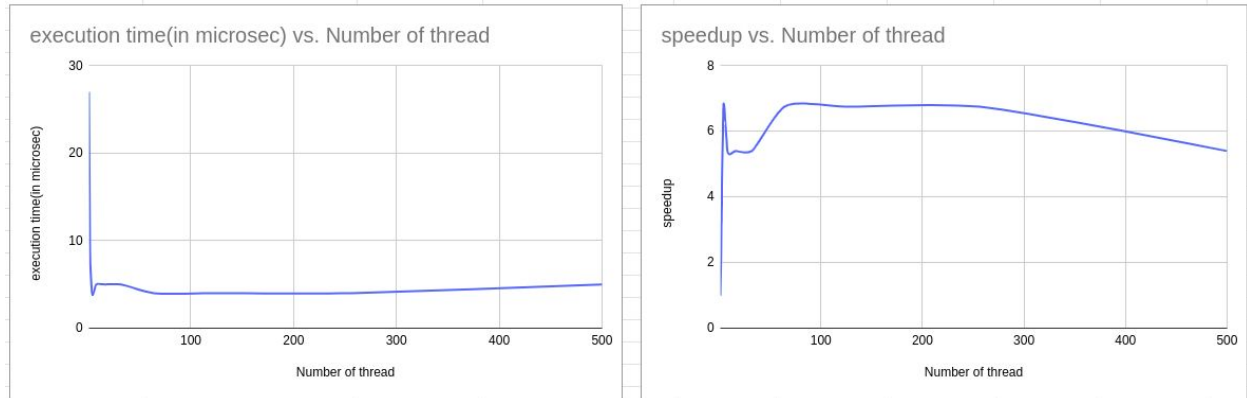
Instead of running the program serially, we can distribute the task of multiplying matrices between threads so that my program could run parallelly and in turn save the execution time.

Therefore, firstly, I converted the matrix in an array and divided the task of multiplying among the threads. For example, if my matrix size is  $4*4$  then I have converted it in an array of size 16. Suppose I have 4 threads then thread1 will execute 1,5,9,13, thread2 will execute 2,6,10,14, thread3 will execute 3,7,11,15 and thread4 will execute 4,8,12,16.  $x$  and  $y$  is the position of element which needs to be calculated.  $k$  is varying as the column of  $A$  or row of  $B$ . In this matrix multiplication is happening.

## Graph and tables

<https://docs.google.com/spreadsheets/d/1DM3hmeEPAiy6AFi0OVakTaHoM5Bg47a1P4glVVslYhU/edit#gid=0>

| Question2        |                             |         |                             |
|------------------|-----------------------------|---------|-----------------------------|
| Number of thread | execution time(in microsec) | speedup | parallelization fraction(f) |
| 1                | 27                          | 1       | 0                           |
| 2                | 9                           | 3       | 1.333333333                 |
| 4                | 4                           | 6.75    | 1.135802469                 |
| 8                | 5                           | 5.4     | 0.9312169312                |
| 16               | 5                           | 5.4     | 0.8691358025                |
| 32               | 5                           | 5.4     | 0.8410991637                |
| 64               | 4                           | 6.75    | 0.8653733098                |
| 128              | 4                           | 6.75    | 0.8585593467                |
| 256              | 4                           | 6.75    | 0.8551924473                |
| 500              | 5                           | 5.4     | 0.8164477102                |



## Calculation of parallelization fraction

$T(1) = 27$  seconds

Here, for  $P = 64$  the execution time is minimum

$T(P) = 4$  seconds

$$\text{Speedup} = \frac{T(1)}{T(P)} = \frac{27}{4} = 6.750.8448150377$$

From Amdahl's Law,

$$\text{Speedup} = \frac{1}{(f/P) + (1-f)} \text{ Where, } f = \text{Parallelization factor } P = \text{Thread Number}$$

$$\text{So, } f = \frac{(1 - T(P)/T(1))}{(1 - (1/P))}$$

Therefore,  $f = 0.8653733098$  which means that approx 86% of the program is parallelizable.