

STORAGE SYSTEM ARCHITECTURES



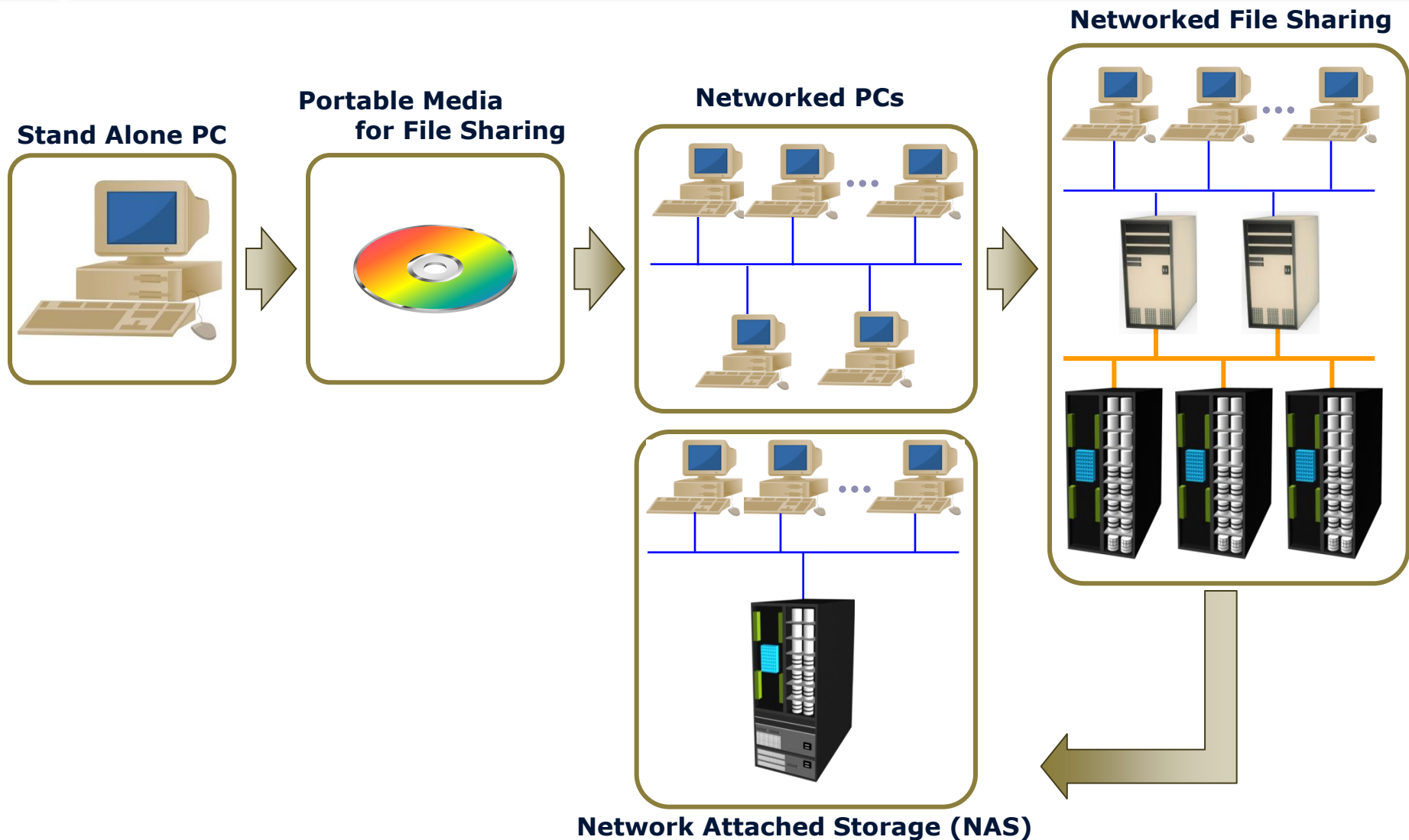


Network-Attached Storage

File Sharing Environment

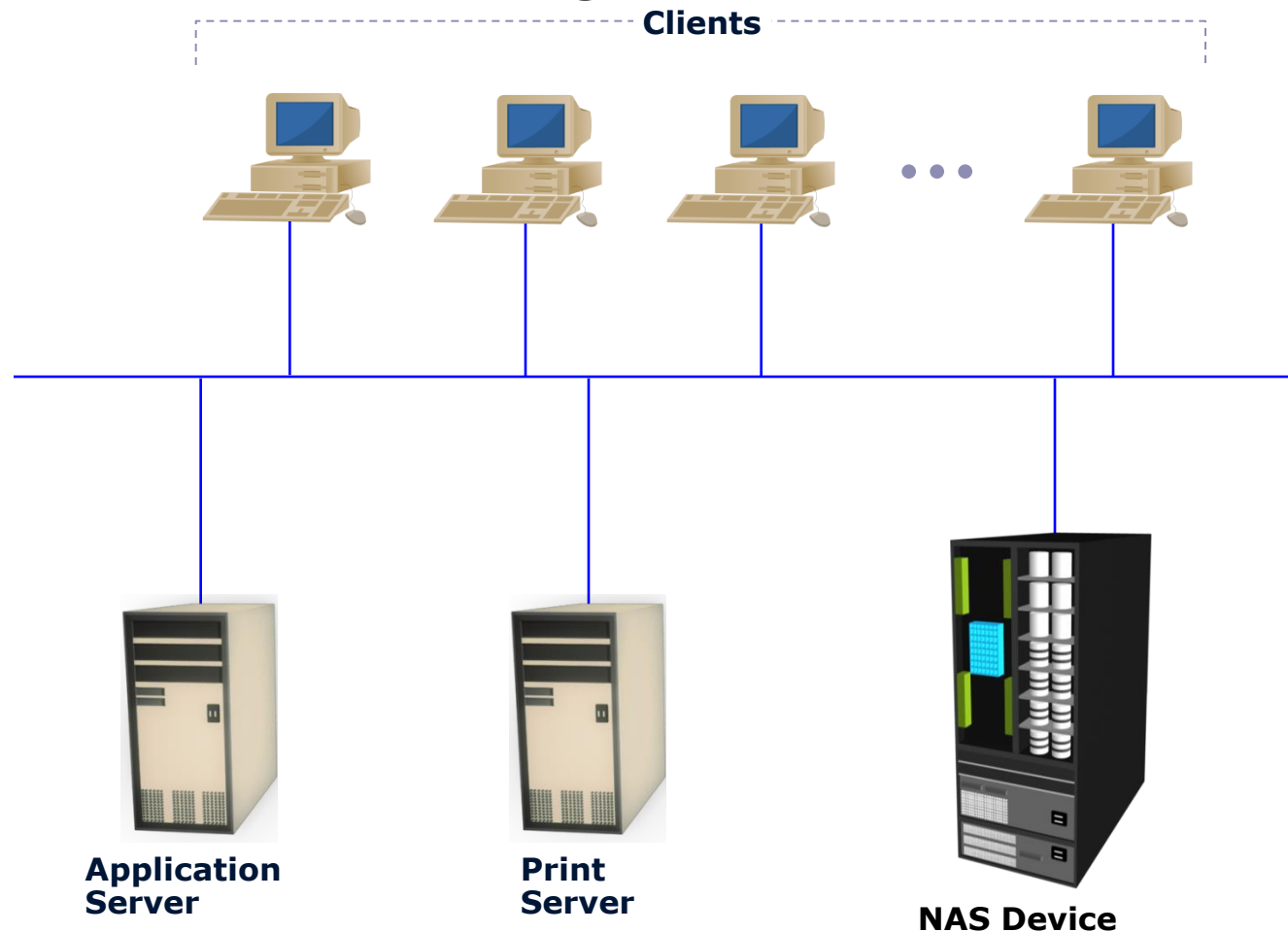
- File system is structured way of storing and organizing data files
- File Sharing
 - ▣ Storing and accessing data files over network
 - ▣ FS must be mounted in order to access files
- Traditional *client/server model*, implemented with file-sharing protocols for remote file sharing
 - ▣ Example: FTP, DNS
- DFS is another examples of file sharing implementation

File Sharing Technology Evolution

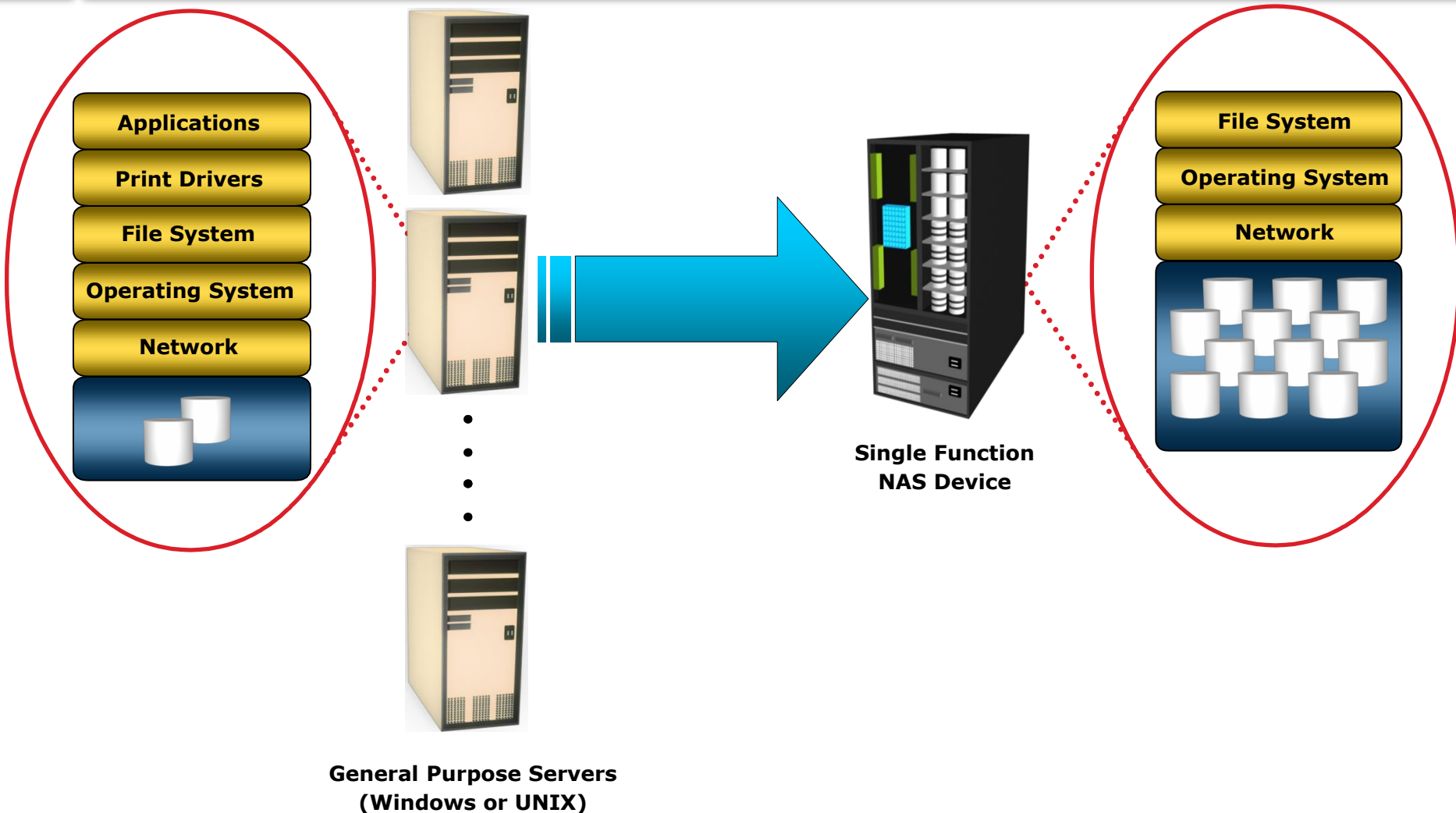


What is NAS ?

- NAS is shared storage on a network infrastructure



General Purpose Servers vs. NAS Devices

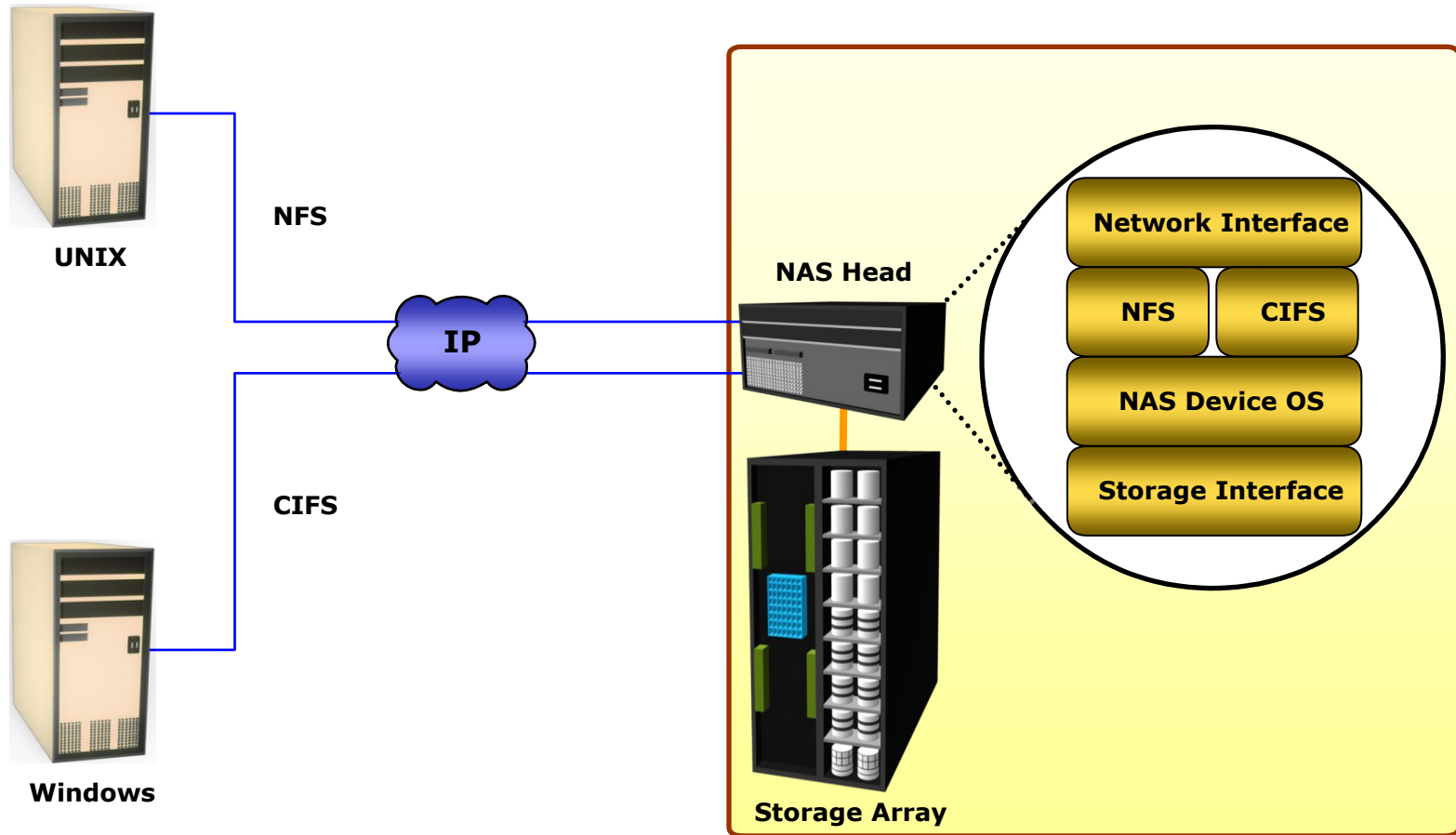


Benefits of NAS



- ❑ Support comprehensive access to information
- ❑ Improves efficiency
- ❑ Improved flexibility
- ❑ Centralizes storage
- ❑ Simplifies management
- ❑ Scalability
- ❑ High availability – through native clustering
- ❑ Provides security integration to environment (user authentication and authorization)

Components of NAS



NAS File Sharing Protocols



- Two common NAS file sharing protocols are:
 - ▣ CIFS – Common Internet File System protocol
 - Traditional Microsoft environment file sharing protocol, based upon the Server Message Block protocol
 - ▣ NFS – Network File System protocol
 - Traditional UNIX environment file sharing protocol

Network File System (NFS)

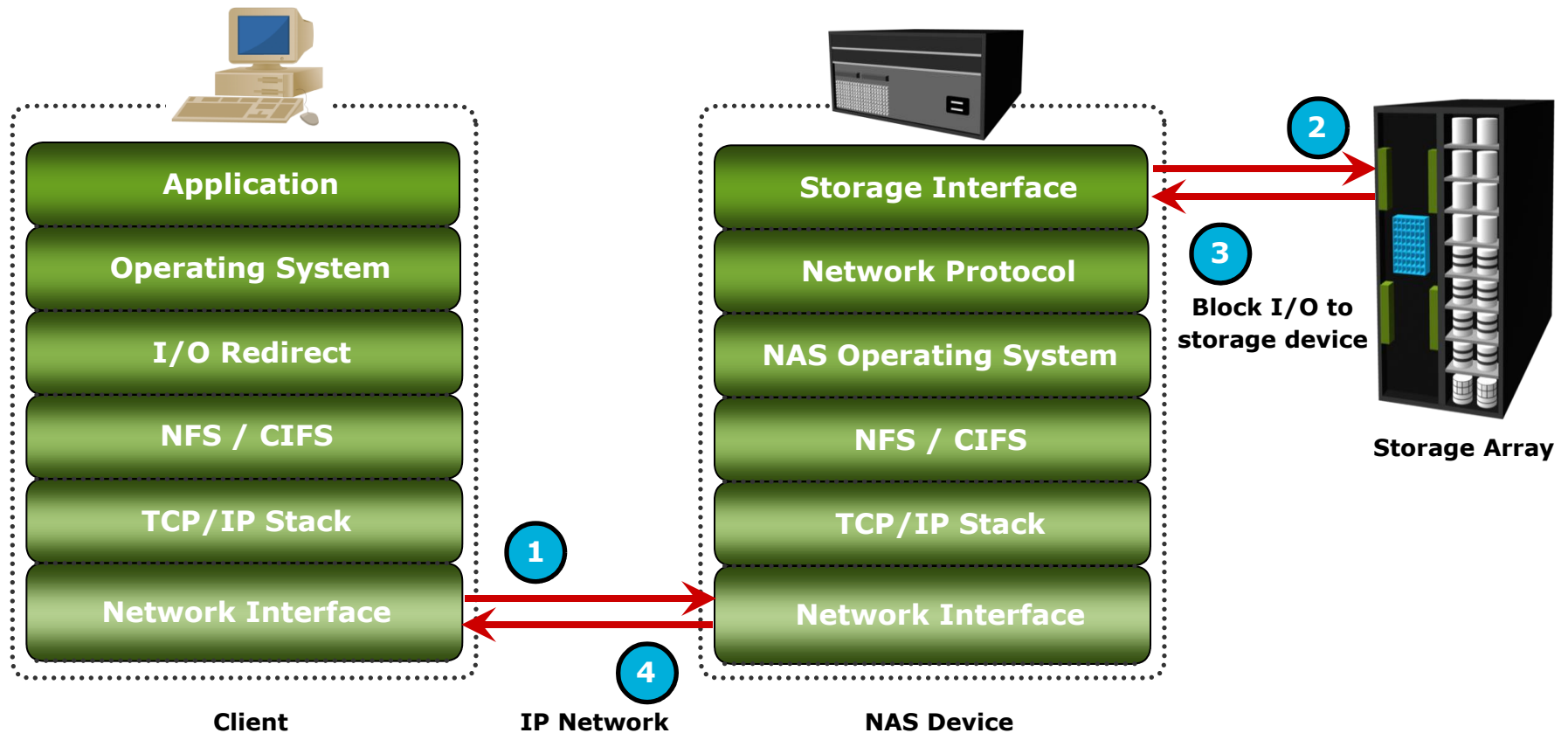


- ❑ Client/server application
- ❑ Uses RPC mechanisms over TCP protocol
- ❑ Mount points grant access to remote hierarchical file structures for local file system structures
- ❑ Access to the mount can be controlled by permissions

NAS File Sharing - CIFS

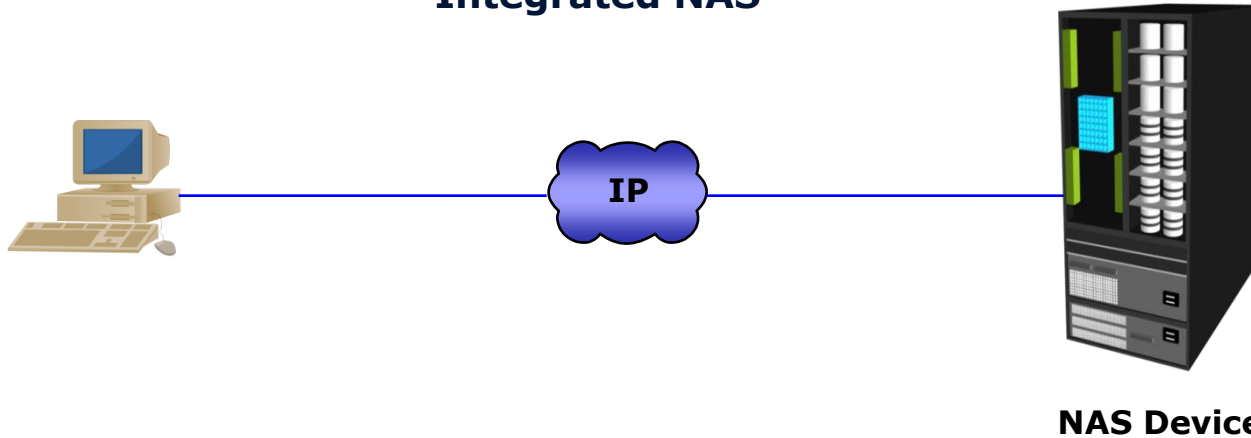
- Common Internet File System
 - ▣ Developed by Microsoft in 1996
 - ▣ An enhanced version of the Server Message Block (SMB) protocol
 - ▣ Stateful Protocol
 - Can automatically restore connections and reopen files that were open prior to interruption
 - ▣ Operates at the Application/Presentation layer of the OSI model
 - ▣ Most commonly used with Microsoft operating systems, but is platform-independent
 - ▣ CIFS runs over TCP/IP and uses DNS (Domain Naming Service) for name resolution

NAS I/O

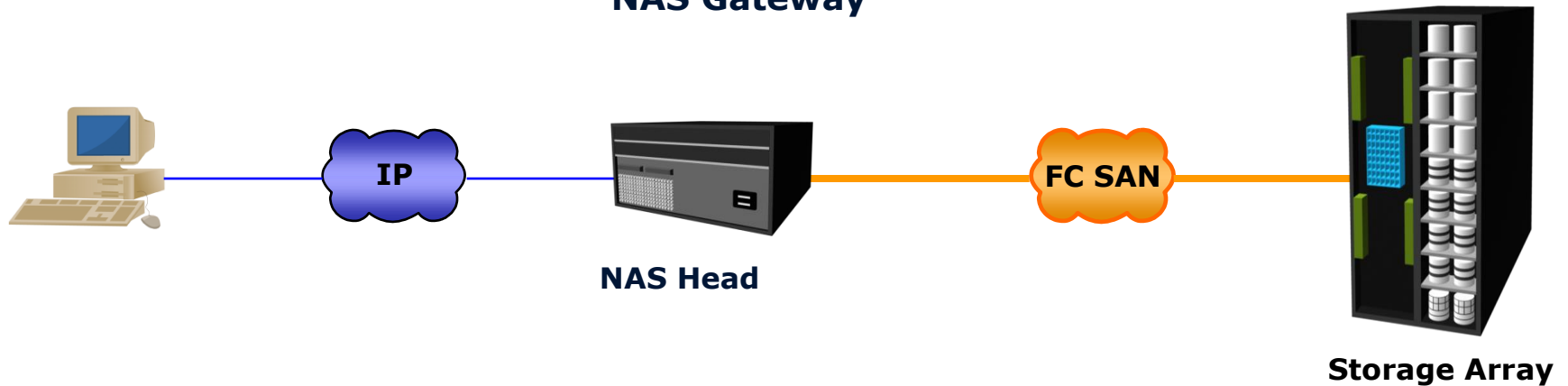


NAS Implementations

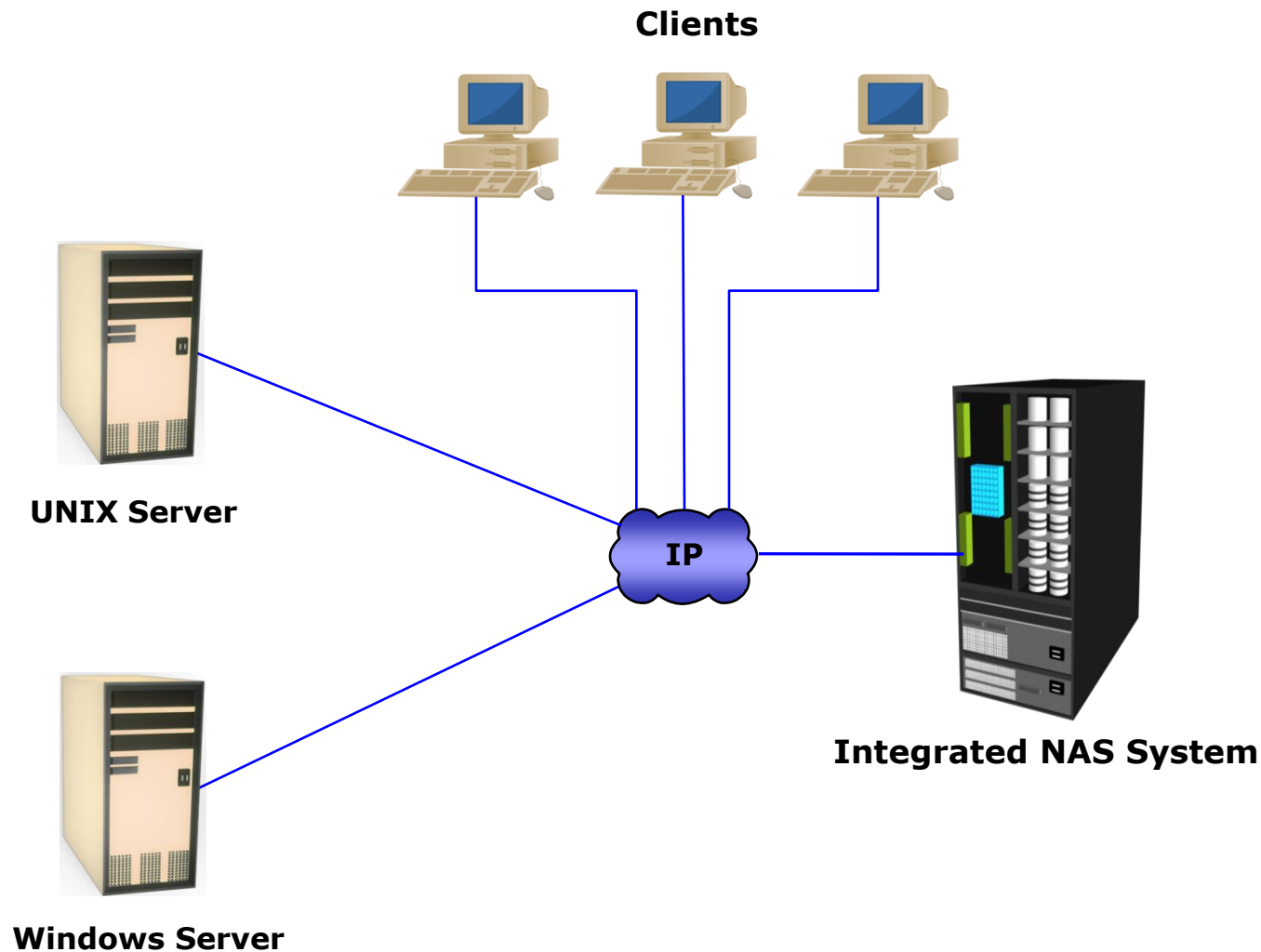
Integrated NAS



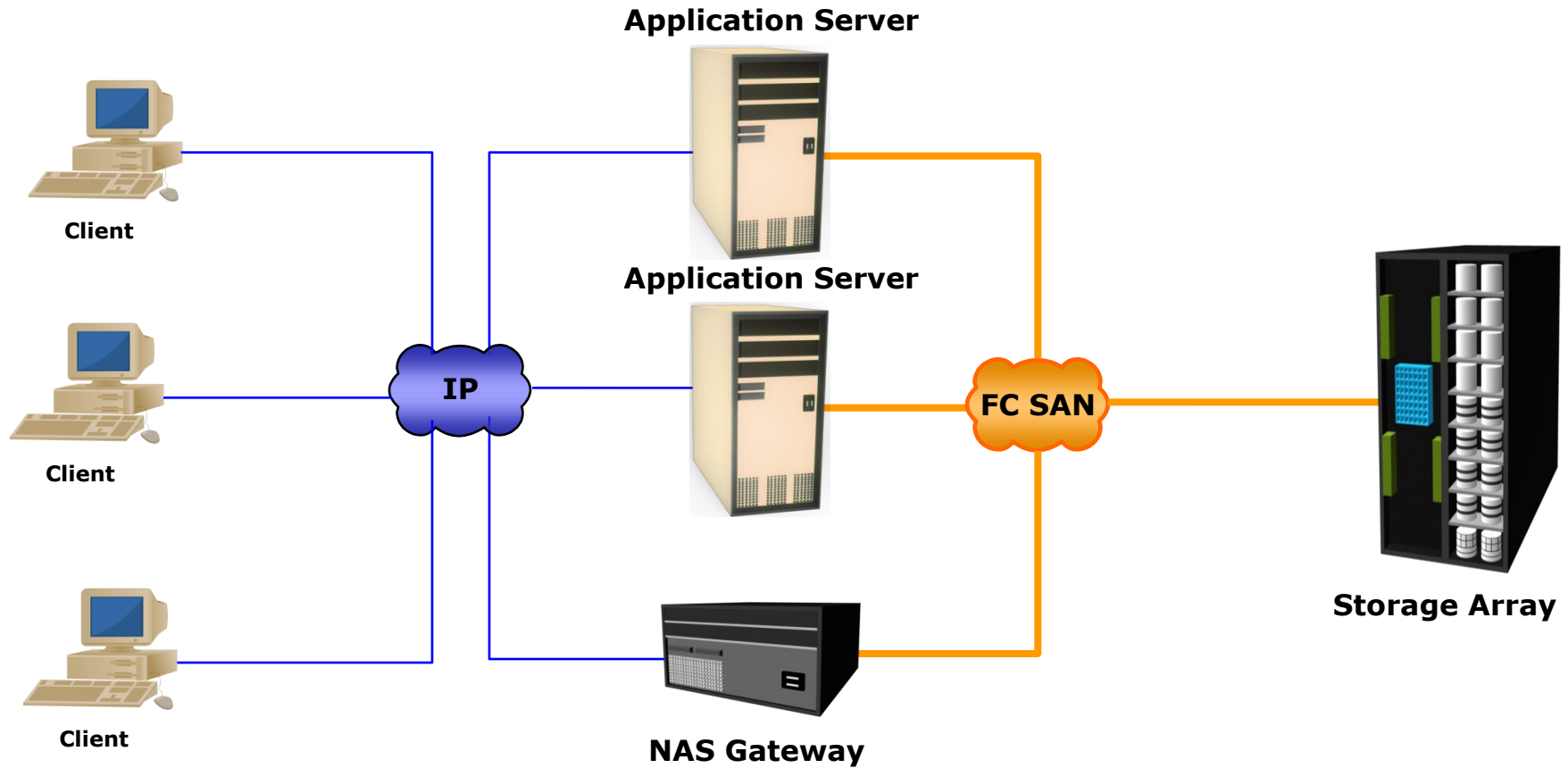
NAS Gateway



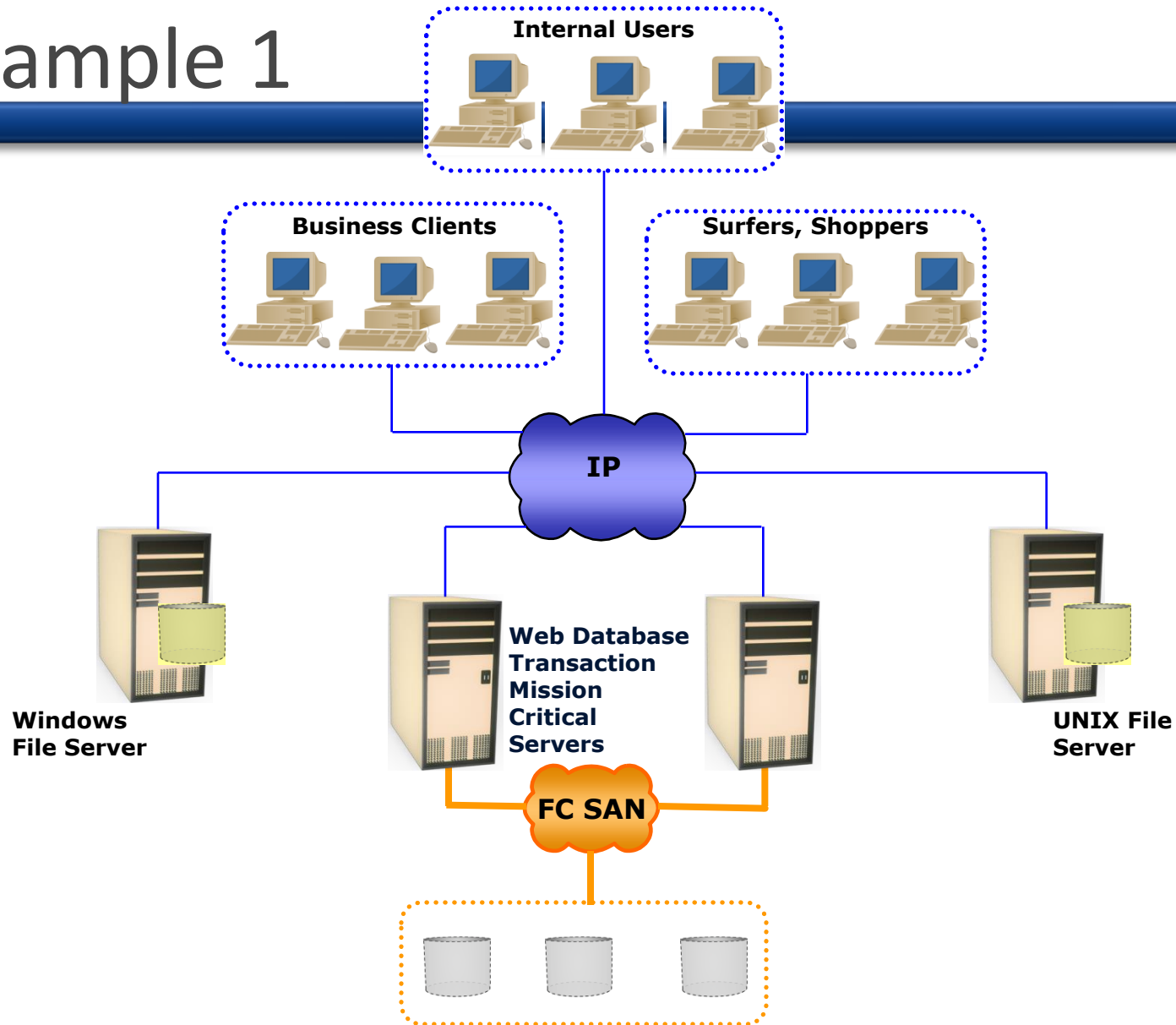
Integrated NAS Connectivity



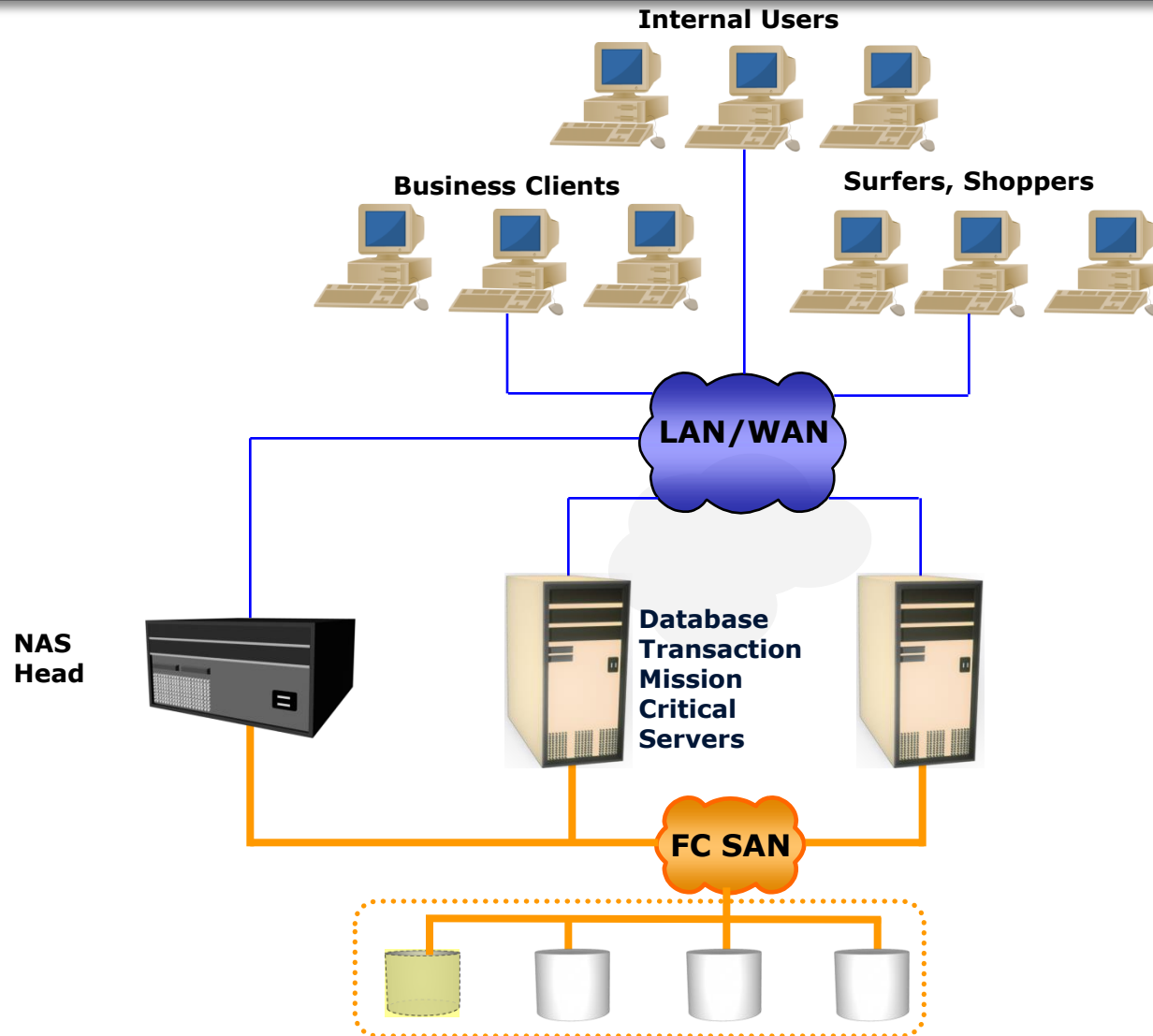
Gateway NAS Connectivity



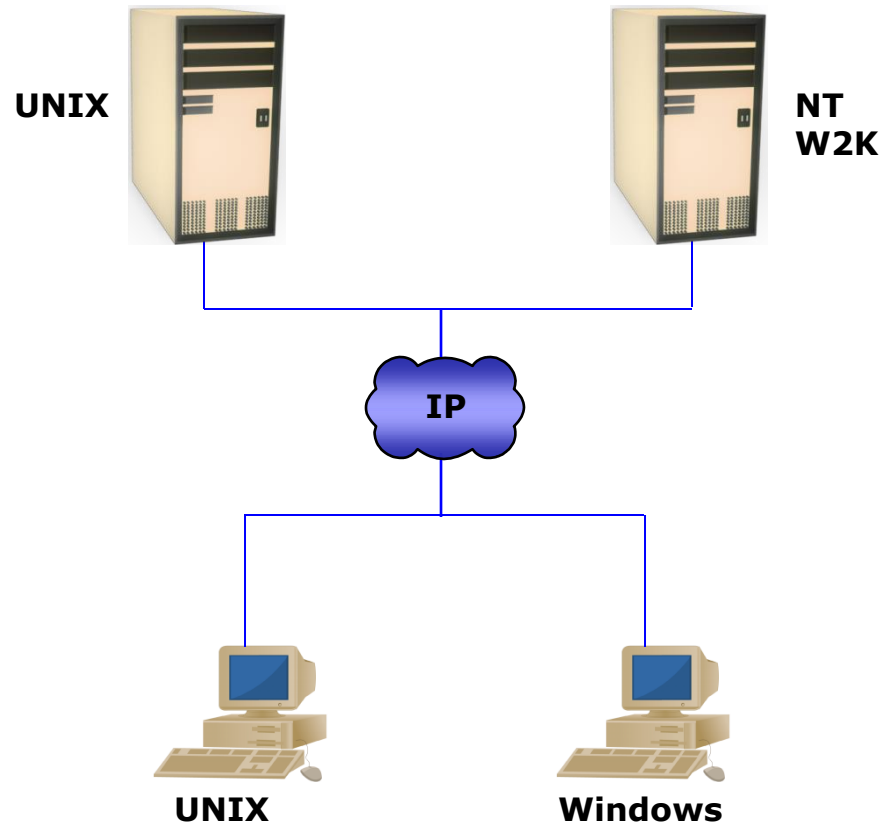
Traditional File Server Environment – Example 1



Storage Consolidation with NAS

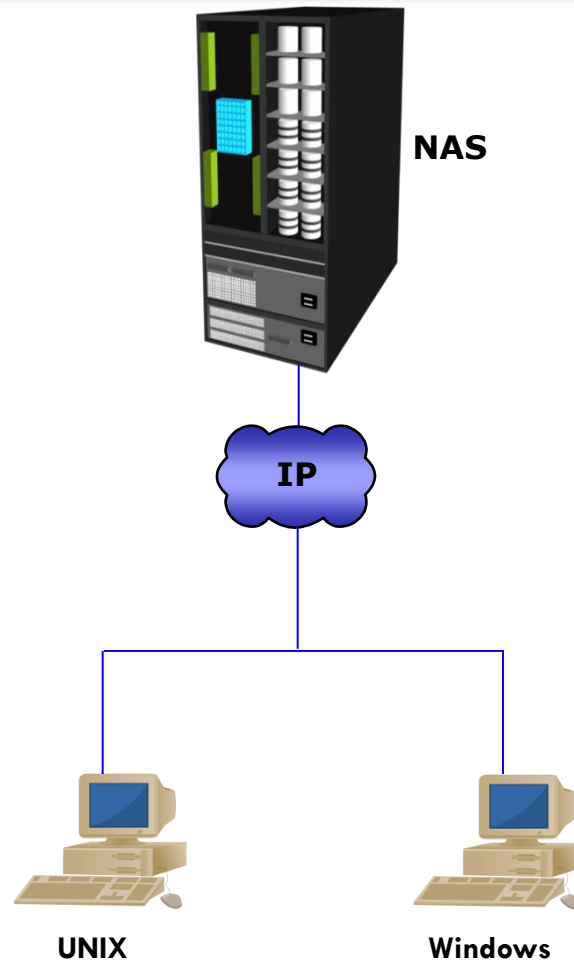


Traditional File Server Environment – Example 2



**General purpose OS serving files via
FTP, CIFS, NFS, HTTP. . .**

Server Consolidation with NAS



**General purpose OS serving files via
FTP, CIFS, NFS, HTTP. . .**

Network Attached Storage

Decreasing Disk Diameters

14" » 10" » 8" » 5.25" » 3.5" » 2.5" » 1.8" » 1.3" » ...
high bandwidth disk systems based on arrays of disks

Network provides
well defined physical
and logical interfaces:
*separate CPU and
storage system!*

**High Performance
Storage Service
on a High Speed
Network**

Network File Services

OS structures
supporting remote
file access

3 Mb/s » 10Mb/s » 50 Mb/s » 100 Mb/s » 1 Gb/s » 10 Gb/s
networks capable of sustaining high bandwidth transfers

Increasing Network Bandwidth

Manufacturing Advantages of Disk Arrays

Conventional:
4 disk
designs

3.5"

5.25"

10"

14"

Disk Product Families

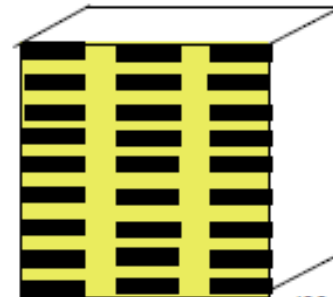
Low End



High End

Disk Array:
1 disk design

3.5"



14" 10" 5.25" 3.5"

Replace Small # of Large Disks with Large # of Small Disks!

	IBM 3390 (K)	IBM 3.5" 0061	x70
<i>Data Capacity</i>	20 GBytes	320 MBytes	23 GBytes
<i>Volume</i>	97 cu. ft.	0.1 cu. ft.	11 cu. ft.
<i>Power</i>	3 KW	11 W	1 KW
<i>Data Rate</i>	15 MB/s	1.5 MB/s	120 MB/s
<i>I/O Rate</i>	600 I/Os/s	55 I/Os/s	3900 I/Os/s
<i>MTTF</i>	250 KHrs	50 KHrs	??? Hrs
<i>Cost</i>	\$250K	\$2K	\$150K

Disk Arrays have potential for

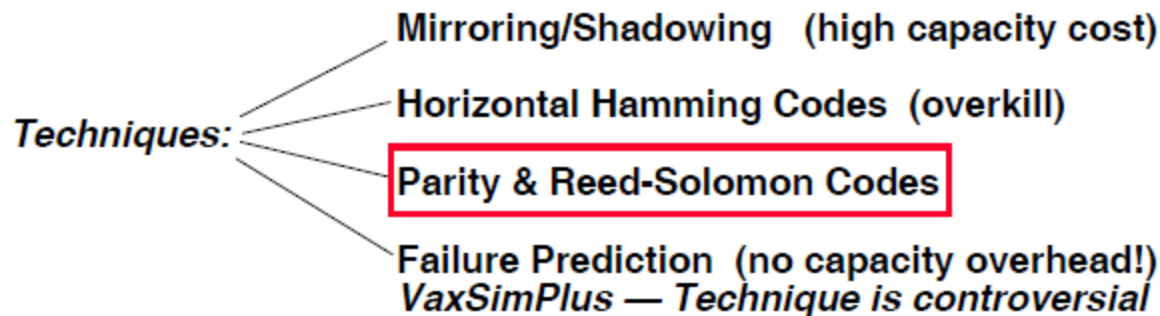
- large data and I/O rates
- high MB per cu. ft., high MB per KW
- reliability?

Array Reliability

- Reliability of N disks = Reliability of 1 Disk \div N
 - ▣ 50,000 Hours \div 70 disks = 700 hours
 - ▣ Disk system MTTF: Drops from 6 years to 1 month!
- Arrays (without redundancy) too unreliable to be useful!

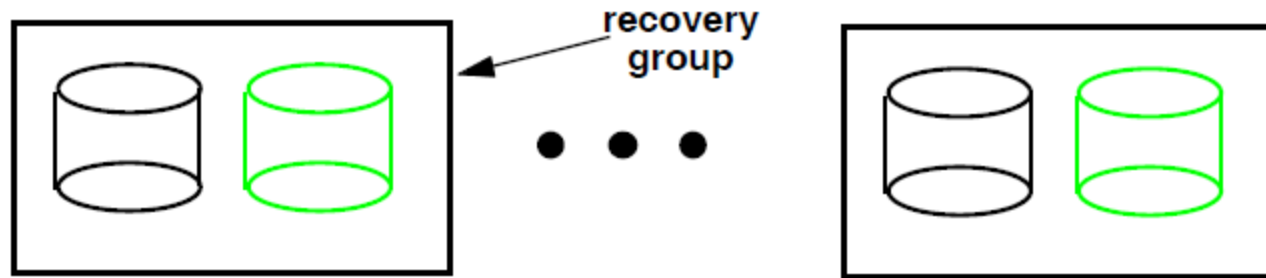
Redundant Arrays of Disks

- Files are "striped" across multiple spindles
- Redundancy yields high data availability
 - ▣ Disks will fail
 - ▣ Contents reconstructed from data redundantly stored in the array
 - Capacity penalty to store it
 - Bandwidth penalty to update



Redundant Arrays of Disks

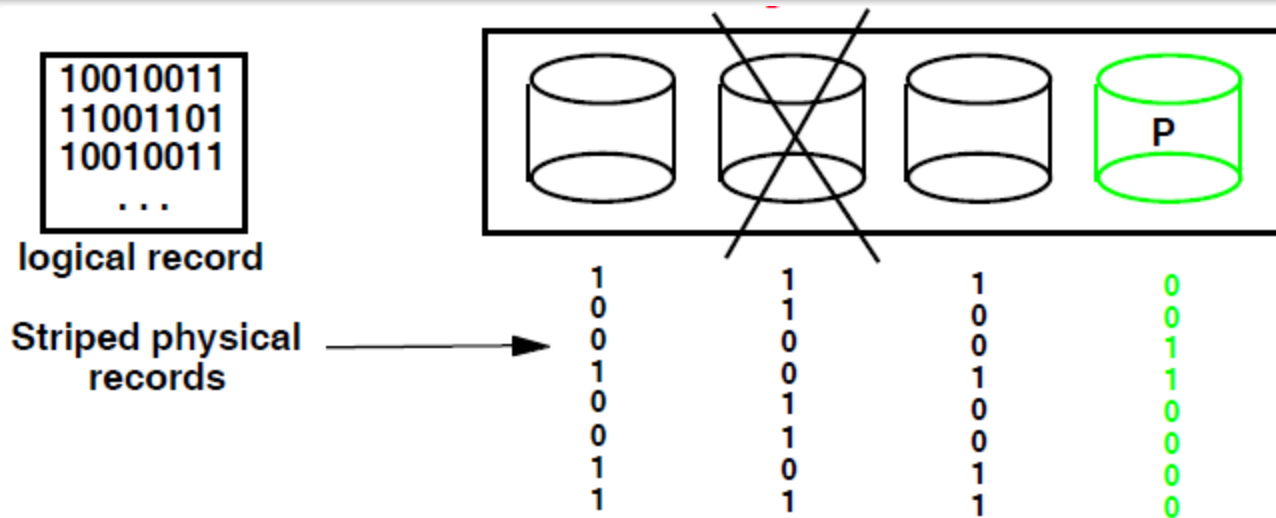
RAID 1: Disk Mirroring/Shadowing



- ❑ Each disk is fully duplicated onto its "shadow"
 - ▣ Very high availability can be achieved
- ❑ Bandwidth sacrifice on write:
 - ▣ Logical write = two physical writes
- ❑ Reads may be optimized
- ❑ Most expensive solution: 100% capacity overhead
 - ▣ *Targeted for high I/O rate , high availability environments*

Redundant Arrays of Disks

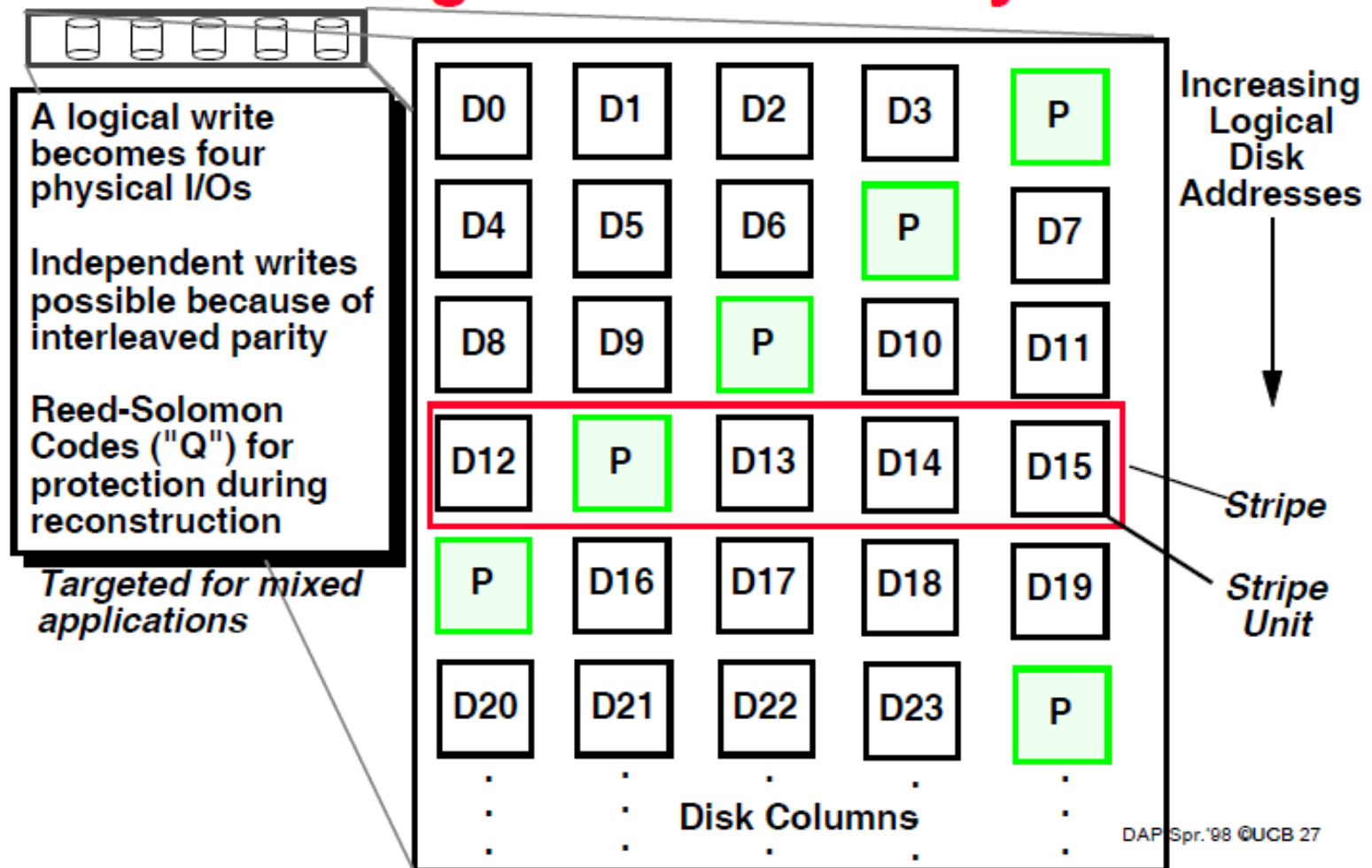
RAID 3: Parity Disk



- Parity computed across recovery group to protect against hard disk failures
 - ▣ 33% capacity cost for parity in this configuration
 - ▣ wider arrays reduce capacity costs, decrease expected availability, increase reconstruction time
- Arms logically synchronized, spindles rotationally synchronized
 - ▣ logically a single high capacity, high transfer rate disk
- *Targeted for high bandwidth applications: Scientific, Image Processing*

Redundant Arrays of Disks RAID 5+:

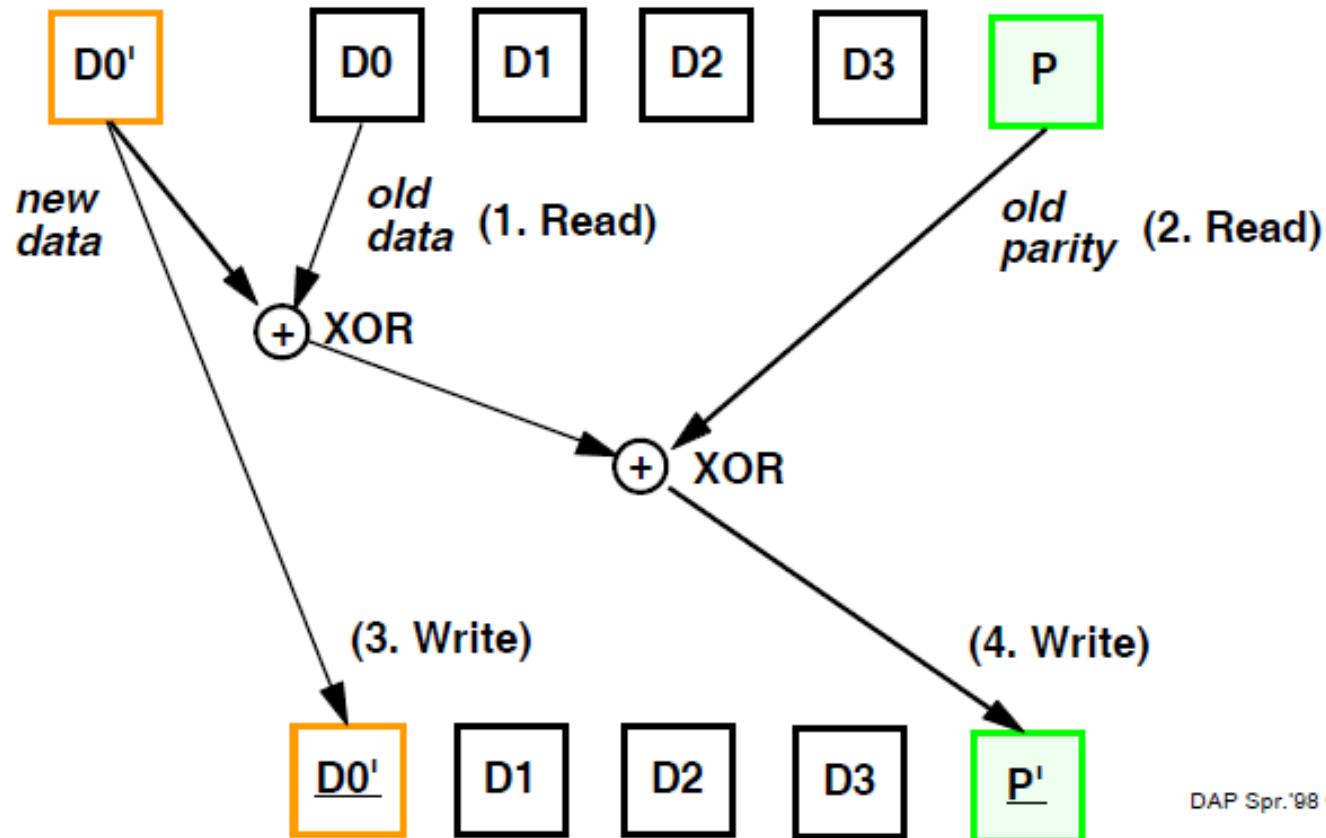
High I/O Rate Parity



Problems of Disk Arrays: Small Writes

RAID-5: Small Write Algorithm

1 Logical Write = 2 Physical Reads + 2 Physical Writes



RAID

- *Disk Mirroring, Shadowing (RAID 1)*
 - ▣ Each disk is fully duplicated onto its "shadow"
 - ▣ Logical write = two physical writes
 - ▣ 100% capacity overhead
- *Parity Data Bandwidth Array (RAID 3)*
 - ▣ Parity computed horizontally
 - ▣ Logically a single high data bw disk
- *High I/O Rate Parity Array (RAID 5)*
 - ▣ Interleaved parity blocks
 - ▣ Independent reads and writes
 - ▣ Logical write = 2 reads + 2 writes
 - ▣ Parity + Reed-Solomon codes

