# Marshmallow SER 502 Team 15

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### **MARSHMALLOW**

This is a Language that has been developed from scratch using C++ and STL

The language has been developed based on the Imperative programming paradigm.

The grammar of the language was made to follow the recursive descent model

The language can support Assignment, Arithmetic operations, Logical operations, Comparison operators, Assignment operations, Identifiers, Print statements, Continue statements, Return statements, Break statements, Conditional statements, Loops statements, Functions

### **Grammar**

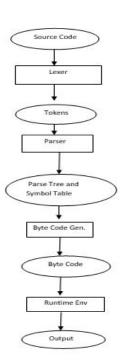
```
marshmallow ::= (stmt)+
letter ::= [a-zA-Z]
digit ::= [0-9]
identifier ::= (letter|"_") (letter | digit | "_")*
funcname ::= identifier
integer ::= digit(digit)*
stmt
       ::= assignment_stmt NEWLINE
          | print_stmt NEWLINE
           | return_stmt NEWLINE
           | break_stmt NEWLINE
           continue_stmt NEWLINE
           | exec_stmt NEWLINE
       | if_stmt
           | while_stmt
           | func_def
```

### **Grammar Contd...**

```
block ::= NEWLINE INDENT (stmt )+ DEDENT
if_stmt ::= "if" expression block
        ( "elif" expression block )*
        ("else" block)?
while_stmt ::= "while" expression block
funcdef ::= "function" funcname "(" parameters ")" block
parameters ::= identifier (',' identifier)*
print ::= "print" expression
assignment_stmt ::=
identifier::= "=" expression
return ::= "return" expression
break ::= "break"
continue ::= "continue"
arguments ::= expression (','expression)*
exec_stmt ::= funcname "(" arguments ")"
```

# **Design flow**

- 1. Source Code
- 2. Lexer
- 3. Tokens
- 4. Parser
- 5. Parse Tree and Symbol Table
- 6. ByteCode Generation
- 7. Byte Code
- 8. Runtime Env
- 9. Output



### **Grammar Contd...**

# Steps involved in creating language

- 1. Lexer: This takes in the program as input and creates tokens out of the program.
  - i. This analyzes the input program to create tokens.
- 2. Parser: Tokens from the Lexer is fed as input to Parser which generates Parse Tree and Symbol table.
  - i. This is a recursive descent parser with look ahead.
  - ii. This part takes care of semantic analysis and generating the parse tree.
  - iii. This generates symbol table which is a doubly linked N-ary tree structure.
  - iv. This throws an error if the given program has syntax errors.
- 3. Intermediate Code: This generates the bytecode based on the parse tree in which is in agreement with the runtime.
  - i. This traverses the parse tree and generates the bytecode using opcodes.
- 4. Runtime Environment: This takes the bytecode and does the execution of the program written. This gives out the output after completing the execution.
  - Stack Model is used for execution.
  - ii. Bytecode is traversed and executed with the stack holding current values.

### SAMPLE PROGRAM EXECUTION

```
a = 5
i = 1
if (a%2 == 0)
  print (a)
  else
  while (i<a)
    print(i)
  i = i+1</pre>
```

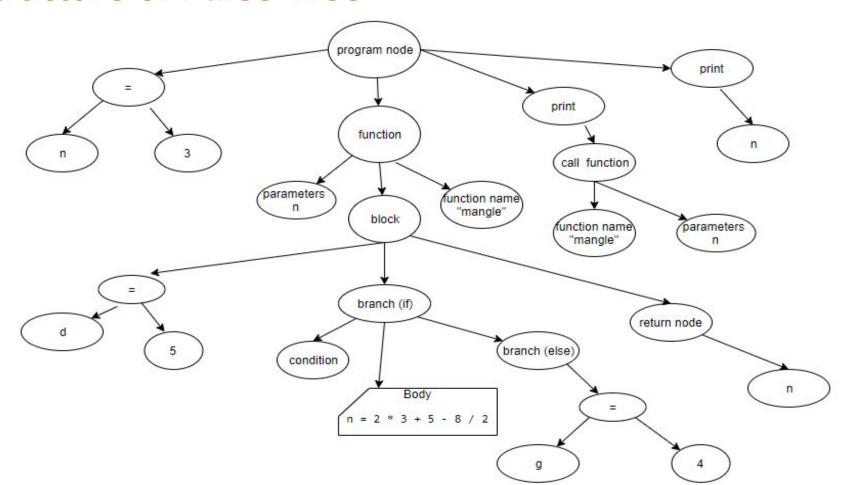
### **LEXER**

```
String[] program = ["a","=","5","NEWLINE","i",","","NEWLINE","INDENT","print","(", "=","1","NEWLINE","if","(","a","%","2","==","0",")","NEWLINE","INDENT","print","(", "a",")","NEWLINE","UEDENT","else,"NEWLINE","INDENT","while","(","i",">","NEWLINE","i","=","i","+","1","NEWLINE","DEDENT","DEDENT","DEDENT","DEDENT","
```

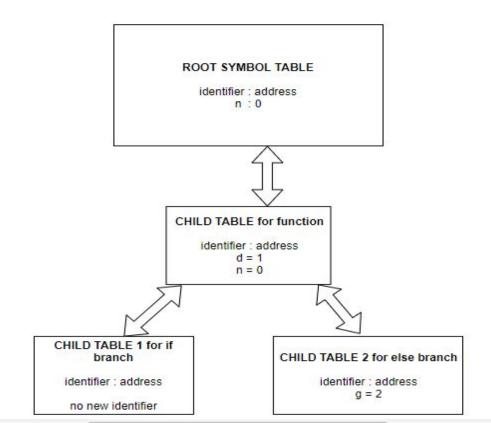
### **SAMPLE PROGRAM**

```
n = 3
function mangle(n)
    d = 5
    if 2 < 3
         n = 2 * 3 + 5 - 8 / 2
    else
         g = 4
return n
print mangle(n)
print n
```

### **Structure of Parse Tree**

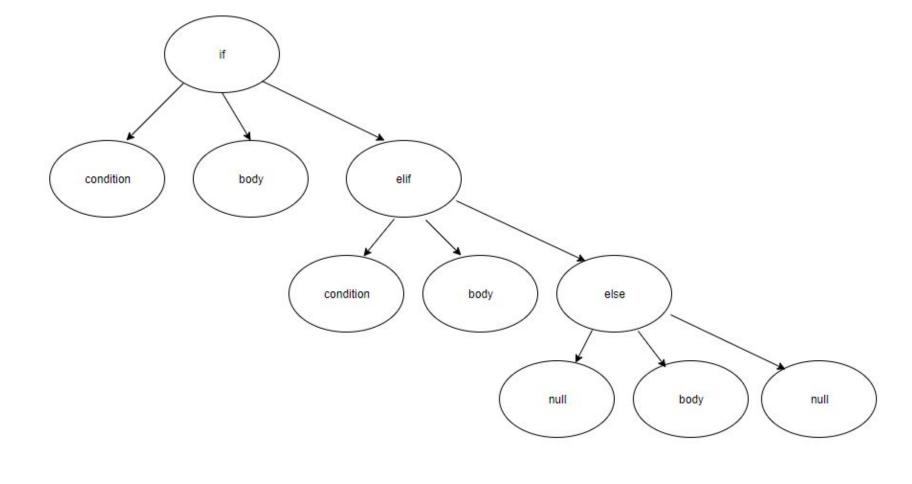


### **SYMBOL TABLE**

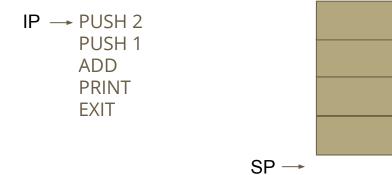


# **ByteCode**

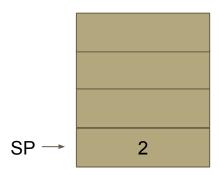
```
0 PUSH 5
2 STORE #0
 4 PUSH 1
6 STORE #1
 8 PUSH 0
10 PUSH 2
12 LOAD #0
 14 MOD
  15 EQ
16 BRF $22
18 LOAD #0
 20 PRINT
 21 BR $40
23 PUSH 10
25 LOAD #1
   27 LT
28 BRF $40
30 LOAD #1
 32 PRINT
33 PUSH 1
35 LOAD #1
  37 ADD
38 STORE #1
  40 EXIT
```

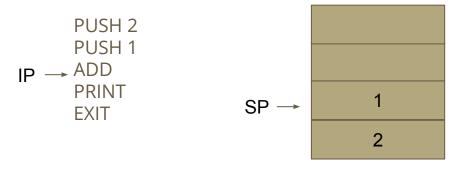


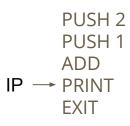
PUSH 2 PUSH 1 ADD PRINT EXIT print 1 + 2





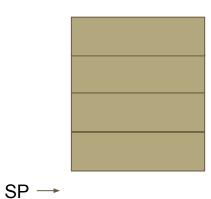








```
PUSH 2
PUSH 1
ADD
PRINT
IP → EXIT
```



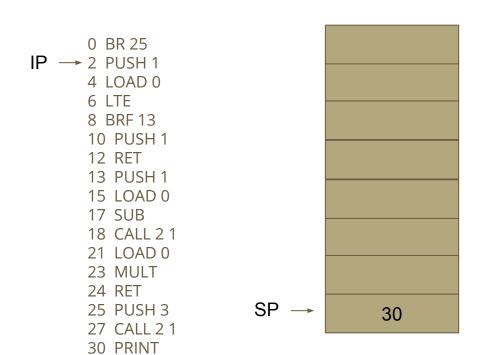
```
0 BR 25
2 PUSH 1
4 LOAD 0
6 LTE
8 BRF 13
10 PUSH 1
12 RET
13 PUSH 1
15 LOAD 0
17 SUB
18 CALL 21
21 LOAD 0
23 MULT
24 RET
25 PUSH 3
27 CALL 21
30 PRINT
```

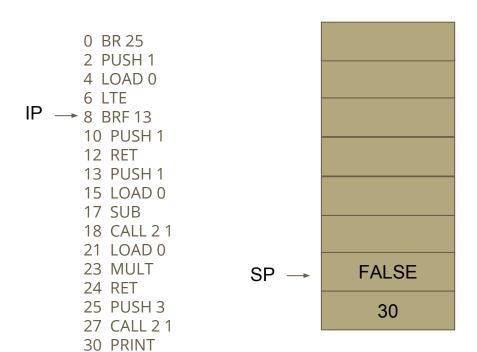
```
Function fact(n)
   if(n \le 1)
      return 1
   return(n *fact (n-1))
fact(3)
```

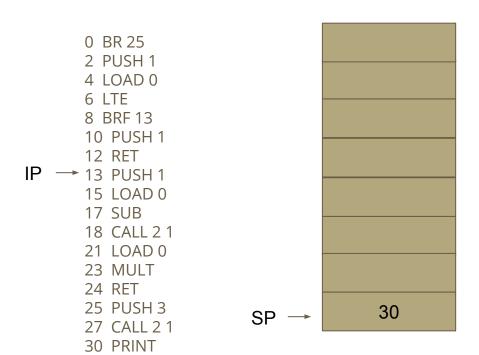
```
IP → 0 BR 25
      2 PUSH 1
      4 LOAD 0
      6 LTE
      8 BRF 13
      10 PUSH 1
      12 RET
      13 PUSH 1
      15 LOAD 0
      17 SUB
      18 CALL 21
      21 LOAD 0
      23 MULT
      24 RET
      25 PUSH 3
      27 CALL 2 1
                         SP →
      30 PRINT
```

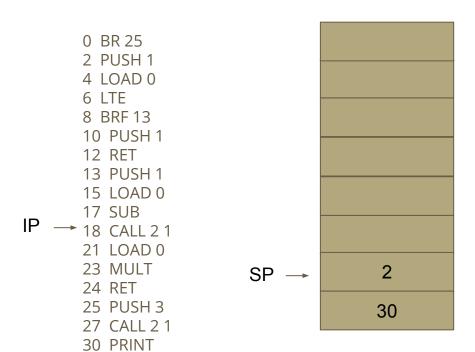
```
0 BR 25
      2 PUSH 1
      4 LOAD 0
      6 LTE
      8 BRF 13
      10 PUSH 1
      12 RET
      13 PUSH 1
      15 LOAD 0
      17 SUB
      18 CALL 21
      21 LOAD 0
      23 MULT
      24 RET
IP → 25 PUSH 3
      27 CALL 2 1
                         SP →
      30 PRINT
```

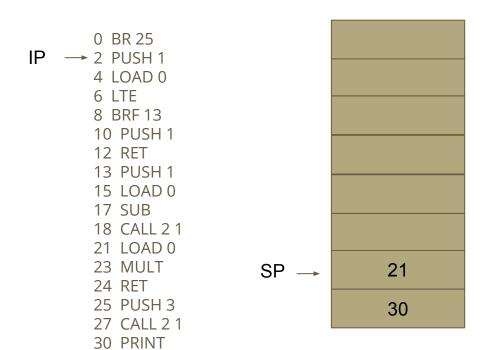
```
0 BR 25
   2 PUSH 1
   4 LOAD 0
   6 LTE
   8 BRF 13
   10 PUSH 1
   12 RET
   13 PUSH 1
   15 LOAD 0
   17 SUB
   18 CALL 21
   21 LOAD 0
   23 MULT
   24 RET
   25 PUSH 3
                                     3
                      SP →
→ 27 CALL 2 1
   30 PRINT
```



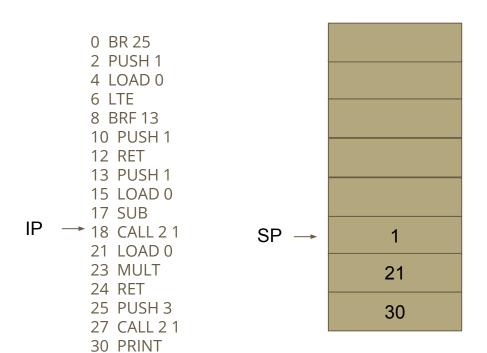




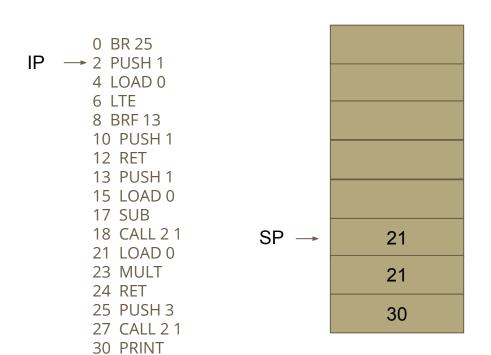




Context2: 0 - 2

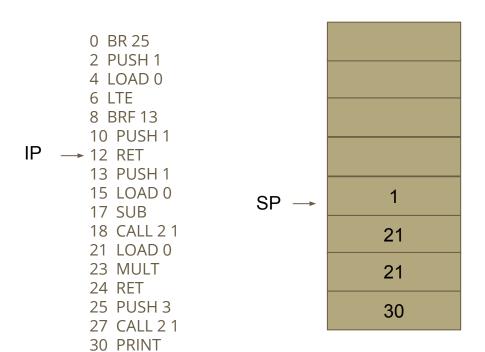


Context2: 0 - 2



Context3: 0 - 1

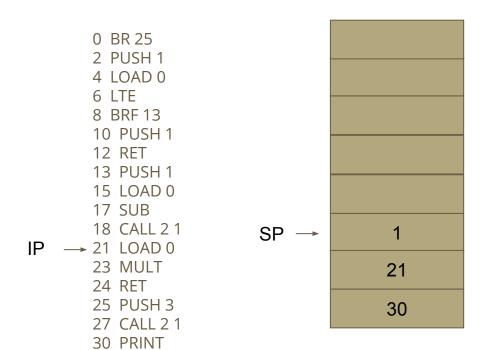
Context2: 0 - 2



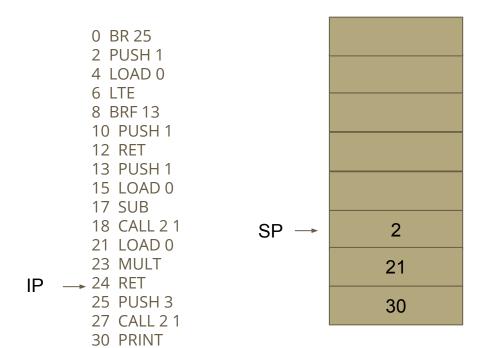
Context3: 0 - 1

Context2: 0 - 2

Context1: 0 - 3



Context2: 0 - 2



Context2: 0 - 2

