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My method $h'(k)$ to convert String name into integer:

```
hh = 1
for x in k:
    hh *= ord(x)
h = hh % self.size
```

quadratic probing

c1	c2	the minimum table size	the average number of steps(<5)
1	1	2539	4.8184
2	0.5	2731	4.7928
3	0.3	2543	4.8784

double hashing

$h''(k)$	the minimum table size	the average number of steps
$g = ((5^{**}0.5) - 1)/2$ $\text{round}(\text{self.size} * (g * h - \text{round}(g * h)))$	2531	4.9052
$3 - (h \% 3)$	3187	4.2088
$1 + (\text{round}(h/\text{self.size}) \% (\text{self.size} - 1))$	2851	4.514

conclusion:

The hypothesis is not true. Double Hashing won't allow us to use smaller tables than Quadratic Probing does. The essential factor for a smaller table is good hash functions. For example, if the second hash function of double hashing is $(3 - (h \% 3))$, the table size will be bigger than the table of quadratic probing with $c1 = 2$, $c2 = 0.5$. Therefore, the collision solution won't have a great impact on the size of the table.

PS: since the table size and the average number of steps are not always in linear relation, it's hard to find the exact smallest table size to achieve the steps < 5. I just can find the approximate one.