

Tutorial-3

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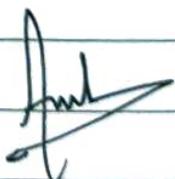
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Tutorial - 3

① Pseudocode for Linear search

```
for ( i=0 to n )
{
    if ( arr[i] == value )
        // element found
}
```

②

void insertion (int arr[], int n) // recursive

```
{ if ( n <= 1 )
```

```
    return;
```

```
insertion ( arr, n-1 );
```

```
int nth = arr[n-1];
```

```
int j = arr[n-2];
```

```
while ( j >= 0 && arr[j] > nth )
```

```
{ arr[j+1] = arr[j];
```

```
    j--;
}
```

```
arr[j+1] = nth;
```

```
for ( i=1 to n )
```

// iterative

```
{ key ← A[i]
```

```
j ← i-1
```

```
while ( j >= 0 and A[j] > key )
```

```
{ A[j+1] ← A[j]
```

```
    j ← j-1
```

```
A[j+1] ← key
```

Insertion sort is online sorting because it doesn't know the whole input, more input can be inserted with the insertion sorting is running.

(3) Complexity

Name	Best	Worse	Average
Selection Sorting	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble Sorting	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion Sorting	$O(n)$	$O(n^2)$	$O(n^2)$
Heap Sorting	$O(n \log(n))$	$O(n \log(n))$	$O(n \log(n))$
Quick Sorting	$O(n \log(n))$	$O(n^2)$	$O(n \log(n))$
Merge Sorting	$O(n \log(n))$	$O(n \log(n))$	$O(n \log(n))$

(4) Inplace Sorting

Bubble

Selection

Insertion

Quick Sort

Heap Sort

Stable Sorting

Merge Sort

Bubble

Insertion

Count

Online Sorting

Insertion

```

(5) int binary (int arr[], int l, int r, int n) // recursive
{
    if (r >= l)
    {
        int mid = l + (r - l) / 2;
        if (arr[mid] == x)
            return mid;
        else if (arr[mid] > x)
            return binary (arr, l, m-1, n);
        else
            return binary (arr, m+1, r, x);
    }
    return -1;
}

```

```

int binary ( int arr[], int l, int r, int n)
{
    while ( l <= r )
    {
        int m = l + (r-l)/2;
        if ( arr[m] == n )
            return m;
        else if ( arr[m] > n )
            r = m-1;
        else
            l = m+1;
    }
    return -1;
}

```

Time complexity of

Binary Search $\Rightarrow O(\log n)$

Linear Search $\Rightarrow O(n)$

(6) Recurrence relation for binary recursive search

$$T(n) = T(n/2) + 1$$

where $T(n)$ is the time required for binary search in an array of size 'n'.

(7) int find (A[], n, K)

{ Sort (A, n)

for (i=0 to n-1)

{ n = binarySearch (A, 0, n-1, K-A[i])

if (n)

return 1

}

return -1

}

$$\begin{aligned}\text{Time complexity} &= O(n \log(n)) + n \cdot O(\log n) \\ &= O(n \log(n))\end{aligned}$$

- ⑧ • QuickSort is the fastest general purpose sort.
 • In most practical situations, quicksort is the method of choice. If stability is important and space is available, merge sort might be best.

- ⑨ A pair $(a[i], a[j])$ is said to be inversion if $a[i] > a[j]$

In $\text{arr[]} = \{7, 21, 31, 8, 10, 1, 20, 6, 4, +5\}$
 total no. of inversion are 31, using merge sort.

- ⑩ The worst case time complexity of quick sort is $O(n^2)$.
 This case occurs when the picked pivot is always an extreme (smallest or largest) element. This happens when input array is sorted or reverse sorted.
 The best case of quick sort is when we will select pivot as a mean element.

- ⑪ Recurrence relation of

$$\text{Merge Sort} \rightarrow T(n) = 2T(n/2) + n$$

$$\text{Quick Sort} \rightarrow T(n) = 2T(n/2) + n$$

- Merge Sort is more efficient and works faster than quick sort in case of larger array size or datasets.
- Worst case complexity for quick sort is $O(n^2)$ whereas $O(n \log(n))$ for merge sort.

(12) Stable selection sort

```
void stableSelection( int arr[], int n )
{ for ( int i=0 ; i<n-1 ; i++ )
    { int min = i;
      for ( int j=i+1 ; j<n ; j++ )
      { if ( arr[min] > arr[j] )
          min = j;
      }
      int key = arr[min];
      while ( min > i )
      { arr[min] = arr[min-1];
        min--;
      }
      arr[i] = key;
    }
}
```

(13) Modified bubble sorting

```
void bubble( int a[], int n )
{ for ( int i=0 ; i<n ; i++ )
    { int swaps = 0;
      for ( int j=0 ; j<n-1-i ; j++ )
      { if ( a[j] > a[j+1] )
          { int t = a[j];
            a[j] = a[j+1];
            a[j+1] = t;
            swaps++;
          }
      }
      if ( swaps == 0 )
        break;
    }
}
```