Fibonacci Seguence

If you have even a passing interest in wethousers, you've house of the Fibonack's Sequence.

 $f_{0} = 0$ $f_{1} = 1$ $f_{2} - f_{1} + f_{0} = 2$ $f_{3} - f_{2} + f_{1} = 3$

etc.

$$\int_{n}^{\infty} f_{n-1} + f_{n-2}$$

The nature of this sequence makes
if a great candidate for rows:-.

de f. Bonacci (n)

if. (n = = 0):

return \$

elf (n==1) :

return 1

elso:

voturn f.bonacci (n-1) + f.bonacci (n-2)

Test:

N	fibonnaci (n)
4	\$ 1
7	Sibonacii (1) +

S. Bonnai (2) + fibrai -> Seems løbe it should work (i) Sæ Libonaci.pu -) and it does !!

We are so clever. Recursionis so cool. We are fxx!xrg smont.

Or, maybe we're not.

What is the computational complexity of this algorithm??

n	* of Calls
4	1
1 2 3 4 5 6 7 8	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
9	u7 + 29 = 76

I'm not sue let's unte a

python program to colculate the # of

calls; and plat it, to see that it

looks like....

See fibonacci - complexity. Pg.

^

Kesult: # Calls ~ 800,000 !!! in general, # calls re Complexity = $O(e^n)$ It is difficult for me to suphosize In bronibly bad this is.

Is there any thing we can do?
Please, Dr. Brosh, save us!!!

OKay, lot's calm down....

When over we have an algorithm

that is $O(e^n)$, the thought is

that we probably write fine something.

That we probably write fine something.

better ... liee,

Von-recursive Fibonaci:

def fibonacci (n):

if
$$n = 0$$
:

veturn ϕ

else:

else:

$$+1 = 0$$
 $+2 = 1$
for i in varye $(2, n+1)$:
$$+3 = +2 + +1 = sun ot truo terms$$

$$+1 = +2 = +3$$

$$+2 = +3$$

$$+2 = +3$$

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$$+2 = +3$$

But wait! There's more.

As you may know, the ratio of Successive terms in the Fibonacci Sequence is interesting, as well.

On 1, 1, 2, 3, 5, 8, 13, 2), ...

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r=1 r=2 r=1.5 r=1.66 r=1.654

ratio seems to be converging ... [1.62-...

Let's use some basic maths: = \$\frac{1}{60003988}...

$$\int_{n} = \int_{n-1}^{n} + \int_{n-2}^{n} \\
 \left(\int_{x_{n-1}}^{x_{n-1}} + \int_{x_{n-1}}^{x_{n-2}} + \int_{x$$

 $\frac{1}{\phi} = 1 + \frac{1}{\phi}$

Quadretiz
$$\Rightarrow b^2 - b - 1 = 0$$
 $b^2 - b - 1 = 0$
 $b^2 - b - 1 = 0$

OKay, So not so magical, perhaps. (3) After a bit more moth (Binet 1843), We find that:

$$\int_{n}^{\infty} = \left(\frac{4}{4}\right)^{n} - \left(\frac{-1}{\phi}\right)^{n}$$

$$= \sqrt{5}$$

New Abonacti Proton.

import math inhi = 1.0 + math. sqrt (5.0)

Let's chech:

2	f franta $f(n)$	f(n)
0	0	0
6	8	8
20	6765	6765
40	102 334 155	102334155
60	1548008755920	
75	211 1485077978055	211 148507797 805¢
		Dops

Mhy? It's coming from the fact

that our calculation of $\phi = \frac{1+05}{2}$

is being trundated at worset in decimal places. This is the maxinom according that one com achieve for storing float in 5 point #'s using 64-bit binomy flouting point representation. So, what to do? -... The really useful modules in python is decimal import decimal as dec The decimal mudule allres us to specify the number of decimal Places, up to relatively arbitrony

precision!

-

```
dec. get context (). prec = 50
onehalf = dec. Decimal (0.5)
This will actually store
    one = dec. Decomal(1)
tus = dec. Decimal (2)
 five = dec. Decimal (5)
Phi = (one + five ** one half)

/ two
      fibonacci (n):
def
        global Phi
                one, two, five, one holf
        global
               round [[Phi ** "
        return
               - (- one/phi) *** n)/
                     five *xx one half
See fibonacci-tornula. Py
```

Summany:

- 1) recursive Fibonnacci -> O(e")
- 2) non-recursive Fibonacci -> O(n)
- 3 formula Fronnacci *-> 0 (1)
 - -) reguires higher level

 floating pt. precision library.

 -> 50, might be slower!!