

Static Taint Analysis for C with LLVM

Xavier Noubissi

Department of Electrical and Computer Engineering
University of Waterloo

Problem: Prevent Software Vulnerabilities

- Format String Attacks
- SQL Injection
- Cross Site Scripting, etc.

Solution: Track Used of Non Trusted Program Input

- Non trusted program input: **Tainted Input**
- Source: origin of tainted input
- Sink: use of tainted input
- Taint Propagation: tracking tainted input

Example

```
1  int main() {
2      int x, b1, b2, y;
3      scanf("%d", &x);
4      b1 = even(x);
5      b2 = odd(3);
6      y = compute(x);
7      return 0;
8  }

10 int compute(int x) {
11     int sum, i;
12     if (x == 2)
13         scanf("%d", &sum);
14     else
15         sum = 0;
16     for (i = 0; i < x; ++ i)
17         sum += i;
18     return sum;
19 }

21 int odd(int x) {
22     if (x == 1)
23         return 0;
24     else
25         return even(x - 1);
26 }

28 int even(int x) {
29     if (x == 0)
30         return 1;
31     else
32         return odd(x - 1);
33 }
```

Figure 1. Motivating Example

Solution: In This Project

- Implicit Taint Propagation: due to **Control Flow**
- Explicit Taint Propagation: due to **Data Flow**

Expected Contributions

- Algorithm to statically detect use of tainted values in C programs
- Handling of interprocedural taint propagation
- Implementation of the algorithm in LLVM

Taint Analysis: Relevant C Program Statements

Statement Type	C Code
COPY	$p = q$
LOAD	$p = *q$
STORE	$*p = q$
ADDROF	$p = \&a$
SOURCE	<i>call gets</i>

Taint Analysis: Transfer Functions

- **COPY** [$p = q$]: **taint p** iff q is tainted
- **LOAD** [$p = *q$]: **taint p** iff it exists $t_q = *q \wedge t_q$ is tainted
- **STORE** [$*p = q$]: **nothing**
- **ADDROF** [$p = \&a$]: **nothing**
- **SOURCE** [$call\ gets(p)$]: **taint all** t_q s.t. $t_q = *p$

Interprocedural, Context-Sensitive Analysis

- Analysis of a callee start with the taint assumptions from the caller

Taint Information from Source

- Developer specify sources and sinks in configuration file
- Analysis do not analyze sources and sinks
- Analysis use annotations for sources: taint propagation

Algorithm Analyze: Implements Interprocedural Analysis

```
input : func : Proc,  
       initDataFlow: Inst  $\rightarrow$  (Var  $\rightarrow 2^{Inst}$ )  
output:  
1  $s_0 \leftarrow \text{first}(f)$   
2  $\text{input}[s_0] \leftarrow \text{initDataFlow}(s_0)$   
3  $\text{worklist} \leftarrow \{s_0\}$   
4 while  $\text{worklist} \neq \emptyset$  do  
5    $i \leftarrow \text{next}(\text{worklist})$   
6    $\text{output}[i] \leftarrow \text{FlowAnalyze}(i)$   
7   foreach  $j \in \text{succs}(i)$  do  
8     if  $\text{output}[i] \not\subseteq \text{input}[j]$  then  
9        $\text{input}[j] \leftarrow \text{input}[j] \cup \text{output}[i]$   
10       $\text{worklist} \leftarrow \text{worklist} \cup \{j\}$   
11    end  
12  end  
13 end
```

Algorithm 1: Analyze

Algorithm Flow: Implements Transfer Functions

```

input : caller : Proc, s : Inst
output:
1 switch TypeOf(s) do
2   case COPY [p = q]
3     if inFlows[q] = ∅ then
4       outFlows[p] ← inFlows[q] ∪ {s};
5     end
6   endsw
7   case LOAD [p = *q]
8     foreach a ∈ ptq(q) do
9       if inFlows[a] = ∅ then
10        outFlows[p] ← inFlows[p] ∪ {s}
11      end
12    end
13  endsw
14  case SOURCE [call func(a0, a1, ..., an)]
15    foreach k ∈ {0, 1, ..., n} do
16      if taint(k) then
17        outFlows[ak] ← inFlows[ak] ∪ {s}
18      end
19    end
20  endsw
21  case CALL [call func(a0, a1, ..., an)]
22    if caller = func then
23      foreach ak, k ∈ {0, 1, ..., n} do
24        fk ← formal(func, ak)
25        if ak ∈ P then
26          foreach b ∈ ptq(ak) do
27            inFlows[fk] ← inFlows[fk] ∪ inFlows[b]
28          end
29        end
30        else if ak ∈ A then
31          tk ← toplevel(ak)
32          foreach b ∈ ptq(tk) do
33            inFlows[fk] ← inFlows[fk] ∪ inFlows[b]
34          end
35        end
36      end
37      Flow(caller, func)
38    end
39  endsw
40  case ADDRESS [p = &a]
41  case STORE [*p = q]
42  case SINK [call func]
43  endsw

```

Implementation

- LLVM infrastructure is ready
- Need to implement the analysis

Future Work

- Make analysis modular
- Better handling of (mutual-) recursive function calls

Conclusion

- Implementation should take less than 1 months
- Need to do evaluation on real world C programs
- We are optimistic about future results

Thank You!

Comments & Questions