How to Grow a TREE from CBASS

Interactive Binary Analysis for Security Professionals

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Outline

- Background
- Interactive Binary Analysis with TREE and CBASS
- Demonstrations
- Conclusions

Interactive Binary Analysis

 Automated binary analyses useful for certain tasks (e.g., finding crashes)

Many binary analyses can't be automated

 Expert experience and heuristics are still key to binary analyses

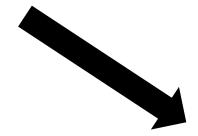
Benefits of Interactive Binary Analysis

Applicable to many security problems

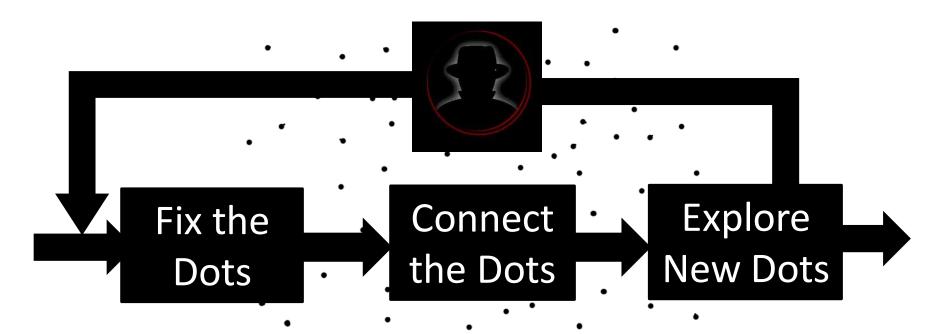
- Our tools increase productivity in:
 - Finding vulnerabilities
 - Analyzing root causes
 - Exploitability and risk assessment

Interactive Analysis Like Connecting Dots

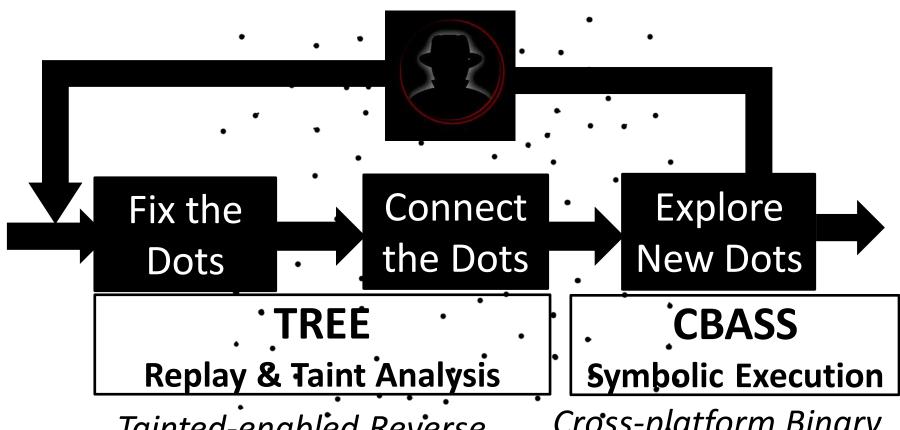




Our Tools are Designed to Help



What Do Our Tools Do?



Tainted-enabled Reverse Engineering Environment Cross-platform Binary Automated Symbolic execution System

Gaps between Research and Interactive Binary Analysis

- Existing research does not support interactive binary analysis
 - No practical tools
 - No uniform trace collection tools
 - No unified Instruction SetArchitecture(ISA) -independentanalysis tools

Bringing Proven Research Techniques to Interactive Binary Analysis

- Our tools use dynamic, trace-based, offline analysis approach
 - Interactive binary analysis [1]
 - Dynamic taint analysis ([2][3][4])
 - Symbolic execution/ SMT solver ([2][5])
 - Trace replay ([6])

Making It Practical

- TREE integrates with IDA Pro now and other mainstream binary analysis environments (later)
- TREE leverages debugging infrastructure to support tracing on multiple platforms
- CBASS uses Intermediate Representation (REIL [6][7])-based approach to support ISA-independent analysis

CBASS Supports Both Automated & Interactive Analysis

TREE
Interactive Analysis

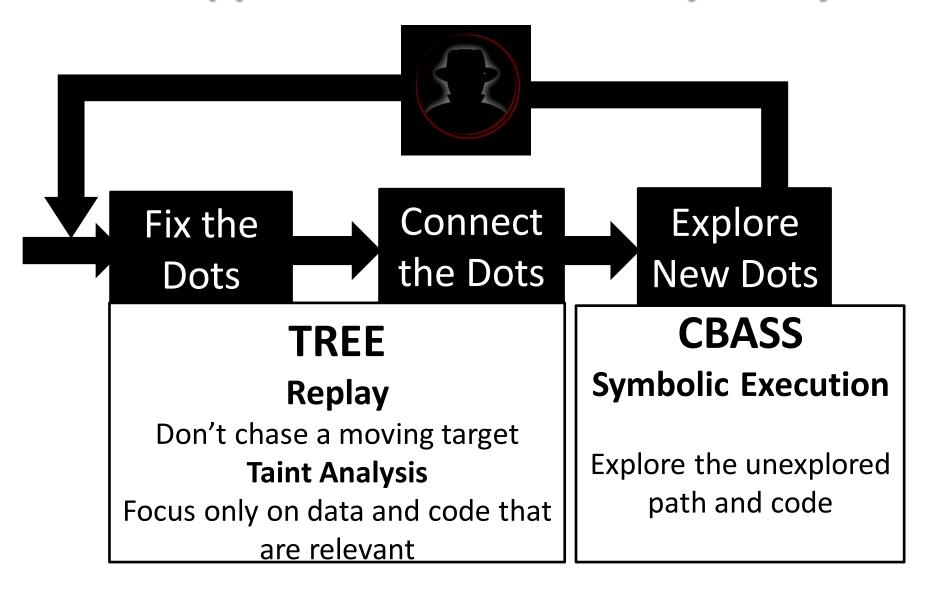
Automated Fuzzer
Automated Analysis

CBASS

IR-based Symbolic Execution Engine

TREE fills gaps for interactive analysis

Tools Support Interactive Binary Analyses



Illustrative Dots in Vulnerability Analysis: A Running Example

```
//INPUT
ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);
                                                 //Vulnerable Function
                                                 void StackOVflow(char *sBig,int num)
//INPUT TRANSFORMATIONS
                                                     char sBuf[8] = \{0\};
//PATH CONDITIONS
                                                     for(int i=0;i<num;i++)
if(sBigBuf[0]=='b') iCount++;
                                                      //Overflow when num>8
if(sBigBuf[1]=='a') iCount++;
if(sBiqBuf[2]=='d') iCount++;
                                                          sBuf[i] = sBig[i];
if(sBigBuf[3]=='!') iCount++;
if(iCount==4)// bad!
         StackOVflow(sBigBuf,dwBytesRead)
                                                      return;
else // Good
         printf("Good!");
```

Our Tools Support

Fixing the Dots (TREE)

Fix the Dots

Reverse engineers don't like moving dots

- Why do the dots move?
 - Concurrency (multi-thread/multi-core)brings non-deterministic behavior
 - ASLR guarantees nothing will be the same

Fix the Dots

- How does TREE work?
 - Generates the trace at runtime
 - Replays it offline
- TREE trace
 - Captures program state = {Instruction, Thread,Register, Memory}
 - Fully automated generation
- TREE can collect traces from multiple platforms
 - Windows/Linux/Mac OS User/Kernel and real devices (Android/ARM, Cisco routers/MIPS, PowePC)

TREE Taint-based Replay vs. Debug-based Replay

- Debug-replay lets you connect the dots
 - Single step, stop at function boundary, Breakpoint

- TREE replay connects dots for you
 - Deterministic replay with taint-point break

Our Tools Support

Connecting the Dots (TREE)

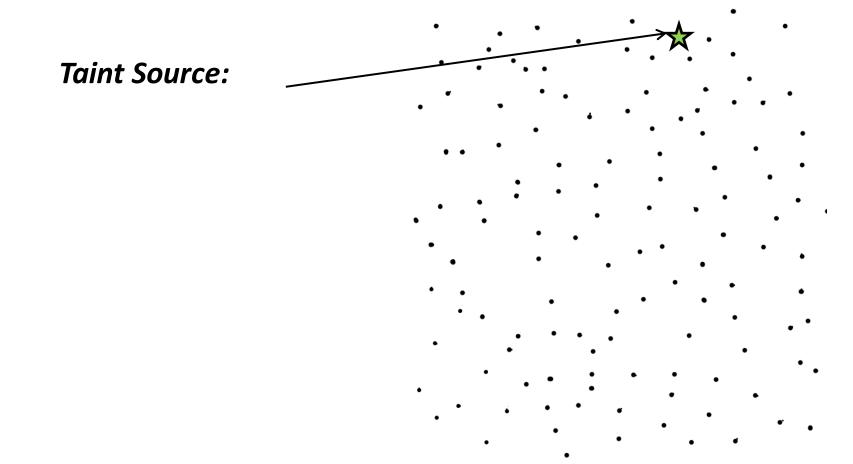
Connecting Dots is Hard

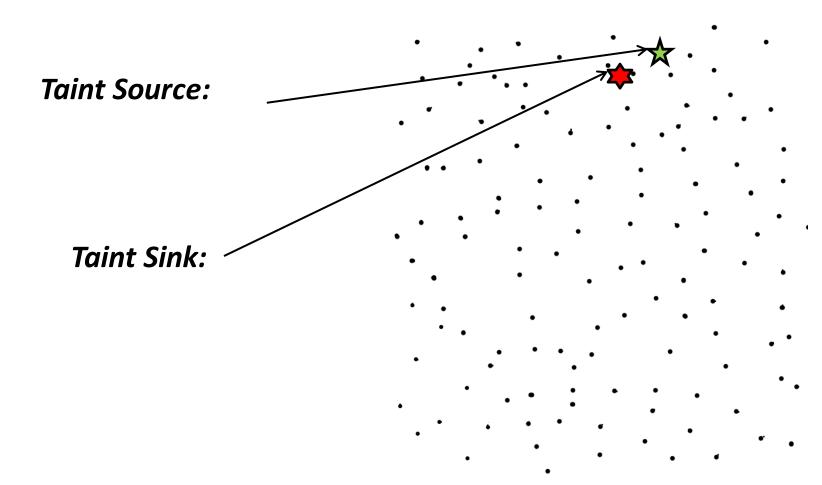
- Basic elements complex in real programs
 - Code size can be thousands (++) of lines
 - Inputs can come from many places
 - Transformations can be lengthy
 - Paths grow exponentially
- Basic elements likely separated by millions of instructions, spatially and temporally
- Multiple protections built in

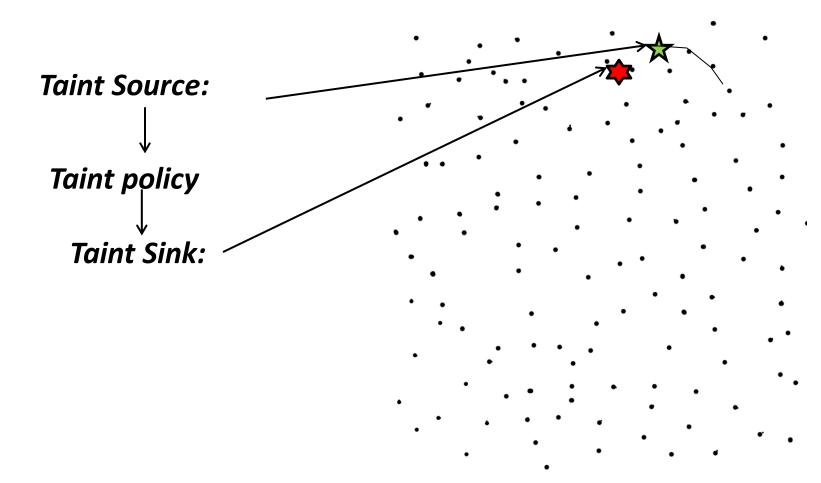
Techniques Help Connect the Dots

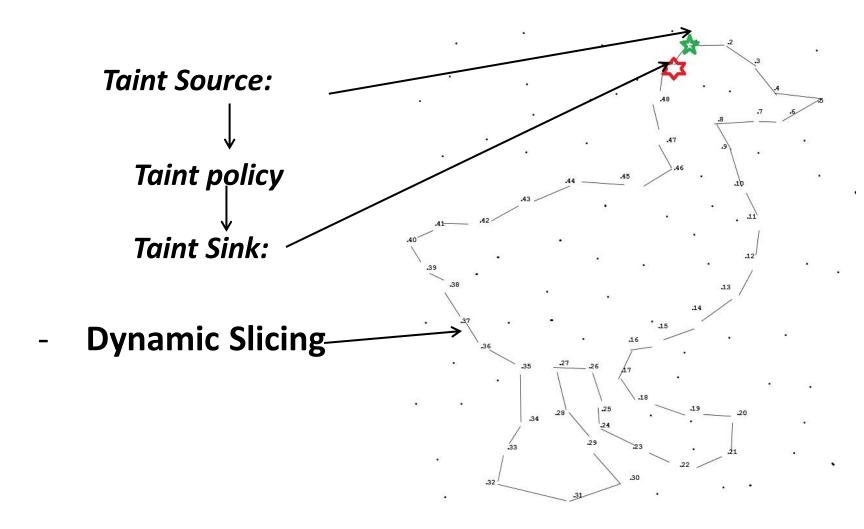
- Dynamic Taint Analysis
 - Basic Definitions
 - Taint source
 - Taint Sink:
 - o Taint Policy:

- Taint-based Dynamic Slicing
 - Taint focused on data
 - Slicing focused on relevant instructions and sequences





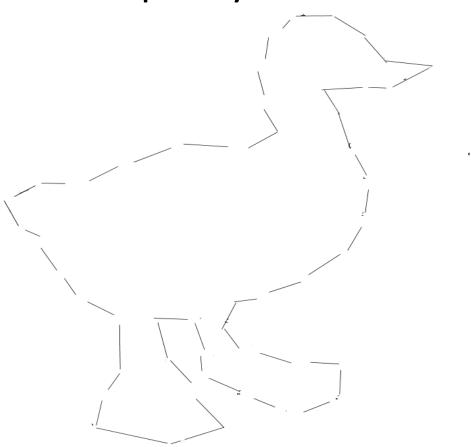




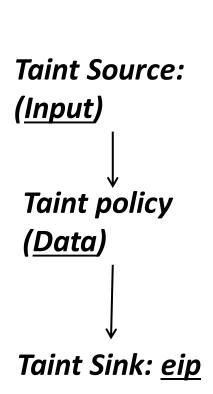
Find the Dots and Slice that Matter

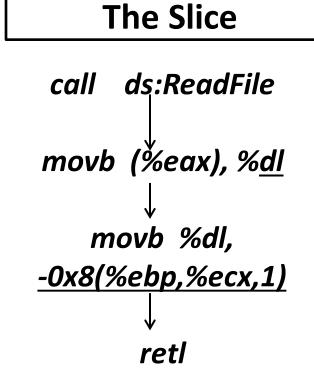
In practice, most dots don't matter – eliminate them quickly to focus on what

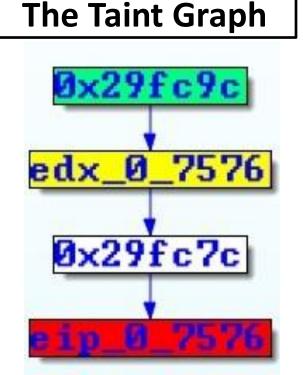
matters



Connecting Dots in Running Example







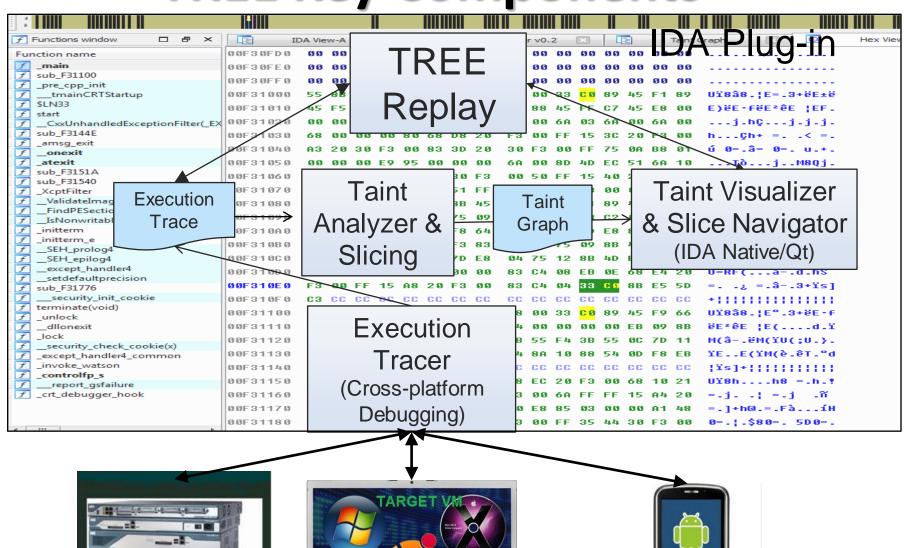
What You Connect is What You Get

- Dots can be connected in different ways
 - Data dependency
 - Address dependency
 - Branch conditions
 - Loop counter

Connect dots in different taint policies

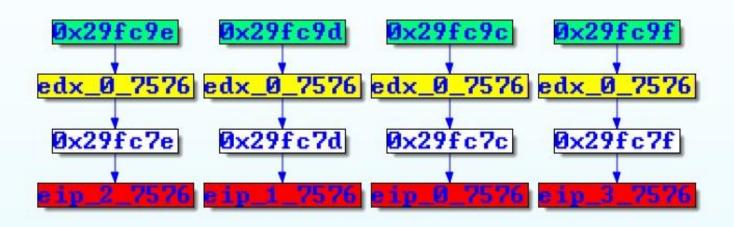


TREE Key Components



CIOSCUD

Taint Graph

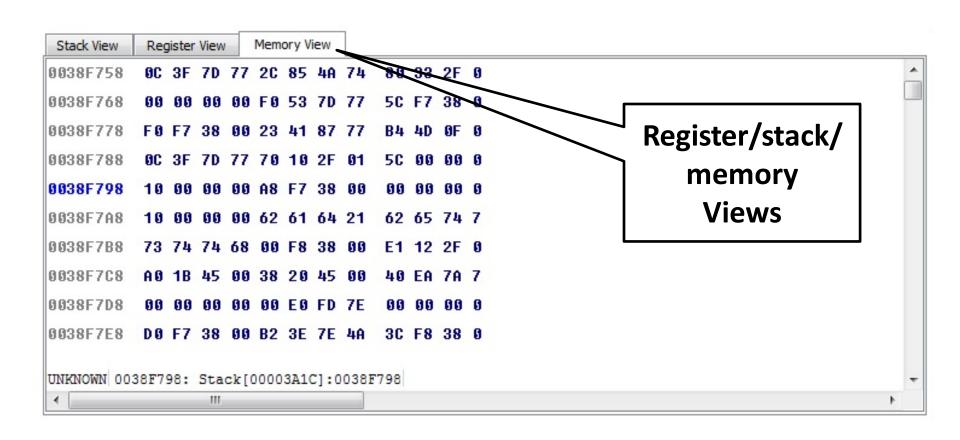


Taint Inter

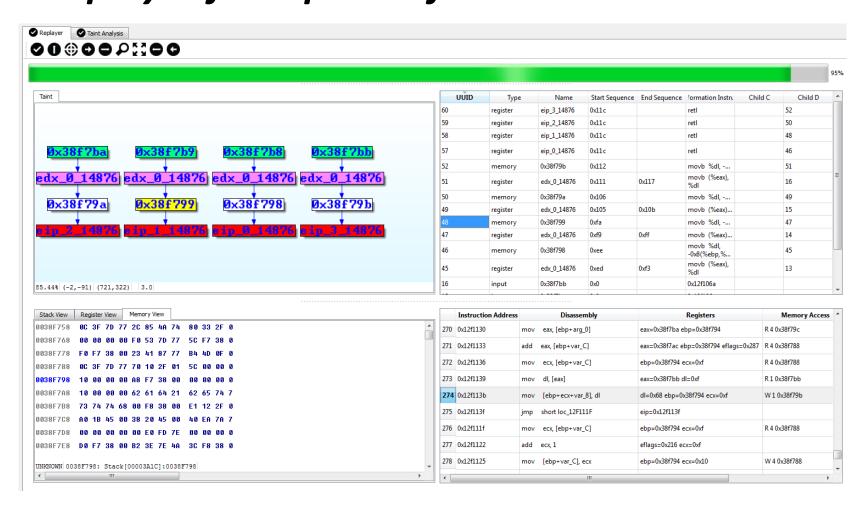
UUID	Туре	Name	Start Sequence	End Sequence	formation Instru	Child C	Child D	^
60	register	eip_3_14876	0x11c		retl		52	
59	register	eip_2_14876	0x11c		retl		50	
58	register	eip_1_14876	0x11c		retl		48	
57	register	eip_0_14876	0x11c		retl		46	
52	memory	0x38f79b	0x112		movb %dl,		51	
51	register	edx_0_14876	0x111	0x117	movb (%eax), %dl		16	Е
50	memory	0x38f79a	0x106		movb %dl,		49	
49	register	edx_0_14876	0x105	0x10b	movb (%eax)		15	
48	memory	0x38f799	0xfa		movb %dl,		47	
47	register	edx_0_14876	0xf9	0xff	movb (%eax)		14	
46	memory	0x38f798	0xee		movb %dl, -0x8(%ebp,%		45	
45	register	edx_0_14876	0xed	0xf3	movb (%eax), %dl		13	
16	input	0x38f7bb	0x0		0x12f106a			
			1					100

Execution
Trace
Table

	Instruction Address	Disassembly	Registers	Memory Access	
270	0x12f1130	mov eax, [ebp+arg_0]	eax=0x38f7ba ebp=0x38f794	R 4 0x38f79c	
271	0x12f1133	add eax, [ebp+var_C]	eax=0x38f7ac ebp=0x38f794 eflags=0x287	R 4 0x38f788	
272	0x12f1136	mov ecx, [ebp+var_C]	ebp=0x38f794 ecx=0xf	R 4 0x38f788	
273	0x12f1139	mov dl, [eax]	eax=0x38f7bb dl=0xf	R 1 0x38f7bb	
274	0x12f113b	mov [ebp+ecx+var_8], dl	dl=0x68 ebp=0x38f794 ecx=0xf	W 1 0x38f79b	
275	0x12f113f	jmp short loc_12F111F	eip=0x12f113f		
276	0x12f111f	mov ecx, [ebp+var_C]	ebp=0x38f794 ecx=0xf	R 4 0x38f788	
277	0x12f1122	add ecx, 1	eflags=0x216 ecx=0xf		
278	0x12f1125	mov [ebp+var_C], ecx	ebp=0x38f794 ecx=0x10	W 4 0x38f788	
4		111			



TREE: The Front-end of Our Interactive Analysis System Replay is focal point of user interaction



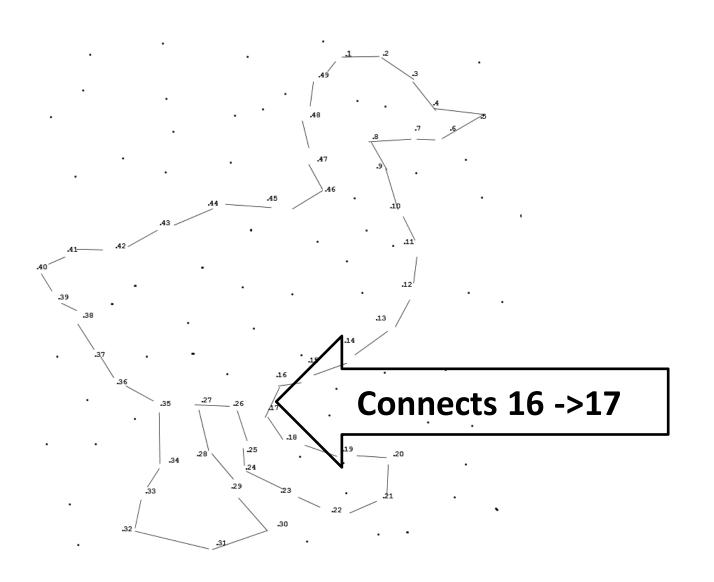
Tree Demo

Using TREE to Analyze a Crash

Our Tools Support

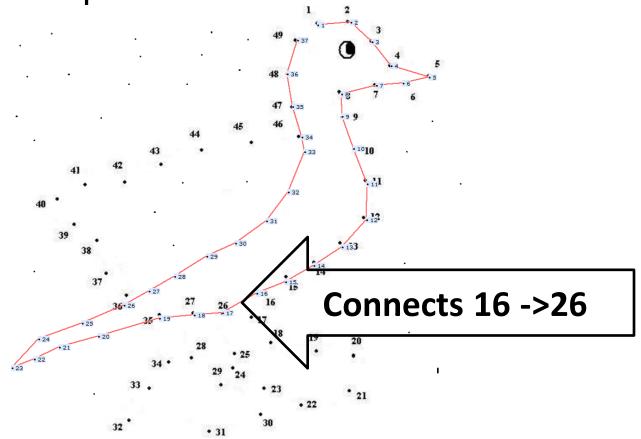
Exploring New Dots

A Key Branch Point for a Duck



The Path for a ...

 Reverse engineers don't just connect dots; they want to explore new dots:



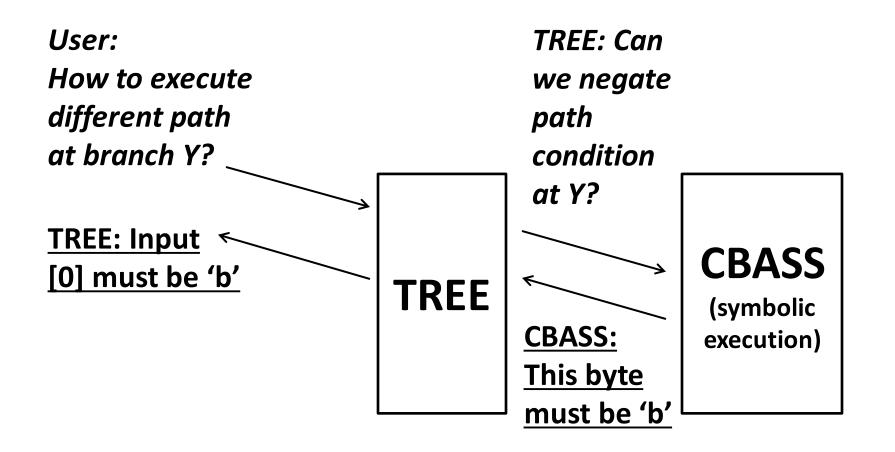
Explore New Dots

 How do you force the program to take a different path to lead to "bad!"?

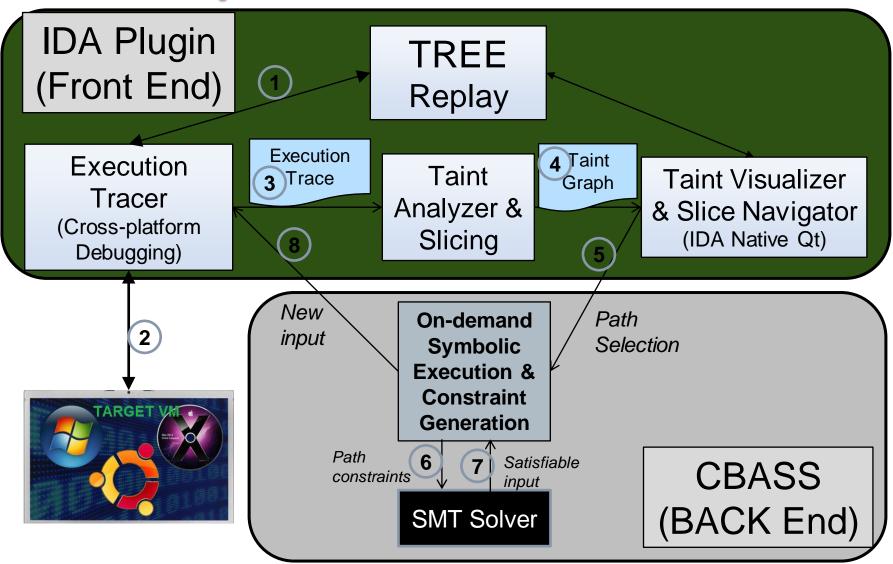
```
//INPUT
ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);
//PATH CONDITION
if(sBigBuf[0]=='b') iCount++;
if(sBigBuf[1]=='a') iCount++;
if(sBigBuf[2]=='d') iCount++;
if(sBigBuf[3]=='!') iCount++;
if(iCount==4) // "bad!" path
         StackOVflow(sBigBuf,dwBytesRead)?
Else // "Good" path
        printf("Good!");
```

Explore New Dots

 User wants execution to take different path at a branch point Y – what input will make that happen?



Explore New Dots Demo



Task 1: Force the Program to Take "bad!" Path

//INPUT

ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);

//INPUT TRANSFORMATION

.

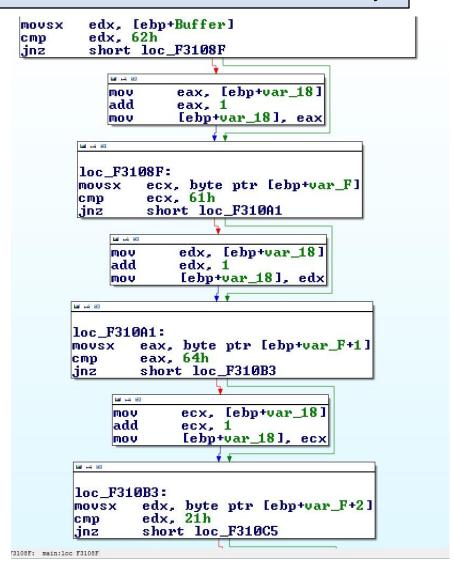
//PATH CONDITION

```
if(sBigBuf[0]=='b') iCount++;
if(sBigBuf[1]=='a') iCount++;
if(sBigBuf[2]=='d') iCount++;
if(sBigBuf[3]=='!') iCount++;
if(iCount==4) // "bad!" path
//Vulnerable Function
```

StackOVflow(sBigBuf,dwBytesRead) else

printf("Good!");

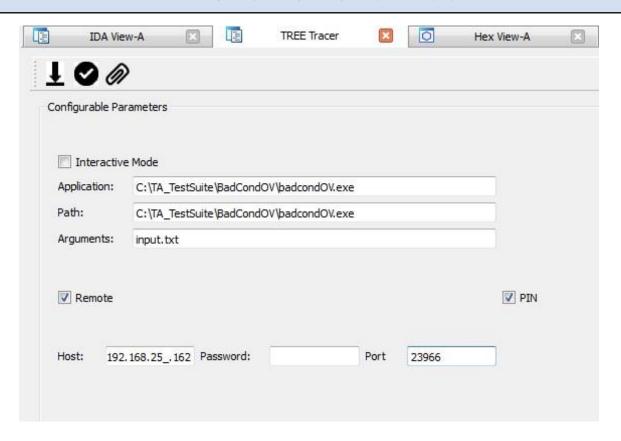
Branch Conditions In Disassembly



1 TREE Pin Trace

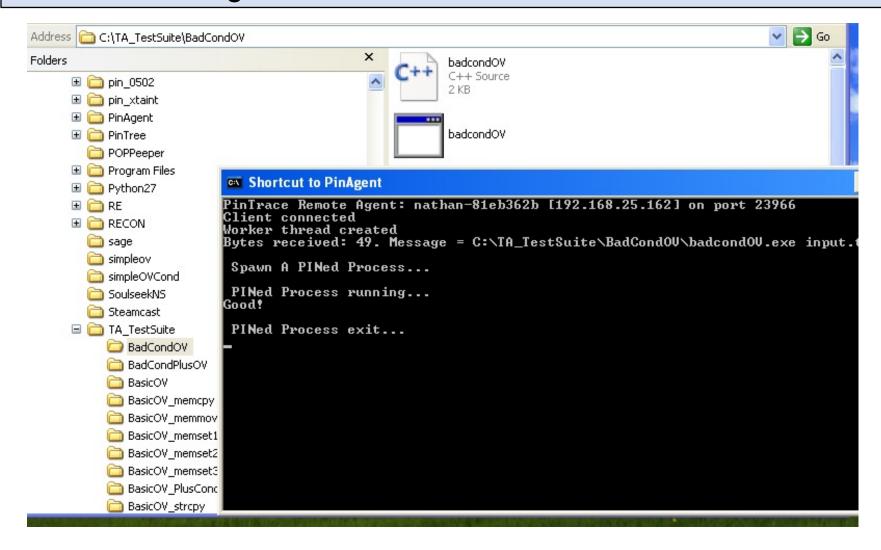
PIN: A popular Dynamic Binary Instrumentation (DBI) Framework

http://software.intel.com/en-us/articles/pin-a-dynamic-binary-instrumentation-tool

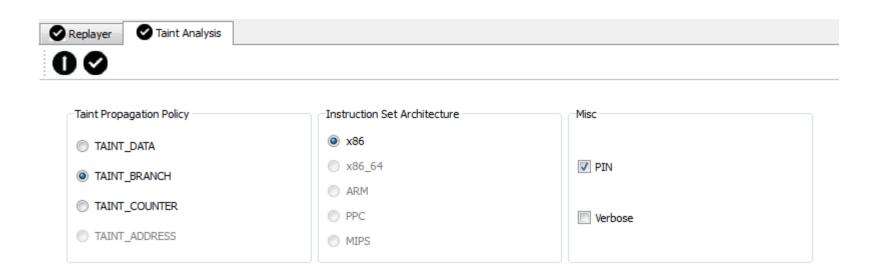


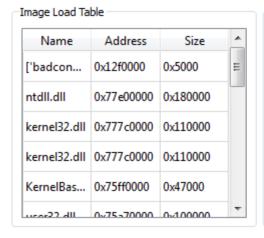
2 TREE Console: Trace Generation

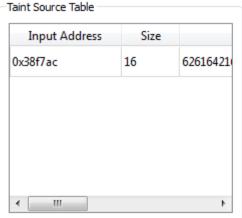
PINAgent: Connects TREE with PIN tracer

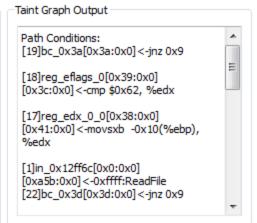


3 TREE: Taint Analysis Configuration

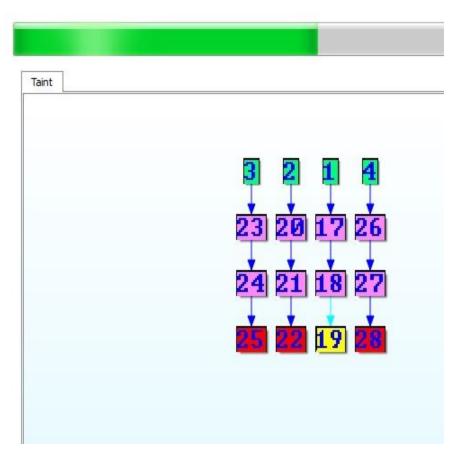






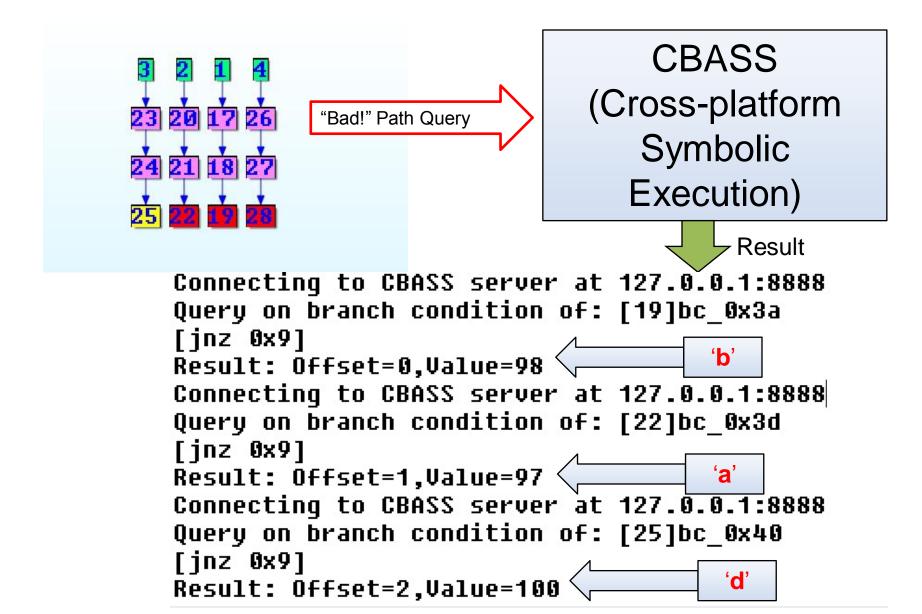


4 TREE: Branch Taint Graph



UÛID	Туре	Name	
1	input	0x12ff6c	
17	register	edx_0_0	
18	register	eflags 0	
19	branch	0x3a	
2	input	0x12ff6d	
20	register	ecx_0_0	
21	register	eflags_0	
22	branch	0x3d	
23	register	eax_0_0	
24	register	eflags_0	
25	branch	0x40	
26	register	edx_0_0	
27	register	eflags_0	
28	branch	0x43	
3	input	0x12ff6e	
4	input	0x12ff6f	

Solution 10 September 5 Negate Tainted Path Condition to Exercise a New ("Bad") Path



On-demand Symbolic Execution (What Happens Behind the Scene)

```
(set-logic QF_AUFBV)

(declare-fun _IN_0x12ff6c_0x0_SEQ0 () (_ BitVec 8))
(declare-fun EXPR_0 () (_ BitVec 32))
(assert (= EXPR_0 (bvsub ((_ sign_extend 24) (bvxor _IN_0x12ff6c_0x0_SEQ0 (_ bv128 8))) (_ bv4294967168 32))))

(assert (= (ite (not (= (ite (not (= (bvand ((_ extract 63 0) (bvsub ((_ sign_extend 32) (bvand ((_ extract 31 0) EXPR_0) (_ bv4294967295 32))) (_ bv98 64))) (_ bv4294967295 64)) (_ bv0 64))) (_ bv0 8)))

(check-sat)
(get-value (_IN_0x12ff6c_0x0_SEQ0))
```

```
RD t5, EMPTY , DWORD edx1]

58:481877 cmp edx, 98

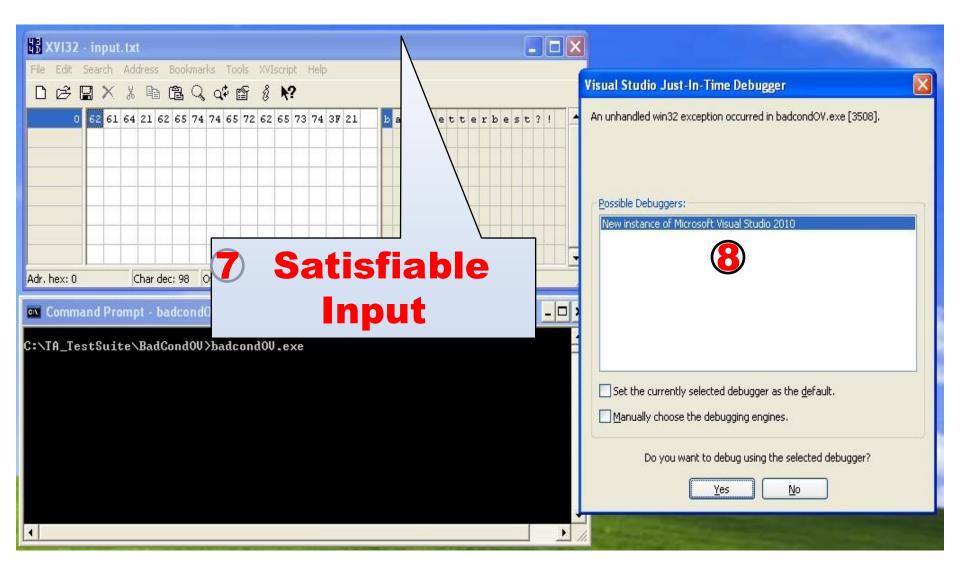
[48187708: and [DWORD edx, DWORD 2147483648, DWORD t8], 48187781: and [DWORD 98, DWORD 2147483648, DWORD t1], 48187782: sub [DWORD edx, DWORD 98, QWORD t2], 48187783: and [QWORD t2], QWORD 2147483648, DWORD t3], 48187784: bab [DWORD t3, DWORD -31, BYTE SF], 48187785: xor [DWORD t8, DWORD t1, DWORD t4], 48187786: xor [DWORD t0, DWORD t3, DWORD t5], 48187787: and [DWORD t4, DWORD t4], 48187786: xor [DWORD t0, DWORD t3, DWORD t5], 48187787: and [DWORD t4, DWORD t5], 48187788: and [QWORD t2], 48187788: bab [DWORD t2], 48187788: bab [DWORD t2], 48187788: bab [DWORD t2], 48187788: and [QWORD t2], 48187788: and [QWORD t2, QWORD 4294967295, DWORD t8], 4818778C: bisz [DWORD t8, EMPTY, BYTE ZF]]

59:48187a jnz loc_48188E

[481878 jnz loc_48188E

[
```

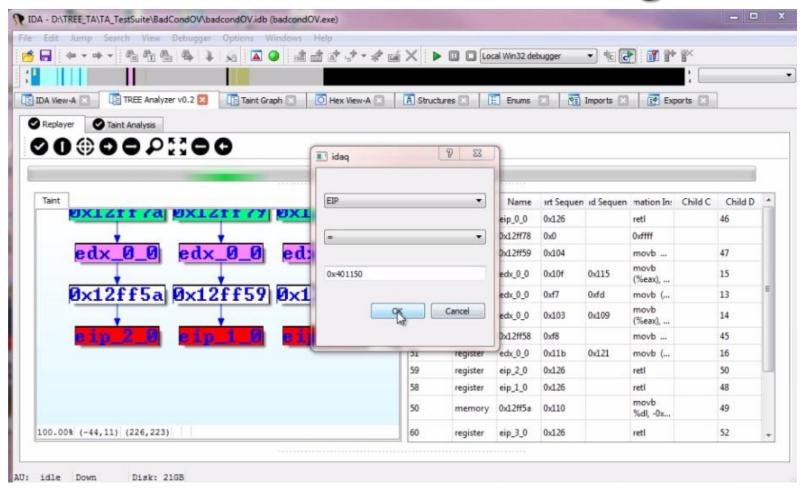
TREE: Re-execute with Satisfiable Input



Task 2: Own the Execution Assume Payload at 0x401150

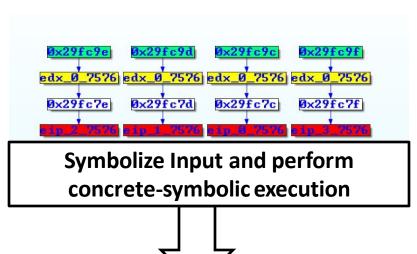
```
.text:00401145
                                 align 10h
.text:00401150
                                 push
                                         ebp
.text:00401151
                                 mov
                                         ebp, esp
                                         1010h
.text:00401153
                                 push
                                         offset aYouHaveBeenHac :
.text:00401158
                                 push
                                         offset aCbassCrossPlat :
.text:0040115D
                                 push
.text:00401162
                                 push
                                         ds:MessageBoxA
                                 call
.text:00401164
.text:0040116A
                                 push
                                         OFFFFFFFF
                                         ds exit
.text:0040116C
                                 call
+av+•00401177
```

TREE Constraint Dialogue



Task 2:

Own the Execution: From Crash to Exploit



Symbolic eip =

(= expr_0 (concat (bvand (bvor _IN_0x12ff6c_0xd_SEQ0 (_ bv0 8)) (_ bv255 8)) (bvand (bvor_IN_0x12ff6c_0xc_SEQ0 (_ bv0 8)) (_ bv255 8))))

Query:

get-value (_IN_0x12ff6c_0xd_SEQ0 _IN_0x12ff6c_0xc_SEQ0 _IN_0x12ff6c_0xe_SEQ0 _IN_0x12ff6c_0xf_SEQ0)

Sat:

(_IN_0x12ff6c_0xd_SEQ0#x11 _IN_0x12ff6c_0xc_SEQ0#x50 _IN_0x12ff6c_0xe_SEQ0#x40 _IN_0x12ff6c_0xf_SEQ0#x00



TREE/CBASS Demo

Using CBASS/TREE to Explore Bad Paths and Refine Exploits

Real World Case Studies

Target Vulnerability	Vulnerability Name	Target Application Mode	Target OS
CVE-2005-4560	Windows WMF	User Mode	Windows
CVE-2207-0038	ANI Vulnerability	User Mode	Windows
OSVDB-2939	AudioCoder Vulnerability	User Mode	Windows
CVE-2011-1985	Win32k Kernel Null Pointer De- reference	Kernel Mode	Windows
CVE-2004-0557	Sound eXchange (SoX) WAV Multiple Buffer Overflow	User Mode	Linux
Compression/ Decompression	Zip on Android	User Mode	Real Device Trace Generation (In Progress)

Highlights from Real World Case Study: Windows WMF Vulnerability (CVE-2005-4560)

- WMF SETABORTPROC Escape Vulnerability
 - http://www.cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2005-4560
 - The Windows Graphical Device Interface library (GDI32.DLL) in Microsoft Windows allows remote attackers to execute arbitrary code via a Windows Metafile (WMF) format image with a crafted SETABORTPROC GDI Escape function call, related to the Windows Picture and Fax Viewer (SHIMGVW.DLL).

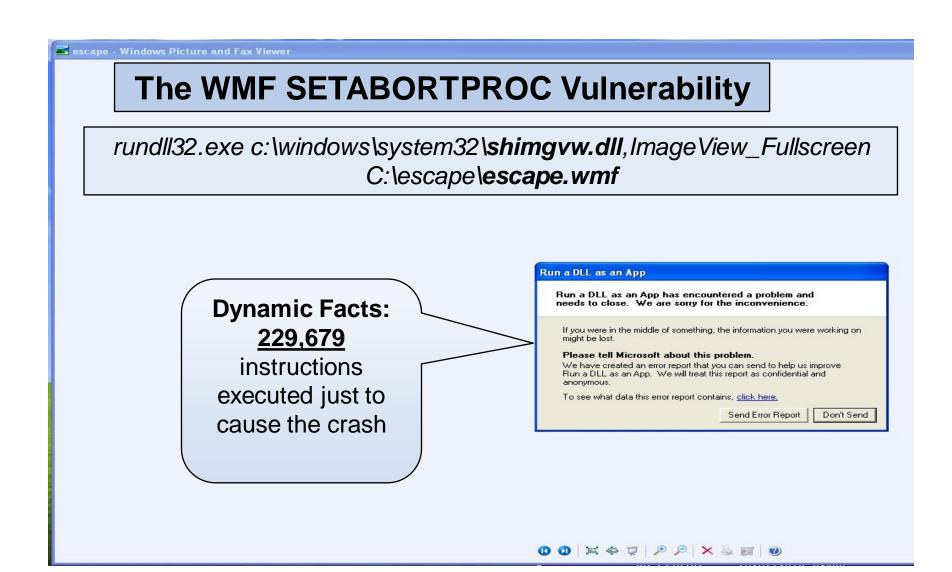
WMF Format

- [MS-WMF]: Windows Metafile Format
 - http://msdn.microsoft.com/en-us/library/cc250370.aspx
- A Simplified One:
 - http://wvware.sourceforge.net/caolan/ora-wmf.html
- Overall WMF File Structure:

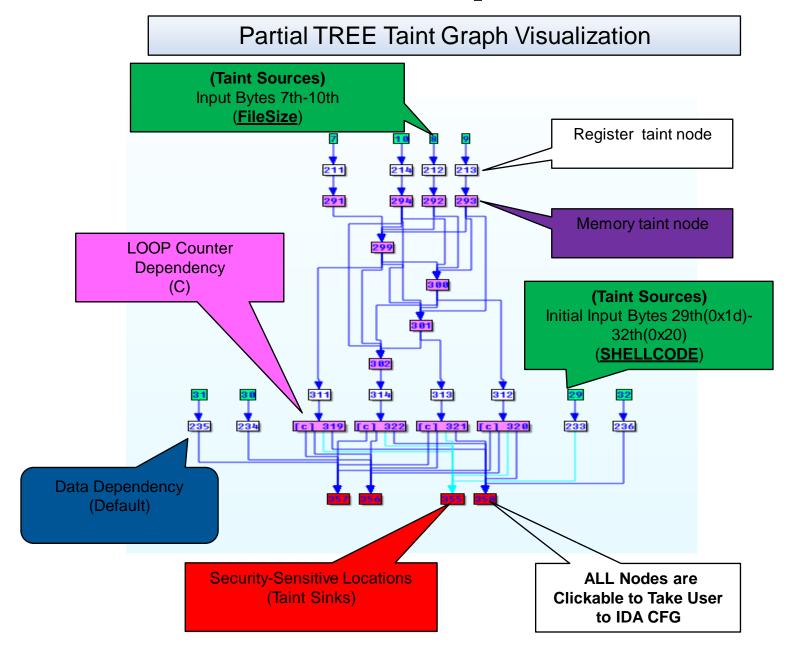
Meta	Meta	Meta	Meta	
Header	Record 1	Record 2	Record 3	

- One type of record is "escape" record
- SETABORTPROC escape allow an application to register a hook function to handle spooler errors

WMF Crash



WMF Taint Graph



WMF File: The Fields & The Vulnerability

Key Structures:

```
typedef struct WindowsMetaHeader
                                                              typedef struct _StandardMetaRecord
WORD FileType; /* Type of metafile (0=memory, 1=disk) */
                                                                 DWORD Size;
WORD HeaderSize: /* Size of header in WORDS (always 9) */
                                                                /* Total size of the record in WORDs */
WORD Version: /* Version of Microsoft Windows used */
                                                                 WORD Function:
                                                                 /* Function number (defined in WINDOWS.H)
DWORD FileSize; /* Total size of the metafile in WORDs */
WORD NumOfObjects; /* Number of objects in the file */
DWORD MaxRecordSize; /* The size of largest record in WORDs */
                                                                 WORD Parameters[]:
WORD NumOfParams; /* Not Used (always 0) */
                                                                /* Parameter values passed to function */
                                                              } WMFRECORD:
} WMFHEAD;
           XVI32 - escape_wmf.NODISPLAY
            File Edit Search Address Bookmarks Tools XVIscript Help
            0 01 00 09 00 00 03 22 00 00 00 63 79 61 6E 69 64 2D 45 01 04 00 00 00 06 26 09 CC CC CC
                 00|00|00|80|03|00|00|0
Shellcode
                                                                                            SetAbortProc
                                                                             Escape
                           Char dec: 0
                                    Overwrite
           Adr. hex: 43
```

WMF Slicing (1)

An Instruction Slice Traced Back from Crash Site to Input

Each node uniquely trace back to one execution event through its sequence number

0x77f330a3 <u>call eax</u> 2 ffd0 0x0 0x3812f Reg(EAX=0xa8b94 ESP=0xb4fb88 EIP=0x77f330a3) W 4 b4fb88

0x77c472e3 rep movsd 2 f3a5 0x0 0xb142 Reg(EDI=0xa8804 eflags=0x10216 ESI=0xa9f8c ECX=0xa) R 4 a9f8c cc_cc_cc W 4 a8804

0x77f2e997 <u>mov ecx, [ebp+arg 8]</u> 3 8b4d10 0x0 0xc5c3 Reg(EBP=0xb4fbf8 ECX=0x7c809a20) R 4 b4fc08 44_0_0

0x77f2e983 mov [ebp+arg 8], eax 3 894510 0x0 0xbd8c Reg(EAX=0x44 EBP=0xb4fbf8) W 4 b4fc08

0x77f2e97f add eax, eax 2 03c0 0x0 0xbd89 Reg(EAX=0x22 eflags=0x246)

0x77f2e949 mov eax, [edi+6] 3 8b4706 0x0 0xbd7d Reg(EAX=0xa8920 EDI=0xa87e8) R 4 a87ee 22_0_00

0x77c472e3 rep movsd 2 f3a5 0x0 0xb13c Reg(EDI=0xa87ec eflags=0x10216 ESI=0xa9f74 ECX=0x10) R 4 a9f74 0_3_22_0 W 4 a87ec

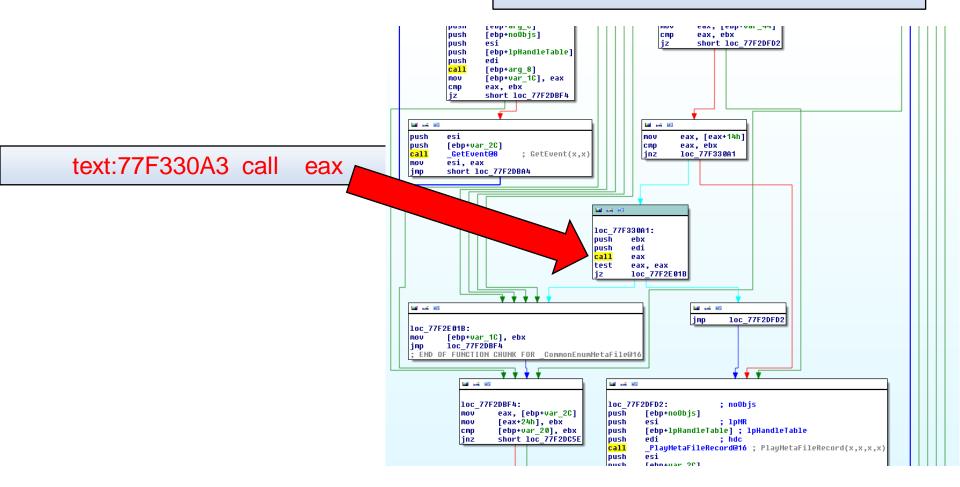
WMF Slicing (2)

An Instruction Slice with Text Helps

Put Instruction In Its Context Helps
More

Module: gdi32.dll

Function: CommonEnumMetaFile



WMF Slicing (3)

An Instruction Slice with Text Helps a Little

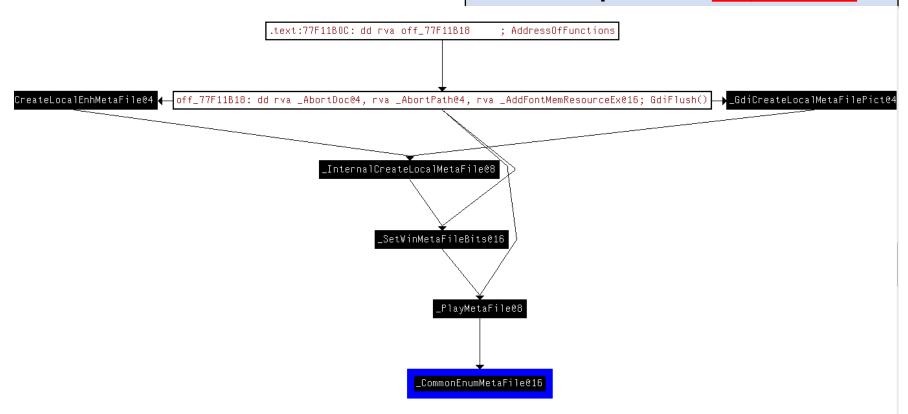
text:77F330A3 call eax

More Context Helps More

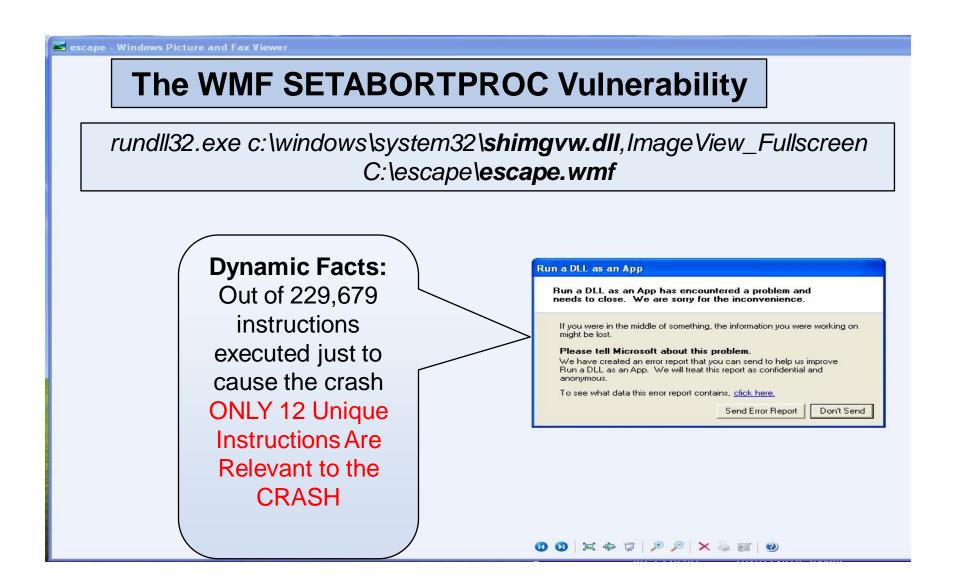
Module: gdi32.dll

Function: CommonEnumMetaFile

Call Graph: caller PlayMetaFile



WMF -- The Relevant Parts



Conclusions

- Our tools support interactive binary analysis, with Replay, Dynamic Taint Analysis, and Symbolic Execution.
- TREE runs on top of IDA Pro and supports cross-platform trace collection, taint analysis and replay.
- CBASS (based on REIL) enables IR-based architecture-independent symbolic execution and can support both automated and interactive analysis.
- YOU drive the tools!

Where You Can Get TREE

- TREE is open source at: http://code.google.com/p/tree-cbass/
 - First version of TREE (Taint Analysis) is released
 - Replay is in Progress
 - CBASS is Following
- Contacts:
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