# Terna Engineering College Computer Engineering Department Program: Sem VIII

**Course: Distributed Computing Lab (CSL802)** 

Faculty: Rohini Patil

# **Experiment No. 7**

**A.1 Aim:** To Implement Chandi–Misra-Haas distributed deadlock detection algorithm.

# PART B (PART B: TO BE COMPLETED BY STUDENTS)

Roll No. 50	Name: AMEY MAHENDRA THAKUR	
Class: BE COMPS B 50	Batch: B3	
Date of Experiment: 25-02-2022	Date of Submission: 25-02-2022	
Grade:		

# **B.1 Software Code written by a student:**

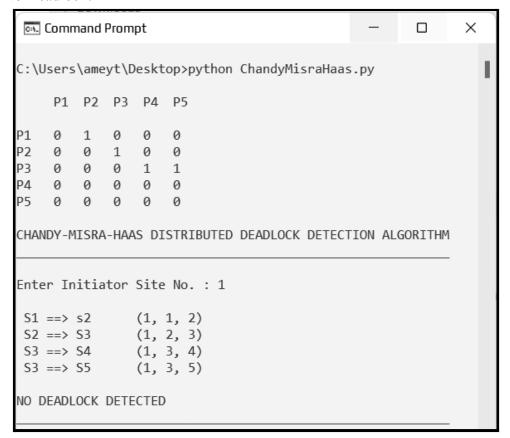
# • ChandyMisraHaas.java

No Deadlock:	Deadlock:
print("" P1 P2 P3 P4 P5	print("" P1 P2 P3 P4 P5
P1 0 1 0 0 0 P2 0 0 1 0 0 P3 0 0 0 1 1 P4 0 0 0 0 0 P5 0 0 0 0 0 "")	P1 0 1 0 0 0 P2 0 0 1 0 0 P3 0 0 0 1 1 P4 1 0 0 0 0 P5 0 0 0 0 0 0 0 0 0 0 0 0 0 0
a = [ [0, 1, 0, 0, 0], $[0, 0, 1, 0, 0],$ $[0, 0, 0, 1, 1],$ $[0, 0, 0, 0, 0],$ $[0, 0, 0, 0, 0]]$	a = [ [0, 1, 0, 0, 0], $[0, 0, 1, 0, 0],$ $[0, 0, 0, 1, 1],$ $[1, 0, 0, 0, 0],$ $[0, 0, 0, 0, 0]]$

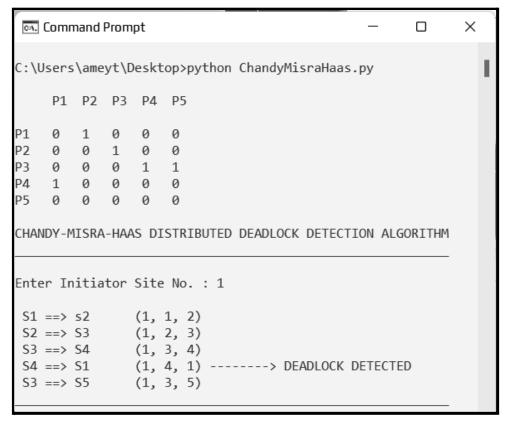
```
flag = 0
def aman(a, i, k):
      end = 5
      for x in range(end):
            if(a[k][x] == 1):
                  if(i == x):
                        print(f S\{k+1\} ==> S\{x+1\} \quad (\{i+1\}, \{k+1\}, \{x+1\}) ----->
                  DEADLOCK DETECTED')
        global flag
        flag = 1
        break
      print(f' S\{k+1\} => S\{x+1\} \quad (\{i+1\}, \{k+1\}, \{x+1\})')
      aman(a,i,x)
print("CHANDY-MISRA-HAAS DISTRIBUTED
                                                DEADLOCK
                                                               DETECTION
ALGORITHM")
print("
                                                               ")
print()
x = 0
end = 5
i = int(input("Enter Initiator Site No.:"))
j = i - 1
print()
for k in range(end):
  if(a[j][k]==1):
    print(f' S\{j+1\} ==> s\{k+1\} \qquad (\{i\}, \{j+1\}, \{k+1\})')
    aman(a,j,k)
if(flag == 0):
  print("\nNO DEADLOCK DETECTED")
print("____
```

## **B.2 Input and Output:**

• No Deadlock:



#### • Deadlock:



#### **B.3 Observations and learning:**

- This is considered an edge-chasing, probe-based algorithm.
- It is also considered one of the best deadlock detection algorithms for distributed systems.
- If a process makes a request for a resource that fails or times out, the process generates a probe message and sends it to each of the processes holding one or more of its requested resources.
- Each probe message contains the following information:
  - the id of the process that is blocked (the one that initiates the probe message);
  - the id of the process is sending this particular version of the probe message;
  - the id of the process that should receive this probe message.
- When a process receives a probe message, it checks to see if it is also waiting for resources.
- If not, it is currently using the needed resource and will eventually finish and release the resource.
- If it is waiting for resources, it passes on the probe message to all processes it knows to be holding resources it has itself requested.
- The process first modifies the probe message, changing the sender and receiver ids.
- If a process receives a probe message that it recognizes as having initiated, it knows there is a cycle in the system and thus, deadlock.
- The advantages of this algorithm include the following:
  - It is easy to implement.
  - Each probe message is of fixed length.
  - There is very little computation.
  - There is very little overhead.
  - There is no need to construct a graph, nor to pass graph information to other sites.
  - This algorithm does not find false (phantom) deadlock.
  - There is no need for special data structures.

#### **B.4 Conclusion:**

Successfully implemented Chandy-Misra-Haas distributed deadlock detection algorithm using python.

#### **B.5 Question of Curiosity:**

Q1: Consider the following statements: A deadlock detection algorithm must satisfy the following two conditions:

**Condition 1**: Progress (No false deadlocks): The algorithm should not report deadlocks that do not exist.

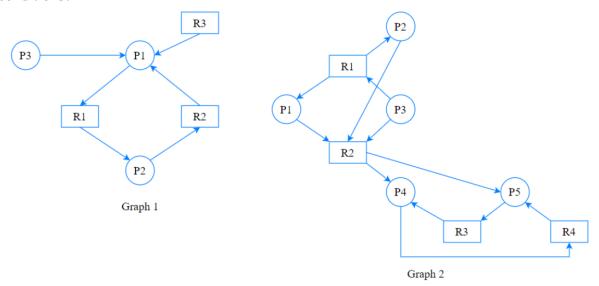
**Condition 2**: Safety (No undetected deadlocks): The algorithm must detect all existing deadlocks infinite time.

#### A. Both conditions are true

- B. Both conditions are false
- C. Only condition 1 is true
- D. Only condition 1 is true

ANS: A. Both conditions are true

Q2: Consider the resource allocation graphs G1 and G2. What can be said about the deadlock conditions?



## A. Both are deadlocked

- B. G1 deadlocked, G2 not
- C. G1 not, G2 deadlocked
- D. None is deadlocked

ANS: A. Both are deadlocked

Q3: Which of the following problem we might face if we invoke the deadlock detection algorithm at the arbitrary interval?

- A. There may be many cycles in the graph and it will not be possible to roll back the deadlocked processes anymore.
- B. There may be many cycles in the graph and we would not be able to tell which of the many deadlocked processes caused the deadlock.
- C. There may be no cycles in the graph and thus it will not be possible to know if any deadlock has happened.
- D. All of the above.

ANS: B. There may be many cycles in the graph and we would not be able to tell which of the many deadlocked processes caused the deadlock

Q4: Consider the following statements related to distributed deadlock detection algorithms:

**Statement 1:** In path-pushing algorithms, distributed deadlocks are detected by maintaining an explicit global WFG. The basic idea is to build a global WFG for each site of the distributed system.

**Statement 2:** In an edge-chasing algorithm, the presence of a cycle in a distributed graph structure is be verified by propagating special messages called probes, along the edges of the graph. These probe messages are different from the request and reply messages.

## A. Both statements are true

- B. Both statements are false
- C. Only statement 1 is true
- D. Only statement 2 is true

ANS: A. Both statements are true