	pef	Threads Process		
Instant Ch. L.	Lest	F - S		
Current State -Pointer	Event info contraction wouting process much		Process States - new state -> process creat	ion.
13/11/2	Swing	to F S	- new state ⇒ process create - ready state → ready home - executing state → cpo - waiting shote	but was tring (
Process ID	-GPR-IR-SP-ACC.	Uses level - implemented by Uncu	Kernel level - implemented by O.S.	Process Context switch
Priority (Scheduling queue)	A counting - wer & Kernel CPU time consumed	- Os doan't recognite - early - CS - leaf	- 0.5 Recognize - complex - CS - more	Prais Skates, Discipaten.
PC	Menery allocation.	- No horolware Folick - whole procur blocker	- Hardware	PCB Operations on Process
O peration on Process	Process termination.	h	Report Stack 30500 -	Threads, types of Threads.
Process ocation - system introducation	EXIT	enterno Heavy	Lyar 3 SSS	Multi threading Multiore processing
-IRPEPCSC -AURICNE -Starting & Batch Tob	3,555-455	S-> Dodgo	Stan Stan Regard - Menory	Amdahl'claw Schedulers, types of schedules.
One to One i more concurry only a parallel kernel - war -> Apr degrande	AMDAHL'S LAW time required to excurse a pg.	Aney Dis	Dynam Multiwan	
)))-m	Miles - time required to execute a 19. On 10. Il process	reation Termination	TO T	
Many to one	(1-F) + F/N	Schooling Long.	Madium	(hat
-managent done in war appare	fract of pg. frag.	Ly loads or memory for exec.	- swap in a out from mounted - speed is in between no state	from R. Q :::::: 2 paringn it to cau - speed V. Jak
Rick	7) MN Apr - many engle les gharing 3) MI Apr. VR.	- new ready	MIN.	- ready to exente state Upo school
AboM of ult >		7. · · · ·		
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Uniprocessor Scheduling	Scheduling Algorithms	Numaricaly Deposity Shed (NP)	,
-improve perf. by keeping	i) first in first Out (FCFO) > NI	DY AT B 10 10 19	, !
OU bury all the time	- allocater the CPO to process	Proces 10 0 3 Pr P3 P4	
	in order of their arrival	P 0 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
DCPU utilization	- Short job (ghas wait until execution	P2 2 0 3 Process P FT TAI WI	
2) Throughput	of long job (1)	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	
	· 2) Shortest Job First (SJF) → NP/P	T FCPS 16 16	
4) Waiting time	conducing the jobs so as to run	P1 P2 P3 P5 P5 GONDAN P4 4 19 19 19 (8	4
5) Response time	SJE > improve tapes	R 2 6 6 1	
6) Fairness	- minimize the average waiting time	Process P FT TAT WI	
Non preemptive	•3) Priority Scheduling	P ₁ 3 0 0	T
- once process is allocated to CPU	- each process has priority - int	12 13 13 14 15 14 15 14 15 14 16 17 17 18 14 16 16 16 16 16 16 16 16 16 16 16 16 16	13
it does not free CPU runtil it	value a singued to it - Smallert int - highest P.	P, 3 13 13 17 Procon P FT TAT WT	
completes its execution	- largest int → lowest P.	P ₁₁ 4 19 10 11 10 3 19 19 19	
Picemptive	-4) Round Robin Scheduling.	Pr 2 (9 1 2 2 1 5	
- Il authors taking away are	- time sharing system	(1) SJF(NP)	
from process during execution	- many processes get CPU on time	P3 P4 P3 P5 P1 P5 2 14 14 9	
Thread Scheduling. 1) User level	shaving baris	INIT (P.)	المالة
2) Kund Level	_ squallest unit time -> quantum	P_{TOLOM} P FT TAT Wi = FR-AF = $(0-0)+(5-1)+(8-6)+(10-9)+(12-11)+$ P_{1} 3 19 19 9	יייי
Multiprocessor Scheduling.	• 5) Mullilevel queue Scheduling.	ρ ₂ 1 1 1 0 = 9	1.
	lategorites the processes into \$.	73 3 7 7	
Process Scheduling	groups.	R5 2 9 9 4.	\perp
-if 1> processor present then echeduling approach.	- Separation between interactive process a batch process.		
then scheduling approach. (6) less imp.	6) Multi level Feedback-Queue Scheduling.	0 1 2 4 9 19 P2 P4 13 P5 P1	201
Ways to und thread Scheduling	- process are not allowed to	Prous P FT TAT WI	100
i) load Sharing	transfer from one queue to another	P1 3' 19 19 9	1
2) Gang Scheduling	- queue -> permanent & cannol be changed - low scheduling cost	, ,	
3) Dedicated Process Assignment 4) Dynamic Scheduling	- not Rexible		
Linux Scheduling > RT T	- ·	1	
Lillon const	3'		
		- Feb	