

<p>File Management</p> <p>File</p> <ul style="list-style-type: none"> - A file is a named collection of related information that is recorded on secondary storage such as magnetic disk, magnetic tape. - File is a sequence of bits, bytes, lines or records whose meaning is defined by File Owner & User <p>Extensions → Types</p> <ul style="list-style-type: none"> GIF → Image HTML → Web page OBJ, O → Object EXE, BIN → Executable TXT, DOC → Text LIB → Library ZIP → Archive PS → Post Script MP3, MP4 → Multimedia PDF → Portable Document Format <p><i>Handwritten note: mega, mega, kabit, kabit</i></p>	<p>File attribute</p> <p>Attribute</p> <ul style="list-style-type: none"> - Protection - Size - Type - Creator - Owner - Password <p>Identifiers</p> <ul style="list-style-type: none"> - Hidden Flag - Read only Flag - ASCII / Binary Flag - Lock Flag - Key length - Time of last access - Time of last change - Maximum Size - System Flags 	<p>File Organization</p> <p>Pile</p> <ul style="list-style-type: none"> - simplest form of file organization - File data is stored in order of arrival - each record in the file comprises one burst of data - idea behind pile is simply to collect mass of data & store it <p>Sequential File</p> <ul style="list-style-type: none"> - most commonly used structure & records <p>Indexed Sequential File</p> <ul style="list-style-type: none"> - Each and every record in file has same length with equal no. of fixed length field in a particular order - as length & location of each field are known, simply the values of fields are required to be stored - it eliminates drawback of sequential file - it contains index to the file to support random access & an overflow file - offers searching ability to search speeding the required record 	<p>File Directories</p> <p>Single level directory system</p> <ul style="list-style-type: none"> - The simplest method is to have one big list of all the files on the disk - The entire system will contain only one directory which is supposed to mention all the files present in file system - The directory contains one entry per each file present in file system - Adv → Simple implementation - Disadv → We cannot have 2 files with same name → Heuristic Searching <p>Two level directory system</p> <ul style="list-style-type: none"> - Separate directory for each user - 1 master directory which contains separate directories dedicated to each user - System doesn't let a user to enter in the other user's directory without permission <p>Tree/Intrinsics</p> <ul style="list-style-type: none"> - Each file has a pathname as → user name / directory name / filename - Different user can have same file name - Efficient searching - Adv → Solve name collision pb - Indpendent user get isolated from each other <p>Hierarchical directory system</p> <ul style="list-style-type: none"> - It is quite general and advantageous for users to group their files together in logical ways - Adv → With this system, in addⁿ to their files, users can access the files of other user by pathname <p>Path names</p> <ul style="list-style-type: none"> → Absolute Path name → Relative Path name <p>Free Space Management</p> <ul style="list-style-type: none"> - A file system is responsible to allocate the free blocks to the file therefore it has to keep track of all the free blocks present in the disk <p>Methods of free spaces</p> <ul style="list-style-type: none"> → Bit Vector - In this approach, the free space is implemented as a bit map vector. It contains the no. of bits where each bit represents each block - If block is empty → Bit is 1 otherwise zero - Initially all blocks are empty <p>Linked List</p> <ul style="list-style-type: none"> - This approach suggests linking together all the free blocks and keeping a pointer in the cache which points to the first free block - All free blocks of list → linked together → pointer <p>Preallocation and dynamic allocation</p> <ul style="list-style-type: none"> - Preallocation policy requires that the max size of file declared at a time of file creation request - For many app^s it is difficult to estimate reliably the max potential size of file - Dynamic allocates a space to the file portions as needed 	<p>Contiguous Allocation</p> <ul style="list-style-type: none"> - Each file occupies a contiguous address space on disk - Assigned disk address is in linear order - Easy to implement - External fragmentation is a major issue <p>Linked list allocation</p> <ul style="list-style-type: none"> - Each file carries list of links to disk blocks - Directory contain link/pointer to first block of file - No external fragmentation <p>Indexed allocation</p> <ul style="list-style-type: none"> - Effectively used in sequential address - Inefficient in use of direct access <p>Indirect allocation</p> <ul style="list-style-type: none"> - Provides solⁿ to problem of contiguous & linked allocation - Index block is created having all pointers of file - Each file has its own index block which stores the address of disk space occupied by file - I-Node (Index-Node) - All types of Unix files are administered by OS by means of index nodes <p>Linux Virtual File System</p> <ul style="list-style-type: none"> - The linux virtual file system or VFS generally is a layer that sits on the top of your actual file system which allows the user to access diff types of file system - Linux of VFS is an interface betⁿ kernel and actual file system - Object oriented principles are used to design VFS <p>4 major object types by VFS</p> <ul style="list-style-type: none"> → Superblock → dentries → Inodes → Files <p>A set of operations are defined for each of the type of objects</p> <ul style="list-style-type: none"> - Each object of one of these types points to 6ⁿ table - The record at addresses of the actual 6ⁿ are kept in 6ⁿ table
<p>File Structure</p> <ul style="list-style-type: none"> - Unstructured Sequence of Bytes - OS treats file as sequence of bytes - Any meaning must be imposed by user level prgm - Both UNIX & WINDOWS use this approach - It provides more flexibility - Any type of content can be kept in file as per convenience - Sequence of fixed length record - File is treated as sequence of fixed length record - Each record has its internal structure - Any read opⁿ on file returns one record and any write opⁿ will append or overwrite one record <p>Tree Structure</p> <ul style="list-style-type: none"> - Files are organized in tree structure - All records not necessarily have same length - Each record has a key field in fixed position - Records in tree are arranged in order to key field, to permit quick searching for a particular key - Main aim is to search record on particular key <p>File types</p> <p>Regular Files</p> <ul style="list-style-type: none"> - It contains user info - The byte sequence, record sequence & tree structure are examples of regular files <p>Directories</p> <ul style="list-style-type: none"> - These are system files for maintaining the structure of file system <p>Character special files</p> <ul style="list-style-type: none"> - related to I/O device such as terminal, printer, serial I/O device <p>Block Special Files</p> <ul style="list-style-type: none"> - These are used to model disks <p>Regular files are normally either ASCII or binary files</p> <ul style="list-style-type: none"> - ASCII file contains lines of text, each line is terminated by either carriage return character or line feed character 	<p>File operations</p> <ul style="list-style-type: none"> - Create - Delete - Open - Close - Read - Write - Append - Seek - Set attribute - Set attribute - Rename <p>File Access</p> <ul style="list-style-type: none"> - Sequential access - Records are accessed in some sequence - info in file is processed in order, one record after the other - such method is the most primitive one - eg: simple usually access files in fashion - direct or random access - It provides accessing the records directly - each record has its own address on the file with by help of which, it can be directly accessed for reading or writing - Records need not be in any sequence with Multiple & need not be in adjacent location on storage medium <p>Indexed Sequential Access</p> <ul style="list-style-type: none"> - This mechanism is built up on basis of sequential access - An index is created for each file which contains pointers to various block - Index is searched sequentially & its pointer is used to access the file directly <p>Directory Operations</p> <ul style="list-style-type: none"> - Create - Delete - Openness - Close dir - Read dir - Rename - Link - Unlink 	<p>Indexed File</p> <ul style="list-style-type: none"> - The sequential file & Indexed sequential file can search the records based on single field of file - Apart from using key field, they cannot search record using other attribute of record - To overcome this difficulty and accomplish the flexibility, a structure containing multiple indexes, one for each type of field is required - Records are accessed only through indexes <p>The Direct or Method File</p> <ul style="list-style-type: none"> - makes use of ability found on disks to access straight any block of a known address - hashing opⁿ on key value using hash function - direct files are frequently used where very fast access is necessary, where there is use of fixed length records and where records are always accessed one by one <p>File Sharing</p> <ul style="list-style-type: none"> - Necessary to share a file among multiple users - Needs to deal with 2 issues <p>① Access Rights</p> <ul style="list-style-type: none"> - None - Knowledge - Execution - Reading - Appending - Updating - Changing protection - Deletion <p>② Simultaneous Access</p> <ul style="list-style-type: none"> - When more than 1 user granted access to open or update a file, the OS or file management system must implement some way to restrict it <p>File Allocation</p> <ul style="list-style-type: none"> - On secondary storage, a file consists of collection of blocks - OS or file management system is responsible for allocating blocks of files - Approach taken for file allocation may influence the approach taken for free space management - Space is allocated to a file as one or more portions 	<p>Free Space Management</p> <ul style="list-style-type: none"> - A file system is responsible to allocate the free blocks to the file therefore it has to keep track of all the free blocks present in the disk <p>Methods of free spaces</p> <ul style="list-style-type: none"> → Bit Vector - In this approach, the free space is implemented as a bit map vector. It contains the no. of bits where each bit represents each block - If block is empty → Bit is 1 otherwise zero - Initially all blocks are empty <p>Linked List</p> <ul style="list-style-type: none"> - This approach suggests linking together all the free blocks and keeping a pointer in the cache which points to the first free block - All free blocks of list → linked together → pointer <p>Preallocation and dynamic allocation</p> <ul style="list-style-type: none"> - Preallocation policy requires that the max size of file declared at a time of file creation request - For many app^s it is difficult to estimate reliably the max potential size of file - Dynamic allocates a space to the file portions as needed <p>Portion Size</p> <ul style="list-style-type: none"> - Choosing a portion size, there is a trade off betⁿ efficiency from the point of view of a single file vs overall system efficiency 	<p>Contiguous Allocation</p> <ul style="list-style-type: none"> - Each file occupies a contiguous address space on disk - Assigned disk address is in linear order - Easy to implement - External fragmentation is a major issue <p>Linked list allocation</p> <ul style="list-style-type: none"> - Each file carries list of links to disk blocks - Directory contain link/pointer to first block of file - No external fragmentation <p>Indexed allocation</p> <ul style="list-style-type: none"> - Effectively used in sequential address - Inefficient in use of direct access <p>Indirect allocation</p> <ul style="list-style-type: none"> - Provides solⁿ to problem of contiguous & linked allocation - Index block is created having all pointers of file - Each file has its own index block which stores the address of disk space occupied by file - I-Node (Index-Node) - All types of Unix files are administered by OS by means of index nodes <p>Linux Virtual File System</p> <ul style="list-style-type: none"> - The linux virtual file system or VFS generally is a layer that sits on the top of your actual file system which allows the user to access diff types of file system - Linux of VFS is an interface betⁿ kernel and actual file system - Object oriented principles are used to design VFS <p>4 major object types by VFS</p> <ul style="list-style-type: none"> → Superblock → dentries → Inodes → Files <p>A set of operations are defined for each of the type of objects</p> <ul style="list-style-type: none"> - Each object of one of these types points to 6ⁿ table - The record at addresses of the actual 6ⁿ are kept in 6ⁿ table