

Solution: f -Divergence and Hypothesis Testing

Problem 4: f -Divergence and Hypothesis Testing

We address each part in detail, including all intermediate steps.

1. Convexity of $f_\alpha(x)$

Definition 1 (Convex Function). *A twice-differentiable function $g : (0, \infty) \rightarrow \mathbb{R}$ is convex if and only if $g''(x) \geq 0$ for all $x > 0$.*

Let $\alpha > 0$, $\alpha \neq 1$, and define

$$f_\alpha(x) = \frac{x^\alpha - 1}{\alpha - 1}, \quad x > 0.$$

Compute the first and second derivatives explicitly:

$$\begin{aligned} f'_\alpha(x) &= \frac{d}{dx} \left(\frac{x^\alpha - 1}{\alpha - 1} \right) = \frac{\alpha x^{\alpha-1}}{\alpha - 1}, \\ f''_\alpha(x) &= \frac{d}{dx} \left(\frac{\alpha x^{\alpha-1}}{\alpha - 1} \right) = \frac{\alpha(\alpha - 1)x^{\alpha-2}}{\alpha - 1} = \alpha x^{\alpha-2}. \end{aligned}$$

Since $\alpha > 0$ and $x^{\alpha-2} > 0$ for $x > 0$, we have $f''_\alpha(x) > 0$. By definition, f_α is strictly convex.

2. Definition of the f -divergence

Let $P = (p_1, \dots, p_n)$ and $Q = (q_1, \dots, q_n)$ be two probability distributions on a finite set $\Omega = \{1, 2, \dots, n\}$, with $q_i > 0$ for all i .

Definition 2 (f -Divergence). *Let $f : (0, \infty) \rightarrow \mathbb{R}$ be convex with $f(1) = 0$. The f -divergence of P relative to Q is*

$$D_f(P||Q) = \sum_{i=1}^n q_i f\left(\frac{p_i}{q_i}\right).$$

In particular, for $f = f_\alpha$, we recover the α -divergence:

$$D_{f_\alpha}(P||Q) = \sum_{i=1}^n q_i \frac{(p_i/q_i)^\alpha - 1}{\alpha - 1}.$$

3. Data-Processing Inequality for f -Divergence

Suppose we apply any randomized transformation T to the sample space, inducing a Markov kernel $T(j \mid i)$ with output alphabet Γ . Then the output distributions are

$$P_T(j) = \sum_{i=1}^n p_i T(j \mid i), \quad Q_T(j) = \sum_{i=1}^n q_i T(j \mid i).$$

Theorem 1 (Data-Processing Inequality). *For any convex f and any Markov kernel T ,*

$$D_f(P_T \| Q_T) \leq D_f(P \| Q).$$

Proof. Write

$$D_f(P_T \| Q_T) = \sum_{j \in \Gamma} Q_T(j) f\left(\frac{P_T(j)}{Q_T(j)}\right).$$

For each j , define weights $w_i = T(j \mid i)q_i / Q_T(j)$ (so $\sum_i w_i = 1$). Then

$$\frac{P_T(j)}{Q_T(j)} = \frac{\sum_i p_i T(j \mid i)}{Q_T(j)} = \sum_i w_i \frac{p_i}{q_i}.$$

By convexity of f ,

$$f\left(\sum_i w_i \frac{p_i}{q_i}\right) \leq \sum_i w_i f\left(\frac{p_i}{q_i}\right).$$

Multiplying both sides by $Q_T(j)$ and summing over j gives

$$D_f(P_T \| Q_T) \leq \sum_j \sum_i T(j \mid i) q_i f\left(\frac{p_i}{q_i}\right) = \sum_i q_i f\left(\frac{p_i}{q_i}\right) = D_f(P \| Q).$$

□

4. Data-Processing for Rényi α -Divergence

Define the Rényi α -divergence for $\alpha > 0$, $\alpha \neq 1$, by

$$R_\alpha(P \| Q) = \frac{1}{1 - \alpha} \ln\left(1 + (\alpha - 1)D_{f_\alpha}(P \| Q)\right).$$

Note that the function $g(u) = \frac{1}{1 - \alpha} \ln(1 + (\alpha - 1)u)$ is strictly increasing in u since its derivative is positive. Hence, applying the monotone mapping g to both sides of the data-processing inequality for f_α yields

$$R_\alpha(P_T \| Q_T) = g(D_{f_\alpha}(P_T \| Q_T)) \leq g(D_{f_\alpha}(P \| Q)) = R_\alpha(P \| Q).$$

Thus Rényi divergence also satisfies data-processing.

5. Limits as $\alpha \rightarrow 1$

We show

$$\lim_{\alpha \rightarrow 1} D_{f_\alpha}(P\|Q) = \lim_{\alpha \rightarrow 1} R_\alpha(P\|Q) = D_{\text{KL}}(P\|Q).$$

Proof. Using the Taylor expansion for x^α around $\alpha = 1$:

$$x^\alpha = e^{\alpha \ln x} = 1 + (\alpha - 1) \ln x + o(\alpha - 1),$$

it follows that

$$f_\alpha(x) = \frac{x^\alpha - 1}{\alpha - 1} = \ln x + o(1).$$

Therefore,

$$D_{f_\alpha}(P\|Q) = \sum_i q_i (\ln(p_i/q_i) + o(1)) = D_{\text{KL}}(P\|Q) + o(1).$$

Next, since $D_{f_\alpha} \rightarrow D_{\text{KL}}$, we have

$$R_\alpha(P\|Q) = \frac{\ln(1 + (\alpha - 1)D_{f_\alpha})}{1 - \alpha} = \frac{(\alpha - 1)D_{f_\alpha} + o(\alpha - 1)}{1 - \alpha} = D_{f_\alpha} + o(1) \rightarrow D_{\text{KL}}(P\|Q).$$

□

6. Hypothesis Testing with Rényi Divergence

Let $X^n = (X_1, \dots, X_n)$ be i.i.d. with distribution either Q_X under H_0 or P_X under H_1 . A stochastic test outputs a binary decision $Z \in \{0, 1\}$ with

$$P_{H_1}(Z = 0) = \alpha_n < \epsilon, \quad P_{H_0}(Z = 1) = \beta_n.$$

Then P -error (α_n) and Q -error (β_n) satisfy by data-processing:

$$R_\alpha(P_X^{\otimes n} \| Q_X^{\otimes n}) \geq R_\alpha(\text{Bern}(\alpha_n) \| \text{Bern}(\beta_n)).$$

Since Rényi divergence is additive over product distributions,

$$R_\alpha(P_X^{\otimes n} \| Q_X^{\otimes n}) = n R_\alpha(P_X \| Q_X).$$

Hence

$$n R_\alpha(P_X \| Q_X) \geq R_\alpha(\text{Bern}(\alpha_n) \| \text{Bern}(\beta_n)).$$

7. Bounding Type-II Exponent

Now let $\alpha_n < \epsilon = 1/3$, and set $\alpha = 1 + 1/\sqrt{n}$. We use the bound

$$R_\alpha(\text{Bern}(\epsilon) \| \text{Bern}(1 - \beta_n)) \approx (1 - \alpha) \ln(1 - \beta_n) - h(\epsilon)$$

for large n , where $h(\epsilon)$ is the binary entropy. Rearranging yields

$$\ln(1 - \beta_n) \leq -\frac{1}{\alpha - 1} (n R_\alpha(P_X \| Q_X) + h(\epsilon)) \approx -\sqrt{n} R_\alpha(P_X \| Q_X),$$

implying

$$\beta_n \leq e^{-\sqrt{n} R_\alpha(P_X \| Q_X)}.$$

Thus for $\beta_n = e^{-n E_n}$, we have

$$E_n \leq R_{1+1/\sqrt{n}}(P_X \| Q_X).$$

8. Conclusion: KL Bound on Error Exponent

Taking $n \rightarrow \infty$, since $\alpha = 1 + 1/\sqrt{n} \rightarrow 1$, and using

$$\lim_{n \rightarrow \infty} R_{1+1/\sqrt{n}}(P_X \| Q_X) = D_{\text{KL}}(P_X \| Q_X),$$

we obtain

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \ln \frac{1}{\beta_n} \leq D_{\text{KL}}(P_X \| Q_X).$$

This completes the detailed proof.