

"""

Binary Search Tree

The Binary Search Tree represents an ordered symbol table of generic key-value pairs. Keys must be comparable. Does not permit duplicate keys.

When associating a value with a key already present in the BST, the previous value is replaced by the new one. This implementation is for an unbalanced BST.

Pseudo Code: <http://algs4.cs.princeton.edu/32bst>

"""

```
class Node(object):
```

```
    """
```

```
    Implementation of a Node in a Binary Search Tree.
```

```
    """
```

```
    def __init__(self, key=None, val=None,
size_of_subtree=1):
        self.key = key
        self.val = val
        self.size_of_subtree = size_of_subtree
        self.left = None
        self.right = None
```

```
class BinarySearchTree(object):
```

```
    """
```

```
    Implementation of a Binary Search Tree.
```

```
    """
```

```
    def __init__(self):
```

```

self.root = None

def _size(self, node):
    if node is None:
        return 0
    else:
        return node.size_of_subtree

def size(self):
    """
    Return the number of nodes in the BST

    Worst Case Complexity: O(1)

    Balanced Tree Complexity: O(1)
    """
    return self._size(self.root)

def is_empty(self):
    """
    Returns True if the BST is empty, False
    otherwise

    Worst Case Complexity: O(1)

    Balanced Tree Complexity: O(1)
    """
    return self.size() == 0

def _get(self, key, node):
    if node is None:
        return None

    if key < node.key:
        return self._get(key, node.left)
    elif key > node.key:
        return self._get(key, node.right)
    else:
        return node.val

```

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def get(self, key):
    """
    Return the value paired with 'key'

    Worst Case Complexity:  $O(N)$ 

    Balanced Tree Complexity:  $O(\lg N)$ 
    """
    return self._get(key, self.root)

def contains(self, key):
    """
    Returns True if the BST contains 'key', False
    otherwise

    Worst Case Complexity:  $O(N)$ 

    Balanced Tree Complexity:  $O(\lg N)$ 
    """
    return self.get(key) is not None

def _put(self, key, val, node):
    # If we hit the end of a branch, create a new
    node
    if node is None:
        return Node(key, val)

    # Follow left branch
    if key < node.key:
        node.left = self._put(key, val,
node.left)
    # Follow right branch
    elif key > node.key:
        node.right = self._put(key, val,
node.right)
    # Overwrite value
    else:

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        node.val = val

        node.size_of_subtree = self._size(node.left)
+ self._size(node.right)+1
        return node

def put(self, key, val):
    """
    Add a new key-value pair.

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    self.root = self._put(key, val, self.root)

def _min_node(self):
    """
    Return the node with the minimum key in the
BST
    """
    min_node = self.root
    # Return none if empty BST
    if min_node is None:
        return None

    while min_node.left is not None:
        min_node = min_node.left

    return min_node

def min_key(self):
    """
    Return the minimum key in the BST

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """

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        min_node = self._min_node()
        if min_node is None:
            return None
        else:
            return min_node.key

def _max_node(self):
    """
    Return the node with the maximum key in the
BST
    """
    max_node = self.root
    # Return none if empty BST
    if max_node is None:
        return None

    while max_node.right is not None:
        max_node = max_node.right

    return max_node

def max_key(self):
    """
    Return the maximum key in the BST

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    max_node = self._max_node()
    if max_node is None:
        return None
    else:
        return max_node.key

def _floor_node(self, key, node):
    """
    Returns the node with the biggest key that is
less than or equal to the

```

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given value 'key'
"""
if node is None:
    return None

if key < node.key:
    # Floor must be in left subtree
    return self._floor_node(key, node.left)

elif key > node.key:
    # Floor is either in right subtree or is
this node
    attempt_in_right = self._floor_node(key,
node.right)
    if attempt_in_right is None:
        return node
    else:
        return attempt_in_right

else:
    # Keys are equal so floor is node with
this key
    return node

def floor_key(self, key):
    """
    Returns the biggest key that is less than or
equal to the given value
    'key'

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    floor_node = self._floor_node(key, self.root)
    if floor_node is None:
        return None
    else:
        return floor_node.key

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def _ceiling_node(self, key, node):
    """
    Returns the node with the smallest key that
    is greater than or equal to
    the given value 'key'
    """
    if node is None:
        return None

    if key < node.key:
        # Ceiling is either in left subtree or is
        this node
        attempt_in_left = self._ceiling_node(key,
        node.left)
        if attempt_in_left is None:
            return node
        else:
            return attempt_in_left
    elif key > node.key:
        # Ceiling must be in right subtree
        return self._ceiling_node(key,
        node.right)
    else:
        # Keys are equal so ceiling is node with
        this key
        return node

def ceiling_key(self, key):
    """
    Returns the smallest key that is greater than
    or equal to the given
    value 'key'

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    ceiling_node = self._ceiling_node(key,

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self.root)
    if ceiling_node is None:
        return None
    else:
        return ceiling_node.key

def _select_node(self, rank, node):
    """
    Return the node with rank equal to 'rank'
    """
    if node is None:
        return None

    left_size = self._size(node.left)
    if left_size < rank:
        return self._select_node(rank - left_size
- 1, node.right)
    elif left_size > rank:
        return self._select_node(rank, node.left)
    else:
        return node

def select_key(self, rank):
    """
    Return the key with rank equal to 'rank'

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    select_node = self._select_node(rank,
self.root)
    if select_node is None:
        return None
    else:
        return select_node.key

def _rank(self, key, node):
    if node is None:

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        return None

    if key < node.key:
        return self._rank(key, node.left)
    elif key > node.key:
        return self._size(node.left) +
self._rank(key, node.right) + 1

    else:
        return self._size(node.left)

def rank(self, key):
    """
    Return the number of keys less than a given
    'key'.

    Worst Case Complexity:  $O(N)$ 

    Balanced Tree Complexity:  $O(\lg N)$ 
    """
    return self._rank(key, self.root)

def _delete(self, key, node):
    if node is None:
        return None
    if key < node.key:
        node.left = self._delete(key, node.left)
    elif key > node.key:
        node.right = self._delete(key,
node.right)

    else:
        if node.right is None:
            return node.left
        elif node.left is None:
            return node.right
        else:
            old_node = node
            node = self._ceiling_node(key,

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node.right)
        node.right =
self._delete_min(old_node.right)
        node.left = old_node.left
        node.size_of_subtree = self._size(node.left)
+ self._size(node.right)+1
        return node

def delete(self, key):
    """
    Remove the node with key equal to 'key'

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    self.root = self._delete(key, self.root)

def _delete_min(self, node):
    if node.left is None:
        return node.right

    node.left = self._delete_min(node.left)
    node.size_of_subtree = self._size(node.left)
+ self._size(node.right)+1
    return node

def delete_min(self):
    """
    Remove the key-value pair with the smallest
key.

    Worst Case Complexity: O(N)

    Balanced Tree Complexity: O(lg N)
    """
    self.root = self._delete_min(self.root)

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def _delete_max(self, node):
    if node.right is None:
        return node.left

    node.right = self._delete_max(node.right)
    node.size_of_subtree = self._size(node.left)
+ self._size(node.right)+1
    return node

def delete_max(self):
    """
    Remove the key-value pair with the largest
key.

    Worst Case Complexity:  $O(N)$ 

    Balanced Tree Complexity:  $O(\lg N)$ 
    """
    self.root = self._delete_max(self.root)

def _keys(self, node, keys):
    if node is None:
        return keys

    if node.left is not None:
        keys = self._keys(node.left, keys)

    keys.append(node.key)

    if node.right is not None:
        keys = self._keys(node.right, keys)

    return keys

def keys(self):
    """
    Return all of the keys in the BST in
ascending order

```

Worst Case Complexity: $O(N)$

Balanced Tree Complexity: $O(N)$
"""

keys = []

return self._keys(self.root, keys)