TP3: Large Scale Graph Learning

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Abstract

The report and the code are due in 2 weeks (deadline $23:59\ 26/12/2018$). You will find instructions on how to submit the report on piazza, as well as the policies for scoring and late submissions. All the code related to the TD must be submitted.

A small preface (especially for those that will not be present at the TD). Some of the experiments that will be presented during this TD makes use of randomly generated datasets. Because of this, there is always the possibility that a single run will be not representative of the usual outcome. Randomness in the data is common in Machine Learning, and managing this randomness is important. Proper experimental setting calls for repeated experiments and confidence intervals. In this case, it will be sufficient to repeat each experiment multiple times and visually see if there is large variations (some experiments are designed exactly to show this variations).

Whenever the instructions require you to complete a file (e.g. exponential_euclidean.m) this means to open the corresponding file, and complete it according to the instructions all the following sections

For each of these file the documentation at the top describes their parameters.

1 Large Scale Semi Supervised Learning

A large part of the algorithms on graph that we have seen so far, depends on matrix representation of the graph (e.g. Laplacian) to compute their results. This means that the execution of these algorithm is strongly influenced by constraint of representing matrices in memory, and the computational complexity of executing operations on them.

We will see later how the computational complexity plays a role, but the largest constraint in real-world systems is memory occupation. In particular, storing in memory all the elements of an $n \times n$ matrix will require around 190 MB for n=5000, but will easily cross the GB threshold at n=12000 and the 10 GB threshold at n=37000. At this point, storing the whole graph in a single machine's memory becomes impossible, and we must either compress it using a graph approximation technique, distribute it across several machine, or store it in a larger memory and process it in multiple passes.

2 Online SSL

Semi-Supervised Learning was introduced as a solution designed for problems where collecting large quantities of Supervised training data (usually labels) is not possible. On the other hand, in SSL scenarios it is usually inexpensive to obtain more samples coming from the same process that generated the labeled samples, but without a label attached. A simple approach is to collect as much additional data as possible, and then run our algorithm on all of it, in a batch or off-line manner. But in other cases it is possible to start the learning process as early as possible, and then refine the solution as the quantity of available information increases. An important example of this approach is stream processing. In this setting, a few labeled examples are provided in advance and set the initial bias of the system while unlabeled examples are gathered online and update the bias continuously. In the online setting, learning is viewed as a repeated game against a potentially adversarial nature. At each step t of this game, we observe an example x_t , and then predict its label f_t . The challenge of the game is that after the game started we do not observe any more true label y_t . Thus, if we want to adapt to changes in the environment, we have to rely on indirect forms of feedback, such as the structure of data. Another difficulty posed by online learning is computational cost. In particular, when t becomes large, a naive approach to hard-HFS, such as recomputing the whole solution from scratch, has prohibitive computational costs. Especially in streaming settings where near real-time requirements are common, it is imperative to have an approach that has scalable time and memory costs. Because most operations related to hard-HFS scale with the number of nodes, one simple but effective approach to scale this algorithm is subsampling, or in other words to compute an approximate solution on a smaller subset of the data in a way that generalizes well to the whole dataset. Several techniques can give different guarantees for the approximation. As you saw in class, incremental k-centers [1] guarantees on the distortion introduced by the approximation allows us to provide theoretical guarantees.

Algorithm 1 Incremental k-centers

- 1: **Input:** an unlabeled \mathbf{x}_t , a set of centroids C_{t-1} , \mathbf{b} , \mathbf{V}
- 2: **if** $(|C_{t-1}| = k+1)$ **then**
- 3: Find two closest centroids c_{add} and c_{rep} . c_{add} will forget the old centroid and will point to the new sample that just arrived, and c_{rep} will take care of representing all nodes that belonged to c_{add}
- 4: c_{add} will now represent the new node, we must update accordingly V and the stored centroids.
- 5: **else**
- 6: $C_t \leftarrow C_{t-1}$
- 7: x_t is added as a new centroid c_t and V is updated accordingly
- 8: end if

There is a lot of practical consideration to keep in mind when implementing this kind of algorithms.

- The labeled nodes are fundamentally different from unlabeled ones. Because of this, it is always a good idea to keep them separate, and never merge them in a centroid. In the implementation this is accomplished with a taboo list **b** that keeps track of nodes that cannot be merged together.
- Ammortized cost matters, especially in streaming application. It is not always possible to stop execution to repartition the centroids, and it is often preferrable to pay a small price at every step to keep execution smooth. In our case, the centroids are updated at every step.

The most important variables for the centroid is V. It allow us to extract labels from the solutions, and to correctly mantain information on the multiplicities of every centroid. Whenever a new node arrives, and we have too many centroids, we choose the two closest centroids c_{add} and c_{rep} . c_{add} will forget the old centroid and will point to the new sample that just arrived, and c_{rep} will take care of representing all nodes that belonged to c_{add} .

2.1. Complete online_ssl_update_centroids using the pseudocode from 1.

Computing the solutions from the quantized graph using hard (constrained) HFS is pretty straightforward.

Algorithm 2 Online HFS with Graph Quantization

- 1: **Input:** t, a set of centroids C_{t-1} , V, \mathbf{b} , labels \mathbf{y}
- 2: Compute the laplacian Δ_t of G using $W_q = V\widetilde{W}_q V$
- 3: Infer labels using hard-HFS
- 4: Predict $\widehat{y}_{t} = \operatorname{sgn}(\mathbf{f}_{u}(t))$.
- 5: With the preceding construction of the centroids, x_t is always present in the reduced graph and does not share the centroid with any other node.
- 2.2. Complete online_ssl_compute_solution.m following the pseudocode from 2

Use create_user_profile to capture a training set of labeled data of your face and someone else. Read the help for cv.VideoCapture to understand how to open a stream. The faces will be preprocessed and saved in the folder data/faces. They will be loaded by online_face_recognition.m (you have to adjust the subjects name in the file).

- 2.3. Complete preprocess_face.m and run online_face_recognition. m. Include some of the resulting frames (not too similar) in the report showing faces correctly being labeled as opposite, and comment on the choices you made during the implementation.
- 2.4. What happens if an unknown person's face is captured by the camera? Modify your code to be able to disregard faces it cannot recognize, and include some of the resulting frames (not too similar) in the report showing unknown faces correctly being labeled as unknown.

3 Large Scale Label Propagation

We will begin by implementing label propagation on a middle scale dataset. In TD2 we have seen how to do this by computing an HFS solution in closed form. The Harmonic property that the HFS solution wants to satisfy is

$$f(x_i) = \frac{\sum_{i \sim j} f(x_j) w_{ij}}{\sum_{i \sim j} w_{ij}}$$

Computing HFS using linear algebra is not particularly easy to parallelize, and does not decompose easily over vertices as required by most distributed

graph frameworks. But the Harmonic property gives us an easy idea for implementing this algorithm iteratively with subsequent sweeps across the graph.

For the implementation, a graph is represented using sparse matrices in Compressed Sparse Column (CSC) format¹. When implementing an algorithm, it is often important to have at least a rough idea of the underlying data representation used by the implementation. For CSC matrices, accessing the matrix column-wise is a much cheaper operation than extracting a row of the matrix. Another common constraint for distributed algorithms is that when the algorithm processes node i, it has access only to the edges incident to node i itself. We will represent this constraint by allowing only operations on single columns of the matrix in iterative_hfs.m, forbidding global operations on the whole matrix (e.g. matrix multiplication).

- 3.1. Complete iterative_hfs.m. Be careful to only use vector operations (e.g. vector-vector multiplication, sum over a single colum), and to respect the column-wise access pattern suggested by CSC.
- 3.2. In normal HFS, regularization is added to the Laplacian to simulate absorption at each step in the label propagation. How can you have a similar regularization in the iterative label propagation?

References

[1] Moses CHARIKAR, Chandra CHEKURI, Tomas FEDER, and Rajeev MOTWANI. Incremental clustering and dynamic information retrieval. SIAM journal on computing, 33(6):1417–1440, 2004.

https://en.wikipedia.org/wiki/Sparse_matrix