

Bit Manipulation

Get Bit:

- This operation *SHIFTS 1* over by *i* bits, creating a value that looks like 00010000. By performing an *AND* with *num*, we clear all bits other than the bit at *i*-th position. Finally, we compare that to 0. If that new value is not zero, then bit *i* must have a 1. Otherwise, bit *i* is a 0.

$\text{num} \& (1 \ll i)$

num:	01101001
1 << 4:	00010000

num & (1 << 4):	00000000

- The following operation isolate the lowest (rightmost) set bit of *num*. The method returns an integer which is the index of the rightmost set bit of *num*.

$\text{num} \& \sim(\text{num}-1)$

num:	01101000
1:	00000001

num-1:	01100111

~(num-1):	10011000
num:	01101000

num & ~(num-1):	00001000

- The following method isolate the lowest (rightmost) unset bit of *num*. . The method returns an integer which is the index of the rightmost unset bit of *num*.

$\sim\text{num} \& (\text{num}+1)$

num:	01100111
1:	00000001

num+1:	01101000
~num:	10011000

~num & (num+1):	00001000

Set Bit:

- Set Bit shifts 1 over by i bits, creating a value like 00010000. By performing an **OR** with num , only the value at bit i will change. All other bits of the mask are zero and will not affect num .

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$num \mid (1 \ll i)$

num:	01101001
$1 \ll 4$:	00010000

$num \mid (1 \ll 4)$:	01111001

Clear Bit:

- This method operates in almost the reverse of setBit. First, we create a number like 11101111 by creating the reverse of it (00010000) and negating it. Then, we perform an **AND** with num . This will clear the i th bit and leave the remainder unchanged.

$num \& \sim(1 \ll 3)$

num:	01101001
$1 \ll 3$:	00001000

$\sim(1 \ll 3)$:	11110111
num:	01101001

$num \& \sim(1 \ll 3)$:	01100001

- The following method clear the lowest set bit in num

$num \& (num-1)$

num:	01101000
1:	00000001

num-1:	01100111
num:	01101000

$num \& (num-1)$:	01100000

- To clear all bits from the most significant bit through i (**inclusive**), we do:

```
num & (1 << i) - 1
```

num:	01101001

1 << 4:	00010000
1:	00000001

(1 << 4) - 1:	00001111
num:	01101001

num & (1 << 4) - 1:	00001001

- To clear all bits from i through 0 (**inclusive**), we do:

```
int clearBitsIThrough0(int num, int i) {
    int mask = ~((1 << (i+1)) - 1);
    return num & mask;
}
```

num:	01101001

1 << 4+1:	00100000
1:	00000001

(1 << 4+1) - 1:	00011111

~(1 << 4+1) - 1:	11100000
num:	01101001

num & ~(1 << 4+1) - 1:	01100000

Update Bit:

- This method merges the approaches of **setBit** and **clearBit**. First, we clear the bit at position i by using a mask that looks like 11101111. Then, we shift the intended value, v , left by i bits. This will create a number with bit i equal to v and all other bits equal to 0. Finally, we Or these two numbers, updating the i th bit if v is 1 and leaving it as 1 otherwise.^{[1][SEP]}

```
int updateBit(int num, int i, int v) {
    int mask = ~(1 << i);
    return (num & mask) | (v << i);
}
```

num:	01101001
v:	00010010

1 << 4:	00010000
~(1 << 4):	11101111
num:	01101001

num & ~(1 << 4):	00011111
v << 4:	00100000

(num & ~(1 << 4)) (v << 4):	00111111

▪ Addition using ONLY Bitwise Operator without using any Arithmetic Operator

<pre>static long add(long x, long y){ while (y != 0) { long carry = x & y; x = x ^ y; y = carry << 1; } return x; }</pre>	X = 0111; Y = 0110; X + Y = 1101		
	0111 (&) 0110 ----- 0110	0111 (^) 0110 ----- 0001	Y = 1100
	1100 (&) 0001 ----- 0000	0001 (^) 1100 ----- 1101	Y = 0000

<pre>static long multiply(long x, long y) { long sum = 0; while (x != 0) { if ((x & 1) != 0) sum = add(sum, y); x >>= 1; y <<= 1; } return sum; }</pre>	Y = 0111; X = 0110; Y × X = 101010 0 1 1 1 (×) 0 1 1 0 ----- 0 0 0 0 0 1 1 1 × 0 1 1 1 × × 0 0 0 0 × × × ----- 0 1 0 1 0 1 0
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<p>let, x is a int & n is wordlength of significant digit in x. Then,</p> <ul style="list-style-type: none">▪ $n = \text{floor}(\log_{10}x) + 1$▪ $\frac{y=x}{x^{n-1}}$: Extract the MSD of x▪ $y = x \% 10$: Extract the LSD of x▪ $x = x \% 10^{n-1}$: Remove MSD of x▪ $x = \frac{x}{10}$: Remove LSD of x.	<p>Example: let $x = 157793$ and $z = 099910$</p> <ul style="list-style-type: none">▪ $n_x = 5 + 1 = 6$; $n_z = 4 + 1 = 5$▪ $\frac{x}{x^{n_x-1}} = 1$; $\frac{z}{z^{n_z-1}} = 9$▪ $x \% 10 = 3$; $z \% 10 = 0$▪ $x = x \% 10^{n-1} = 57793$;▪ $z = z \% 10^{n-1} = 9910$▪ $x = \frac{x}{10} = 15779$; $z = \frac{z}{10} = 9991$
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- $1111\ 1111\ 1111\ 1111 = 0xFFFF = 65535 = 2^{16} - 1$
- $1\ 0000\ 0000\ 0000\ 0000 = 0x10000 = 65536 = 2^{16}$
- $1010\ 1010 = 0xAA$
- $0101\ 0101 = 0x55$
- $\sim((1 \ll 4) - 1) = 111110000$

Bit Facts and Trick:

<i>Bitwise: XOR (^)</i>	<i>Bitwise: AND (&)</i>	<i>Bitwise: OR ()</i>
$X \wedge X = 0s:$ <pre> 0 1 1 0 (XOR) 0 1 1 0 ----- 0 0 0 0 </pre>	$X \& X = X:$ <pre> 0 1 1 0 (AND) 0 1 1 0 ----- 0 1 1 0 </pre>	$X X = X:$ <pre> 0 1 1 0 (OR) 0 1 1 0 ----- 0 1 1 0 </pre>
$X \wedge (\sim X) = 1s:$ <pre> 0 1 1 0 (XOR) 1 0 0 1 ----- 1 1 1 1 </pre>	$X \& (\sim X) = 0s:$ <pre> 0 1 1 0 (AND) 1 0 0 1 ----- 0 0 0 0 </pre>	$X \wedge (\sim X) = 1s:$ <pre> 0 1 1 0 (OR) 1 0 0 1 ----- 1 1 1 1 </pre>
$X \wedge (0s) = X:$ <pre> 0 1 1 0 (XOR) 0 0 0 0 ----- 0 1 1 0 </pre>	$X \& (0s) = 0s:$ <pre> 0 1 1 0 (AND) 0 0 0 0 ----- 0 0 0 0 </pre>	$X \wedge (0s) = X:$ <pre> 0 1 1 0 (OR) 0 0 0 0 ----- 0 1 1 0 </pre>
$X \wedge (1s) = \sim X:$ <pre> 0 1 1 0 (XOR) 1 1 1 1 ----- 1 0 0 1 </pre>	$X \& (1s) = X:$ <pre> 0 1 1 0 (AND) 1 1 1 1 ----- 0 1 1 0 </pre>	$X \wedge (1s) = 1s:$ <pre> 0 1 1 0 (OR) 1 1 1 1 ----- 1 1 1 1 </pre>