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### Chapter 1. Solutions to Lambda Calculi with Types

**Problem 1.** (Exercise 3.1.13) \*Exercise Statement\*

Solution: \*Solution!\*

**Problem 2.** (Exercise 4.1.20) \*Exercise Statement\*

Problem 3. (Exercise 4.2.8) \*Exercise Statement\*

Problem 4. (Exercise 5.1.16) \*Exercise Statement\*

# Chapter 2. Solutions to Domain-Theoretic Foundations of Functional Programming

### 2.1 PCF and its Operational Semantics

**Problem 1.** (Page 14) \*Problem Statement\*

**Problem 2.** (Page 16) (Lemma 2.1.) The evaluation relation  $\Downarrow$  is deterministic, i.e. whenever  $M \Downarrow V$  and  $M \Downarrow W$  then  $V \equiv W$ 

**Solution:** We prove this by induction on the structure of the derivation. **Base cases.** 

- By the rules of the BigStep semantics for PCF, the lemma for the following base cases is trivial:
  - $-M \equiv x$ , then  $x \downarrow x$ . So V and W can only be x; thus,  $V \equiv W \equiv x$ .
  - $-M \equiv \lambda x : \sigma.M$ , then  $\lambda x : \sigma.M \downarrow \lambda x : \sigma.M$ .
  - $-M\equiv 0$ , then  $0\downarrow 0$ .

#### Inductive Steps.

- If  $M \equiv succ(M)$ , then it must be derived by the rule  $\frac{M \Downarrow \underline{n}}{succ(M) \Downarrow \underline{n+1}}$ . Then we would have  $V \equiv \underline{n+1}$  and  $W \equiv \underline{m+1}$  since the successor rule is the only way to derive succ(M). By IH, we know that  $\underline{n} = \underline{m}$ , thus  $\underline{n+1} = \underline{m+1}$ , and hence V = W.
- If  $M \equiv M(N)$ . The derivation for M(N) must be of the form  $\frac{M \Downarrow \lambda x: \sigma.E \ E[N/x] \Downarrow V}{M(N) \Downarrow V}$ . A second derivation for M(N) must use the same rule. i.e.,  $\frac{M \Downarrow \lambda x: \sigma'.E' \ E'[N/x] \Downarrow W}{M(N) \Downarrow W}$ . But then by IH, we would have  $\lambda x: \sigma.E \equiv \lambda x: \sigma'.E'$ . So  $\sigma \equiv \sigma'$  and E = E'. Now, we have  $E[N/x] \Downarrow V$  and  $E[N/x] \Downarrow W$ . By the IH on the sub-derivation for E[N/x], we conclude  $V \equiv W$ .
- If  $M \equiv pred(M)$ , then the rules are  $\frac{M \Downarrow \underline{0}}{pred(M) \Downarrow \underline{0}}$  and  $\frac{M \Downarrow \underline{n+1}}{pred(M) \Downarrow \underline{0}}$ . For the derivation  $pred(M) \Downarrow V$ , we must have a sub-derivation for  $M \Downarrow \underline{x}$  for some numeral  $\underline{x}$ . Similarly, for  $pred(M) \Downarrow W$ , we must have a sub-derivation for  $M \Downarrow \underline{y}$  for some numeral  $\underline{y}$ . By the IH on the sub-derivation for M, we can conclude that  $\underline{x} \equiv y \equiv k$ .

Let's examine k. If  $k \equiv \underline{0}$ , then both derivations must be  $\frac{M \Downarrow \underline{0}}{pred(M) \Downarrow \underline{0}}$ . Thus,  $V \equiv W \equiv \underline{0}$ . Same argument is valid for the case  $k \equiv n+1$ .

• The other cases  $(Y_{\sigma}, \text{ and both cases of } ifz)$  can be proved likewise.

**Problem 3.** (Page 16) \*Problem Statement\*

**Problem 4.** (Page 17) \*Problem Statement\*

**Problem 5.** (Page 17) \*Problem Statement\*

**Problem 6.** (Page 19) \*Problem Statement\*

### 2.2 The Scott Model of PCF

Problem 1. (Page 26) \*Problem Statement\*

Problem 2. (Page 26) \*Problem Statement\*

Problem 3. (Page 27) \*Problem Statement\*

Problem 4. (Page 30) \*Problem Statement\*

Problem 5. (Page 33) \*Problem Statement\*

Problem 6. (Page 34) \*Problem Statement\*

### 2.3 Milner's Context Lemma

Problem 1. (Page 44) \*Problem Statement\*

# 2.4 Logical Relations

Problem 1. (Page 52) \*Problem Statement\*

**Problem 2.** (Page 54) \*Problem Statement\*