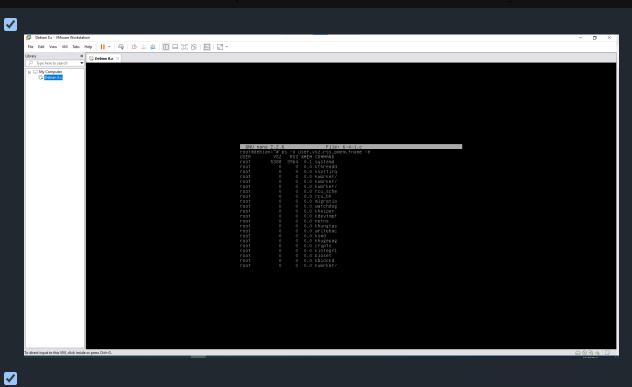
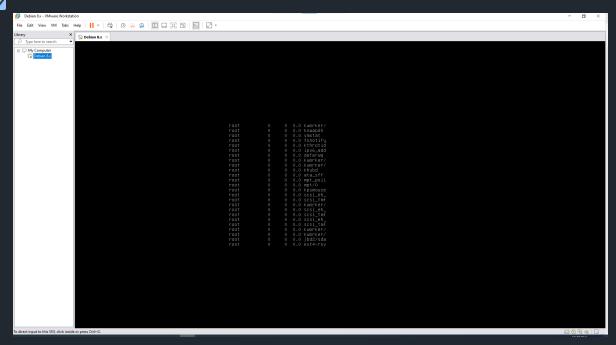
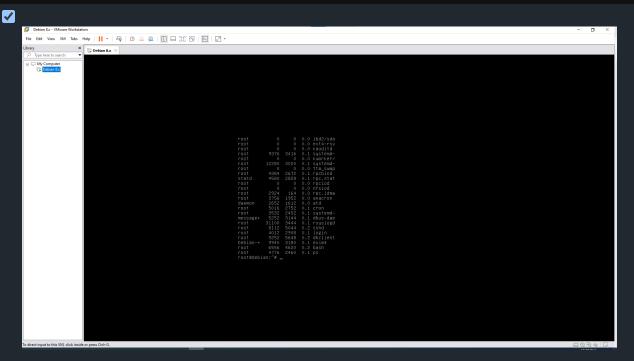


The malloc function returns a void pointer to the reserved space. The return value is NULL if not enough storage is available, or if size was specified as zero.

✓ Using ps



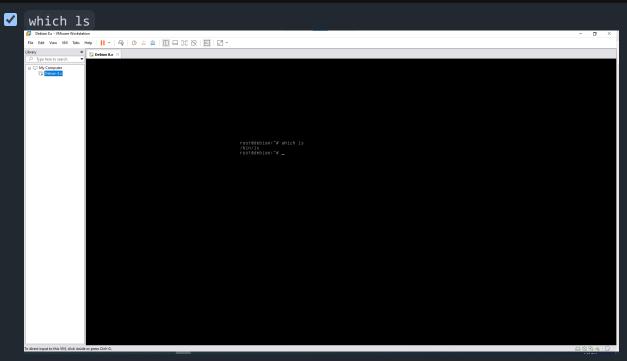


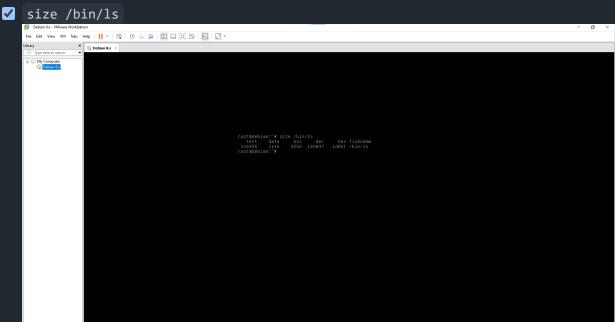


✓ Columns are:

- user: This column displays the user ID or username of the person who owns and runs the process. Example Output: root, user1. Alias: uname.
- vsz: This column shows the virtual memory size of the process in kilobytes. It's the total memory size that a process can access, such as the memory of the shared libraries, the allocated memory such as stack and heap, and swapped out memory.

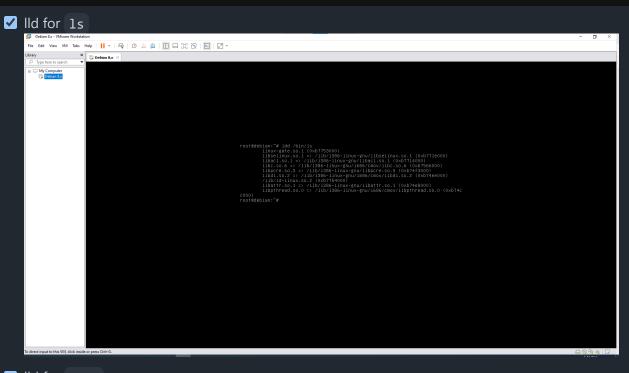
 Example Output: 1048576 (for 1 GB). Alias: vsize.
- rss: This column shows the resident set size, which is the non-swapped physical memory a task has used (in kilobytes). It doesn't show all of the information about the memory, like loaded libraries, heap, or stack. It shows the shared libraries if their pages are in memory. Example Output: 2048 (for 2 MB).
- pmem: This column shows the percentage of resident set size (RSS) to the physical memory used by the process. Example Output: 0.1
- fname: This column displays the first 8 bytes of the process's executable file name. The output in this column may contain spaces. Example Output: bash, python
- Getting started with memory segments

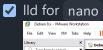


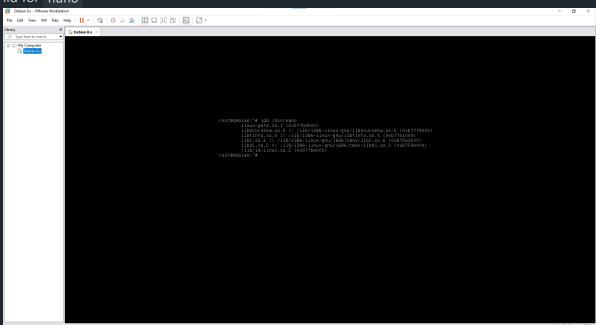


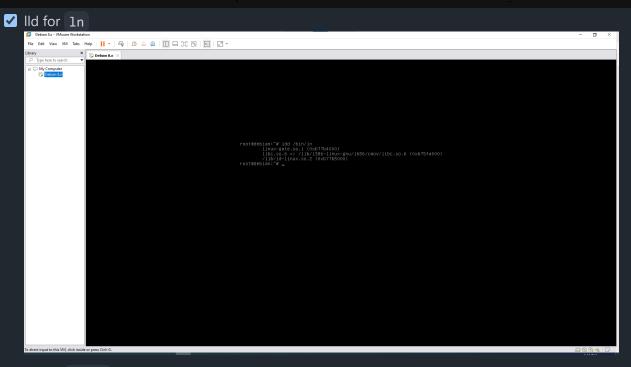
As you see, 116455 bytes from memory are allocated to text part, 1140 bytes for initialized data, 3252 bytes for bss, and 120847 bytes for dec.

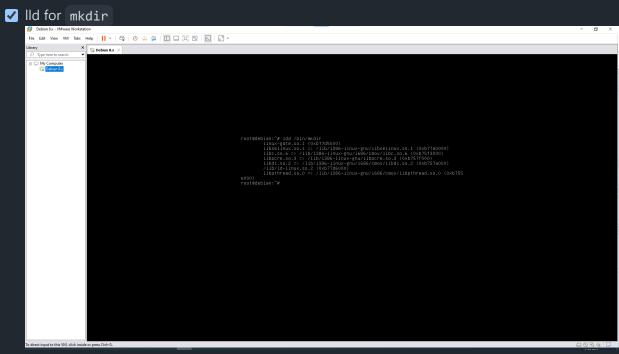
- ✓ This command doesn't show stack and heap size.
- ✓ Getting started with memory sharing



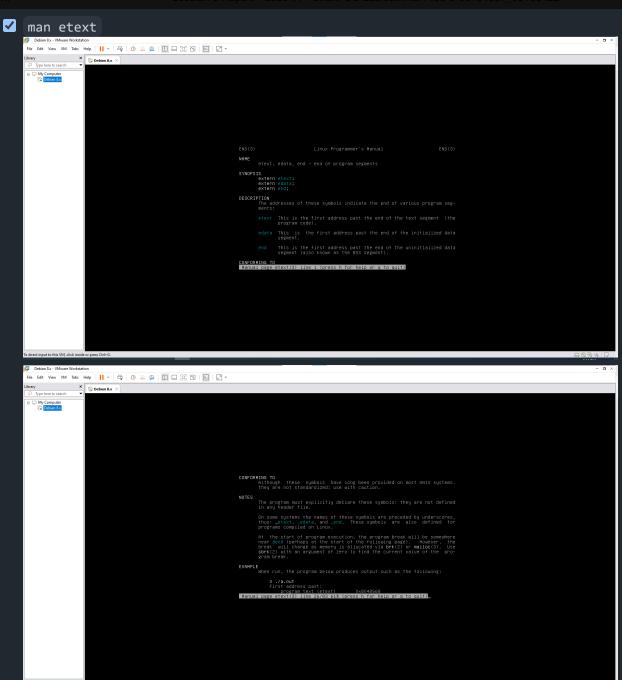


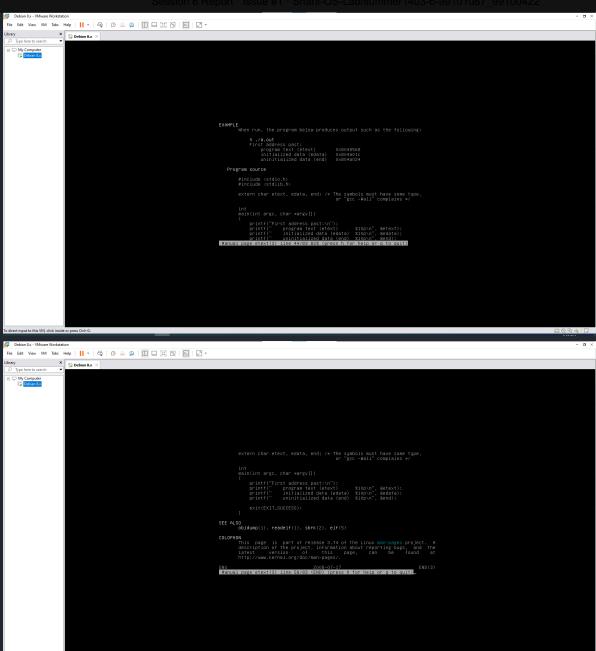




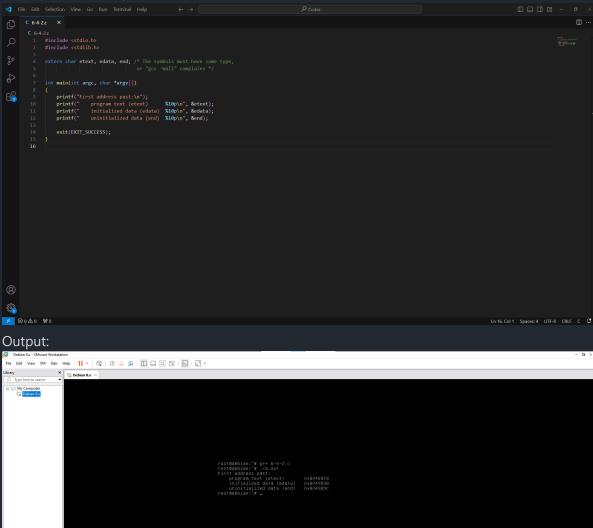


✓ Getting started with addresses





✓ Code (in file 6-4-2.c):



As you see, the addresses are exactly the same as what we expected.

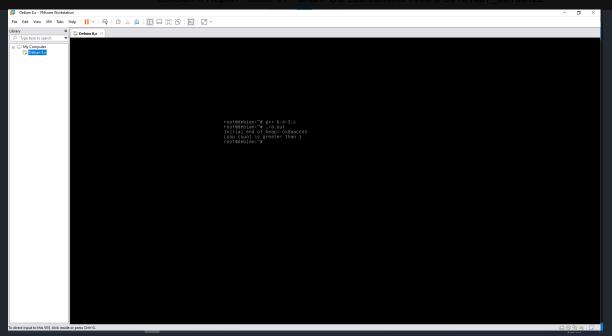
- Description about etext, the meaning of the comment, why using the term symbol, why using extern:
 - ✓ The symbols are:
 - etext: This is the first address past the end of the text segment (the program code). The text segment contains the compiled machine code of the program.
 - edata: This is the first address past the end of the initialized data segment
 - end: This is the first address past the end of the uninitialized data segment (also known as the BSS segment).
 - Meaning of the comment: The '-Wall' flag enables all compiler's warning messages, which helps in writing cleaner and more error-free code. In C, declaring symbols with extern

- requires specifying a type. Without a type, the compiler would issue warnings or errors. Here, char is chosen as the type because it is a basic type that simply gives these symbols a valid type without any specific size or alignment constraints.
- Why use the term symbol: These symbols () are used to refer to specific locations in memory and do not hold values themselves. They are part of the program's memory map, which is why they are called symbols rather than variables.
- Why use extern: The extern keyword is used to declare these symbols without defining them. This informs the compiler that these symbols are defined elsewhere, typically by the linker. These symbols are defined by the linker, not by the C code, so using extern tells the compiler that their actual addresses will be resolved during the linking stage. Also, by declaring them as extern char, the code can take the address of these symbols (using &etext, &edata, &end) to print or otherwise.
- sbrk analysis:

Code (in file 6-4-3.c):

```
| Tile tild Selection | View Go Run Terminal Help | Code | Piccolas | Code | Co
```

Output:



As you see, the loop count is not equal to 1. That's because memory allocators often request larger chunks of memory from the operating system and then manage these chunks internally. When malloc is called, it might be allocating memory from a pre-existing pool rather than extending the heap. This means the end of the heap doesn't change with every call to malloc. In the above code, the size of malloc was 1 MB. If we change it to 1 KB, the address of end of heap will change and the loop count will be 1. Its output will be:

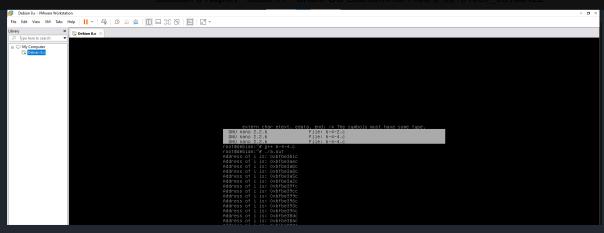
We could write code like this to see if malloc changes the address of end of heap:

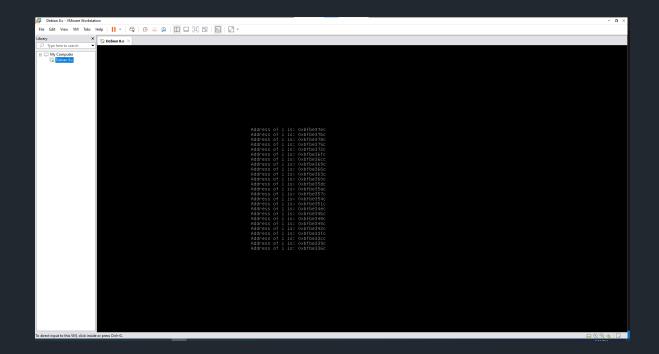
```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

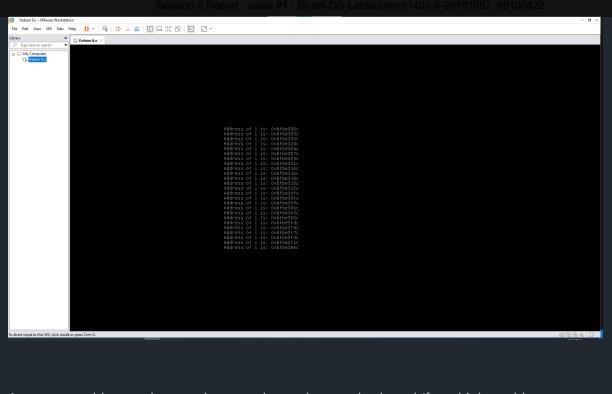
int main() {
    printf("End of heap before malloc: %p\n", sbrk(0));
    malloc(1024 * 1024);
    printf("End of heap after malloc: %p\n", sbrk(0));
    return 0;
}
```

✓ stack growth analysis: Code (in file 6-4-4.c):

Output:





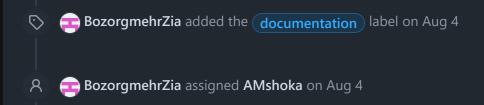


As you see, addresses decrease because the stack grows backward (from higher address to lower address).

The codes are:

Codes.zip







Projects

None yet

(3)

