

Selection sort

Implementation:—

```
void selection sort (int arr[], int n)
```

```
{
  for (int i = 0; i < (n-1); i++)
```

```
{
  for (int j = i+1; j < n; j++)
```

```
{
  if (arr[i] > arr[j])
```

```
    swap (arr[i], arr[j]);
```

```
}
```

```
}
```

```
return;
```

Analysis:

```
for (int i = 0; i < (n-1); i++)
```

```
{
```

```
for (int j = i+1; j < n; j++)
```

```
{
```

```
if (arr[i] > arr[j])
```

```
    swap (arr[i], arr[j]);
```

```
}
```

```
}
```

• cost time

 $C_1 (n-1)$ $C_2 \frac{n(n+1)}{2}$ $C_3 \frac{n(n+1)}{2}$ $C_4 \frac{n(n+1)}{2}$

\therefore Time function $f(t) = c_1(n-1) + c_2 \frac{n(n+1)}{2} + c_3 \frac{n(n+1)}{2} + c_4 \frac{n(n+1)}{2}$

Best case occurs when the array is already sorted

1	2	3	4	5
---	---	---	---	---

Value of c_4 will be 0, so, time function will be

$$f(t) = \left(\frac{c_2+c_3}{2}\right) * n^2 + \left(\frac{c_2+c_3+2c_1}{2}\right) * n - c_1$$

This function look like $f(t) = an^2 + bn + c$ which is quadric function.

\therefore worst case will be occurred when the array is reversly sorted

5	4	3	2	1
---	---	---	---	---

$$f(t) = \left(\frac{c_4+c_2+c_3}{2}\right) * n^2 + \left(\frac{c_4+c_2+c_3+2c_1}{2}\right) * n - c_1$$

look like $f(t) = (an^2 + bn + c)$; which is quadric function.

Time complexity:

$$f(t) = an^2 + bn + c \text{ [Best case]}$$

$$O(n^2)$$

Worst case:

$$f(t) = an^2 + bn + c$$

$$O(n^2)$$

Implementation:

Insertion sort

void insertion (int arr[], int n)

```
{
    int i, j, x;
    for (i = 0; i < n; i++)
    {
        x = arr[i];
        j = i - 1;
        while (j > 0 && arr[j] > x)
        {
            arr[j + 1] = arr[j];
            j--;
        }
        arr[j + 1] = x;
    }
}
```

Analysis

void insertion (int arr[], int n) cost time

```
{
    int i, j, x;
    for (i = 0; i < n; i++)
```

$C_1 \quad n$

```
{
    x = arr[i];
```

```
    j = i - 1;
```

```
    while (j > 0 && arr[j] > x)
```

$C_2 \quad n-1$

$C_3 \quad n-1$

```
    arr[j + 1] = arr[j]; j--;
```

$C_4 \quad \frac{n(n+1)}{2}$

$C_5 \quad \frac{n(n+1)}{2}$

arr[j+1] = x;

c_6

$n-1$

∴ Time function:

$$f(t) = c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 \frac{n(n+1)}{2} + c_5 \frac{n(n+1)}{2} + c_6 n(n+1)$$

Best case occurs when array is already sorted

1 2 3 5 7

when c_4 will execute for $(n-1)$ time and c_5 will execute for 0 time.

So, function will be,

$$f(t) = c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 (n-1) + c_5 * 0 + c_6 (n-1)$$

$$= n(c_1 + c_2 + c_3 + c_4 + c_6) - (c_2 + c_3 + c_4 + c_6)$$

∴ This look like $y = ax - b$, That's linear eqⁿ worst case occurs when the array is neverly sorted.

7 5 3 2 1

∴ Time function will be

$$f(t) = c_1 n + c_2 + c_3 (n-1) + c_4 \frac{n(n+1)}{2} + c_5 \frac{n(n+1)}{2} + c_6 (n-1)$$

$$= \frac{c_1 + c_2}{2} n^2 + \frac{2(c_1 + c_2 + c_3 + c_6)}{2} n - (c_2 + c_3 + c_4 + c_6)$$

\therefore look like $y = ax^2 + bx + c$; This is a quadratic eqn.

Time complexity:

Best case

we know, $y = ax + b$ [\therefore Best case]

So, $O(n)$

Worst case:

we know, $y = ax^2 + bx + c$ [\therefore worst case]

So, $O(n^2)$