# Exercise 3

## Compilation 0368:3133

Due 27/12/2017

## 1 Introduction

We continue our journey of building a compiler for the invented object oriented language RioMare. Remember that the entire specification of RioMare appears inside the relevant folder of the course website. In order to make this document self contained, all the information needed to complete the third exercise is brought here again.

# 2 Programming Assignment

The third exercise implements a semantic analyzer that recursively scans the AST produced by CUP, and checks if it contains any semantic errors. The input for the semantic analyzer is a (single) text file containing a RioMare program, and the output is a (single) text file indicating whether the input program is semantically valid or not. In addition to that, whenever the input program is valid semantically, the semantic analyzer will add meta data to the abstract syntax tree, which is needed for later phases (code generation and optimization). The added meta data content will not be checked in exercises 3, but the best time to design and implement this addition is exercise 3.

# 3 The RioMare Semantics

This section describes the semantics of RioMare, and provides a multitude of legal and illegal example programs.

### 3.1 Types

The RioMare programming language defines two native types: integers and strings. In addition, it is possible to define a class by specifying its data members and methods. Also, given an existing type T, one can define an array of T's. Note, that defining classes and arrays is only possible in the uppermost (global) scope. The exact details follow.

#### 3.1.1 Classes

Classes contain data members and methods, and can only be defined in the uppermost (global) scope. They can refer to/extend only previously defined classes, to ensure that the class hierarchy has a tree structure. Following the same concept, a method M1 can not refer to a method M2, whenever M2 is defined after M1 in the class. In contrast to all that, a method M can refer to a data member d, even if d is defined after M in the class. Table 1 summarizes these facts.

```
CLASS Son EXTENDS Father
   {
        int bar;
                                          ERROR
   CLASS Father
        void foo() { PrintInt(8); }
   }
3
   CLASS Edge
        Vertex u;
        Vertex v;
   }
                                          ERROR
   CLASS Vertex
        int weight;
   CLASS UseBeforeDef
                                          ERROR
        void foo() { bar(8); }
        void bar(int i) { PrintInt(i); }
   CLASS UseBeforeDef
        void foo() { PrintInt(i); }
                                          OK
        int i;
   }
```

Table 1: Referring to classes, methods and data members

**Methods overloading** is *illegal* in RioMare, with the obvious exception of overriding a method in a derived class. Similarly, it is illegal to define a variable with the same name of a previously defined variable (shadowing), or a previously defined method. Table 2 summarizes these facts.

```
CLASS Father
       int foo() { return 8; }
                                          ERROR
   CLASS Son EXTENDS Father
       void foo() { PrintInt(8); }
   CLASS Father
       int foo(int i) { return 8; }
                                          OK
   CLASS Son EXTENDS Father
       int foo(int i) { return i; }
   CLASS IllegalSameName
                                          ERROR
       void foo() { PrintInt(8); }
       void foo(int i) { PrintInt(i); }
3
   CLASS Father
       int foo;
                                          ERROR
   CLASS Son EXTENDS Father
       string foo;
   }
```

Table 2: Method overloading and variable shadowing are both illegal in Ri-oMare.

In addition, if class Son is derived from class Father, then any place in the program that semantically allows an expression of type Father, should semantically allow an expression of type Son. For example,

```
ARRAY IntArray EQ int[]
void F(IntArray A){ PrintInt(A[8]); }
void main(){ F(NULL); }
```

Table 3: NULL sent instead of an integer array is semantically allowed.

**Arrays** can only be defined in the uppermost (global) scope. They are defined with respect to some previously defined type, as in the following example:

```
ARRAY IntArray EQ int[]
```

Defining an integer matrix, for example, is possible as follows:

```
ARRAY IntArray EQ int[] ARRAY IntMat EQ IntArray[];
```

In addition, any place in the program that semantically allows an expression of type array, should semantically allow NULL instead. For instance,

```
ARRAY IntArray EQ int[]
void F(IntArray A){ PrintInt(A[8]); }
void main(){ F(NULL); }
```

Table 4: NULL sent instead of an integer array is semantically allowed.

### 3.2 Asignments

Assigning an expression to a variable is clearly legal whenever the two have the same type. In addition, if class Son is derived from class Father, then any place in the program that semantically allows an expression of type Father, should semantically allow an expression of type Son. Assigning NULL to array and class variables is legal, but assigning NULL to string variables is illegal. Last, Table 5 summarizes these facts

Binary operations are performed only between integers and strings. A variable is a recursive entity, represented as a sub tree in the AST. Assigning a value to a variable, Table 6 summarizes the context free grammar of RioMare. You will need to feed this grammar to CUP, and make sure there are no shift-reduce conflicts.

# 4 Input

The input for this exercise is a single text file, the input RioMare program.

```
CLASS Father
{
    int foo() { return 8; }
                                   OK
}
Father f := NULL;
CLASS Father { ... }
CLASS Son EXTENDS Father { ... }
void foo(Father f) { ... }
                                   OK
Father f := new Son;
foo(new Son);
ARRAY gradesArray = int[];
ARRAY IDsArray
                 = int[];
                                   ERROR
IDsArray := new int[8];
gradesArray := new int[8];
gradesArray := IDsArray;
string s := NULL;
                                   ERROR
```

Table 5: Assignments in RioMare.

# 5 Output

The output is a *single* text file that contains a *single* word. Either OK when the input program is correct semantically, or otherwise ERROR(*location*), where *location* is the line number of the *first* error that was encountered.

## 6 Submission Guidelines

The skeleton code for this exercise resides (as usual) in subdirectory EX3 of the course repository. COMPILATION/EX3 should contain a makefile building your source files to a runnable jar file called COMPILER (note the lack of the .jar suffix). Feel free to use the makefile supplied in the course repository, or write a new one if you want to. Before you submit, make sure that your exercise compiles and runs on the school server: nova.cs.tau.ac.il. This is the formal running environment of the course.

**Execution parameters** compiler receives 2 input file names:

InputRioMareProgram.txt OutputStatus.txt

```
Program
                  dec^+
           ::=
                   funcDec | varDec | classDec | arrayDec
\operatorname{dec}
            ::=
varDec
            ::=
                  ID ID [ ASSIGN exp ] ';'
                  ID ID '(' [ ID ID [ ',' ID ID ]* ] ')' '{' stmt [ stmt ]* '}'
\operatorname{funcDec}
                  CLASS ID [ EXTENDS ID ] '{' cField [ cField ]* '}'
{\it classDec}
            ::=
                   ARRAY ID = ID '['\ ']'
arrayDec
           ::=
exp
            ::=
                   var
                  '(' exp ')'
                 exp BINOP exp
                  [ var '.' ] ID '(' [ exp [ ',' exp ]* ] ')'
                  ['-'] INT | NIL | STRING | NEW ID | NEW ID '[' exp ']'
var
            ::=
                   var '.' ID
            ::=
                   var '[' exp ']'
            ::=
\operatorname{stmt}
            ::=
                  varDec
                  var ASSIGN exp ';'
            ::= RETURN [exp]';
            ::= IF'(' \exp')''\{' \text{ stmt } [ \text{ stmt } ]^*'\}'
                  WHILE '(' exp ')' '{' stmt [ stmt ]* '}'
            ::=
                  [var '.'] ID '(' [exp [',' exp ]*]')' ';'
                  varDec | funcDec
cField
                  + | - | * | / | < | > | =
BINOP
INT
                   [1-9][0-9]^* \mid 0
```

Table 6: Context free grammar for the RioMare programming language.