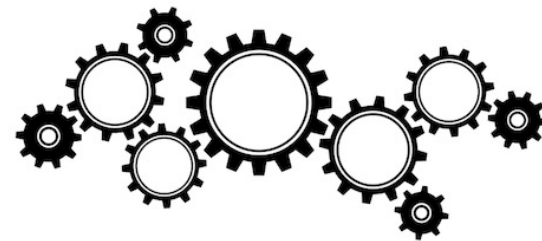
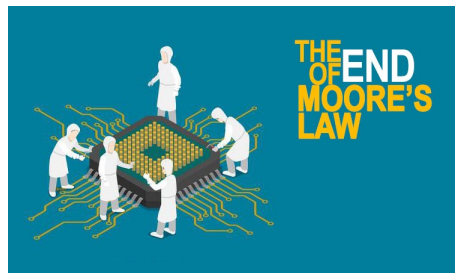


# Innovating IT Infrastructure with Computational Storage – Can you believe (or achieve) the hype?

Tong Zhang, Chief Scientist @ ScaleFlux

# The Rise of Computational Storage



Homogeneous Computing


Heterogenous Computing

Compute 

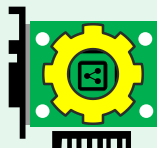
FPGA/GPU/TPU



End of Moore's Law

Networking 

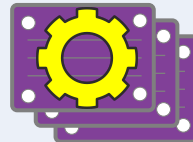
SmartNICs



10 → 100-400Gb/s

Storage 

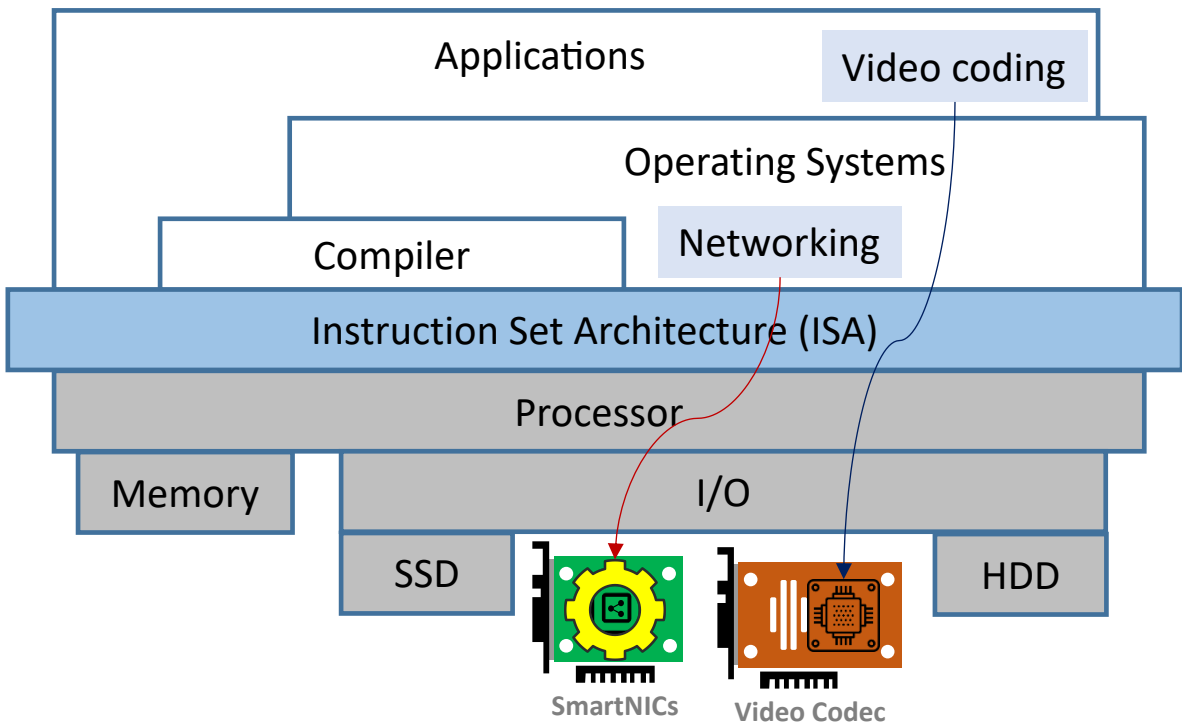
Computational Storage



Fast & Big Data Growth

Domain Specific Compute

# Beauty of Abstraction



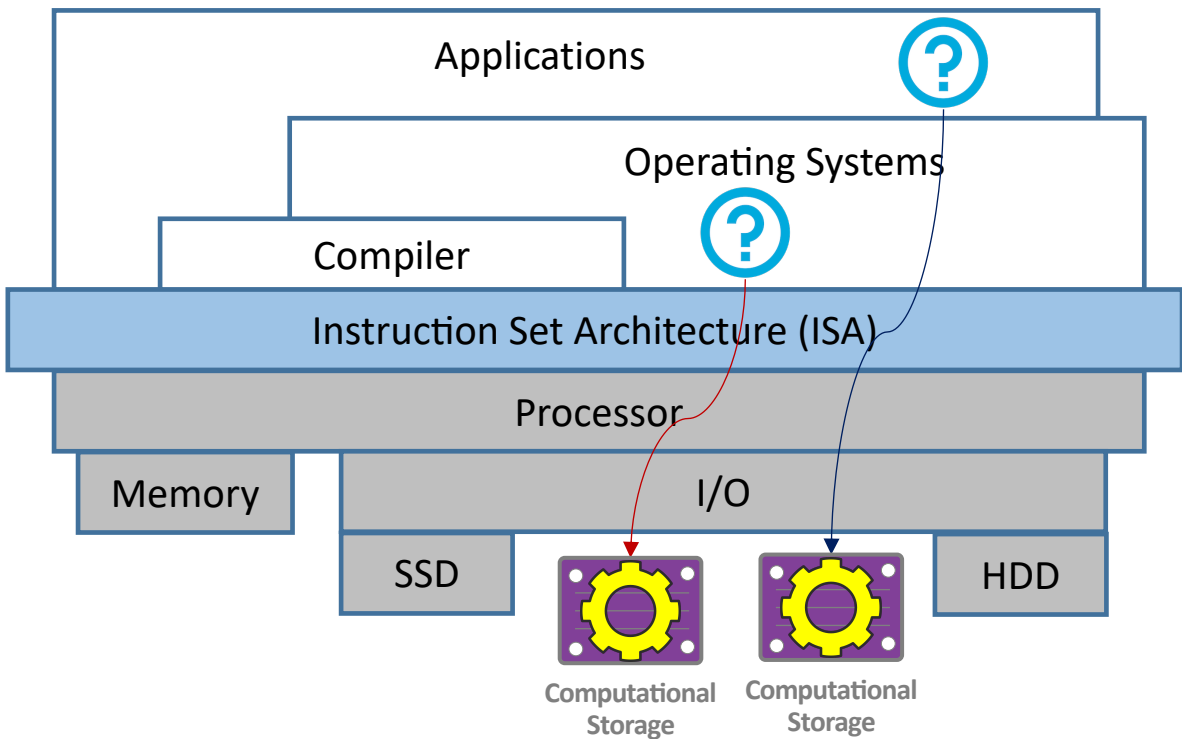
Break the principle of abstraction



- x **Cost:** Modify one (or more) layers
- x **Cost:** Enhance cross-layer interface
- x **Cost:** Vendor lock-in
- ✓ **Benefit:** Improve the system performance/efficiency



# Beauty of Abstraction



Break the principle of abstraction

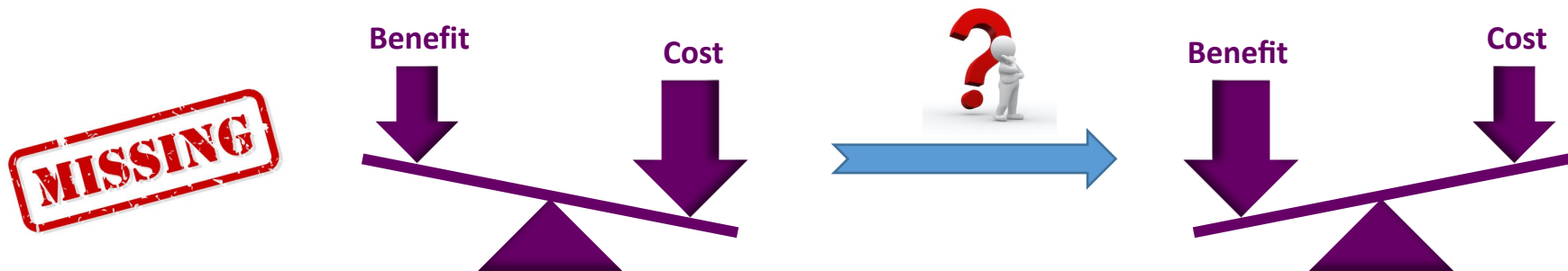


- x **Cost**: Modify one (or more) layers
- x **Cost**: Enhance cross-layer interface
- x **Cost**: Vendor lock-in
- ✓ **Benefit**: Improve the system performance/efficiency



# Computational Storage

- ❑ Academia started to explore this simple concept since 20 years ago
  - Coined many terms: “Intelligent RAM”, “Active Disk”, “Intelligent SSD”, ...
  - Not surprisingly, very impressive performance gains were demonstrated **on papers** for applications such as linear algebra, multimedia/graph processing, database, ...



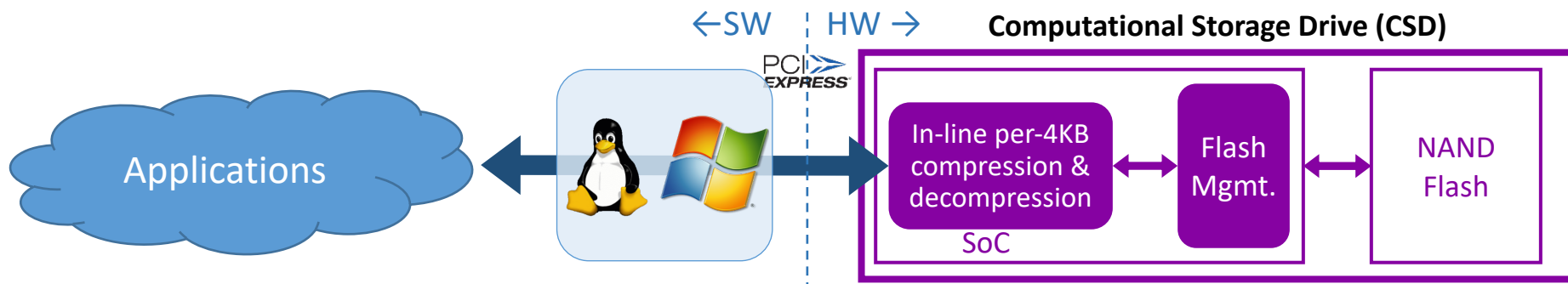
- x **Cost:** Modify one (or more) layers
- x **Cost:** Enhance cross-layer interface
- x **Cost:** Vendor lock-in

**VS**

- ✓ **Benefit:** Improve the system performance/efficiency

# Computational Storage

- ❑ A perfect low-hanging fruit: In-storage **transparent** compression

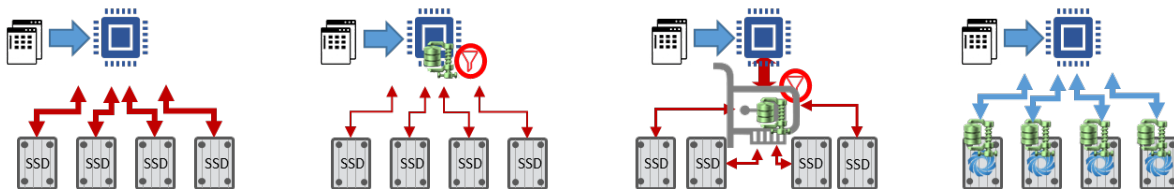


- ~~x **Cost:** Modify one (or more) layers~~
- ~~x **Cost:** Enhance cross layer interface~~
- ~~x **Cost:** Vendor lock in~~

**VS**

- ✓ **Benefit:** Improve the system performance/efficiency

# Comparing Compression Options



	No Compression	Host-Based	Offload Card	CSD
No CPU Overhead	✓	✗	✓	✓
Reduced \$/User GB	✗	✓	✓	✓
Performance scales with capacity	✓	✗	✗	✓
Transparent App Integration	-	✗	✗	✓
Zero App Latency	✓	✗	✗	✓
No incremental power usage	✓	✗	✗	✓
No incremental physical footprint	✓	✓	✗	✓

**Scalable CSD-based compression reduces Cost/GB without choking the CPU**

# Seamless Deployment

- Zero changes to application source code
- >2x storage cost reduction at zero host CPU usage
- Representative use cases: Relational database (e.g., MySQL, Postgre, Oracle), key-value store (e.g., Aerospike), data analytics (e.g., ClickHouse, Apache Kudu)

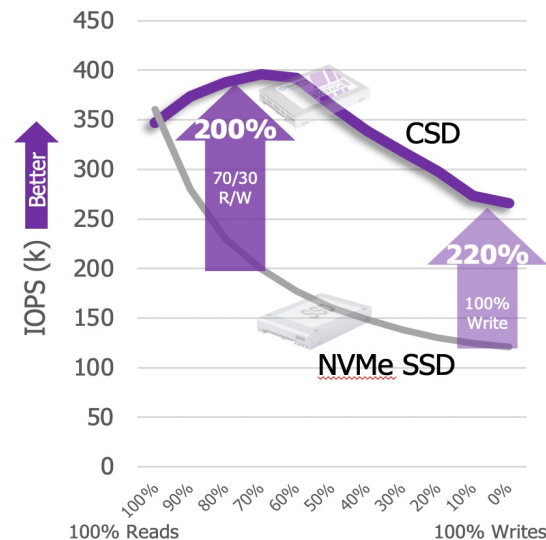


- ❑ >50% storage cost saving
- ❑ 50% TPS improvement
- ❑ 38% latency reduction



- ❑ >50% storage cost saving
- ❑ 35% TPS improvement
- ❑ 20% latency reduction

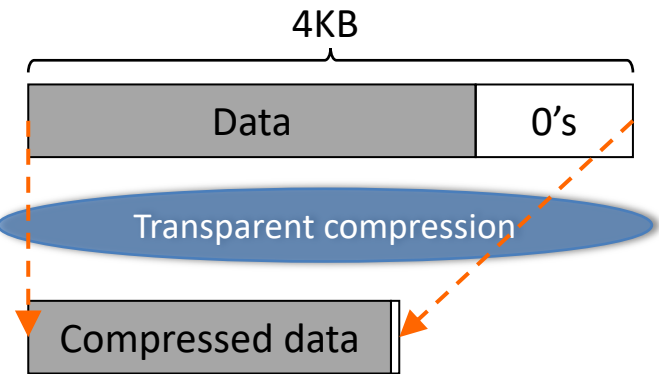
FIO: 8K Random R/W IOPS



- ❑ >50% storage cost saving
- ❑ 50% TPS improvement
- ❑ 98% tail latency reduction



# System-level Innovations Enabled by Transparent Compression



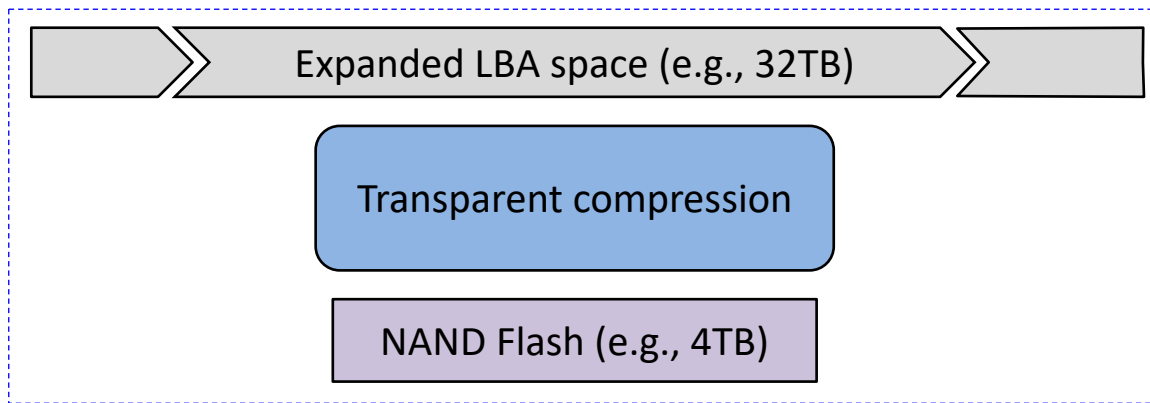
**Unnecessary** to completely fill each 4KB sector with user data



OS/Applications can employ **sparse** data structure without sacrificing storage cost



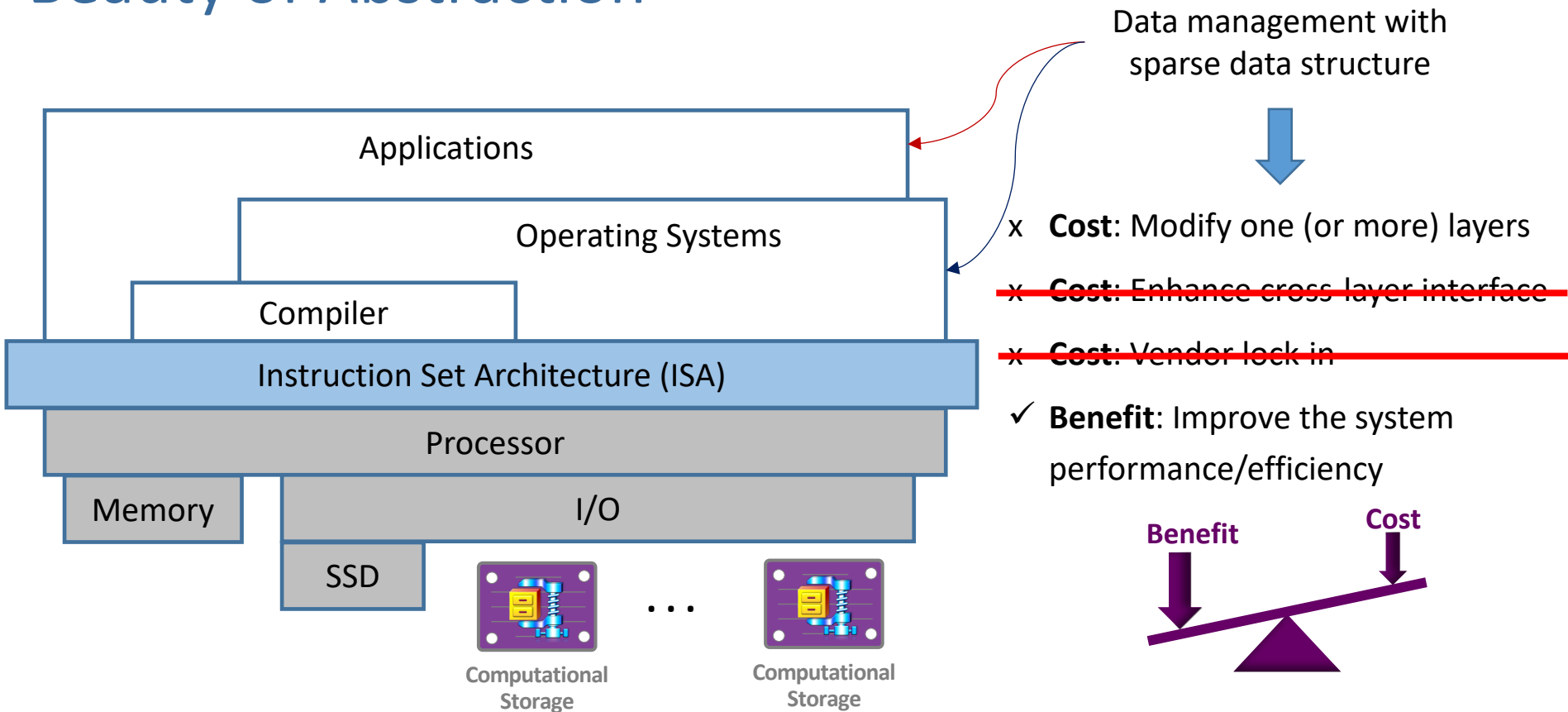
Sparse data structure → cost/performance benefits



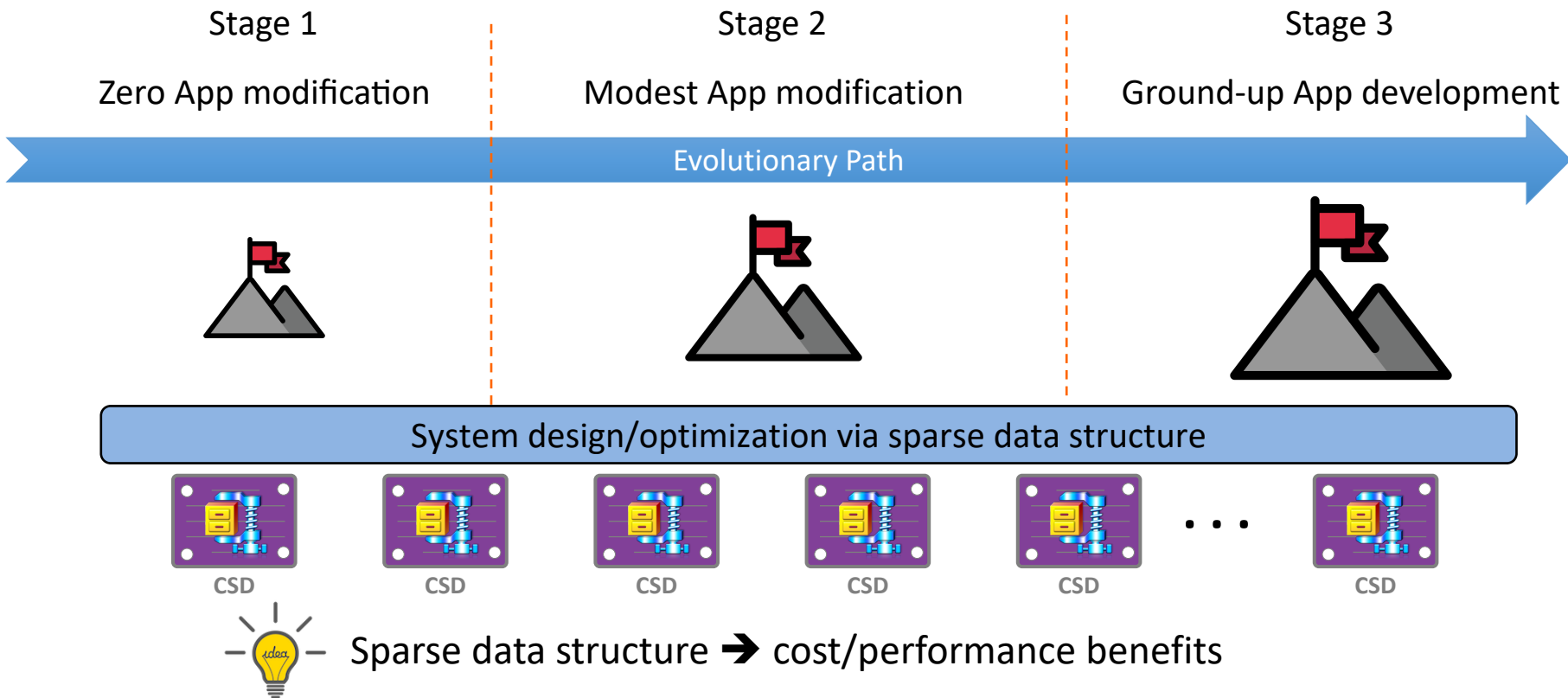
**Unnecessary** to use all the LBAs



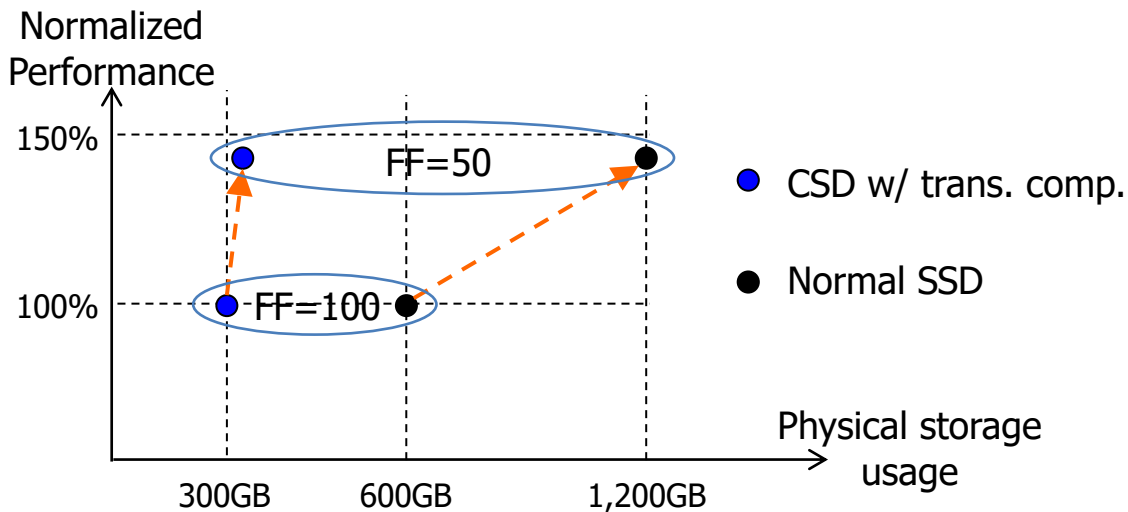
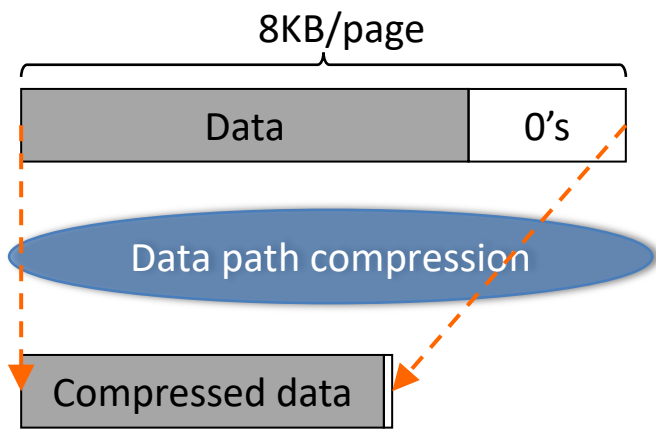
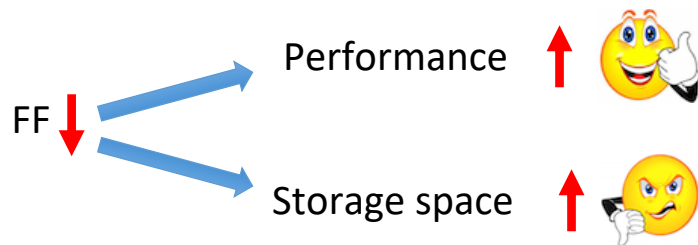
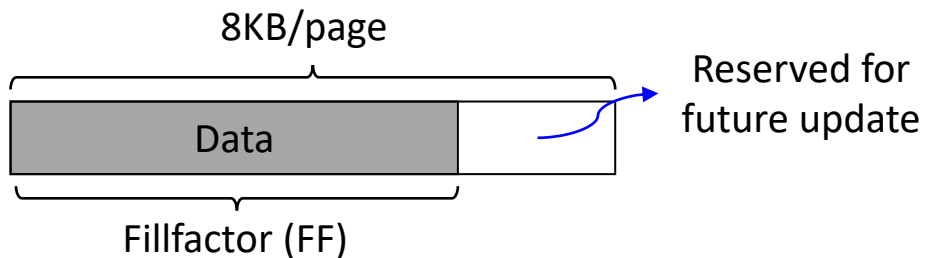
# Beauty of Abstraction



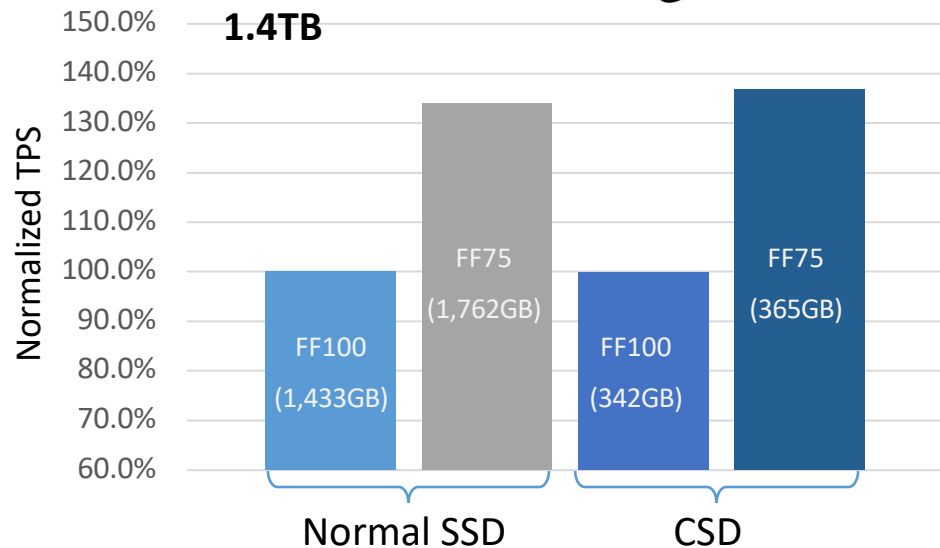
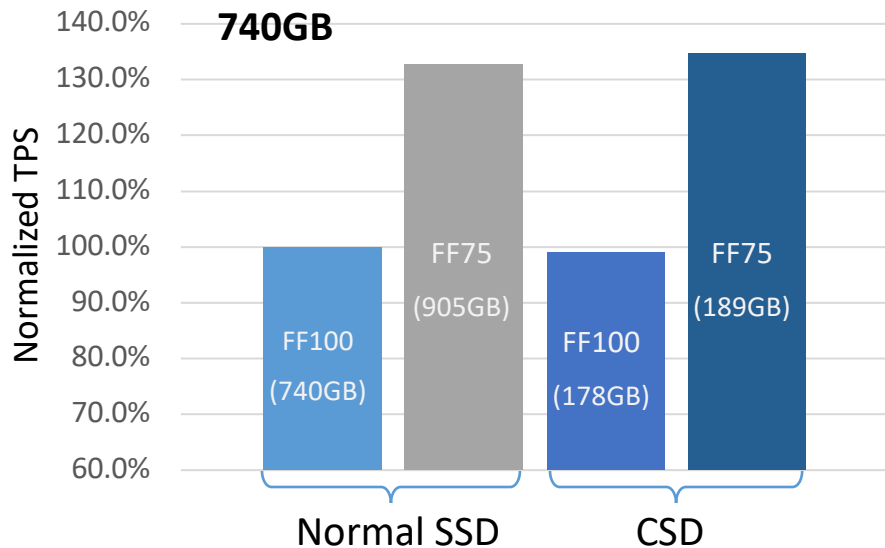
# System-level Innovations Enabled by Transparent Compression



# Stage-1 Example: PostgreSQL



# Stage-1 Example: Experiments (Sysbench-TPCC)



Fillfactor	Logical size (GB)	Drive	Physical size (GB)	Comp Ratio
100	740	Normal SSD	740	1.00
		CSD	178	4.12
75	905	Normal SSD	905	1.00
		CSD	189	4.75

Fillfactor	Logical size (GB)	Drive	Physical size (GB)	Comp Ratio
100	1,433	Normal SSD	1,433	1.00
		CSD	342	4.19
75	1,762	Normal SSD	1,762	1.00
		CSD	365	4.82

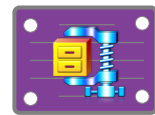
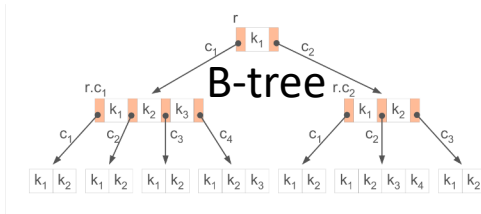
# Stage-2 Example: Reduce B-tree Write Amplification

## ❑ Log-structured merge tree (LSM-tree)

- Significant recent interest: RocksDB, Cassandra, HBase, ScyllaDB, ...

Advantages of LSM-tree

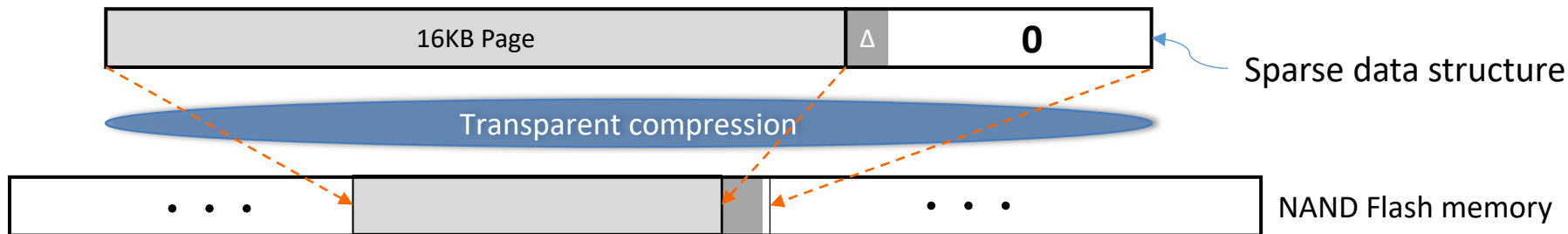
- ~~➤ Less storage space usage~~
- ~~➤ Much smaller write amplification~~



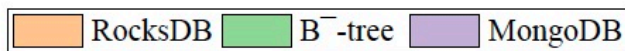
Computational Storage

## ❑ Localized page modification logging

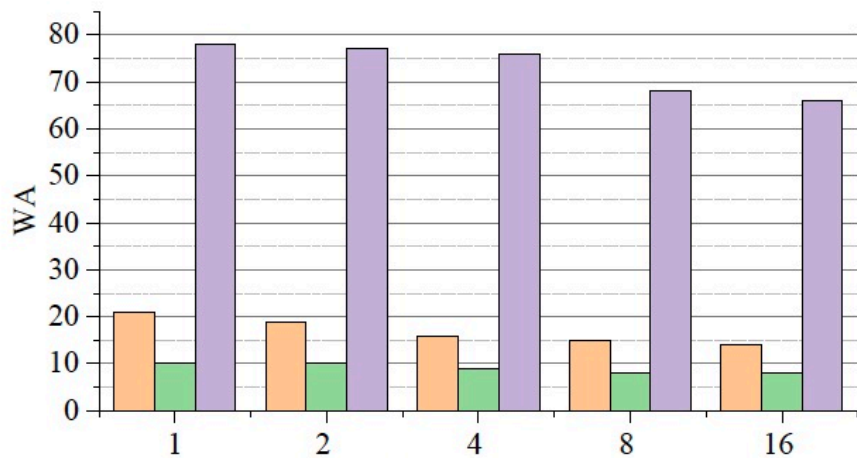
- Very simple idea: Log page modification instead of re-writing the entire page every time



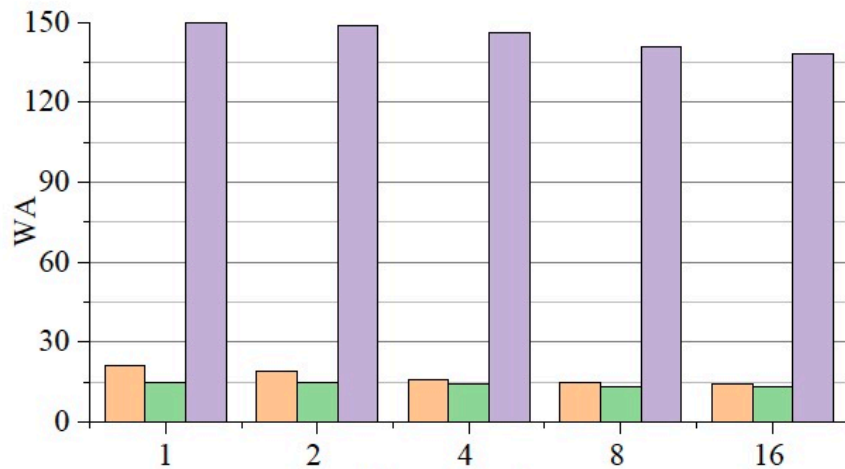
# Stage-2 Example: Reduce B-tree Write Amplification



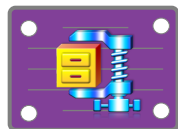
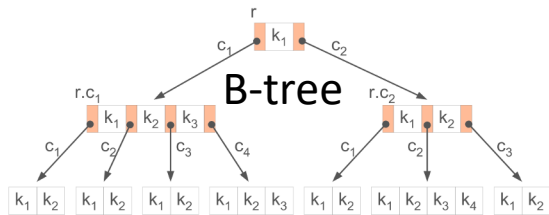
150GB dataset & 1GB cache



(a) 128B record, 8KB page



(b) 128B record, 16KB page



Computational Storage

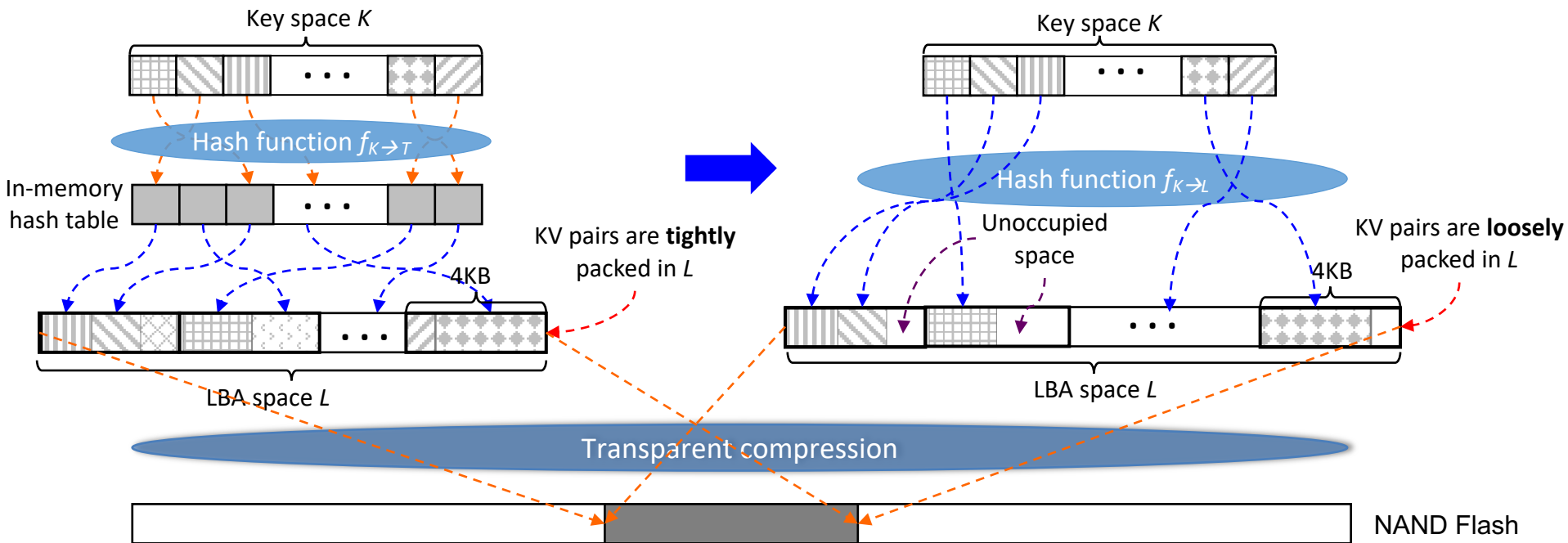


Make LSM-tree (e.g., RocksDB)  
**much less appealing**

# Stage-3 Example: Table-less Hash-based KV Store

## ❑ Very simple idea

- Hash *key space* directly onto *logical storage space* → eliminate the in-memory hash table
- Transparent compression eliminates the “unoccupied space” from physical storage space





# Stage-3 Example: Table-less Hash-based KV Store

## ❑ Eliminate in-memory hash table

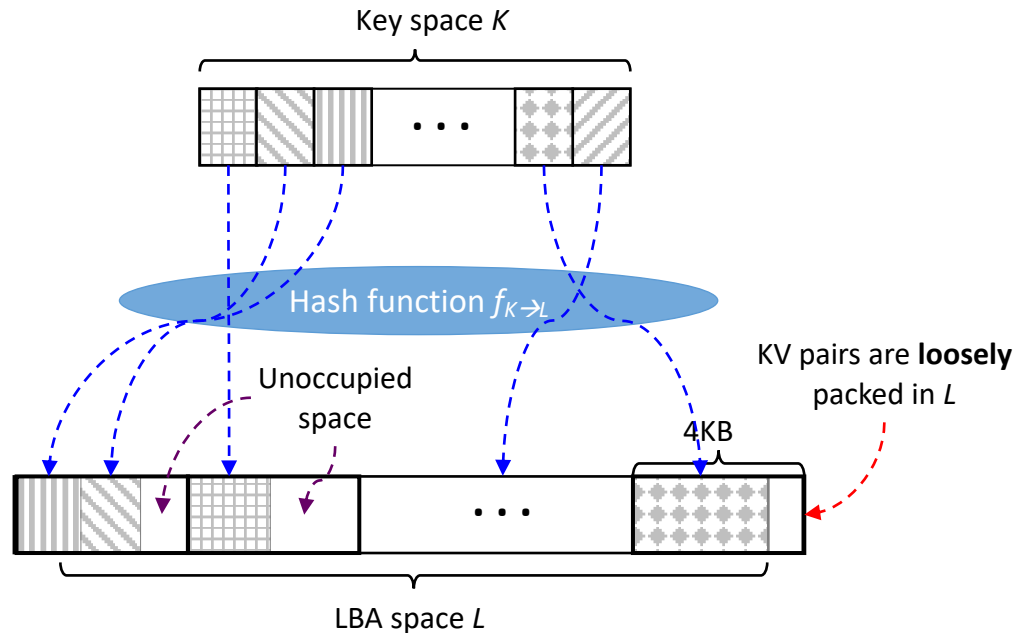
- ✓ Very small memory footprint
- ✓ High operational parallelism
- ✓ Short data access data path
- ✓ Very simple code base

## ❑ Under-utilize logical storage space

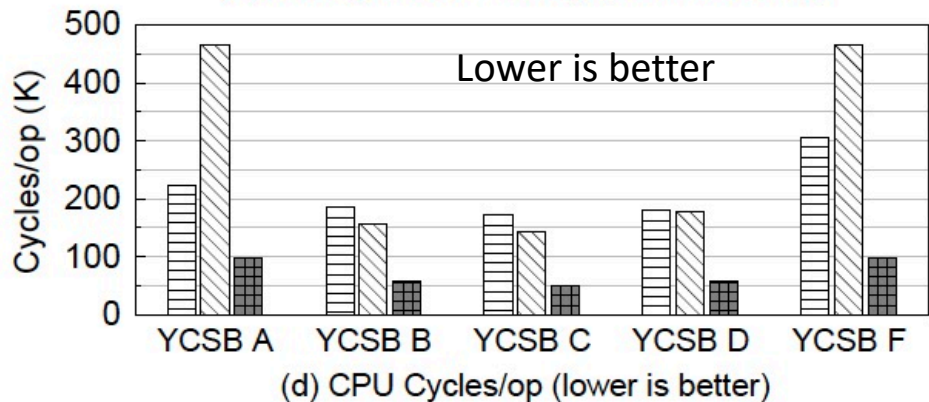
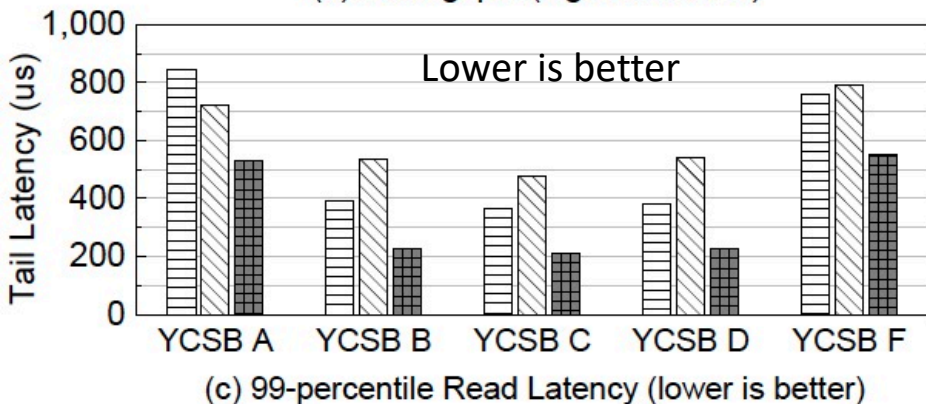
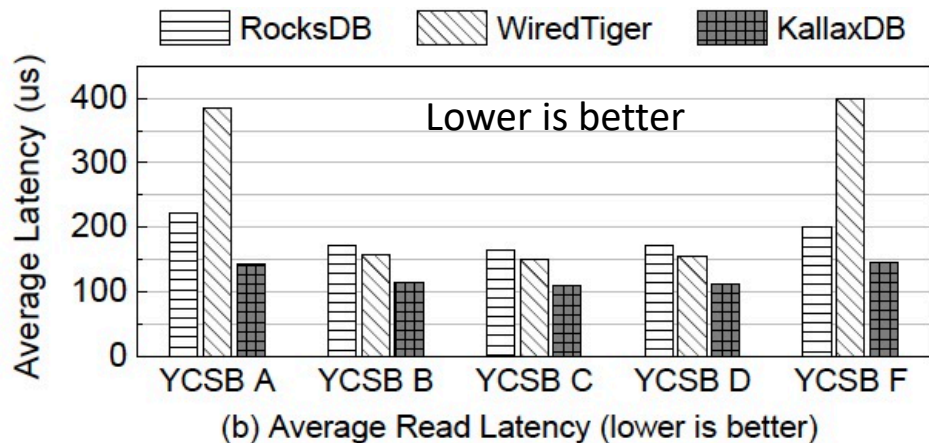
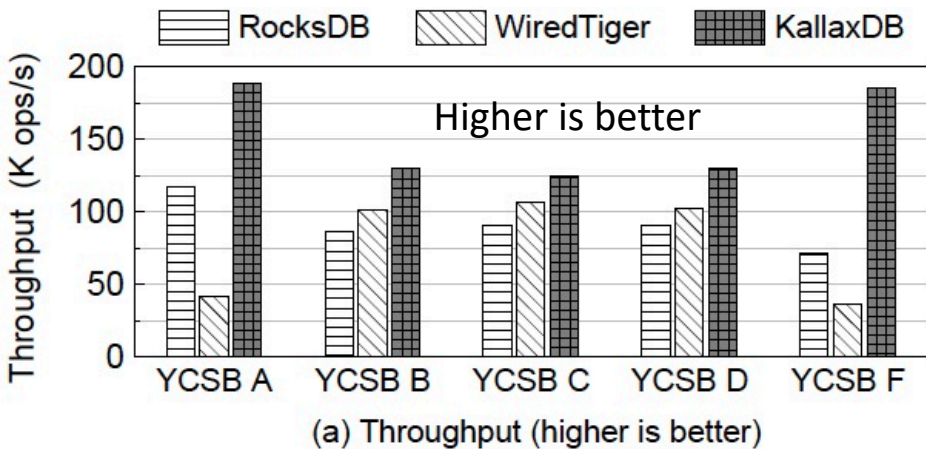
- ✓ Obviate frequent background operations (e.g., GC and compaction)



High performance, low memory cost, and low CPU usage

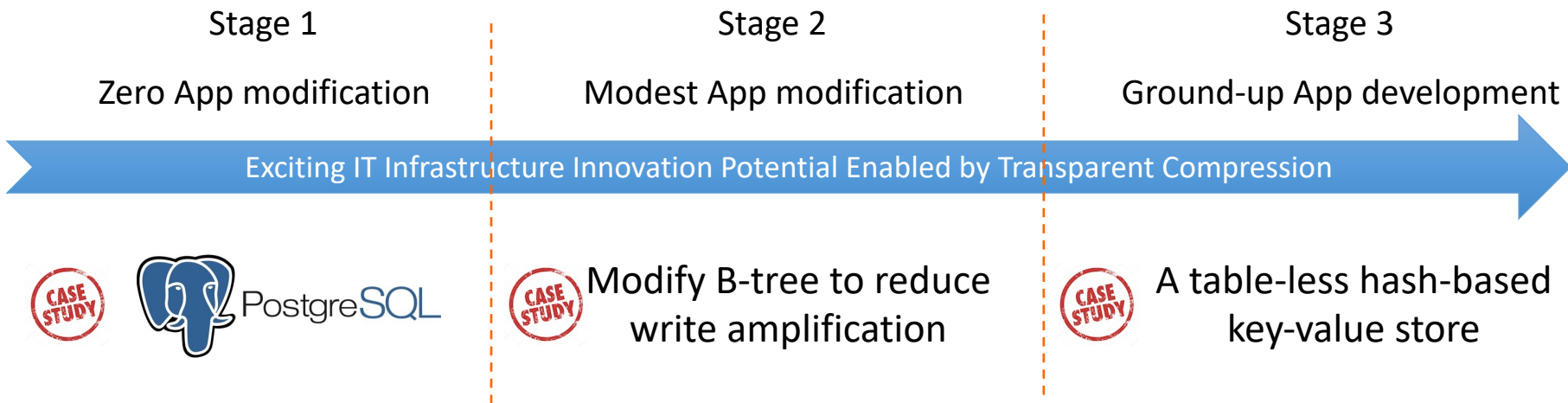


# Stage-3 Example: Results (400-byte KV, 400GB dataset)



# Conclusion

- ❑ Yes, computational storage indeed has a real potential to innovate IT infrastructure
- ❑ The best starting point: Computational storage drive with transparent compression



- ❑ Beyond performance gain, two additional HUGE benefits to end customers:
  - ✓ Avoid modifying interface across applications ↔ OS/filesystem ↔ driver ↔ HW
  - ✓ Avoid computational storage drive vendor lock-in