#### Data Mining and Web Algorithms

- Why web search algorithms (information Retrieval systems)
- Different Retrieval Systems
  - Boolean Retrieval
  - Ranked Retrieval
  - Link Analysis
  - Ranking algorithms
- Web Crawlers
- Web Caching Algorithms
- Recommendation Systems

## Data Mining and Web Algorithm - Web Search engines are actually a specific type of information retrieval systems. > Information Retrieval (FR) is finding material (usually documents) of an instructured ratio (usually for) that satisfier an information need from wither large collections." -> In general, on IR combe 7 E-mail . Starch - searching your laptor - Lyal knowledge rebeval - Audio and Visual Retieval Systems - Nieb Search Systems Usually IR confdes instructioned data / semi-Strictured data. Data Streducol Un-structural Some-strectural (Data base Table) (twitter, Comments) ( ISOM, e-mail messages) (XML)

-> An IR Problem !-

En We wond to search the plays of shakespax's which contains the words Brutu and Causar

and not Calpunia.

IR Models:- Apply "gep" Command of unix and

Toption I:- Apply "gep" Command of unix and

you will get the answer.

Proble: but for larged set of documents and different bund of query such as "near countymen", it is with grep.

Banie Tems: (i) Corpun: a large reporting of documents
(ii) Information need (query): A topic about
which I want to called the dates
(information)

(iii) Revence: Some of the documents in confun that may contain what I want to search.

Option ? - Boolean Refrard Model :-

- Each downent or guny is treated an a"bag" of words or terms". Word sequence is not considered. Fiven a collection of documents D, let (5)  $V = L + 1.11_1 - - + t_1 p_1^2$  be the distinctor words!

terms in the collection. V is Called

the vocabulary.

A weight wij 70 is a provated with each km tio of document dy ED. For a firm that does not appear in document dy, wij = 0.

dy = (wij 1 wzj 1 - w 1 v 1 j)

dy = (wij 1 wzj 1 - w 1 v 1 j)

dzument '

matrix Es

		Julin	1 Pho	Hamlet	Orthello
Irtems	Almony S Cleopatra	(atsar (D2)	Tempert (D3)	(64)	(01-)
	1	1	0	0	0
Antony		1	0	1	0
Bruten	1		0	1	1
Caesar	1	- 1	-		
alapuma	0	1	0	.0	0
mercy	1	0	1	1	1
0					

Ques .- Brutu AND CAESAR AND HOT CALPURNIA Solute: 11010 AND 11011 AND 10111 => 100100 [Thi inplies DI and D4)

Problem :- Building a term-document matrix in above fastion has too many zero and its . Here matrix in externely space on it has few non-zero entre. (Menoy Wastage Problem)

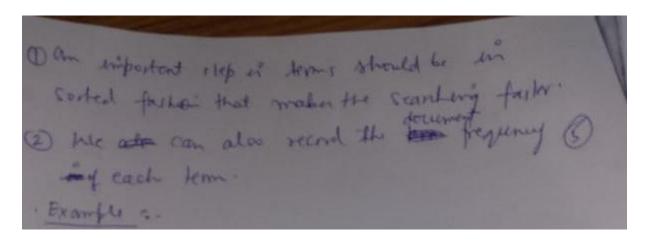
Problem 2: - Retrieval results on usually pour belowse tem-fequency is not considered!

oblion 33 - Invested Index:

- An indra always maps back from tems to the parts of the document where they occur.
- It is a stemdard term in IR or web search.
- -> An inverted wider han two therigs
  - (i) Dictionary (Vocabulary): Set of tems
  - (ii) Posting list ! for each fem, we have a list that records which documents the term occur in , (which is Called posting). This lost is

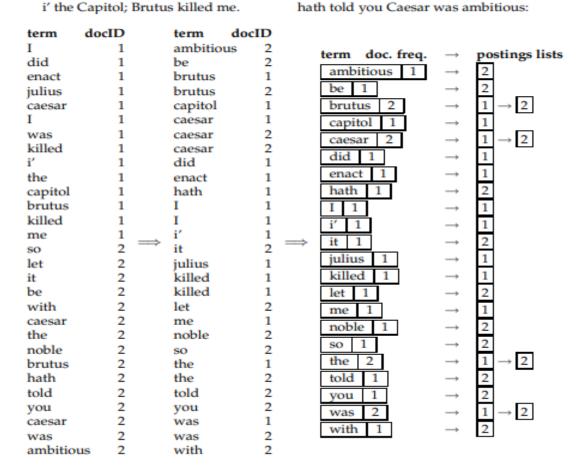
Called posting hist (inverted list). Dictionary in best som - memory dist and portings -1/1/2/4 and kept standisk Calpurnia -> [2] ---

Didtonam



Doc 1 I did enact Julius Caesar: I was killed

Doc 2 So let it be with Caesar. The noble Brutus



Incressing Boolean Busines: D: Brutu AND Calparmy

(1) locate Brutu in Dictionary

(2) Rehere its posting:

(3) horate calparma in the Dicharany

(4) Rehere its posting;

(5) Introcest them [(1) & 4)]

Brutu: 1124

Calparma: 2

: quy result is 2 (Dournat 2 has Both the ten)

#### **Problems with Boolean retrieval models**

- 1. Works for Boolean queries that match or do not to match.
- 2. Most of the users are not capable to write Boolean queries.
- 3. In the case of large document collections, the resulting number of matching documents can far exceed the number a human user could possibly shift through.
  - Most users don't want to wade through 1000s of results.
  - This is particularly true of web search.

Accordingly, it is essential for a search engine to **rank** the documents matching a query. To do this, the search engine computes, for each matching document, a score with respect to the query at hand.

- 4. Boolean queries using AND gives very few and OR gives a lot. With a ranked list of documents it does not matter how large the retrieved set is.
- 5. Queries in web search is usually free text.

#### Ranked retrieval

# Binary term-document incidence matrix

	Antony and Cleopatra	Julius Caesar	The Tempest	Hamlet	Othello	Macbeth
Antony	1	1	0	0	0	1
Brutus	1	1	0	1	0	0
Caesar	1	1	0	1	1	1
Calpurnia	0	1	0	0	0	0
Cleopatra	1	0	0	0	0	0
mercy	1	0	1	1	1	1
worser	1	0	1	1	1	0

Each document is represented by a binary vector  $\in \{0,1\}^{|\mathcal{N}|}$ 

It doesn't consider term frequency (how many times a term occurs) in document.

#### Term-document count matrices

Consider the number of occurrences of a term in a document:

	Antony and Cleopatra	Julius Caesar	The Tempest	Hamlet	Othello	Macbeth
Antony	157	73	0	0	0	0
Brutus	4	157	0	1	0	0
Caesar	232	227	0	2	1	1
Calpurnia	0	10	0	0	0	0
Cleopatra	57	0	0	0	0	0
mercy	2	0	3	5	5	1
worser	2	0	1	1	1	0

#### Vector Space Model

- 1. Documents are treated as bag of words model.
  - John is quicker than Mary and Mary is quicker than John have the same vectors
  - This is called the bag of words model.
- 2. Each document is represented as a vector. Let the vector derived from document d, with one component in the vector for each dictionary term is denoted by V(d). The set of documents in a collection then may be viewed as a set of vectors in a vector space, in which there is one axis for each term.
- 3. However, the term weights are no longer 0 or 1. Each term weight is computed based on some variations of **term-frequency(tf)** or **tf-idf** (**inverse document frequency**) scheme.

#### Term frequency tf

- A document that mentions a query term more often has more to do with that query and therefore should receive more score.
- In this scheme each term in a document are assigned a weight depending on the number of occurrences of term in the document.
- Computing Score b/w a query term 't' and a document 'd', is based on the weight of 't' in 'd'.
- The simplest approach is to assign the weight to be equal to the number of occurrences of term 't' in document 'd'.

#### Term frequency tf

- We want to use tf when computing query-document match scores. But how?
- Raw term frequency is not what we want:
  - A document with 10 occurrences of the term is more relevant than a document with 1 occurrence of the term.
  - But not 10 times more relevant.
- Relevance does not increase proportionally with term frequency.

## Log-frequency weighting

The log frequency weight of term t in d is

$$w_{t,d} = \begin{cases} 1 + \log_{10} \mathsf{tf}_{t,d}, & \text{if } \mathsf{tf}_{t,d} > 0 \\ 0, & \text{otherwise} \end{cases}$$

- $0 \to 0, 1 \to 1, 2 \to 1.3, 10 \to 2, 1000 \to 4$ , etc.
- Score for a document-query pair: sum over terms t in both q and d:

$$score = \sum_{t \in q \cap d} (1 + \log tf_{t,d})$$

 The score is 0 if none of the query terms is present in the document.

#### Document frequency (df)

- Frequent terms are less informative than rare terms
- Consider a query term that is frequent in the collection (e.g., high, increase, line)
- A document containing such a term is more likely to be relevant than a document that doesn't
- But it's not a sure indicator of relevance.
- For frequent terms, we want high positive weights for words like high, increase, and line
- But lower weights than for rare terms.
- We will use document frequency (df) to capture this.

## idf weight

- df<sub>t</sub> is the <u>document</u> frequency of t: the number of documents that contain t
  - df, is an inverse measure of the informativeness of t
  - $df_t \leq N$
- We define the idf (inverse document frequency) of t
   by
   AC 1 2 (N/AC)

 $idf_t = \log_{10} (N/df_t)$ 

• We use log (N/df<sub>t</sub>) instead of N/df<sub>t</sub> to "dampen" the effect of idf.

Will turn out the base of the log is immaterial.

#### tf-idf weighting

 The tf-idf weighting scheme assigns to term 't' a weight in document 'd' given by

$$tf - idf_{t,d} = (1 + \log tf_{t,d}) \times \log_{10}(N/df_t)$$
 [1]

- Highest when 't' occurs many times within a small number of documents.
- Lower when the term occurs fewer times in a document, or occurs in many documents
- Lowest when the term occurs in virtually all documents.

#### The vectors space model for scoring

 To compensate for the effect of document length, the standard way of quantifying the similarity between two documents d1 and d2 is to compute the cosine similarity of their vector representations V (d1) and V (d2)

$$\mathrm{sim}(d_1,d_2) = \frac{\vec{V}(d_1) \cdot \vec{V}(d_2)}{|\vec{V}(d_1)||\vec{V}(d_2)|}, \quad ----(3)$$

- where the numerator represents the dot product of the vectors V (d1) and V (d2) and,
- The denominator is the product of their Euclidean lengths.

## The vectors space model for scoring

• The dot product  $x \cdot y$  of two vectors is defined as

$$\sum_{i=1}^{M} x_i y_i$$
.

Let V (d) denote the document vector for d, with M components  $V_1(d)$  . . . $V_M(d)$ . The Euclidean length of d is defined to be

 $\sqrt{\sum_{i=1}^{M} \vec{V}_{i}^{2}(d)}$ .

We can then rewrite equation (3) as

$$sim(d_1, d_2) = \vec{v}(d_1) \cdot \vec{v}(d_2).$$
 ----(4)

#### **Queries as vectors**

- we can also view a query as a vector.
- we can use the cosine similarity between the query vector and a document vector as a measure of the score of the document for that query.
- The resulting scores can then be used to select the topscoring documents for a query. Thus we have

$$score(q, d) = \frac{\vec{V}(q) \cdot \vec{V}(d)}{|\vec{V}(q)||\vec{V}(d)|}$$

#### **Example**

## TERM VECTOR MODEL BASED ON wi = tfi\*IDFi

Query, Q: "gold silver truck"

D₁: "Shipment of gold damaged in a fire"

D2: "Delivery of silver arrived in a silver truck"

D<sub>3</sub>: "Shipment of gold arrived in a truck" D = 3; IDF = log(D/df<sub>i</sub>)

		Co	unts	, tf <sub>i</sub>				Weights, w <sub>i</sub> = tf <sub>i</sub> *IDF <sub>i</sub>			
Terms	Q	$D_1$	D <sub>2</sub>	D <sub>3</sub>	dfi	D/df <sub>i</sub>	IDFi	Q	$D_1$	$D_2$	D <sub>3</sub>
а	0	1	1	1	3	3/3 = 1	0	0	0	0	0
arrived	0	0	1	1	2	3/2 = 1.5	0.1761	0	0	0.1761	0.1761
damaged	0	1	0	0	1	3/1 = 3	0.4771	0	0.4771	0	0
delivery	0	0	1	0	1	3/1 = 3	0.4771	0	0	0.4771	0
fire	0	1	0	0	1	3/1 = 3	0.4771	0	0.4771	0	0
gold	1	1	0	1	2	3/2 = 1.5	0.1761	0.1761	0.1761	0	0.1761
in	0	1	1	1	3	3/3 = 1	0	0	0	0	0
of	0	1	1	1	3	3/3 = 1	0	0	0	0	0
silver	1	0	2	0	1	3/1 = 3	0.4771	0.4771	0	0.9542	0
shipment	0	1	0	1	2	3/2 = 1.5	0.1761	0	0.1761	0	0.1761
truck	1	0	1	1	2	3/2 = 1.5	0.1761	0.1761	0	0.1761	0.1761

### **Similarity Analysis**

$$\begin{aligned} |D_1| &= \sqrt{0.4771^2 + 0.4771^2 + 0.1761^2 + 0.1761^2} = \sqrt{0.5173} = 0.7192 \\ |D_2| &= \sqrt{0.1761^2 + 0.4771^2 + 0.9542^2 + 0.1761^2} = \sqrt{1.2001} = 1.0955 \\ |D_3| &= \sqrt{0.1761^2 + 0.1761^2 + 0.1761^2 + 0.1761^2} = \sqrt{0.1240} = 0.3522 \end{aligned}$$

$$\therefore |\mathbf{D}_i| = \sqrt{\sum_i w_{i,j}^2}$$

$$|Q| = \sqrt{0.1761^2 + 0.4771^2 + 0.1761^2} = \sqrt{0.2896} = 0.5382$$

$$\therefore |Q| = \sqrt{\sum_{i} w_{Q, j}^2}$$

Next, we compute all dot products (zero products ignored)

$$Q \bullet D_1 = 0.1761 * 0.1761 = 0.0310$$

$$Q \bullet D_2 = 0.4771 * 0.9542 + 0.1761 * 0.1761 = 0.4862$$

$$Q \bullet D_3 = 0.1761 * 0.1761 + 0.1761 * 0.1761 = 0.0620$$

$$\therefore Q \bullet D_i = \sum_i w_{Q,j} w_{i,j}$$

Now we calculate the similarity values

Cosine 
$$\theta_{D_1} = \frac{Q \bullet D_1}{|Q|^* |D_1|} = \frac{0.0310}{0.5382 * 0.7192} = 0.0801$$

Cosine 
$$\theta_{D_2} = \frac{Q \bullet D_2}{|Q|^* |D_2|} = \frac{0.4862}{0.5382 *1.0955} = 0.8246$$

Cosine 
$$\theta_{D_3} = \frac{Q \bullet D_3}{|Q| * |D_3|} = \frac{0.0620}{0.5382 * 0.3522} = 0.3271$$

$$\therefore$$
 Cosine  $\theta_{D_i} = Sim(Q, D_i)$ 

$$\therefore \operatorname{Sim}(Q, D_i) = \frac{\sum_{i}^{W} Q_{i,j} W_{i,j}}{\sqrt{\sum_{j}^{W} W_{Q,j}^2} \sqrt{\sum_{i}^{W} W_{i,j}^2}}$$

 Finally we sort and rank the documents in descending order according to the similarity values

Rank 1: Doc 2 = 0.8246

Rank 2: Doc 3 = 0.3271

Rank 3: Doc 1 = 0.0801

GIrpmation Retrieval (document Retrieval) invertigate relevant document in a small and trusted set. (Eg newspaper ash In , patents etc) -> Brutin bleb Search, bleb is huge and full of Unbrusted documents, random things, web span etc - Heb can be thought of an a graph consiling of o Modin: - hlebpager o Edgu: - Hyperlinks Page de Hypertine Mode B - Web challings:--(i) heb contains many rouse of information. Who to toust? Trick: - Trustworthy pages may point to each other. [Hyperlinks] - (ii) What is but answer to query "hews paper - No single answer

- Hence, by looking at web-streeting

We can rank the web pager.

Link analysi: 
To analysi: 
do link analysis. for link analysis.

I Different Appraches in this lourse are 
o Page Fank

O Hubs and authorities [HITS]

Pegikank r.- [

- Pegikank algorithm determines the importance of

Web pages based on link structure

- It is a probability distribution used to

represent the likelihood that a peron

randomly clicking on links will arrive

at any particular page.

- Central part of Croogle's Search engine in

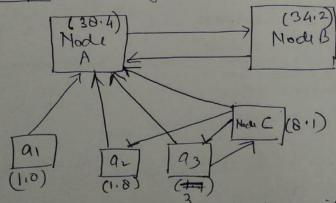
laggerank developed by Larry Page and

Sergey Brin at standford University.

Bari Idea for Page Pank

- Consder links as votes
  - o lage in more important if it ton more links.
  - o Whether incorning outgoing links.
  - o Iscoming likes as actually more unifortent and conveys that that webpage is more tousted!
  - o Are all the links are same? It defends on the importance of other pages incoming to it.

Example: - Web lager on a graph with lage Ranks



1) Mode A har high page Rank on it has more incoming a likes.

2) But Hode B has also higher rank as compared to Hode a Which also has single incorning link (formag) as it has getting link (thicoming) from Mode A. 3) Hence nit is recurire. Simple Recurrire fromulation can be done by making their statements - Each links vot (Rome) is proportional to importance of its source Page. -7 21 page j' with importance orgo tan noutlinks, each, link gets my - Hence, Page jois own & 7,0 = 3p + 1B importance is the sum of the votes on its links. -> A "vole" from an important page is more worth. -> A page is important if it is possited by other important page. Hence rank of page j'" can be defined as ry° = \( \frac{\sigma\_i^o}{\dio} \), where dio => outdegree

of rock i° regularly on any page other that
there going to page that it points to

C A (B)
(B2)

Vier com go dirately to "c" rather than going to BI &BZ pomited by web page A.

and hence include some probability of going to page of that page (going to page of and page (going to page of and hence we can finulate the following equation.

Organal Paghonk algorithm ! --> 2t in an iterative algorithm that storts with green. PR(A) = (1-d) + d [ PR(TI) + PR(T2) PEN)

Where, Hypobability for rondom sufer jumping to a page

> PR(A) = Page Rank of Page A vialrays (1-d). -> PR(Ti) = Pagi Rank of Pager Ti which like to page A ( informed link ) -> C(Ti) = No of outboard links on page Ti. -> d = demping factor lier BIW 0 81. - According to lage and Brim (who have proported) d = 0.85. Eg No-of Nodu = 3 Let un assure inlital are 1: PR(A) = (1-d) + d[ PR(C) ((C) = (015) + (015) [ RACO] PR(B) = (1-a) + a \[ \frac{PR(A)}{C(A)} \] = (1-05) + (01) PR(A)

$$PR(c) = (0.5) + (0.5) \left[ \frac{PR(A)}{C(A)} + \frac{PR(B)}{C(B)} \right]$$

$$= 0.5 + 0.5 \left[ \frac{PR(A)}{2} + \frac{PR(B)}{C(B)} \right] - (3)$$
Let
$$PR(A) = P(B) = P(C) = 1 \text{ initial } \text{ foth initials}$$

$$PR(A) = 0.5 + 0.5 = 1$$

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Sum of Pagetark of all the page = Mocof web pager = 1+1+5+ 1.125 = 2.07533 It has to repeated (iteratively) until we get the Same (or almost same) value of page Pank.

Another notation of pagerank no: -

$$PR(A) = \frac{(1-d)}{N} + d \left[ \frac{PR(T)}{C(T)} + \frac{PR(T)}{C(T)} - \frac{PR(T)}{C(T)} \right]$$

Then sum of all pages will be one.

#### **References:**

- [1] Manning, Christopher D., Prabhakar Raghavan, and Hinrich Schütze. Introduction to information retrieval. Cambridge university press, 2008.
- [2] Leskovec, Jure, Anand Rajaraman, and Jeffrey David Ullman. Mining of massive data sets. Cambridge university press, 2020.