## Homework 1

**Problem 1.** Prove IP2 by induction on the property  $Q(x) = {}^{\alpha}P(k)$  holds for all k < x."

*Proof.* Note Q(0) is vacuously true since there are no natural numbers less than 0. Suppose Q(n) is true. Then P(k) holds for all k < n. Now if P(n) is true, then P(k) holds for all k < n + 1 and thus Q(n + 1) is true. Then Q holds for any natural number n which means P holds for the same set of numbers. This proves IP2 holds.

**Problem 2.** Prove that the relation <, as we defined it in class, is transitive on  $\mathbb{N}$ . That is, show that for all  $k, m, n \in \mathbb{N}$ , if k < m and m < n then k < n.

*Proof.* Let P(n) be the statment "For all k < m and m < n, k < n". Note that P(0) is vacuously true since there are no natural numbers less than 0. Suppose P(n) is true. Choose k < n + 1 and m < k. If m < n then we know k < n by our inductive hypothesis. It remains to show the case k = n. Suppose k = n and m < n. Then  $m \subseteq n$ . But note that  $n + 1 = n \cup \{n\}$ . Thus  $m \subseteq n + 1$  and so m < n + 1. By the principle of induction, P(n) holds for all  $n \in \mathbb{N}$ .

**Problem 3.** Prove that  $(\mathbb{N}, <)$  is a well ordered set.

*Proof.* Let  $A \subseteq \mathbb{N}$  be a subset with no least element. Let  $B = \mathbb{N} \setminus A$  be the set of natural numbers not in A. Note that  $0 \in B$  because 0 is less than every natural number and so it would be the least element of A. Also, if  $n \in B$ , then  $n+1 \in B$  as well, otherwise n+1 would be a least element of A. But then  $B = \mathbb{N}$  and so  $A = \emptyset$ . Thus all nonempty subsets of  $\mathbb{N}$  have least elements.

**Problem 4.** Prove that there is no function  $f: \mathbb{N} \to \mathbb{N}$  such that for all  $n \in \mathbb{N}$ , f(n) > f(n+1).

*Proof.* Suppose such a function f exists. We can show that each element of  $f(\mathbb{N})$  is distinct. Suppose f(k) < f(n) for all k < n. Then note that f(n) < f(n+1) and by the transitivity of < on  $\mathbb{N}$ , we have f(k) < f(n+1) for each k < n+1. Thus by the second version of the induction principle, we know every element of  $f(\mathbb{N})$  is distinct.

Let f(0) = k. Note that there are only k natural numbers less than k. Consider f(k+1). We know  $f(k+1) < f(k), f(k-1), \ldots, f(1), f(0)$ . Since each of  $f(0), f(1), \ldots, f(k)$  is distinct, by the pigeon hole principle, one of these must be equal to f(k+1). This is a contradiction and so f cannot exist.  $\square$ 

**Problem 5.** Verify that the definition we gave in class for  $\vDash$  is unambiguous for each wff A.

*Proof.* Let S be a set of sentence symbols and  $M \subseteq S$  be a model. First suppose that A has length 1. Then either  $A \in M$  or  $A \notin M$  and so either  $M \models A$  or  $M \not\models A$ . Now suppose that for all wffs B of length n, we have  $M \models B$ . Then by definition  $M \not\models A = (\neg(B))$ . Also, if C is a wff of length n, then  $M \models A = ((B) \land (C))$ . Since these are the only two ways of making a wff, by the principle of induction there is no ambiguity in the symbol  $\models$  for any wff A.