

Space Complexity

Input space : num1, num2 : 2 units

Auxilliary space : res : 1 unit

Total space = Input space + Auxilliary space

= 2 + 1

= 3 units

Irrespective of type and value of data, space required is constant space (k)

Total space \propto constant space

Total space $\propto k$

Total space $\propto K * 1$

Total space = 1

To indicate complexities, we need to use one notataion ie 'order of' notation $o()$

Space complexity = $o(1)$

```
Algorithm to find sum of two variable
int sum(int num1, int num2)
{
    int res = num1 + num2;
    return res;
}
```

Space Complexity

```
Algorithm to find sum of array elements
int sumofarray(int arr[], int size)
{
    int sum = 0;
    for(int i = 0 ; i < size ; i++)
        sum += arr[i]
    return sum;
}
```

size = n
Input space : n unit
Auxilliary space : size, sum, i : 3 units

Total space \propto Input space + Auxilliray

Total space $\propto n + 3$

if $n \gg \gg \gg \gg \gg 3$

Total space $\propto n$

Space complexity = $o(n)$

Space complexity : $o(1)$, $o(n)$, $o(n^2)$, $o(n^3)$,.....

```
Algorithm to display matrix
void display_matrix(int mat[][], int row, int col)
{
    int i,j;
    for(i = 0 ; i < row ; i++)
    {
        for(j = 0 ; j < col ; j++)
            printf("%d", mat[i][j]);
        printf("\n");
    }
}
```

size $n \times n$

Input space : n^2 units
Auxilliray space : row, col, i, j : 4 units

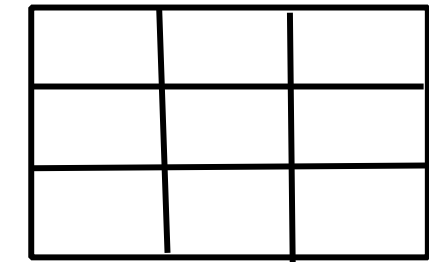
Total space \propto Input space + Auxilliary space
Total space $\propto n^2 + 4$

if $n \gg \gg \gg \gg$

Total space $\propto n^2$

Space complexity = $o(n^2)$

Space Complexity



$3 \times 3 = 9 = 9$ units

size $m \times n$

Input space : $m * n$ units
Auxilliray space : row, col, i, j : 4 units

Total space \propto Input space + Auxilliary space
Total space $\propto m * n + 4$

if $m, n \gg \gg \gg \gg$

Total space $\propto m * n$

Space complexity = $o(m * n)$

Time Complexity

```
statement;
```

Time = 1 unit

Time complexity = $O(1)$

```
for(int i = 0 ; i < n ; i++)  
{  
    statements;  
}
```

Condition will be checked = $n + 1$ times
statement executes = n times

Total time = $n + 1 + n$

Total time = $2n + 1$

if $n \gg \gg \gg \gg \gg$

Total time = $2n$

Total time $\propto n$

Time complexity = $O(n)$

```
for(int i = n ; i >= 0 ; i--)  
{  
    statements;  
}
```

Time Complexity

```
for(int i = 0 ; i < n ; i++)  
{  
    for(int j = 0 ; j < n ; j++)  
    {  
        statements  
    }  
}
```

Outer condition checked : n times
Inner condition checked : n * n times
statements will execute : n * n times

Total time = $n + n^2 + n^2$
 = $n + 2n^2$
 = $2n^2$

Total time $\propto n^2$

Time complexity = $O(n^2)$

Time Complexity

```
for(int i = n ; i > 0 ; i=i/2)
{
    statement;
}
```

$i = n, n/2, n/4, n/8, \dots$

$i = n/2^0, n/2^1, n/2^2, n/2^3, \dots, n/2^{\text{itr}}$

for $i=1$, last time condition will be true

$$n/2^{\text{itr}} = 1$$

$$2^{\text{itr}} = n$$

$$\log 2^{\text{itr}} = \log n$$

$$\text{itr} \log 2 = \log n$$

$$\text{itr} = \log n / \log 2$$

$$\text{itr} = (1/\log 2) \log n$$

$$\text{itr} \propto \log n$$

Time complexity = $O(\log n)$

Time complexity : $O(\log \log n), O(\log n), O(n \log n), O(1), O(n), O(n^2), O(n^3), \dots$