



# **Sunbeam Institute of Information Technology**

## **Pune and Karad**

### **PreCAT**

## **Module – Data Structures**

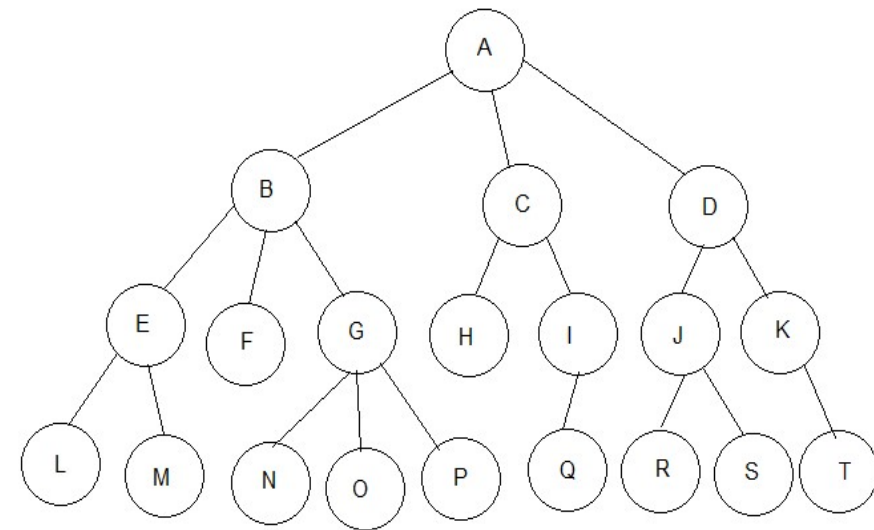
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# Tree : Terminologies

- **Tree** is a **non linear** data structure in which one specially designated node is called as “**root**”.
- **Root** is a **starting point** of the tree.
- Remaining elements can be partitioned into  $m$  disjoint subsets where each of subset is a tree.
- All elements are connected in **Hierarchical manner**.
- Every element of a tree is called as **node** of the tree.
- **Parent node**:- having other child nodes connected
- **Child node**:- immediate descendant of a node
- **Leaf node**:-
  - Terminal node of the tree.
  - Leaf node does not have child nodes.
- **Ancestors**:- all nodes in the path from root to that node.
- **Descendants**:- all nodes accessible from the given node
- **Siblings**:- child nodes of the same parent



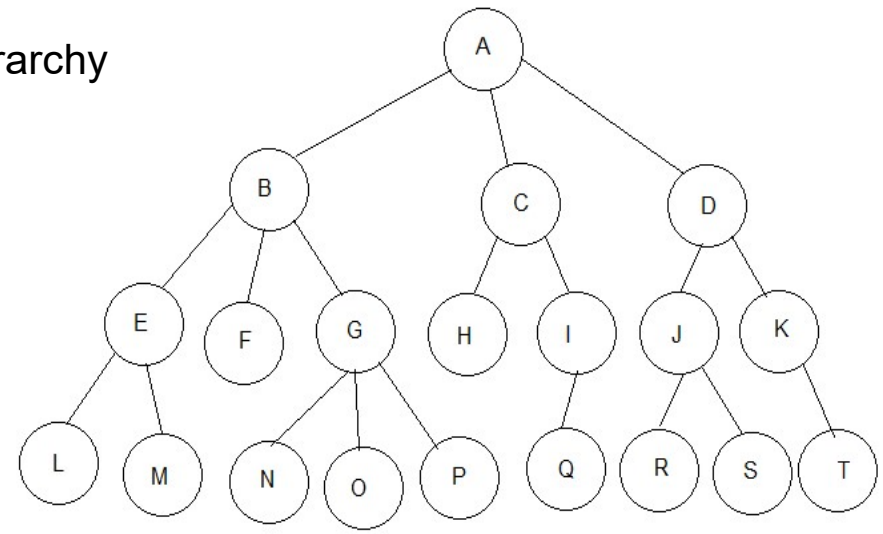
# Tree : Terminologies

- **Degree of a node** :- number of child nodes for any given node.
- **Degree of a tree** :- Maximum degree of a child in tree.
- **Level of a node** :- indicates position of the node in tree hierarchy
  - Level of child = Level of parent + 1
  - Level of root = 1

- **Height of node** :- level of given node
- **Depth of node** :- level of node - 1
- **Height of a tree** :- Maximum height of a node
- **Depth of a tree** :- Maximum depth of a node
- Tree with zero nodes (ie empty tree) is called as

“**Null tree**”. Height of Null tree is 0.

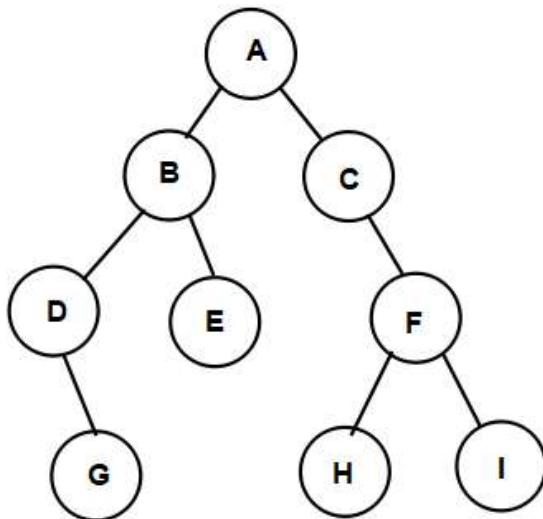
- Tree can grow up to any level and any node can have any number of Childs.
- That's why operations on tree becomes un efficient.
- Restrictions can be applied on it to achieve efficiency and hence there are different types of trees.



# Tree : Types

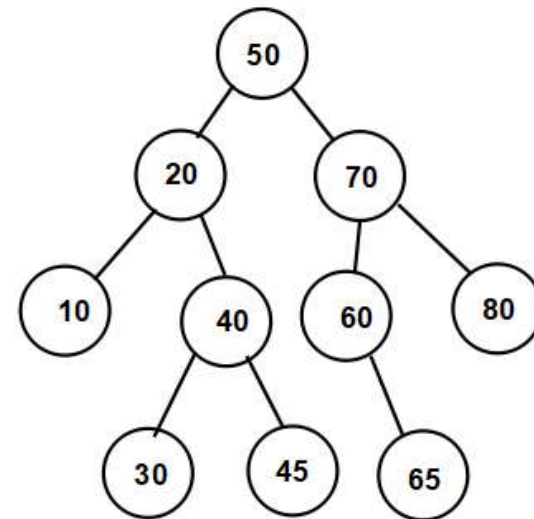
- **Binary Tree**

- Tree in which each node has maximum two child nodes
- Binary tree has degree 2. Hence it is also called as 2- tree



- **Binary Search Tree**

- Binary tree in which left child node is always smaller and right child node is always greater or equal to the parent node.
- Searching is faster
- Time complexity :  $O(h)$        $h$  – height of tree



# Tree : Types

- **Complete Binary Tree**

- Binary tree in which all leaf nodes are at same level.

- **Strictly Binary Tree**

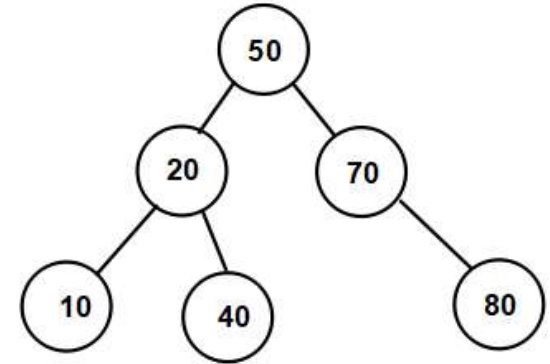
- Binary tree in which each non leaf node has exact two child nodes.

- **Full Binary Tree**

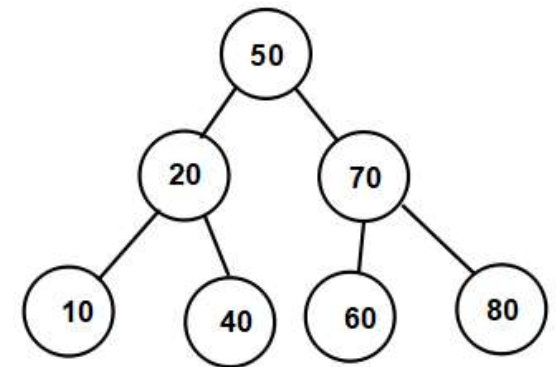
- Binary tree with its full capacity for the given height.
- In other words, adding one more node will increase height of the tree.
- It is always complete as well as strictly binary tree.
- Number of elements =  $2^h - 1$

- **Almost Complete Binary Tree**

- The binary tree which follows two conditions
  - All leaf nodes are at level  $h$  or  $h-1$ .
  - All leaf nodes at last level ( $h$ ) are aligned to left as much as possible.



Complete Binary Tree



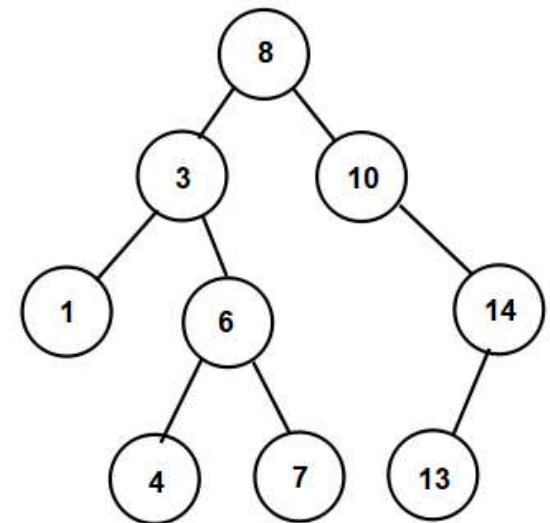
Strictly Binary Tree



# Binary Search Tree : Traversal

- **Pre-Order:-** V L R
- **In-order:-** L V R
- **Post-Order:-** L R V
- The traversal algorithms can be implemented easily using recursion.
- Non-recursive algorithms for implementing traversal needs stack to store node pointers.

- **Pre-Order :-** 8 3 1 6 4 7 10 14 13
- **In-Order :-** 1 3 4 6 7 8 10 13 14
- **Post-Order :-** 1 4 7 6 3 13 14 10 8



# BST - Types

- **Skewed Binary tree**

- In binary tree if only left or right links are used, then tree grows in only one direction and such tree is called as skewed binary tree.
  - Left skewed binary tree
  - Right skewed binary tree
- Tree has maximum height that is same as number of elements
- Time complexity of searching is  $O(n)$ . (Like linked list)

- **Balanced BST**

- If nodes in BST are arranged so that its height is kept as less as possible, is called as balanced BST.
- Balance Factor = Height of left subtree – Height of right subtree
- In balanced BST, balance factor of each node is -1, 0 or +1
- A tree can be balanced by applying series of left or right rotations.

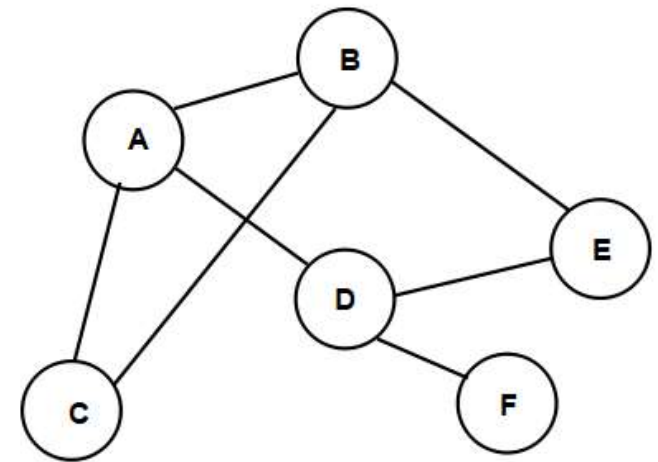
- **AVL Tree**

- AVL tree is a self balancing binary search tree.
- Node are balanced on each insert and delete operation.
- Difference between heights of left and right subtrees can not be more than 1 for all nodes.



# Graph : Terminologies

- **Graph** is a non linear data structure having set of vertices (nodes) and set of edges (arcs).
  - $G = \{V, E\}$   
Where V is a set of vertices and E is a set of edges
  - **Vertex (node)** is an element in the graph  
 $V = \{A, B, C, D, E, F\}$
  - **Edge (arc)** is a line connecting two vertices  
 $E = \{(A,B), (A,C), (B,C), (B,E), (D, E), (D,F), (A,D)\}$
- Vertex A is set be **adjacent** to B, if and only if there is an edge from A to B.
- **Degree of vertex** :- Number of vertices adjacent to given vertex
- **Path** :- Set of edges connecting any two vertices is called as path between those two vertices.
  - Path between A to D =  $\{(A, B), (B, E), (E, D)\}$
- **Cycle** :- Set of edges connecting to a node itself is called as cycle.
  - $\{(A, B), (B, E), (E, D), (D, A)\}$
- **Loop** :- An edge connecting a node to itself is called as loop. Loop is smallest cycle.





# Graph : Types

- **Simple Graph**
  - Graph not having multiple edges between adjacent nodes and no loops.
- **Complete Graph**
  - Simple graph in which node is adjacent with every other node.
  - Number of Edges =  $n(n-1)/2$  where,  $n$  – number of vertices
- **Connected Graph**
  - Simple graph in which there is some path exist between any two vertices
- **Weighted Graph**
  - A graph in which edge is associated with a number (ie weight)
- **Directed Graph (Di-graph)**
  - A graph in which each edge has some direction





Thank you!

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