



Sunbeam Institute of Information Technology
Pune and Karad
PreCAT

Module – Data Structures

Trainer - Devendra Dhande

Email – devendra.dhande@sunbeaminfo.com



Algorithm Analysis

- Analysis is done to determine how much resources it require.
- Resources such as time or space
- There are two measures of doing analysis of any algorithm
 - Space Complexity
 - Unit space to store the data into the memory (Input space) and additional space to process the data (Auxiliary space)
 - e.g. Algorithm to find sum of all array elements.
int arr[n] – n units of input space
sum, index, size – 3 units of auxiliary space
Total space required = input space + auxiliary space = $n + 3 = n$ units
 - Time Complexity
 - Unit time required to complete any algorithm
 - Approximate measure of time required to complete algorithm
 - Depends on loops in the algorithm
 - Also depends on some external factors like type of machine, no of processed running on machine.
 - That's why we can not find exact time complexity.
- Method used to calculate complexities, is “**Asymptotic Analysis**”



Asymptotic Analysis

- It is a mathematical way to calculate complexities of an algorithm.
- It is a study of change in performance of the algorithm, with the change in the order of inputs.
- It is not exact analysis
- Few mathematical notations are used to denote complexities.
- These notations are called as “Asymptotic notations” and are
 - Omega notation (Ω)
 - Represents lower bound of the running algorithm
 - It is used to indicate the best case complexity of an algorithm
 - Big – Oh notation (O)
 - Represents upper bound of the running algorithm
 - It is used to indicate the worst case complexity of an algorithm
 - Theta notation (Θ)
 - Represents upper and lower bound of the running time of an algorithm (tight bound)
 - It is used to indicate the average case complexity of an algorithm



Time Complexity

Statement;

constant

```
for(i=0; i< n; i++)  
{  
    statements;  
}
```

Linear

```
for(i=0; i< n; i++)  
{  
    for(j=0; j< n; j++)  
    {  
        statements;  
    }  
}
```

Quadratic

```
for(i=n; i>0; i/=2)  
{  
    statement  
}
```

Logarithmic



Searching Algorithms : Time Complexity

Linear Search :

	No of Comparisons		Running Time	Time Complexity
Best Case	1	Key found at very first position	$O(1)$	$O(1)$
Average Case	$n/2$	Key found at in between position	$O(n/2) = O(n)$	$O(n)$
Worst Case	n	Key found at last position or not found	$O(n)$	$O(n)$

Binary Search :

	No of Comparisons		Running Time	Time Complexity
Best Case	1	Key found in very first iteration	$O(1)$	$O(1)$
Average Case	$\log n$	Key found at non-leaf position	$O(\log n)$	$O(\log n)$
Worst Case	$\log n$	if either key is not found or key is found at leaf position	$O(\log n)$	$O(\log n)$



Sorting Algorithms : Comparisons

- Selection sort algorithm is too simple, but performs poor and no optimization possible.
- Bubble sort can be improved to reduce number of iterations.
- Insertion sort performs well if number of elements are too less. Good if adding elements and resorting.
- Quick sort is stable if number of elements increased. However worst case performance is poor.
- Merge sort also perform good, but need extra auxiliary space.

Algorithm	Best Case	Average Case	Worst Case
Selection sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble sort	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$





Thank you!

Devendra Dhande

<devendra.dhande@sunbeaminfo.com>

