Functional Design of MirrorMe

Rieks Joosten (rieks.joosten@tno.nl)

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Chapter 1

Introduction

This document¹ defines the functionality of an information system called 'MirrorMe'. It defines the database and the business services of MirrorMe by means of business rules². Those rules are listed in chapter 2, ordered by theme.

The diagnosis in chapter 3 is meant to help the authors identify shortcomings in their Ampersand script.

 $^{^1{\}rm This}$ document was generated at 15-7-2017 on 20:12:38, using Ampersand-v3.8.6 [HEAD:49946cb70], build time: 08-Jul-17 09:11:28 Coordinated Universal Time.

 $^{^2\}mathrm{Rule}$ based design characterizes the Ampers and approach, which has been used to produce this document.

Chapter 2

Shared Language

This chapter describes functional requirements for 'MirrorMe' in natural language. It contains definitions and agreements. The purpose of this chapter is to create shared understanding among stakeholders. All definitions and agreements have been numbered for the sake of traceability.

2.1 Arguments

The structure of legal arguments has been studied by scholars such as Toulmin~[?], Verheij (name some more). To help legal professionals construct such arguments, the structure of arguments must be formalized. For this purpose, we can consider each (legal) case as a set of statements. Each statement can be considered true or

Definition 1: a set of TTexts that are controlled by a single authority and *Scope* (together) serve a particular purpose

If people make claims, they do observations, and they will reason about them. In order to decide whether a claim is true, statements are organised in a structure that represents the argument. The concept "statement" is used to represent claims, observations, and all other utterances that can be considered true or false. Statements and Claims are modeled/implemented as TTexts.

Definition 2: the collection of one template string, one instance thereof, and *TText* one value within some scope

Let us treat a statement made by an individual as a claim of that individual that the statement is true. For this reason, each individual who makes the claim can be registered in the relation "claimant". The claimant 'claims' (a) the validity of the ttTemplate and (b) the veracity of ttValue. Thus, the claimaint is solely authorized to change ttValue and ttTemplate.

- **Agreement 1** A claimant is the party that makes a claim, especially one that is legally cognizable.
- **Agreement 2** the Actor that is the claimant of all toplevel TTexts.
- **Agreement 3** The application can register a moment as the time a case has been created.
- **Agreement 4** A statement that needs another statement to be true, is registered in the relation requires.

Case Management 2.2

This pattern contains some basic administration for cases.

Als iemand inlogt in het systeem moet diens 'context' worden geactiveerd, d.w.z. de gegevens over de persoon die het systeem nodig heeft om te kunnen berekenen wat hij/zij wel en niet mag doen, en welke gegevens van het systeem daarbij mogen worden gebruikt. Om zulke computations te kunnen maken wordt een aantal zaken geregistreerd en aan één persoon gekoppeld.

Definition 3: een verzameling gegevens die (een deel van) de gebruikerscontext van één gebruiker binnen het systeem beschrijft

Account

Een persoon gebruikt een gegevensruimte (en heet dan 'user') door met een browser (bijv. Chrome of Firefox) het systeem te benaderen dat de gegegensruimte beheert. Als meerdere personen een gegevensruimte delen, moet het systeem de context van elk van hen kunnen onderscheiden, bijvoorbeeld om:

- de interactie 'klein' te houden, d.w.z. alleen gegevens te laten zien die voor hem/haar relevant zijn;
- ervoor te zorgen dat een user niet ziet wat hij niet mag zien;
- te kunnen bijhouden welke persoon, of welk(e) organisatie(onderdeel) verantwoordelijk is voor een zekere transactie;
- automatisch gegevens betreffende de user of zijn context aan transacties toe te kunnen voegen

We gebruiken de term 'SESSION' of 'session' om de verzameling van gegevens betreffende één (actieve) user mee aan te geven. Deze term correspondeert met de gelijknamige term browsers gebruiken, nl. 'een verbinding (door de browser) met een webservice (die een URL heeft)'. Het systeem houdt één session bij voor elke actieve user, d.w.z. voor elke browser die het systeem benadert. Merk op dat dit in het bijzonder geldt als de user in verschillende tabbladen van dezelfde browser het systeem benadert - er is dan toch maar één session (en één user).

Definition 4: een verzameling van gegevens die de context beschrijven waarin één persoon het systeem gebruikt

SESSION

Agreement 5 A SESSION may be linked to an Account, which specifies the (user)context of the session.

Placeholder Extraction and TText Hierar-2.3 chy creation Service

This service ensures that - ttTemplateParsedText=ttTemplate AND ttTemplatePlaceholders contains all TTNames that are mentioned in the ttTemplate of a TText.

The idea is that the specification of a TText is parsed to see if it contains names of TTexts. Such names are recognized by the fact that they are surrounded by square brackets ([and]).

So, - ttTemplateParsedText stores the text that has been parsed (in PHP) to produce the contents of ttTemplatePlaceholders. This implies that whenever ttTemplateParsedText is empty for some TText, ttTemplatePlaceholders is empty for that same atom. - whenever ttTemplateParsedText differs from ttTemplate, - it is first removed, and ttTemplatePlaceholders are discarded - then ttTemplate is being parsed, and ttTemplateParsedText and ttTemplatePlaceholders are filled again.

All of this happens in the same Scope, i.e. in the scope to which the TTexts belong. Note that for this to work, all TTexts that are mentioned in a ttTemplate must exist.

Definition 5: a label used to identify a TText (within a scope)

TTName

Definition 6: a sequence of words

TTPhrase

- **Agreement 6** Statements that are in a given scope are considered valid within that scope.
- **Agreement 7:** TTexts that have no placeholders detected, yet have a ttTemplate, must have been parsed
- **Agreement 8:** Whenever a template phrase changes, the parsed text must be deleted so that the template is parsed again
- **Agreement 9:** TTexts whose template text is not parsed, must not specify placeholders that are detected

2.4 Placeholder Replace-By-TTValue Service

This service ensures that the ttInstance phrase of a TText is the ttTemplate phrase of that TText, where every occurrence of a reference to another TText in that template (i.e. occurences of [<TTName>], called 'placeholder's) is replaced by the ttValue of that other TText, provided that the ttValue is populated for that TText.

Definition 7: (the representation of) a value of a TText

TTValue

The property ttInstanceResetReq is used to start reconstructing an instance phrase of a TText.

- **Agreement 10** If a TText has the property ttInstanceResetReq, this means that its instance phrase needs to be (re)constructed from scratch.
- **Agreement 11** SRC TText has been used to replace placeholders in the TGT TText instance phrase
- **Agreement 12** References to this TText (i.e.: placeholders) exist, and have been replaced with the specified TTValue
- Agreement 13: If a TText has a value, and is used in a TText that is not being re-initialized, then it must appear in the ttValueIsUsedInInstanceOfTText of the TText.
- **Agreement 14:** When the value of a TText differs from the value used in replacements, or has become void, then the instance-texts of all TTexts in which the replacements took place must be discarded and recreated.

2.5 TText Reset Service

This service relinquishes any depencies that the ttInstance of a TText has on other TTexts (as registered in ttValueIsUsedInInstanceOfTText). This condition will be realized by setting the property ttInstanceResetReq for that TText.

Agreement 15: If a TText has a template phrase and no instance phrase, the TText must be reset/initialized, thus allowing its instance phrase to be constructed

Agreement 16: Resetting a TText means that the registration of replaced placeholders (for that TText) must be reset (cleared/deleted).

Before reconstructing an instance phrase of a TText, the ExecEngine must have discarded all administration related to that TText.

Agreement 17: A TText that needs to be (re-)initialized and does not use values of TTexts in its UsedValue, must have the specification-text as its UsedValue, which completes the (re-)initialization.

2.6 Validity

In order to talk about true and false statements in a precise way, we need the idea of contexts.

Consider the statement "John has blond hair". If this statement is known to be true in some context c, the tuple ("John has blond hair", c) can exist in the relation true. However, if this tuple is not in the relation true, it does not follow that John does not have blond hair. The truth of the TText value (as decided by its claimant) is accepted by the owner of the Scope

Agreement 18 A statement that is considered true within a context is registered in the relation true.

The relation false is dual to true.

Agreement 19 A statement that is considered false within a context is registered in the relation false.

For reconstructing events, it can be necessary to administer the moment an observation or a claim is made. For this reason, we use the relation "observed".

Agreement 20 The application can register a moment as the time a statement has been made.

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, given that all its placeholders have been instantiated.

Agreement 21: A statement can be true or false in a context only if it is valid within that context.

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, if all its placeholders have been instantiated.

Agreement 22: A statement cannot be both true and false in the same context.

2.7 Wetsartikelen

Chapter 3

Diagnosis

This chapter provides an analysis of the Ampersand script of 'MirrorMe'. This analysis is intended for the author(s) of this script. It can be used to complete the script or to improve possible flaws.

MirrorMe does not specify which roles may change the contents of which relations.

MirrorMe assigns rules to roles. The following table shows the rules that are being maintained by a given role.

Rule	ExecEngine	SYSTEM
caseAuthor	×	
TText template parsing -	×	
extract placeholders		
TText template parsing -	×	
delete parsed templates		
TText template parsing -	×	
delete placeholders used in		
template		
Compute transitive closure	×	
of 'ttIsUsedBy'		
Create TTexts for	×	
placeholders (if necessary)		
When the name of a TText	×	
is in the list of placeholders		
of another TText, the first		
TText is said to be used by		
the second TText		
A (first) TText is only used	×	
by a (second) TText if the		
name of the first is in the list		
of placeholders of the second		
When a TText is used by	×	
another TText, it inherits		
the latter's Scope		
Register that the value of a	×	
TText has been used to		
replace placeholders		

Rule	ExecEngine	SYSTEM
After a value update, all TTexts that used the value must be reset	×	
When a TText value is NOT used in an instance of (another) TText, then it may not say that its value is used	×	
to replace a placeholder Request TText instance phrase reset if the template phrase does, and the instance phrase does not	×	
exist Resetting a TText implies that the relations with other TTexts in which it is	×	
actually used for replacement, are broken Resetting a TText is complete if other TTexts are not used by this TText	×	

Concepts Scope, Scope, TTPhrase, TTName, TTValue, Actor, Claim, LegalThing, Moment, RechtsGrond, URL, Regeling, Bijlage, Aanwijzing, Artikel, Lid, Sub, BasisWettenBestand, Titel, Aantal, Domein, Regelingsoort, Concept, Processlow, and Tekst remain without a purpose.

```
The relations object[Claim * LegalThing], ttTemplateParsedText[TText * TTPhrase],
ttTemplatePlaceholders[TText*TTName], ttIsUsedBy[TText], ttIsUsedByCopy[TText],
ttIsUsedByStar[TText], url[RechtsGrond*URL], regeling[RechtsGrond*Regeling],
bijlage[RechtsGrond*Bijlage],
                                                                                                                                  aanwijzing[RechtsGrond*Aanwijzing],
artikel[RechtsGrond*Artikel], lid[RechtsGrond*Lid], sub[RechtsGrond*Sub],
bwb[Regeling * BasisWettenBestand], titel[Regeling * Titel], artikelen[Regeling * Aantal],
url[Regeling * URL], gedelegeerdUit[Regeling], regelgeving[Domein * Regeling],
soort[Regeling * Regelingsoort], onderwerp[Regeling * Concept], procesflow[Regeling * Procesflow],
afkorting[Regeling*Tekst], class[TText*Concept], legalGround[TText*LegalGround],
evidence[TText*Evidence], autoLoginAccount[Account], loginUserid[SESSION*UserID],
loginPassword[SESSION * Password],
                                                                                                                                                                                       logoutRequest[SESSION],
isoLevelGTE[ISOLevel],
                                                                                                          sessionReqdISOLevel[SESSION*ISOLevel],
session AuthISOLevel[SESSION*ISOLevel], moduleName[Module*ModuleName], moduleName[ModuleName], moduleName[Modu
module VsnMajor [Module*Module VsnMajor], module VsnMinor [Module*Module VsnMinor], module VsnMinor [Module VsnMinor], module Module [Module VsnMinor], module Module [Module VsnMinor], module [Mo
personOrg[Person*Organization], uidUserid[UID*UserID], uidIssuer[UID*IdP],
siamcompFirstName[SIAMPersonRefComputation*FirstName],
siamcompLastName[SIAMPersonRefComputation*LastName], siamcompPersonRef[SIAMPersonRefComputation*LastName]
loginCreateAccount[SESSION], registerPassword[SESSION*Password],
registerOrgRef[SESSION*OrgRef], loginReq[SESSION], loginISOLevel[SESSION*ISOLevel], loginISOLevel[SESSION*ISOLevel], loginISOLevel[SESSION*ISOLevel], loginISOLevel[SESSION*ISOLevel], loginISOLevel[SESSION*ISOLevel], loginISOLevel[SESSION*ISOLevel[SESSION*ISOLevel]], loginISOLevel[SESSION*ISOLevel[SESSION*ISOLevel]], loginISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel]]], loginISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOLevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*ISOlevel[SESSION*Isolevel[SESSION*ISOlevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Isolevel[SESSION*Iso
sessionDev[SESSION], ttTrace[TText*Assignment], asmVar[Assignment*TText],
                                                                                                                                  asmPOT[Assignment*PointOfTime],
asmVal[Assignment*TTValue],
asmHasPred[Assignment], compVar[Computation * TText], compArg[Computation * Assignment],
compRes[Computation * TTValue],
                                                                                                                                                                 sessionLoginAssist[SESSION],
sessionLogoutAssist[SESSION], sessionSRI[SESSION], sessionSIA[SESSION],
```

```
scope ID[Scope*Scope ID], scope Owner[Scope*Account], scope Descr[Scope*Scope Descr], \\ ttName[TText*TTName], ttValue[TText*TTValue], ttTemplate[TText*TTPhrase], \\ ttInstance[TText*TTPhrase], ttICO[TText], ttICCO[TText], ttDescr[TText*TTPhrase], \\ and ttOwner[TText*Account] all lack both a purpose and a meaning.
```

```
The purpose of relations owner[Scope * Actor], created[Scope * Moment],
                 ttValueUsedToReplacePlaceholders[TText*TTValue],
requires[TText],
ttValueIsUsedInInstanceOfTText[TText],
                                             ttScope[TText * Scope],
accUserid[Account*UserID], accPassword[Account*Password], sessionAccount[SESSION*Account]
accIsInitialized[Account], accIsActive[Account], accDeactivateReg[Account],
sessionUserid[SESSION*UserID], orgAbbrName[Organization*OrgAbbrName],
orgFullName[Organization*OrgFullName], accOrg[Account*Organization],
sessionOrg[SESSION*Organization], personFirstName[Person*FirstName], \\
personLastName[Person*LastName],
                                        accPerson[Account * Person],
sessionPerson[SESSION*Person], accAllowedRoles[Account*Role],
accDefaultRoles[Account*Role], sessionAllowedRoles[SESSION*Role],
sessionActiveRoles[SESSION*Role], accUID[Account*UID], personRef[Person*PersonRef],
accPersonRef[Account*PersonRef], sessionPersonRef[SESSION*PersonRef],
accOrgRef[Account * OrgRef],
                               sessionOrgRef[SESSION*OrgRef],
scopeIII[Scope], and scopeIsaCC[Scope] is not documented.
```

Relations claimant, object, created, requires, ttIsUsedByStar, observed, url, regeling, bijlage, aanwijzing, artikel, lid, sub, bwb, titel, artikelen, url, gedelegeerdUit, regelgeving, soort, onderwerp, procesflow, and afkorting are not used in any rule.

- fig. 3.1 shows a conceptual diagram with all relations declared in 'Arguments'.
- fig. 3.2 shows a conceptual diagram with all relations declared in 'Placeholder Extraction and TText Hierarchy creation Service'.
- fig. 3.3 shows a conceptual diagram with all relations declared in 'Validity'.
- fig. 3.4 shows a conceptual diagram with all relations declared in 'Wetsartikelen'.

Rules are defined without documenting their purpose:

- caseAuthor line 155:5, file C:\Afstuderen\repos\ampersandmodels\MirrorMe\MirrorMe.adl
- $TText\ template\ parsing\ -\ extract\ placeholders$ line 91:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- $TText\ template\ parsing\ -\ delete\ parsed\ templates$ line 103:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- TText template parsing delete placeholders used in template line 109:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- $TTexts\ can\ only\ use\ TTexts\ that\ are\ in\ the\ same\ ttScope$ line 122:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- Compute transitive closure of 'ttIsUsedBy' line 129:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- Create TTexts for placeholders (if necessary)

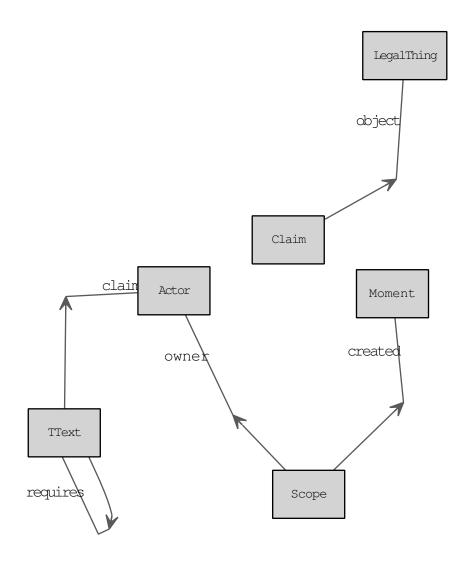


Figure 3.1: Concept diagram of relations in Arguments

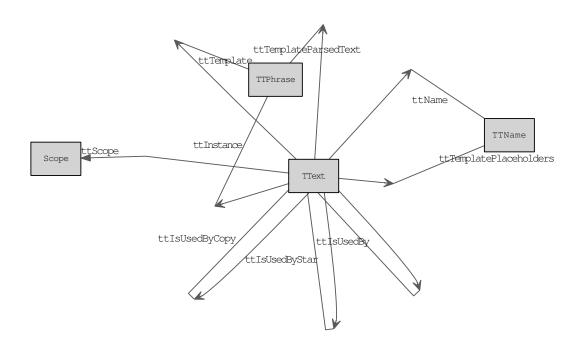


Figure 3.2: Concept diagram of relations in Placeholder Extraction and TText Hierarchy creation Service

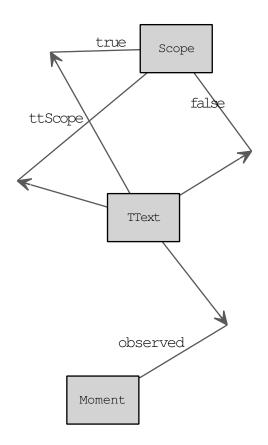


Figure 3.3: Concept diagram of relations in Validity

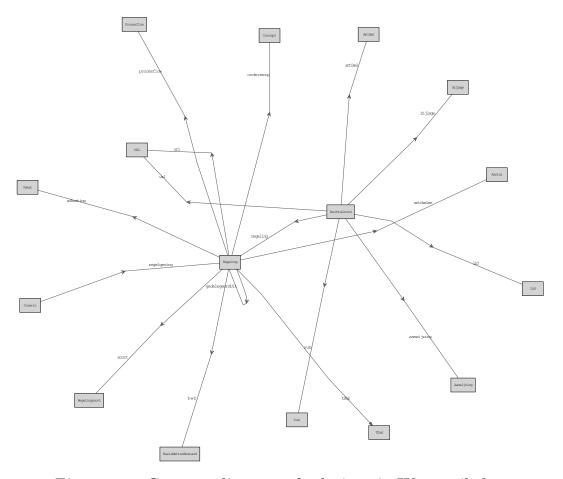


Figure 3.4: Concept diagram of relations in Wetsartikelen

- line 133:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- When the name of a TText is in the list of placeholders of another TText, the first TText is said to be used by the second TText
 - line 140:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- A (first) TText is only used by a (second) TText if the name of the first is in the list of placeholders of the second
 - line 144:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- When a TText is used by another TText, it inherits the latter's Scope line 150:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- $ttValueUsedToReplacePlaceholders\ integrity$ line 19:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- Register that the value of a TText has been used to replace placeholders line 22:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- After a value update, all TTexts that used the value must be reset line 38:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- When a TText value is NOT used in an instance of (another) TText, then it may not say that its value is used to replace a placeholder line 50:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- Request TText instance phrase reset if the template phrase does, and the instance phrase does not exist
 - line 35:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- Resetting a TText implies that the relations with other TTexts in which it is actually used for replacement, are broken
- line 41:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc Rules are defined, the meaning of which is documented by means of computer generated language:
 - caseAuthor line 155:5, file C:\Afstuderen\repos\ampersandmodels\MirrorMe\MirrorMe.adl
 - TTexts can only use TTexts that are in the same ttScope line 122:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
 - Compute transitive closure of 'ttIsUsedBy' line 129:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
 - Create TTexts for placeholders (if necessary) line 133:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc

- When the name of a TText is in the list of placeholders of another TText, the first TText is said to be used by the second TText
 - line 140:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- A (first) TText is only used by a (second) TText if the name of the first is in the list of placeholders of the second
 - line 144:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- When a TText is used by another TText, it inherits the latter's Scope line 150:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcExtract.svc
- $ttValueUsedToReplacePlaceholders\ integrity$ line 19:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- When a TText value is NOT used in an instance of (another) TText, then it may not say that its value is used to replace a placeholder line 50:1, file C:\Afstuderen\repos\ampersandmodels\TText\TText_PlcReplace.svc
- ttValueIsUsedInInstanceOfTText integrity

 $line~31:1, file~C: \\ Afstuderen \\ repos \\ ampersand models \\ TText \\ TText \\ PlcExtract.svc$

The table below shows for each theme (i.e. process or pattern) the number of relations and rules, followed by the number and percentage that have a reference. Relations declared in multiple themes are counted multiple times.

		With			Entire	
Theme	Relation	n s referen	ce %	Rules	context	%
Arguments	5	1	20%	0	0	-
Case Management	0	0	-	1	0	0%
Placeholder Extraction	5	0	0%	9	0	0%
and TText Hierarchy						
creation Service						
Placeholder	1	0	0%	4	0	0%
Replace-By-TTValue						
Service						
TText Reset Service	2	0	0%	4	0	0%
Validity	4	0	0%	2	0	0%
Wetsartikelen	17	0	0%	0	0	-
Entire context	110	1	0%	99	0	0%

The table below shows the signal rules per role.

role	in pattern	rule
ExecEngine	Case Management	caseAuthor
ExecEngine	Placeholder Extraction and TText Hierarchy	TText template parsing - extract placeholders
ExecEngine	creation Service Placeholder Extraction and TText Hierarchy creation Service	TText template parsing - delete parsed templates

role	in pattern	rule
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	TText template parsing - delete placeholders used in template
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	Compute transitive closure of 'ttIsUsedBy'
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	Create TTexts for placeholders (if necessary)
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	When the name of a TText is in the list of placeholders of another TText, the first TText is said to be used by the second TText
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	A (first) TText is only used by a (second) TText if the name of the first is in the list of placeholders of the second
ExecEngine	Placeholder Extraction and TText Hierarchy creation Service	When a TText is used by another TText, it inherits the latter's Scope
ExecEngine	Placeholder Replace-By-TTValue Service	Register that the value of a TText has been used to replace placeholders
ExecEngine	Placeholder Replace-By-TTValue Service	After a value update, all TTexts that used the value must be reset
ExecEngine	Placeholder Replace-By-TTValue Service	When a TText value is NOT used in an instance of (another) TText, then it may not say that its value is used to replace a placeholder
ExecEngine	TText Reset Service	Request TText instance phrase reset if the template phrase does, and the instance phrase does not exist
ExecEngine	TText Reset Service	Resetting a TText implies that the relations with other TTexts in which it is actually used for replacement, are broken

role	in pattern	rule
ExecEngine	TText Reset Service	Resetting a TText is complete if other TTexts are not used by
ExecEngine	MirrorMe	this TText A TText that is not in a Scope may not exist
ExecEngine	MirrorMe	Account activation/initialization
ExecEngine	MirrorMe	Activate 'Developer'
ExecEngine	MirrorMe	Activate roles in a session
ExecEngine	MirrorMe	Activation of session Organization
ExecEngine	MirrorMe	Active accounts are initialized
ExecEngine	MirrorMe	Add arguments to computations
ExecEngine	MirrorMe	Add to Assignments history
ExecEngine	MirrorMe	Assign computation result to TText (provided actual
ExecEngine	MirrorMe	arguments were used) Auto DelPair sessionSIA
ExecEngine	MirrorMe	Auto activate auto-login accounts
ExecEngine	MirrorMe	Auto create SIAMPersonRefComputati
ExecEngine	MirrorMe	Auto maintain 'accPersonRef' relation
ExecEngine	MirrorMe	Auto maintain 'personRef' relation
ExecEngine	MirrorMe	Automatically depopulate personOrg
ExecEngine	MirrorMe	Automatically login
ExecEngine	MirrorMe	Automatically populate personOrg
ExecEngine	MirrorMe	By default, a session authentication level is '1'
ExecEngine	MirrorMe	Clear allowed session roles
ExecEngine	MirrorMe	Clear the OrgRef in a session
ExecEngine	MirrorMe	Create Account upon request
ExecEngine	MirrorMe	Create Assignment for TTexts that have a (new) value

role	in pattern	rule
ExecEngine	MirrorMe	Create Assignment for TTexts that have no value
ExecEngine	MirrorMe	Create computations
ExecEngine	MirrorMe	Create/Update copied ttDescr
ExecEngine	MirrorMe	Create/Update copied ttOwner
ExecEngine	MirrorMe	Deactivate 'Developer' role
ExecEngine	MirrorMe	Deactivate roles in a session
ExecEngine	${\rm MirrorMe}$	Deactivation of session Organization
ExecEngine	${\rm MirrorMe}$	Default roles must be allowed roles
ExecEngine	MirrorMe	Delete computations (that have results) of which one or more arguments have changed
ExecEngine	${\rm MirrorMe}$	Delete obsolete Assignments
ExecEngine	${\rm MirrorMe}$	Delete obsolete Computations
ExecEngine	${\rm MirrorMe}$	Determine the 'userid' in a session
ExecEngine	${\rm MirrorMe}$	Determine the OrgRef in a session
ExecEngine	${\rm MirrorMe}$	Determine the acting person in a session
ExecEngine	${\rm MirrorMe}$	Determine the personRef in a session
ExecEngine	${\rm MirrorMe}$	Disable automated login
ExecEngine	MirrorMe	Every TText with scope and name can be used as a binding
ExecEngine	${\rm MirrorMe}$	Initialize copied ttDescr
ExecEngine	${\rm MirrorMe}$	Initialize copied ttOwner
ExecEngine	MirrorMe	InsPair sessionAuthISOLevel (MPTrx specific)
ExecEngine	${\rm MirrorMe}$	Limit history size to 2 predecessors
ExecEngine ExecEngine	MirrorMe MirrorMe	Logout Registering the first authenticated session account

role	in pattern	rule
ExecEngine	MirrorMe	Reset login help
ExecEngine	MirrorMe	SURttName
ExecEngine	MirrorMe	SURttPhrase
ExecEngine	MirrorMe	Service U/PW Login
		$\operatorname{request}$
ExecEngine	MirrorMe	Service deactivation
		request
ExecEngine	MirrorMe	Sessions with an
		inactive sessionaccount
		may not exist
ExecEngine	MirrorMe	Set allowed session
		roles
ExecEngine	MirrorMe	Set default Scope
		ownership
ExecEngine	MirrorMe	Set default TText
		ownership
ExecEngine	MirrorMe	The userid of a
		customer is its (initial)
		personref
ExecEngine	MirrorMe	U/PW Login
ExecEngine	MirrorMe	Undetermine the
		'userid' in a session
ExecEngine	MirrorMe	Undetermine the acting
		person in a session
ExecEngine	MirrorMe	Undetermine the
_		personRef in a session
ExecEngine	MirrorMe	Unowned TTexts that
		are used by an owned
		TText, will be assigned
_		the same owner
ExecEngine	MirrorMe	Update arguments of
		computations that do
B B :	3.6	not have a result
ExecEngine	MirrorMe	Update deleted ttDescr
ExecEngine	MirrorMe	Update deleted
		ttOwner

The population in this script does not specify any work in progress.

The population in this script violates no rule.

Chapter 4

Conceptual Analysis

This chapter defines the formal language, in which functional requirements of 'MirrorMe' can be analysed and expressed. The purpose of this formalisation is to obtain a buildable specification. This chapter allows an independent professional with sufficient background to check whether the agreements made correspond to the formal rules and definitions.

4.1 Theme: 'Arguments'

The structure of legal arguments has been studied by scholars such as Toulmin~[?], Verheij (name some more). To help legal professionals construct such arguments, the structure of arguments must be formalized. For this purpose, we can consider each (legal) case as a set of statements. Each statement can be considered true or

fig. 4.1 shows a conceptual diagram of this pattern.

4.2 Theme: 'Case Management'

This pattern contains some basic administration for cases.

fig. 4.2 shows a conceptual diagram of this pattern.

4.3 Theme: 'Placeholder Extraction and TText Hierarchy creation Service'

This service ensures that - ttTemplateParsedText=ttTemplate AND - ttTemplatePlaceholders contains all TTNames that are mentioned in the ttTemplate of a TText.

The idea is that the specification of a TText is parsed to see if it contains names of TTexts. Such names are recognized by the fact that they are surrounded by square brackets ([and]).

So, - ttTemplateParsedText stores the text that has been parsed (in PHP) to produce the contents of ttTemplatePlaceholders. This implies that whenever

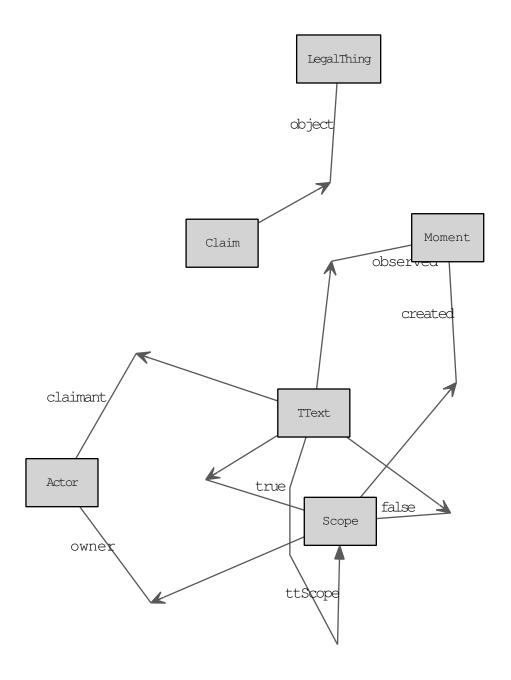


Figure 4.1: Concept diagram of the rules in Arguments

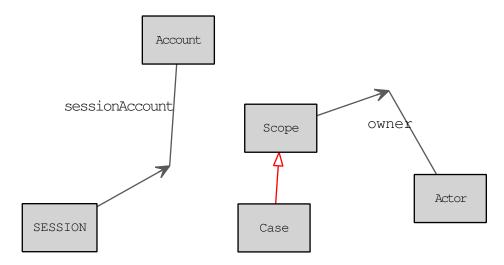


Figure 4.2: Concept diagram of the rules in Case Management

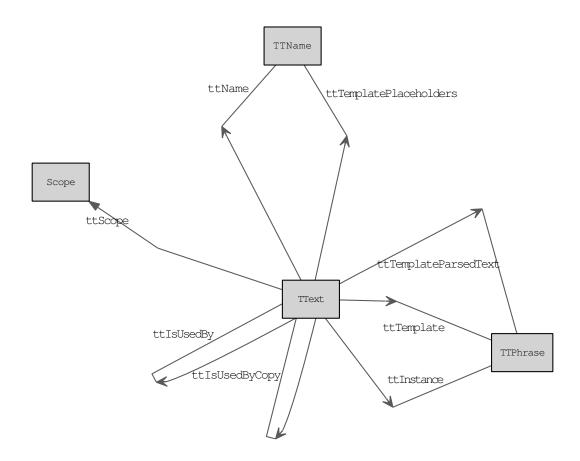


Figure 4.3: Concept diagram of the rules in Placeholder Extraction and TText Hierarchy creation Service

ttTemplateParsedText is empty for some TText, ttTemplatePlaceholders is empty for that same atom. - whenever ttTemplateParsedText differs from ttTemplate, - it is first removed, and ttTemplatePlaceholders are discarded - then ttTemplate is being parsed, and ttTemplateParsedText and ttTemplatePlaceholders are filled again.

All of this happens in the same Scope, i.e. in the scope to which the TTexts belong. Note that for this to work, all TTexts that are mentioned in a ttTemplate must exist.

fig. 4.3 shows a conceptual diagram of this pattern.

4.3.1 Rules

This section itemizes the rules with a reference to the shared language of stakeholders for the sake of traceability.

agreement ¿lst:SharedLangrule:TTexts_32can_32only_32use_32TTexts_32that_32are_32in has been made:

Using relations relation ¿eq:ConceptualAnalysisdeclaration:ttIsUsedBy_91TText_42TText(ttIsUsedBy), relation ¿eq:ConceptualAnalysisdeclaration:ttScope_91TText_42Scope_93?(ttScope), this is formalized as

an undocumented agreement: ttIsUsedBy ⊢ ttScope; ttScope ig. 4.4 shows a conceptual diagram of this rule.

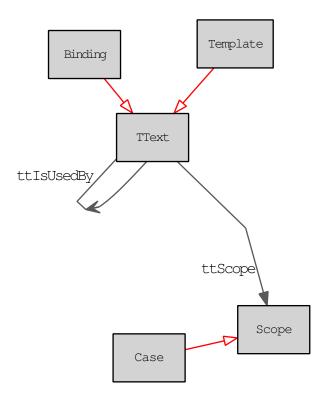


Figure 4.4: Concept diagram of rule TTexts can only use TTexts that are in the same ttScope

4.4 Theme: 'Placeholder Replace-By-TTValue Service'

This service ensures that the ttInstance phrase of a TText is the ttTemplate phrase of that TText, where every occurrence of a reference to another TText in that template (i.e. occurences of [<TTName>], called 'placeholder's) is replaced by the ttValue of that other TText, provided that the ttValue is populated for that TText.

fig. 4.5 shows a conceptual diagram of this pattern.

4.4.1 Rules

This section itemizes the rules with a reference to the shared language of stakeholders for the sake of traceability.

agreement ; lst: Shared Langrule: ttValue Used To Replace Placeholders 32 integrity?

has been made:

Using relations relation ¿eq:ConceptualAnalysisdeclaration:ttValueUsedToReplacePlacehol (ttValueUsedToReplacePlaceholders), relation ¿eq:ConceptualAnalysisdeclaration:ttValue_917 (ttValue), this is formalized as

an undocumented agreement: ttValueUsedToReplacePlaceholders \(\text{ttValue} \)

fig. 4.6 shows a conceptual diagram of this rule.

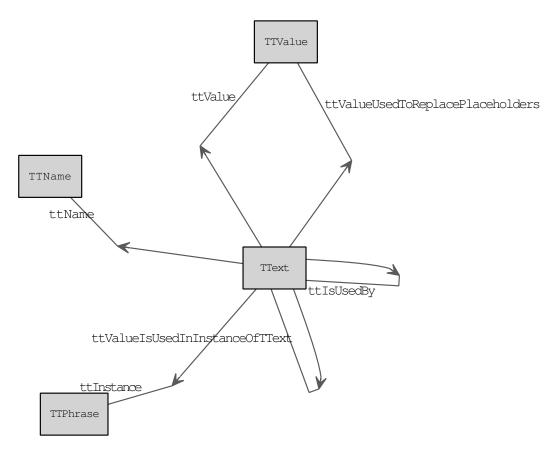
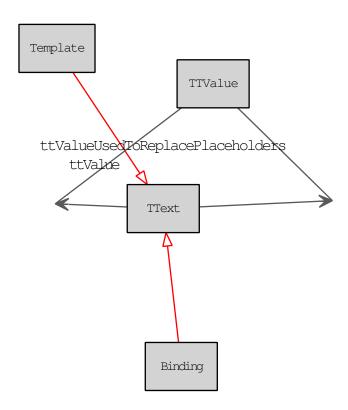


Figure 4.5: Concept diagram of the rules in Placeholder Replace-By-TTValue Service



 ${\bf Figure~4.6:~Concept~diagram~of~rule~ttValueUsedToReplacePlaceholders~integrity}$

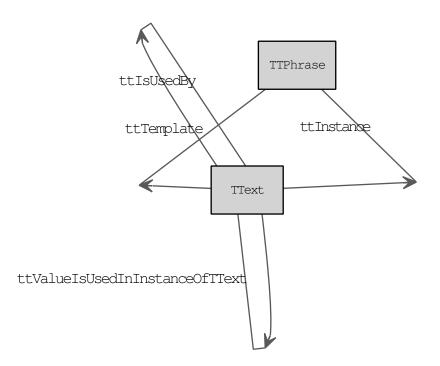


Figure 4.7: Concept diagram of the rules in TText Reset Service

4.5 Theme: 'TText Reset Service'

This service relinquishes any depencies that the ttInstance of a TText has on other TTexts (as registered in ttValueIsUsedInInstanceOfTText). This condition will be realized by setting the property ttInstanceResetReq for that TText.

fig. 4.7 shows a conceptual diagram of this pattern.

4.5.1 Rules

This section itemizes the rules with a reference to the shared language of stakeholders for the sake of traceability.

agreement ; lst:SharedLangrule:ttValueIsUsedInInstanceOfTText_32integrity?

has been made:

Using relations relation ¿eq:ConceptualAnalysisdeclaration:ttValueIsUsedInInstanceOfTTex (ttValueIsUsedInInstanceOfTText), relation ¿eq:ConceptualAnalysisdeclaration:ttIsUsedBy_9 (ttIsUsedBy), this is formalized as

an undocumented agreement: $ttValueIsUsedInInstanceOfTText \vdash ttIsUsedBy$

fig. 4.8 shows a conceptual diagram of this rule.

4.6 Theme: 'Validity'

In order to talk about true and false statements in a precise way, we need the idea of contexts.

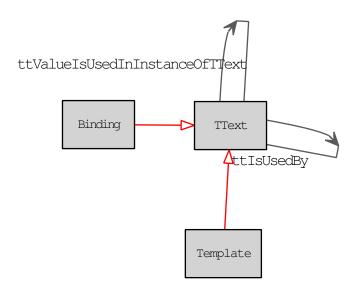


Figure 4.8: Concept diagram of rule ttValueIsUsedInInstanceOfTText integrity

fig. 4.9 shows a conceptual diagram of this pattern.

4.6.1 Rules

This section itemizes the rules with a reference to the shared language of stakeholders for the sake of traceability.

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, given that all its placeholders have been instantiated.

Therefore agreement **¿lst:SharedLangrule:TrueFalseValid?** exists:

A statement can be true or false in a context only if it is valid within that context.

Using relations relation ¿eq:ConceptualAnalysisdeclaration:true_91TText_42Scope_93? (true), relation ¿eq:ConceptualAnalysisdeclaration:false_91TText_42Scope_93? (false), relation ¿eq:ConceptualAnalysisdeclaration:ttScope_91TText_42Scope_93? (ttScope), this is formalized as

Rule 21: true \cup false \vdash ttScope

fig. 4.10 shows a conceptual diagram of this rule.

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, if all its placeholders have been instantiated.

Therefore agreement ; lst: SharedLangrule: Inconsistency? exists:

A statement cannot be both true and false in the same context.

Using relations relation ¿eq:ConceptualAnalysisdeclaration:true_91TText_42Scope_93? (true), relation ¿eq:ConceptualAnalysisdeclaration:false_91TText_42Scope_93? (false), this is formalized as

Rule 22: : true \cap false $\vdash \overline{V_{[\text{TText*Scope}]}}$

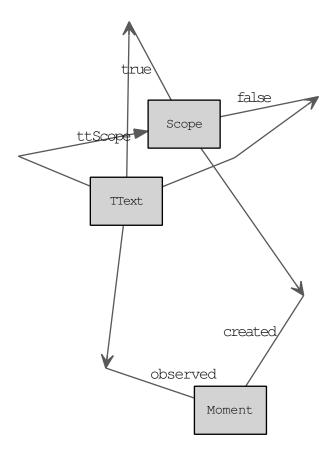


Figure 4.9: Concept diagram of the rules in Validity

fig. 4.11 shows a conceptual diagram of this rule.

4.7 Theme: 'Wetsartikelen'

fig. 4.12 shows a conceptual diagram of this pattern.

The definitions of concepts can be found in the glossary.

4.7.1 Declared relations

This section itemizes the declared relations with properties and purpose.

Let us treat a statement made by an individual as a claim of that individual that the statement is true. For this reason, each individual who makes the claim can be registered in the relation "claimant". The claimant 'claims' (a) the validity of the ttTemplate and (b) the veracity of ttValue. Thus, the claimaint is solely authorized to change ttValue and ttTemplate.

For this purpose, the following relation has been defined

Relation 1: : claimant[TText * Actor]

A claimant is the party that makes a claim, especially one that is legally cognizable.

The following univalent relation has been defined

http://www.thefreedictionar

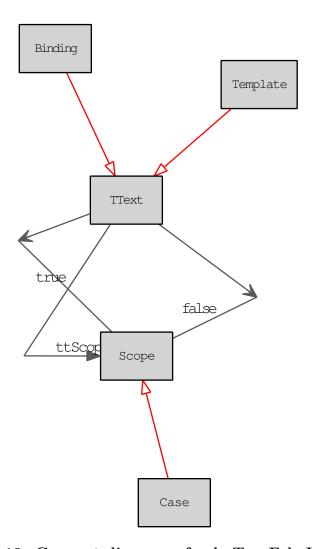


Figure 4.10: Concept diagram of rule TrueFalseValid

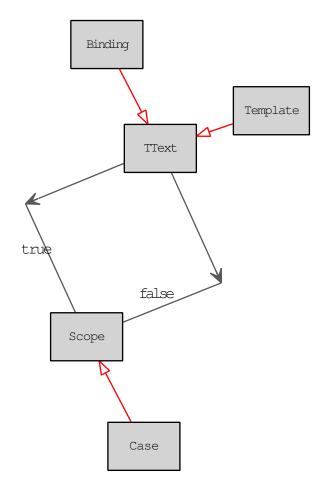


Figure 4.11: Concept diagram of rule Inconsistency

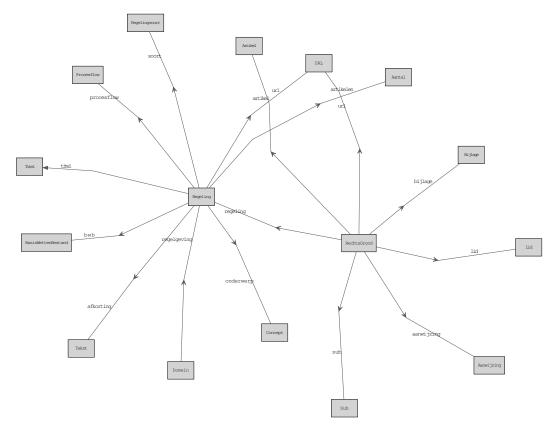


Figure 4.12: Concept diagram of the rules in Wetsartikelen ${\bf e}$

an undocumented relation: object[Claim * LegalThing]

(No meaning has been specified)

The following univalent relation has been defined

Relation 2: : owner[Scope * Actor]

the Actor that is the claimant of all toplevel TTexts.

The following univalent relation has been defined

Relation 3: : created[Scope * Moment]

The application can register a moment as the time a case has been created.

The following relation has been defined

Relation 4: : requires[TText * TText]

A statement that needs another statement to be true, is registered in the relation requires.

The following univalent relation has been defined

an undocumented relation: ttTemplateParsedText[TText * TTPhrase]

(No meaning has been specified)

The following relation has been defined

an undocumented relation: ttTemplatePlaceholders[TText * TTName]

(No meaning has been specified)

The following irreflexive and antisymmetric relation has been defined

an undocumented relation: ttIsUsedBy[TText * TText]

(No meaning has been specified)

The following relation has been defined

an undocumented relation: ttIsUsedByCopy[TText * TText]

(No meaning has been specified)

The following irreflexive, antisymmetric, and transitive relation has been defined

an undocumented relation: ttIsUsedByStar[TText * TText]

(No meaning has been specified)

The following univalent relation has been defined

Relation 12: : ttValueUsedToReplacePlaceholders[TText * TTValue]

References to this TText (i.e.: placeholders) exist, and have been replaced with the specified TTValue

The property ttInstanceResetReq is used to start reconstructing an instance phrase of a TText.

For this purpose, the following symmetric, antisymmetric, univalent, and injective relation has been defined

Relation 10: : ttInstanceResetReq[TText * TText]

If a TText has the property ttInstanceResetReq, this means that its instance phrase needs to be (re)constructed from scratch.

The following irreflexive and antisymmetric relation has been defined

Relation 11: : ttValueIsUsedInInstanceOfTText[TText * TText]

SRC TText has been used to replace placeholders in the TGT TText instance phrase

Consider the statement "John has blond hair". If this statement is known to be true in some context c, the tuple ("John has blond hair", c) can exist in the relation true. However, if this tuple is not in the relation true, it does not follow that John does not have blond hair. The truth of the TText value (as decided by its claimant) is accepted by the owner of the Scope

For this purpose, the following univalent relation has been defined

Relation 18: : true[TText * Scope]

A statement that is considered true within a context is registered in the relation true.

The relation false is dual to true.

For this purpose, the following univalent relation has been defined

Relation 19: : false[TText * Scope]

A statement that is considered false within a context is registered in the relation false.

The following function has been defined

Relation 6: : ttScope[TText * Scope]

Statements that are in a given scope are considered valid within that scope.

For reconstructing events, it can be necessary to administer the moment an observation or a claim is made. For this reason, we use the relation "observed".

For this purpose, the following univalent relation has been defined

Relation 20: : observed[TText * Moment]

The application can register a moment as the time a statement has been made.

The following relation has been defined

an undocumented relation: url[RechtsGrond * URL]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: regeling[RechtsGrond * Regeling]

(No meaning has been specified)

The following univalent relation has been defined an undocumented relation: bijlage[RechtsGrond * Bijlage] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: aanwijzing[RechtsGrond * Aanwijzing] (No meaning has been specified) The following relation has been defined an undocumented relation: artikel[RechtsGrond * Artikel] (No meaning has been specified) The following relation has been defined an undocumented relation: lid[RechtsGrond * Lid] (No meaning has been specified) The following relation has been defined an undocumented relation: sub[RechtsGrond * Sub] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: bwb[Regeling * BasisWettenBestand] (No meaning has been specified) The following function has been defined an undocumented relation: titel[Regeling * Titel] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: artikelen[Regeling * Aantal] (No meaning has been specified) The following relation has been defined an undocumented relation: url[Regeling * URL] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: gedelegeerdUit[Regeling * Regeling] (No meaning has been specified) The following relation has been defined an undocumented relation: regelgeving[Domein * Regeling] (No meaning has been specified) The following univalent relation has been defined

an undocumented relation: soort[Regeling * Regelingsoort]

(No meaning has been specified)

The following relation has been defined

an undocumented relation: onderwerp[Regeling * Concept]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: processflow[Regeling * Processflow]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: afkorting[Regeling * Tekst]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: class[TText * Concept]

(No meaning has been specified)

The following relation has been defined

an undocumented relation: legalGround[TText * LegalGround]

(No meaning has been specified)

The following relation has been defined

an undocumented relation: evidence[TText * Evidence]

(No meaning has been specified)

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: autoLoginAccount[Account * Account]

(No meaning has been specified)

The following injective function has been defined

Relation 23: : accUserid[Account * UserID]

An Account registers a pseudonym for whatever the Account applies to

The following univalent relation has been defined

Relation 24: : accPassword[Account * Password]

An Account registers a password for whatever the Account applies to

The following univalent relation has been defined

Relation 5: : sessionAccount[SESSION * Account]

A SESSION may be linked to an Account, which specifies the (user)context of the session.

The following symmetric, antisymmetric, univalent, and injective relation has been defined

Relation 25: : accIsInitialized[Account * Account]

An account may have the property that it is initialized (e.g. some default values have been set)

The following symmetric, antisymmetric, injective, and surjective function has been defined

Relation 26: : accIsActive[Account * Account]

An Account may have the property that it is active, meaning that it can be used to login

The following symmetric, antisymmetric, univalent, and injective relation has been defined

Relation 27: : accDeactivateReq[Account * Account]

A request may exist to deactivate the account (as soon as it is no longer a sessionAccount)

The following univalent relation has been defined

Relation 28: : sessionUserid[SESSION * UserID]

In a SESSION, the userid (i.e. the name/text that identifies the Account that is used to login) may be known.

The following univalent relation has been defined

an undocumented relation: loginUserid[SESSION * UserID]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: loginPassword[SESSION * Password]

(No meaning has been specified)

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: logoutRequest[SESSION * SESSION]

(No meaning has been specified)

The following reflexive, antisymmetric, and transitive relation has been defined

an undocumented relation: isoLevelGTE[ISOLevel * ISOLevel]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: sessionReqdISOLevel[SESSION * ISOLevel]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: sessionAuthISOLevel[SESSION * ISOLevel]

(No meaning has been specified)

The following function has been defined

an undocumented relation: moduleName[Module * ModuleName]

(No meaning has been specified)

The following function has been defined

an undocumented relation: moduleVsnMajor[Module*ModuleVsnMajor]

(No meaning has been specified)

The following function has been defined

an undocumented relation: moduleVsnMinor[Module*ModuleVsnMinor]

(No meaning has been specified)

The following injective and univalent relation has been defined

Relation 29: : orgAbbrName[Organization * OrgAbbrName]

An organization may have a short name by which it can be identified

The following injective and univalent relation has been defined

Relation 30: : orgFullName[Organization * OrgFullName]

An organization has a long (full) name by which it may be identified

The following univalent relation has been defined

Relation 31: : accOrg[Account * Organization]

An Account may contain (web)service-specific attributes regarding a specific Organization

The following univalent relation has been defined

Relation 32: : sessionOrg[SESSION * Organization]

In a session, the (accountable) Organization may be known.

The following relation has been defined

an undocumented relation: personOrg[Person * Organization]

(No meaning has been specified)

The following function has been defined

Relation 33: : personFirstName[Person * FirstName]

The first name of a Person that is registered.

The following function has been defined

Relation 34: : personLastName[Person * LastName]

The last name of a Person that is registered.

The following univalent relation has been defined

Relation 35: : accPerson[Account * Person]

An Account may contain (web)service-specific attributes regarding a specific Person

The following univalent relation has been defined

Relation 36: : sessionPerson[SESSION * Person]

In a SESSION, the user (i.e. the Person that acts 'at the other side') may be known.

The following relation has been defined

Relation 37: : accAllowedRoles[Account * Role]

An Account registers the Roles that MAY be activated in a SESSION to which the Account is assigned

The following relation has been defined

Relation 38: : accDefaultRoles[Account * Role]

An Account registers the Roles that ARE activated in a SESSION once the Account is assigned

The following relation has been defined

Relation 39: : sessionAllowedRoles[SESSION * Role]

Within a SESSION, a Role may be activated

The following relation has been defined

Relation 40: : sessionActiveRoles[SESSION * Role]

A SESSION has activated a Role

In order to enable the suspension of sessions, the first sessionaccount must be remembered since it must be used to verify that the user that logs in to a suspended session in order to reactivate it, is the same as the user that did the original login.

For this purpose, the following univalent relation has been defined

Relation 41: : firstSessionAccount[SESSION * Account]

The Account that is authenticated first in a session, must be remembered.

The following function has been defined

an undocumented relation: uidUserid[UID * UserID]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: uidIssuer[UID * IdP]

(No meaning has been specified)

The following injective relation has been defined

Relation 42: : accUID[Account * UID]

An Account registers a UID for whatever the Account applies to

The following univalent relation has been defined

Relation 43: : personRef[Person * PersonRef]

A textstring to refer to this person

The following univalent relation has been defined

Relation 44: : accPersonRef[Account * PersonRef]

An Account may have a textstring that refers to the Person for which the account has attributes

The following univalent relation has been defined

Relation 45: : sessionPersonRef[SESSION * PersonRef]

In a SESSION, a textstring that refers to the user (i.e. the Person that acts 'at the other side') may be known.

The following univalent relation has been defined

an undocumented relation: siamcompFirstName[SIAMPersonRefComputation* FirstName]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: siamcompLastName[SIAMPersonRefComputation* LastName]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: siamcompPersonRef[SIAMPersonRefComputation* PersonRef]

(No meaning has been specified)

The following univalent relation has been defined

Relation 46: : accOrgRef[Account * OrgRef]

An Account may have a textstring that refers to the Organization to which the Account is related

The following univalent relation has been defined

Relation 47: : sessionOrgRef[SESSION * OrgRef]

In a SESSION, a textstring that refers to the Organization ('at the other side') may be known.

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: loginCreateAccount[SESSION *
SESSION]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: registerPassword[SESSION * Password]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: registerOrgRef[SESSION * OrgRef]

(No meaning has been specified) The following symmetric, antisymmetric, univalent, and injective relation has been defined an undocumented relation: loginReq[SESSION * SESSION] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: loginISOLevel[SESSION * ISOLevel] (No meaning has been specified) The following symmetric, antisymmetric, univalent, and injective relation has been defined an undocumented relation: sessionDev[SESSION * SESSION] (No meaning has been specified) The following injective function has been defined an undocumented relation: ttTrace[TText * Assignment] (No meaning has been specified) The following function has been defined an undocumented relation: asmVar[Assignment * TText] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: asmVal[Assignment * TTValue] (No meaning has been specified) The following univalent relation has been defined an undocumented relation: asmPOT[Assignment * PointOfTime] (No meaning has been specified) The following univalent and injective relation has been defined an undocumented relation: asmHasPred[Assignment * Assignment] (No meaning has been specified) The following injective function has been defined an undocumented relation: compVar[Computation * TText] (No meaning has been specified) The following relation has been defined an undocumented relation: compArg[Computation * Assignment] (No meaning has been specified)

an undocumented relation: compRes[Computation * TTValue]
(No meaning has been specified)

The following univalent relation has been defined

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: sessionLoginAssist[SESSION*SESSION]

(No meaning has been specified)

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: sessionLogoutAssist[SESSION *
SESSION]

(No meaning has been specified)

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: sessionSRI[SESSION * SESSION]

(No meaning has been specified)

The following symmetric, antisymmetric, univalent, and injective relation has been defined

an undocumented relation: sessionSIA[SESSION * SESSION]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: scopeID[Scope * ScopeID]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: scopeOwner[Scope * Account]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: scopeDescr[Scope * ScopeDescr]

(No meaning has been specified)

The following irreflexive and antisymmetric relation has been defined

Relation 48: : scopeIII[Scope * Scope]

The structured content of a (SRC) Scope may be included in that of a (TGT) Scope

The following symmetric, antisymmetric, univalent, and injective relation has been defined

Relation 49: : scopeIsaCC[Scope * Scope]

A Scope may have the property of being a carbon copy of another Scope

The following univalent relation has been defined

an undocumented relation: ttName[TText * TTName]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: ttValue[TText * TTValue]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: ttTemplate[TText * TTPhrase]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: ttInstance[TText * TTPhrase]

(No meaning has been specified)

The following univalent, irreflexive, and antisymmetric relation has been defined

an undocumented relation: ttICO[TText * TText]

(No meaning has been specified)

The following univalent, irreflexive, and antisymmetric relation has been defined

an undocumented relation: ttICCO[TText * TText]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: ttDescr[TText * TTPhrase]

(No meaning has been specified)

The following univalent relation has been defined

an undocumented relation: ttOwner[TText * Account]

(No meaning has been specified)

Chapter 5

Data structure

This chapter contains the result of the data analysis. It is structured as follows:

We start with the classification model, followed by a list of all relations, that are the foundation of the rest of the analysis. Finally, the logical and technical data model are discussed.

5.1 Classifications

A number of concepts is organized in a classification structure. This is shown in fig. 5.1.

5.2 Rules

In this section an overview of all rules is given. For every rule, the formal expression is shown. The process rules are given first, followed by the invariants.

5.2.1 Process rules

Process rules are rules that are signalled.

Process rule: caseAuthor

 $V_{\text{[Case*SESSION]}}$; "_SESSION"; sessionAccount \vdash owner; $V_{\text{[Actor*Account]}}$

Process rule: TText template parsing - extract placeholders

TTexts that have no placeholders detected, yet have a ttTemplate, must have been parsed

 $(I_{[TText]} - ttTemplatePlaceholders; ttTemplatePlaceholders "); ttTemplate <math>\vdash ttTemplateParsedText$

Process rule: TText template parsing - delete parsed templates

Whenever a template phrase changes, the parsed text must be deleted so that the template is parsed again

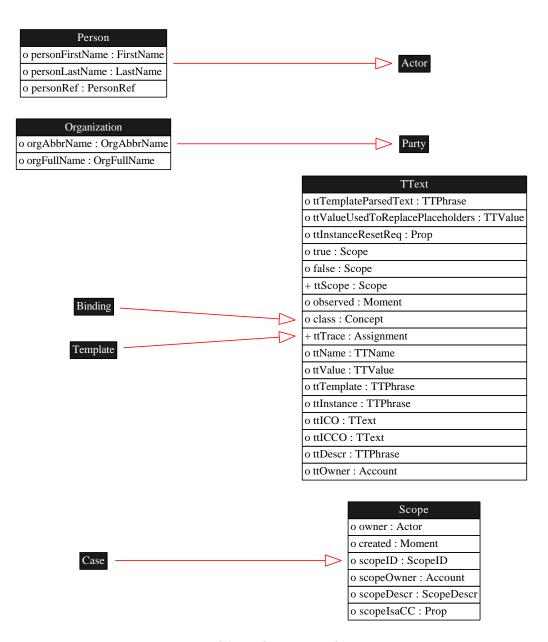


Figure 5.1: Classification of MirrorMe

$\begin{array}{ll} \textbf{Process rule:} \ \textit{TText template parsing - delete placeholders used in} \\ \textit{template} \end{array}$

TTexts whose template text is not parsed, must not specify placeholders that are detected

 $\frac{(I_{[\text{TText}]} - \text{ttTemplateParsedText}; \text{ttTemplateParsedText}^{\smile}); \text{ttTemplatePlaceholders} \vdash \overline{V_{[\text{TText}*TTName]}}$

Process rule: Compute transitive closure of 'ttIsUsedBy'

ttIsUsedBy = ttIsUsedByCopy

Process rule: Create TTexts for placeholders (if necessary)

ttScope; $ttTemplatePlaceholders \vdash ttScope$; ttName

Process rule: When the name of a TText is in the list of placeholders of another TText, the first TText is said to be used by the second TText

ttName; $ttTemplatePlaceholders <math>\cap ttScope$; $ttScope \vdash ttIsUsedBy$

Process rule: A (first) TText is only used by a (second) TText if the name of the first is in the list of placeholders of the second

 $ttIsUsedBy \vdash ttName; ttTemplatePlaceholders \lq \cap ttScope; ttScope \lq$

Process rule: When a TText is used by another TText, it inherits the latter's Scope

 $ttIsUsedBy; ttScope \vdash ttScope$

Process rule: Register that the value of a TText has been used to replace placeholders

If a TText has a value, and is used in a TText that is not being re-initialized, then it must appear in the ttValueIsUsedInInstanceOfTText of the TText.

 $(I_{[TText]} \cap ttValue; ttValue); ttIsUsedBy; (I_{[TText]} - ttInstanceResetReq) \vdash ttValueIsUsedInInstanceOfTText$

Process rule: After a value update, all TTexts that used the value must be reset

When the value of a TText differs from the value used in replacements, or has become void, then the instance-texts of all TTexts in which the replacements took place must be discarded and recreated.

 $(I_{[TText]} \cap (ttValueUsedToReplacePlaceholders; \overline{I_{[TTValue]}}; ttValue \subset \cup (ttValue; ttValue \subset))); ttValueIsUsedInInttValueIsUsedInInstanceOfTText; ttInstanceResetReq$

Process rule: When a TText value is NOT used in an instance of (another) TText, then it may not say that its value is used to replace a placeholder

 $ttValueUsedToReplacePlaceholders \vdash ttValueIsUsedInInstanceOfTText; V_{[TText*TTValue]}$

Process rule: Request TText instance phrase reset if the template phrase does, and the instance phrase does not exist

If a TText has a template phrase and no instance phrase, the TText must be reset/initialized, thus allowing its instance phrase to be constructed

 $(I_{[\text{TText}]} - \text{ttInstance}; \text{ttInstance}); \text{ttTemplate} \vdash \text{ttInstanceResetReq}; \text{ttTemplate}$

Process rule: Resetting a TText implies that the relations with other TTexts in which it is actually used for replacement, are broken

Resetting a TText means that the registration of replaced placeholders (for that TText) must be reset (cleared/deleted).

tt Value Is
Used InInstance OfTText; tt Instance Reset Req
 $\overline{V_{[\text{TText*TText}]}}$

Process rule: Resetting a TText is complete if other TTexts are not used by this TText

Before reconstructing an instance phrase of a TText, the ExecEngine must have discarded all administration related to that TText.

A TText that needs to be (re-)initialized and does not use values of TTexts in its UsedValue, must have the specification-text as its UsedValue, which completes the (re-)initialization.

 $(ttInstanceResetReq-ttValueIsUsedInInstanceOfTText\@ifnext{"}; ttValueIsUsedInInstanceOfTText\@ifnext{"}; ttValueIsUsedInInstanceOfTText\@ifnext{"}; ttValueIsUsedInInstanceOfTText\@ifnext{"}; ttValueIsUsedInInstanceOfTText\@ifnext{"}; ttTemplate\@ifnext{"} ttInstance\@ifnext{"}; ttValueIsUsedInInstanceOfTText\@ifnext{"}; t$

Process rule: A TText that is not in a Scope may not exist

TTexts are destroyed when they are not assigned a ttScope, or their ttScope ceases to exist

 $I_{[\texttt{TText}]} \vdash \texttt{ttScope}; \texttt{ttScope} \check{\ }$

Process rule: Account activation/initialization

 $I_{[Account]} \vdash accIsInitialized$

Process rule: Activate 'Developer' role

sessionDev ⊢ sessionActiveRoles; "Developer"; sessionActiveRoles

Process rule: Activate roles in a session

 $(I_{\rm [SESSION]} - {\rm sessionActiveRoles}; {\rm sessionActiveRoles} \); {\rm sessionAccount}; {\rm accDefaultRoles} \vdash {\rm sessionActiveRoles}$

Process rule: Activation of session Organization

The Organization of the user of the session may be known.

 $sessionAccount; accOrg \vdash sessionOrg$

Process rule: Active accounts are initialized

 $accIsActive \vdash accIsInitialized$

Process rule: Add arguments to computations

If the computation should uses (the value of) a TText as an argument (because it is used by the TText for which the computation computes its result), then the latest assignment of that TText should become an argument unless there is an ambiguity (i.e. the computation already has an argument that is an assignment for that TText).

 $compVar; ttIsUsedBy\ \ddot{}; ttTrace \cap (compArg; asmVar; asmVar\ \ddot{}) \vdash compArg$

Process rule: Add to Assignments history

 $asmVar \vdash asmHasPred$; $asmVar \cup ttTrace$

Process rule: Assign computation result to TText (provided actual arguments were used)

 $compVar \cong ; compRes \vdash ttValue$

Process rule: Auto DelPair sessionSIA

sessionSIA $\vdash V_{\text{[SESSION*Account]}}; (I_{\text{[Account]}} - \text{accIsActive}); V_{\text{[Account*SESSION]}}$

Process rule: Auto activate auto-login accounts

 $autoLoginAccount \vdash accIsActive$

Process rule: Auto create SIAMPersonRefComputation

personFirstName; personLastName | siamcompFirstName; siamcompLastName

Process rule: Auto maintain 'accPersonRef' relation

 $accPerson; personRef \vdash accPersonRef$

Process rule: Auto maintain 'personRef' relation

 $(personFirstName; siamcompFirstName \lq \cap personLastName; siamcompLastName \lq); siamcompPersonRef \vdash personRef$

Process rule: Automatically depopulate personOrg

 $personOrg \vdash accPerson \ \ ; accOrg$

Process rule: Automatically login

autoLoginAccount \cap accIsActive \cap $V_{[Account*SESSION]};$ ("_SESSION" - sessionAccount; sessionAccount"); $V_{[SESSION*Account]} \vdash \overline{V_{[Account*Account]}}$

Process rule: Automatically populate personOrg

accPerson \(\); accOrg \(\) personOrg

Process rule: By default, a session authentication level is '1'

 $I_{[SESSION]} \cap sessionAccount; sessionAccount \vdash sessionAuthISOLevel; sessionAuthISOLevel \vdash$

Process rule: Clear allowed session roles

 $sessionAllowedRoles \vdash sessionAccount; accAllowedRoles$

Process rule: Clear the OrgRef in a session

 $(I_{\text{[SESSION]}} \cap \text{sessionAccount}; \text{sessionAccount}); \text{sessionOrgRef} \vdash \text{sessionAccount}; \text{accOrgRef}$

Process rule: Create Account upon request

 $loginCreateAccount \cap (sessionAccount; sessionAccount \check{}) \cap loginUserid; loginUserid \check{}\cap \\ registerPassword; registerPassword \check{}\cap sessionPersonRef; sessionPersonRef \check{}\vdash \\ (loginUserid; accUserid \check{}\cap registerPassword; accPassword \check{}\cap sessionPersonRef; accPersonRef \check{}); V_{[Account*SESS]})$

Process rule: $Create\ Assignment\ for\ TTexts\ that\ have\ a\ (new)\ value$

ttValue ⊢ ttTrace; asmVal

Process rule: Create Assignment for TTexts that have no value

 $I_{[\texttt{TText}]} - \texttt{ttValue}; \texttt{ttValue} \succeq \texttt{ttTrace}; (I_{[\texttt{Assignment}]} - \texttt{asmVal}; \texttt{asmVal} \succeq); \texttt{ttTrace} \succeq \texttt{ttTrace}; (I_{[\texttt{Assignment}]} - \texttt{asmVal}; \texttt{asmVal} \succeq); \texttt{ttTrace} \succeq \texttt{ttTrace}$

Process rule: Create computations

 $I_{[\text{TText}]} \cap \text{ttIsUsedBy} \subset \text{compVar} \subset \text{compVar} \subset \text{compVar}$

Process rule: Create/Update copied ttDescr

 $ttICCO; ttDescr \vdash ttDescr$

Process rule: Create/Update copied ttOwner

 $ttICCO; ttOwner \vdash ttOwner$

Process rule: Deactivate 'Developer' role

sessionActiveRoles; "Developer" \(\) sessionDev; sessionActiveRoles; "Developer"

Process rule: Deactivate roles in a session

sessionActiveRoles \vdash sessionAccount; $V_{[Account*Role]}$

Process rule: Deactivation of session Organization

The Organization of the user of the session may be known.

 $sessionOrg \vdash sessionAccount; accOrg$

Process rule: Default roles must be allowed roles

 $accDefaultRoles \vdash accAllowedRoles$

Process rule: Delete computations (that have results) of which one or more arguments have changed

 $\frac{(I_{[\texttt{Computation}]} \cap \texttt{compRes}; \texttt{compRes}^{\smile}); \texttt{compArg}; (I_{[\texttt{Assignment}]} - \texttt{asmVar}; \texttt{ttTrace}) \vdash \overline{V_{[\texttt{Computation}*Assignment]}}$

Process rule: Delete obsolete Assignments

 $I_{[Assignment]} \vdash asmVar; asmVar$

Process rule: Delete obsolete Computations

 $I_{[\text{Computation}]} \vdash \text{compVar}; \text{compVar}$

Process rule: Determine the 'userid' in a session

 $sessionAccount; accUserid \vdash sessionUserid$

Process rule: Determine the OrgRef in a session

 $sessionAccount; accOrgRef \vdash sessionOrgRef$

Process rule: Determine the acting person in a session

 $sessionAccount; accPerson \vdash sessionPerson$

Process rule: Determine the personRef in a session

 $sessionAccount; accPersonRef \vdash sessionPersonRef$

Process rule: Disable automated login

 $\underline{\text{autoLoginAccount}} \cap V_{[\text{Account*SESSION}]}; ("_\text{SESSION}" \cap \text{sessionAccount}; \text{sessionAccount}"); V_{[\text{SESSION*Account}]} \vdash V_{[\text{Account*Account}]}$

Process rule: Every TText with scope and name can be used as a binding

 $I_{[\text{TText}]} \cap \text{ttScope}; \text{ttScope} \subset \text{ttName}; \text{ttName} \subset I_{[\text{Binding}]}$

Process rule: Initialize copied ttDescr

 $(I_{[\text{TText}]} - \text{ttDescr}; \text{ttDescr}); \text{ttICO}; \text{ttDescr} \vdash \text{ttDescr}$

Process rule: Initialize copied ttOwner

 $(I_{[\text{TText}]} - \text{ttOwner}; \text{ttOwner}); \text{ttICO}; \text{ttOwner} \vdash \text{ttOwner}$

Process rule: InsPair sessionAuthISOLevel (MPTrx specific)

We set sessionAuthISOLevel to whatever value is specified at login

 $(I_{[SESSION]} \cap sessionAccount; sessionAccount); loginISOLevel \vdash sessionAuthISOLevel$

Process rule: Limit history size to 2 predecessors

asmHasPred; asmHasPred; asmHasPred $\vdash \overline{V_{[Assignment*Assignment]}}$

Process rule: Logout

A request to logout in a session must be processed in that session.

"_SESSION"; logoutRequest $\vdash \overline{V_{\text{[SESSION*SESSION]}}}$

Process rule: Registering the first authenticated session account

Any account may be used to log into a session for the first time.

("_SESSION" — firstSessionAccount; firstSessionAccount \Bigcirc); sessionAccount \Bigcirc firstSessionAccount

Process rule: Reset login help

sessionLoginAssist \cap sessionAccount; sessionAccount $\ \vdash \overline{V_{\text{[SESSION*SESSION]}}}$

Process rule: SURttName

If a TTName is no longer used in a TText, it is registered no longer.

 $I_{[\text{TTName}]} \vdash \text{ttName} \ \ddot{}; \text{ttName}$

Process rule: SURttPhrase

If a TTPhrase is no longer used in a TText, it is registered no longer.

Process rule: Service U/PW Login request

When a user is authenticated, the corresponding Account will become the sessionAccount (provided it is active).

 $loginReq; (loginUserid; accUserid \lq \cap loginPassword; accPassword \lq); accIsActive \vdash sessionAccount$

Process rule: Service deactivation request

 $accDeactivateReq \vdash sessionAccount$; sessionAccount

Process rule: Sessions with an inactive sessionaccount may not exist

 $sessionAccount \vdash sessionAccount; accIsActive$

Process rule: Set allowed session roles

 $sessionAccount; accAllowedRoles \vdash sessionAllowedRoles$

Process rule: Set default Scope ownership

Ownership of a Scope is assigned by default to the Account that has created the Scope

 $(I_{[\text{Scope}]} - \text{scopeOwner}; \text{scopeOwner} \overset{\smile}{\smile}); V_{[\text{Scope*SESSION}]}; "_\text{SESSION}"; \text{sessionAccount} \vdash \text{scopeOwner}$

Process rule: Set default TText ownership

Ownership of a TText is automatically assigned to the Account that has created the TText

 $(I_{[\texttt{TText}]} - \texttt{ttOwner}; \texttt{ttOwner} \check{}); V_{[\texttt{TText*SESSION}]}; "_\texttt{SESSION}"; \texttt{sessionAccount} \vdash \texttt{ttOwner}$

Process rule: The userid of a customer is its (initial) personnef

 $(I_{[Account]} - accUserid; accUserid); accPersonRef \vdash accUserid; V_{[UserID*PersonRef]})$

Process rule: U/PW Login

When a user is authenticated, the corresponding Account will become the sessionAccount (provided it is active).

(login Userid; acc
Userid $\ \cap\$ login Password; acc Password $\);$ acc Is
Active $\ \ \, \vdash$ session Account

Process rule: Undetermine the 'userid' in a session

 $sessionUserid \vdash sessionAccount; accUserid$

Process rule: Undetermine the acting person in a session

 $sessionPerson \vdash sessionAccount; accPerson$

Process rule: Undetermine the personRef in a session

 $(I_{[SESSION]} \cap sessionAccount; sessionAccount); sessionPersonRef \vdash sessionAccount; accPersonRef$

Process rule: Unowned TTexts that are used by an owned TText, will be assigned the same owner

 $(I_{[\text{TText}]} - \text{ttOwner}; \text{ttOwner}); \text{ttIsUsedBy}; \text{ttOwner} \vdash \text{ttOwner}$

Process rule: Update arguments of computations that do not have a result

 $\frac{(I_{[\text{Computation}]} - \text{compRes}; \text{compRes}); \text{compArg}; (I_{[\text{Assignment}]} - \text{asmVar}; \text{ttTrace}) \vdash \overline{V_{[\text{Computation}*Assignment]}}$

Process rule: Update deleted ttDescr

 $ttICCO; (I_{[TText]} - ttDescr; ttDescr) \vdash (I_{[TText]} - ttDescr; ttDescr); ttICCO$

Process rule: Update deleted ttOwner

 $ttICCO; (I_{[TText]}-ttOwner; ttOwner) \vdash (I_{[TText]}-ttOwner; ttOwner); ttICCO$

5.2.2 Invariants

Invariants are rules that are enforced by the database. It is guaranteed that violations cannot occur in the database.

Invariant: TTexts can only use TTexts that are in the same ttScope $ttIsUsedBy \vdash ttScope$; ttScope

 $Invariant:\ ttValueUsedToReplacePlaceholders\ integrity$

 $ttValueUsedToReplacePlaceholders \vdash ttValue$

Invariant: $ttValueIsUsedInInstanceOfTText\ integrity$

 $ttValueIsUsedInInstanceOfTText \vdash ttIsUsedBy$

Invariant: TrueFalseValid

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, given that all its placeholders have been instantiated.

A statement can be true or false in a context only if it is valid within that context.

 $true \cup false \vdash ttScope$

Invariant: Inconsistency

Every statement (TText) that is in scope is considered valid. This means that it can be true or false, if all its placeholders have been instantiated.

A statement cannot be both true and false in the same context.

 $\mathsf{true} \cap \mathsf{false} \vdash \overline{V_{[\mathsf{TText}*\mathsf{Scope}]}}$

Invariant: An account can only be created for users that are not logged in

Invariant: An account may only be created if it has not been previously registered

 $loginCreateAccount; loginUserid; accUserid \vdash \overline{V_{[SESSION*Account]}}$

Invariant: At least one account must exist that is active

 $V_{[Account*Account]}; accIsActive; V_{[Account*Account]}$

Invariant: Authenticate user

loginUserid¨; loginPassword ⊢ accUserid¨; accPassword

Violations of this rule will result in an error message for the user:

• <violation message should be printed here>

Invariant: AutoLoginAccounts must be active

 $autoLoginAccount \vdash accIsActive$

Invariant: Different computation arguments (assignments) must refer to different TTexts

 $\operatorname{compArg}^{\smile}; \operatorname{compArg} - I_{[\operatorname{Assignment}]} \vdash \operatorname{asmVar}; \overline{I_{[\operatorname{TText}]}}; \operatorname{asmVar}^{\smile}$

Invariant: Every assignment has a successor or represents the current value of a TText

 $V_{\text{[ONE*Assignment]}}; \text{asmHasPred} \cup V_{\text{[ONE*TText]}}; \text{ttTrace}$

Invariant: ISOLevels (LoAs) must be in the range [1-5]

 $I_{[\mathrm{ISOLevel}]} = "1" \cup "2" \cup "3" \cup "4" \cup "5"$

Invariant: If one assignment precedes another, they affect the same TText

 $asmVar \subset smHasPred; asmVar \vdash I_{[TText]}$

Invariant: Integrity of session accounts

At most one user may log into any session so as to guarantee that it is unambiguously clear which account was used for doing things in the system.

At most one account may be activated in any session.

 $("_SESSION" \cap firstSessionAccount; firstSessionAccount"); sessionAccount \vdash firstSessionAccount$

Violations of this rule will result in an error message for the user:

• <violation message should be printed here>

Invariant: The trace of each TText starts with an assignment that has no successors

(ttTrace; asmHasPred)

Invariant: The user must be authenticated with at least the required ISOLevel

sessionAuthISOLevel; sessionRegdISOLevel \vdash isoLevelGTE

Invariant: The value of a TText is per definition the latest assignment value

ttValue = ttTrace; asmVal

Invariant: This file expects to load SIAM version 2.x

"SIAM"; moduleVsnMajor ⊢ moduleVsnMajor; "2"

Violations of this rule will result in an error message for the user:

• <violation message should be printed here>

Invariant: You must specify a password to create an account

 $loginCreateAccount \vdash registerPassword; registerPassword$

Invariant: You must specify a username to create an account

 $loginCreateAccount \vdash loginUserid; loginUserid$

Invariant: UNI object/Claim*LegalThing/

object[Claim*LegalThing] is univalent

object; $\overline{I_{[\text{LegalThing}]}}$; object $\vdash \overline{I_{[\text{Claim}]}}$

Invariant: UNI owner/Scope*Actor/

owner[Scope*Actor] is univalent

owner; $\overline{I_{[Actor]}}$; owner $\vdash \overline{I_{[Scope]}}$

Invariant: UNI created/Scope *Moment/

 ${\tt created[Scope*Moment]} \ is \ univalent$

created; $\overline{I_{[Moment]}}$; created $\vdash \overline{I_{[Scope]}}$

Invariant: $UNI\ ttTemplateParsedText/TText*TTPhrase$

ttTemplateParsedText[TText*TTPhrase] is univalent

Invariant: $IRF\ ttIsUsedBy[TText*TText]$

 $ttIsUsedBy[TText*TText] \ is \ irreflexive$

 $\mathsf{ttIsUsedBy} \vdash \overline{I_{[\mathsf{TText}]}}$

 ${\bf Invariant:}\ ASY\ ttIsUsedBy[TText*TText]$

ttIsUsedBy[TText*TText] is antisymmetric

 $\mathbf{ttIsUsedBy} ^{\smallsmile} \cap \mathbf{ttIsUsedBy} \vdash I_{[\mathbf{TText}]}$

Invariant: $IRF\ ttIsUsedByStar[TText*TText]$

 $ttIsUsedByStar[TText*TText] \ is \ irreflexive$

 $ttIsUsedByStar \vdash \overline{I_{[TText]}}$

Invariant: $ASY\ ttIsUsedByStar[TText*TText]$

 $ttIsUsedByStar[TText*TText] \ is \ antisymmetric$

 $\mathbf{ttIsUsedByStar}^{\smile} \cap \mathbf{ttIsUsedByStar} \vdash I_{[\mathsf{TText}]}$

Invariant: $TRN \ ttIsUsedByStar[TText*TText]$

 $ttIsUsedByStar[TText*TText] \ is \ transitive$ $ttIsUsedByStar; ttIsUsedByStar \vdash ttIsUsedByStar$

$Invariant: \ UNI\ ttValue\ Used\ ToReplacePlaceholders[TText*TTValue]$

$$\label{eq:tvalue} \begin{split} &\text{ttValueUsedToReplacePlaceholders}[\text{TText*TTValue}] \text{ is univalent} \\ &\text{ttValueUsedToReplacePlaceholders}; \\ &\overline{I_{[\text{TTValue}]}}; \\ &\text{ttValueUsedToReplacePlaceholders} &\hookrightarrow \\ &\vdash \\ &\overline{I_{[\text{TText}]}} \end{split}$$

Invariant: SYM ttInstanceResetReq[TText*TText]

 $ttInstanceResetReq[TText*TText] \ is \ symmetric$ $ttInstanceResetReq = ttInstanceResetReq ^{\sim}$

Invariant: ASY ttInstanceResetReq[TText*TText]

tt InstanceResetReq
[TText*TText] is antisymmetric tt InstanceResetReq
 \cap tt InstanceResetReq
 $\vdash I_{\texttt{[TText]}}$

$Invariant: \ UNI \ ttInstanceResetReq/TText*TText/$

tt InstanceResetReq[TText*TText] is univalent tt InstanceResetReq; $\overline{I_{[\text{TText}]}}$; tt InstanceResetReq $\ \vdash \overline{I_{[\text{TText}]}}$

$Invariant: \ INJ \ ttInstanceResetReq[TText*TText]$

$$\begin{split} & \text{ttInstanceResetReq}[\text{TText*TText}] \text{ is injective} \\ & \text{ttInstanceResetReq} \vdash \overline{I_{[\text{TText}]}}; \\ & \text{ttInstanceResetReq} \vdash \overline{I_{[\text{TText}]}} \end{split}$$

Invariant: $IRF\ ttValueIsUsedInInstanceOfTText/TText*TText/$

tt Value Is Used In
Instance OfTText[TText*TText] is irreflexive tt Value Is Used In
Instance OfTText $\vdash \overline{I_{[\text{TText}]}}$

$Invariant: \ ASY \ ttValue Is Used In Instance Of TText[TText*TText]$

tt Value Is Used In
Instance OfTText \top \cap tt Value Is Used In
Instance OfTText $\vdash I_{[\text{TText}]}$

${\bf Invariant:}\ \ UNI\ true[TText*Scope]$

 $\begin{aligned} & \text{true}[\text{TText*Scope}] \text{ is univalent} \\ & \text{true}; \overline{I_{[\text{Scope}]}}; \text{true}^{\smile} \vdash \overline{I_{[\text{TText}]}} \end{aligned}$

${\bf Invariant:}\ \ UNI\ false[TText*Scope]$

false[TText*Scope] is univalent false; $\overline{I_{[Scope]}}$; false $\vdash \overline{I_{[TText]}}$

Invariant: $UNI\ ttScope/TText*Scope/$

ttScope[TText*Scope] is univalent ttScope; $\overline{I_{[Scope]}}$; ttScope $\overset{\smile}{\vdash}$ $\overline{I_{[TText]}}$

Invariant: $TOT \ ttScope[TText*Scope]$

$$\begin{split} & ttScope[TText*Scope] \text{ is total} \\ & I_{[TText]} \vdash ttScope; ttScope \check{\ } \end{split}$$

Invariant: $UNI\ observed/TText*Moment/$

 $\begin{aligned} & \text{observed}[\text{TText*Moment}] \text{ is univalent} \\ & \text{observed}; \overline{I_{[\text{Moment}]}}; \text{observed}^{\smallsmile} \vdash \overline{I_{[\text{TText}]}} \end{aligned}$

 $Invariant: \ UNI \ regeling[RechtsGrond*Regeling]$

$$\begin{split} & \operatorname{regeling}[\operatorname{RechtsGrond*Regeling}] \text{ is univalent} \\ & \operatorname{regeling}; \overline{I_{[\operatorname{Regeling}]}}; \operatorname{regeling} & \vdash \overline{I_{[\operatorname{RechtsGrond}]}} \end{split}$$

Invariant: $UNI\ bijlage[RechtsGrond*Bijlage]$

$$\begin{split} & \text{bijlage[RechtsGrond*Bijlage] is univalent} \\ & \text{bijlage; } \overline{I_{\text{[Bijlage]}}}; \\ & \text{bijlage} \vdash \overline{I_{\text{[RechtsGrond]}}} \end{split}$$

Invariant: UNI aanwijzing/RechtsGrond*Aanwijzing/

aanwijzing[RechtsGrond*Aanwijzing] is univalent aanwijzing; $\overline{I_{[\text{Aanwijzing}]}}$; aanwijzing $\ \vdash \overline{I_{[\text{RechtsGrond}]}}$

 ${\bf Invariant:}\ \ UNI\ bwb[Regeling*BasisWettenBestand]$

bwb[Regeling*BasisWettenBestand] is univalent bwb; $\overline{I_{[\text{BasisWettenBestand}]}}$; bwb $^{\smile} \vdash \overline{I_{[\text{Regeling}]}}$

 $Invariant: \ UNI \ titel[Regeling*Titel]$

$$\begin{split} & \text{titel}[\text{Regeling*Titel}] \text{ is univalent} \\ & \text{titel}; \overline{I_{[\text{Titel}]}}; \text{titel} & \vdash \overline{I_{[\text{Regeling}]}} \end{split}$$

Invariant: $TOT\ titel[Regeling*Titel]$

titel[Regeling*Titel] is total $I_{\text{[Regeling]}} \vdash \text{titel}; \text{titel}^{\sim}$

 $Invariant: \ UNI \ artikelen [Regeling*Aantal]$

artikelen [Regeling*Aantal] is univalent artikelen; $\overline{I_{[\text{Aantal}]}}$; artikelen $\overset{\smile}{=} \overline{I_{[\text{Regeling}]}}$

Invariant: UNI gedelegeerdUit/Regeling*Regeling/

gedelegeerdUit[Regeling*Regeling] is univalent gedelegeerdUit; $\overline{I_{[\text{Regeling}]}}$; gedelegeerdUit $^{\smile} \vdash \overline{I_{[\text{Regeling}]}}$

Invariant: UNI soort/Regeling*Regelingsoort/

 $\begin{aligned} & \operatorname{soort}[\operatorname{Regeling*Regelingsoort}] \text{ is univalent} \\ & \operatorname{soort}; \overline{I_{[\operatorname{Regelingsoort}]}}; \operatorname{soort}^{\smile} \vdash \overline{I_{[\operatorname{Regeling}]}} \end{aligned}$

Invariant: UNI processlow/Regeling*Processlow/

$$\begin{split} & \operatorname{procesflow}[\operatorname{Regeling}^*\operatorname{Procesflow}] \text{ is univalent} \\ & \operatorname{procesflow}; \overline{I_{[\operatorname{Procesflow}]}}; \operatorname{procesflow}^{\smile} \vdash \overline{I_{[\operatorname{Regeling}]}} \end{split}$$

Invariant: UNI afkorting[Regeling*Tekst]

afkorting [Regeling*Tekst] is univalent afkorting; $\overline{I_{[\text{Tekst}]}}$; afkorting $\overset{\smile}{=} \overline{I_{[\text{Regeling}]}}$

Invariant: $UNI\ class[TText*Concept]$

 $\begin{aligned} & \operatorname{class}[\operatorname{TText}^*\operatorname{Concept}] \text{ is univalent} \\ & \operatorname{class}; \overline{I_{[\operatorname{Concept}]}}; \operatorname{class}^{\smile} \vdash \overline{I_{[\operatorname{TText}]}} \end{aligned}$

Invariant: SYM autoLoginAccount/Account*Account/

Invariant: $ASY\ autoLoginAccount[Account*Account]$

 ${\bf autoLoginAccount[Account*Account] \ is \ antisymmetric}$ ${\bf autoLoginAccount} ' \cap {\bf autoLoginAccount} \vdash I_{[{\bf Account}]}$

Invariant: UNI autoLoginAccount/Account*Account/

autoLoginAccount[Account*Account] is univalent autoLoginAccount; $\overline{I_{[Account]}}$; autoLoginAccount $\vdash \overline{I_{[Account]}}$

Invariant: INJ autoLoginAccount/Account*Account]

 $autoLoginAccount[Account^*Account] is injective \\ autoLoginAccount^{\smile}; \overline{I_{[Account]}}; autoLoginAccount \vdash \overline{I_{[Account]}}$

Invariant: INJ accUserid/Account*UserID/

accUserid[Account*UserID] is injective accUserid $\overline{I}_{[Account]}$; accUserid $\overline{I}_{[UserID]}$

Invariant: UNI accUserid/Account*UserID/

accUserid[Account*UserID] is univalent accUserid; $\overline{I_{[\text{UserID}]}}$; accUserid $\ \vdash \overline{I_{[\text{Account}]}}$

Invariant: $TOT\ acc Userid[Account*UserID]$

$$\begin{split} & \operatorname{accUserid}[\operatorname{Account}^*\operatorname{UserID}] \text{ is total} \\ & I_{[\operatorname{Account}]} \vdash \operatorname{accUserid}; \operatorname{accUserid}^{\smile} \end{split}$$

Invariant: $UNI\ accPassword/Account*Password/$

$$\begin{split} & \operatorname{accPassword}[\operatorname{Account}^*\operatorname{Password}] \text{ is univalent} \\ & \operatorname{accPassword}; \overline{I_{[\operatorname{Password}]}}; \operatorname{accPassword}^{\smile} \vdash \overline{I_{[\operatorname{Account}]}} \end{split}$$

Invariant: $UNI\ sessionAccount[SESSION*Account]$

$$\begin{split} & \operatorname{sessionAccount}[\operatorname{SESSION*Account}] \text{ is univalent} \\ & \operatorname{sessionAccount}; \overline{I_{[\operatorname{Account}]}}; \operatorname{sessionAccount} ^{\smile} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: $SYM\ acc Is Initialized | Account^*Account|$

 ${\it accIsInitialized[Account*Account] is symmetric} \\ {\it accIsInitialized} = {\it accIsInitialized} \check{\ }$

Invariant: ASY accIsInitialized/Account*Account/

 $\label{eq:accIsInitialized} $$\operatorname{accIsInitialized} \cap \operatorname{accIsInitialized} \vdash I_{[\operatorname{Account}]}$$

Invariant: UNI accIsInitialized/Account*Account/

$$\begin{split} & \text{accIsInitialized[Account*Account] is univalent} \\ & \text{accIsInitialized;} \overline{I_{[\text{Account}]}}; \text{accIsInitialized} ^{\smile} \vdash \overline{I_{[\text{Account}]}} \end{split}$$

Invariant: INJ accIsInitialized/Account*Account]

$$\begin{split} & accIsInitialized[Account*Account] \ is \ injective \\ & accIsInitialized\ \ \ ; \overline{I_{[Account]}}; accIsInitialized\ \ \ \ \overline{I_{[Account]}} \end{split}$$

Invariant: SYM accIsActive/Account*Account/

$$\begin{split} & \operatorname{accIsActive}[\operatorname{Account}^* \operatorname{Account}] \text{ is symmetric} \\ & \operatorname{accIsActive} = \operatorname{accIsActive}^{\smile} \end{split}$$

Invariant: ASY accIsActive/Account*Account/

accIsActive[Account* Account] is antisymmetric accIsActive $' \cap$ accIsActive $\vdash I_{[Account]}$

Invariant: UNI accIsActive/Account*Account/

accIsActive[Account*Account] is univalent accIsActive; $\overline{I_{[Account]}}$; accIsActive $\vdash \overline{I_{[Account]}}$

Invariant: INJ accIsActive/Account*Account/

$$\begin{split} & \text{accIsActive}[\text{Account*Account}] \text{ is injective} \\ & \text{accIsActive}^{\smile}; \overline{I_{[\text{Account}]}}; \text{accIsActive} \vdash \overline{I_{[\text{Account}]}} \end{split}$$

Invariant: $SYM\ accDeactivateReq[Account*Account]$

$$\label{eq:accDeactivateReq} \begin{split} & accDeactivateReq[Account*Account] \ is \ symmetric \\ & accDeactivateReq = accDeactivateReq \ \end{split}$$

Invariant: $ASY\ accDeactivateReq[Account*Account]$

accDeactivateReq[Account*Account] is antisymmetric accDeactivateReq $\ \cap$ accDeactivateReq $\ \vdash I_{[Account]}$

Invariant: $UNI\ accDeactivateReq[Account*Account]$

$$\begin{split} & accDeactivateReq[Account*Account] \text{ is univalent} \\ & accDeactivateReq; \overline{I_{[Account]}}; accDeactivateReq ^{\smile} \vdash \overline{I_{[Account]}} \end{split}$$

Invariant: INJ accDeactivateReq[Account*Account]

$$\begin{split} & \text{accDeactivateReq[Account*Account] is injective} \\ & \text{accDeactivateReq} \vdash \overline{I_{[\text{Account}]}}; \\ & \text{accDeactivateReq} \vdash \overline{I_{[\text{Account}]}} \end{split}$$

Invariant: UNI sessionUserid[SESSION*UserID]

sessionUserid[SESSION*UserID] is univalent sessionUserid; $\overline{I_{[\text{UserID}]}}$; sessionUserid $\ \vdash \overline{I_{[\text{SESSION}]}}$

Invariant: UNI loginUserid[SESSION*UserID]

$$\begin{split} & \text{loginUserid[SESSION*UserID] is univalent} \\ & \text{loginUserid;} \overline{I_{[\text{UserID}]}}; \\ & \text{loginUserid}^{\smile} \vdash \overline{I_{[\text{SESSION}]}} \end{split}$$

Invariant: UNI loginPassword/SESSION*Password/

$$\begin{split} & \operatorname{loginPassword[SESSION*Password]} \text{ is univalent} \\ & \operatorname{loginPassword;} \overline{I_{[Password]}}; \operatorname{loginPassword}^{\smile} \vdash \overline{I_{[SESSION]}} \end{split}$$

Invariant: SYM logoutRequest/SESSION*SESSION

$$\begin{split} & logoutRequest[SESSION*SESSION] \ is \ symmetric \\ & logoutRequest = logoutRequest \ \ \ \ \end{split}$$

Invariant: ASY logoutRequest[SESSION*SESSION]

 $logoutRequest[SESSION*SESSION] is antisymmetric \\ logoutRequest <math>\vdash I_{[SESSION]}$

Invariant: UNI logoutRequest/SESSION*SESSION]

$$\begin{split} & \operatorname{logoutRequest[SESSION*SESSION]} \text{ is univalent} \\ & \operatorname{logoutRequest;} \overline{I_{[SESSION]}}; \operatorname{logoutRequest} ^{\smile} \vdash \overline{I_{[SESSION]}} \end{split}$$

Invariant: INJ logoutRequest/SESSION*SESSION|

$$\begin{split} & \operatorname{logoutRequest[SESSION*SESSION]} \text{ is injective} \\ & \operatorname{logoutRequest}\check{\ \ }; \overline{I_{[\operatorname{SESSION}]}}; \operatorname{logoutRequest} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: RFX isoLevelGTE/ISOLevel*ISOLevel

isoLevelGTE[ISOLevel*ISOLevel] is reflexive $I_{\texttt{[ISOLevel]}} \vdash \texttt{isoLevelGTE}$

Invariant: ASY isoLevelGTE/ISOLevel*ISOLevel

isoLevelGTE[ISOLevel*ISOLevel] is antisymmetric isoLevelGTE \cap isoLevelGTE \vdash $I_{[ISOLevel]}$

Invariant: TRN isoLevelGTE/ISOLevel*ISOLevel]

isoLevelGTE[ISOLevel*ISOLevel] is transitive isoLevelGTE; isoLevelGTE \vdash isoLevelGTE

$Invariant: \ UNI \ session ReqdISOL evel [SESSION*ISOL evel]$

$$\begin{split} & \operatorname{sessionReqdISOLevel[SESSION*ISOLevel]} \text{ is univalent} \\ & \operatorname{sessionReqdISOLevel}; \overline{I_{[\mathrm{ISOLevel}]}}; \operatorname{sessionReqdISOLevel} & \vdash \overline{I_{[\mathrm{SESSION}]}} \end{split}$$

Invariant: $UNI\ sessionAuthISOLevel/SESSION*ISOLevel/$

sessionAuthISOLevel[SESSION*ISOLevel] is univalent sessionAuthISOLevel; $\overline{I_{[\text{ISOLevel}]}}$; sessionAuthISOLevel ${}^{\smile}$ \vdash $\overline{I_{[\text{SESSION}]}}$

Invariant: UNI moduleName/Module*ModuleName/

$$\begin{split} & \operatorname{moduleName}[\operatorname{Module*ModuleName}] \text{ is univalent} \\ & \operatorname{moduleName}; \overline{I_{[\operatorname{ModuleName}]}}; \operatorname{moduleName} & \vdash \overline{I_{[\operatorname{Module}]}} \end{split}$$

Invariant: TOT moduleName/Module*ModuleName/

$$\begin{split} & \operatorname{moduleName}[\operatorname{Module*ModuleName}] \text{ is total} \\ & I_{[\operatorname{Module}]} \vdash \operatorname{moduleName}; \operatorname{moduleName} \check{} \end{split}$$

Invariant: UNI moduleVsnMajor/Module*ModuleVsnMajor/

$$\begin{split} & \operatorname{moduleVsnMajor}[\operatorname{Module}*\operatorname{ModuleVsnMajor}] \text{ is univalent} \\ & \operatorname{moduleVsnMajor}; \overline{I_{[\operatorname{ModuleVsnMajor}]}}; \operatorname{moduleVsnMajor} & \vdash \overline{I_{[\operatorname{Module}]}} \end{split}$$

$Invariant: \ TOT \ module VsnMajor [Module *Module VsnMajor]$

$$\begin{split} & \text{moduleVsnMajor}[\text{Module*ModuleVsnMajor}] \text{ is total} \\ & I_{[\text{Module}]} \vdash \text{moduleVsnMajor}; \text{moduleVsnMajor} \\ & \end{split}$$

$Invariant: \ UNI \ module VsnMinor [Module *Module VsnMinor]$

$$\begin{split} & \operatorname{moduleVsnMinor}[\operatorname{Module*ModuleVsnMinor}] \text{ is univalent} \\ & \operatorname{moduleVsnMinor}; \overline{I_{[\operatorname{ModuleVsnMinor}]}}; \operatorname{moduleVsnMinor} & \vdash \overline{I_{[\operatorname{Module}]}} \end{split}$$

Invariant: TOT moduleVsnMinor/Module*ModuleVsnMinor/

moduleVsnMinor[Module*ModuleVsnMinor] is total $I_{[Module]} \vdash moduleVsnMinor; moduleVsnMinor$

Invariant: INJ orgAbbrName[Organization*OrgAbbrName]

$$\begin{split} & \text{orgAbbrName}[\text{Organization*OrgAbbrName}] \text{ is injective} \\ & \text{orgAbbrName}^{\smile}; \overline{I_{[\text{Organization}]}}; \text{orgAbbrName} \vdash \overline{I_{[\text{OrgAbbrName}]}} \end{split}$$

Invariant: UNI orgAbbrName[Organization*OrgAbbrName]

 $\label{eq:orgAbbrName} orgAbbrName[Organization*OrgAbbrName] is univalent \\ orgAbbrName; \overline{I_{[OrgAbbrName]}}; orgAbbrName <math>\ \vdash \overline{I_{[Organization]}}$

Invariant: INJ orgFullName[Organization*OrgFullName]

orgFullName[Organization*OrgFullName] is injective orgFullName $\[\] \overline{I_{[Organization]}};$ orgFullName $\[\] \overline{I_{[OrgFullName]}}$

Invariant: UNI orgFullName[Organization*OrgFullName]

orgFullName[Organization*OrgFullName] is univalent orgFullName; $\overline{I_{[\text{OrgFullName}]}}$; orgFullName $\vdash \overline{I_{[\text{Organization}]}}$

Invariant: UNI accOrg/Account*Organization

$$\begin{split} & \operatorname{accOrg}[\operatorname{Account}^*\operatorname{Organization}] \text{ is univalent} \\ & \operatorname{accOrg}; \overline{I_{[\operatorname{Organization}]}}; \operatorname{accOrg}^{\smile} \vdash \overline{I_{[\operatorname{Account}]}} \end{split}$$

Invariant: UNI sessionOrg/SESSION*Organization/

sessionOrg[SESSION*Organization] is univalent sessionOrg; $\overline{I_{[Organization]}}$; sessionOrg $^{\smile} \vdash \overline{I_{[SESSION]}}$

Invariant: UNI personFirstName/Person*FirstName/

 $\begin{aligned} & \operatorname{personFirstName}[\operatorname{Person*FirstName}] \text{ is univalent} \\ & \operatorname{personFirstName}; \overline{I_{[\operatorname{FirstName}]}}; \operatorname{personFirstName} & \vdash \overline{I_{[\operatorname{Person}]}} \end{aligned}$

Invariant: TOT personFirstName/Person*FirstName/

personFirstName[Person*FirstName] is total $I_{[Person]} \vdash personFirstName; personFirstName \overset{\smile}{}$

Invariant: UNI personLastName/Person*LastName/

 $personLastName[Person*LastName] is univalent \\ personLastName; \overline{I_{[LastName]}}; personLastName \cite{Green} + \overline{I_{[Person]}}$

Invariant: $TOT\ personLastName[Person*LastName]$

$$\begin{split} & personLastName[Person*LastName] \text{ is total} \\ & I_{[Person]} \vdash personLastName; personLastName \\ & \\ \end{split}$$

Invariant: UNI accPerson/Account*Person/

$$\begin{split} & accPerson[Account*Person] \text{ is univalent} \\ & accPerson; \overline{I_{[Person]}}; accPerson \\ & \vdash \overline{I_{[Account]}} \end{split}$$

Invariant: UNI sessionPerson[SESSION*Person]

 $sessionPerson[SESSION*Person] is univalent \\ sessionPerson; \overline{I_{[Person]}}; sessionPerson \\ ^{\backsim} \vdash \overline{I_{[SESSION]}}$

${\bf Invariant:}\ \ UNI\ first Session Account [SESSION*Account]$

 $\begin{aligned} & \text{firstSessionAccount}[\text{SESSION*Account}] \text{ is univalent} \\ & \text{firstSessionAccount}; \overline{I_{[\text{Account}]}}; \\ & \text{firstSessionAccount} & \vdash \overline{I_{[\text{SESSION}]}} \end{aligned}$

${\bf Invariant:}\ \ UNI\ uidUserid[UID*UserID]$

$$\begin{split} & \text{uidUserid}[\text{UID*UserID}] \text{ is univalent} \\ & \text{uidUserid}; \overline{I_{[\text{UserID}]}}; \text{uidUserid}^{\smile} \vdash \overline{I_{[\text{UID}]}} \end{split}$$

$Invariant: \ TOT \ uidUserid[UID*UserID]$

uidUserid[UID*UserID] is total $I_{\text{[UID]}} \vdash \text{uidUserid}; \text{uidUserid}^{\sim}$

${\bf Invariant:}\ \ UNI\ uidIssuer[UID*IdP]$

uidIssuer[UID*IdP] is univalent uidIssuer; $\overline{I_{[\text{IdP}]}}$; uidIssuer $\ \vdash \overline{I_{[\text{UID}]}}$

Invariant: $INJ \ acc \ UID / Account * UID /$

accUID[Account*UID] is injective

 $\text{accUID}^{\smile}; \overline{I_{[\text{Account}]}}; \text{accUID} \vdash \overline{I_{[\text{UID}]}}$

Invariant: UNI personRef/Person*PersonRef/

 $personRef[Person*PersonRef] \ is \ univalent$

 $\operatorname{personRef}; \overline{I_{[\operatorname{PersonRef}]}}; \operatorname{personRef} \vdash \overline{I_{[\operatorname{Person}]}}$

Invariant: UNI accPersonRef/Account*PersonRef/

accPersonRef[Account*PersonRef] is univalent

 $\mathbf{accPersonRef}; \overline{I_{[\mathsf{PersonRef}]}}; \mathbf{accPersonRef} \overset{\smile}{\smile} \overline{I_{[\mathsf{Account}]}}$

Invariant: UNI sessionPersonRef[SESSION*PersonRef]

 $session Person Ref[SESSION*Person Ref] \ is \ univalent$

 $sessionPersonRef; \overline{I_{[PersonRef]}}; sessionPersonRef \sim \overline{I_{[SESSION]}}$

$Invariant: \ UNI\ siam compFirstName [SIAMPersonRefComputation*FirstName]$

 $siamcompFirstName[SIAMPersonRefComputation*FirstName] \ is \ univalent$

 $siamcompFirstName; \overline{I_{[FirstName]}}; siamcompFirstName \lq \vdash \overline{I_{[SIAMPersonRefComputation]}}$

$Invariant: \ UNI\ siam compLastName [SIAMPersonRefComputation*LastName]$

 $siamcompLastName[SIAMPersonRefComputation*LastName] \ is \ univalent$

$Invariant: \ UNI\ siam compPersonRef[SIAMPersonRefComputation*PersonRef]$

 $siamcompPersonRef[SIAMPersonRefComputation*PersonRef]\ is\ univalent$

 $siamcompPersonRef; \overline{I_{[PersonRef]}}; siamcompPersonRef \sim \overline{I_{[SIAMPersonRefComputation]}}$

Invariant: UNI accOrgRef/Account*OrgRef/

 ${\it accOrgRef[Account*OrgRef]} \ is \ univalent$

 $\mathbf{accOrgRef}; \overline{I_{[\mathsf{OrgRef}]}}; \mathbf{accOrgRef} \vdash \overline{I_{[\mathsf{Account}]}}$

$Invariant: \ UNI \ sessionOrgRef[SESSION*OrgRef]$

 ${\bf sessionOrgRef[SESSION*OrgRef]} \ {\bf is} \ {\bf univalent}$

 $\operatorname{sessionOrgRef}; \overline{I_{[\operatorname{OrgRef}]}}; \operatorname{sessionOrgRef} \widetilde{\smile} \overline{I_{[\operatorname{SESSION}]}}$

$Invariant: \ SYM \ loginCreateAccount[SESSION*SESSION]$

 $loginCreateAccount[SESSION*SESSION] \ is \ symmetric$

 $loginCreateAccount = loginCreateAccount \cite{count}$

Invariant: ASY loginCreateAccount[SESSION*SESSION]

$$\label{eq:count} \begin{split} & \operatorname{loginCreateAccount}[\operatorname{SESSION*SESSION}] \text{ is antisymmetric} \\ & \operatorname{loginCreateAccount}^{\smile} \cap \operatorname{loginCreateAccount} \vdash I_{[\operatorname{SESSION}]} \end{split}$$

Invariant: UNI loginCreateAccount[SESSION*SESSION]

 $\begin{aligned} & \operatorname{loginCreateAccount}[\operatorname{SESSION}^*\operatorname{SESSION}] \text{ is univalent} \\ & \operatorname{loginCreateAccount}; \overline{I_{[\operatorname{SESSION}]}}; \operatorname{loginCreateAccount}^{\smile} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{aligned}$

Invariant: $INJ\ loginCreateAccount[SESSION*SESSION]$

$$\begin{split} & \operatorname{loginCreateAccount}[\operatorname{SESSION}^*\operatorname{SESSION}] \text{ is injective} \\ & \operatorname{loginCreateAccount}^{\smile}; \overline{I_{[\operatorname{SESSION}]}}; \operatorname{loginCreateAccount} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: UNI registerPassword[SESSION*Password]

$$\begin{split} & \operatorname{registerPassword[SESSION*Password]} \text{ is univalent} \\ & \operatorname{registerPassword;} \overline{I_{[\operatorname{Password}]}}; \operatorname{registerPassword}^{\smile} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: UNI registerOrgRef/SESSION*OrgRef/

$$\begin{split} & \operatorname{registerOrgRef[SESSION*OrgRef]} \text{ is univalent} \\ & \operatorname{registerOrgRef;} \overline{I_{[\operatorname{OrgRef}]}}; \operatorname{registerOrgRef} & \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: SYM loginReq[SESSION*SESSION]

$$\begin{split} & \operatorname{loginReq}[\operatorname{SESSION*SESSION}] \text{ is symmetric} \\ & \operatorname{loginReq} = \operatorname{loginReq} \widetilde{} \end{split}$$

Invariant: ASY loginReq[SESSION*SESSION]

 $\begin{aligned} & \operatorname{loginReq[SESSION*SESSION]} \text{ is antisymmetric} \\ & \operatorname{loginReq}^{\smile} \cap \operatorname{loginReq} \vdash I_{[\operatorname{SESSION}]} \end{aligned}$

Invariant: UNI loginReq[SESSION*SESSION]

 $\begin{aligned} & \operatorname{loginReq[SESSION*SESSION]} \text{ is univalent} \\ & \operatorname{loginReq;} \overline{I_{[SESSION]}}; \operatorname{loginReq} ^{\smile} \vdash \overline{I_{[SESSION]}} \end{aligned}$

Invariant: INJ loginReq/SESSION*SESSION

$$\begin{split} & \operatorname{loginReq[SESSION*SESSION]} \text{ is injective} \\ & \operatorname{loginReq}^{\smile}; \overline{I_{[SESSION]}}; \operatorname{loginReq} \vdash \overline{I_{[SESSION]}} \end{split}$$

Invariant: UNI loginISOLevel/SESSION*ISOLevel/

$$\begin{split} & \text{loginISOLevel}[\text{SESSION*ISOLevel}] \text{ is univalent} \\ & \text{loginISOLevel}; \overline{I_{\text{[ISOLevel]}}}; \text{loginISOLevel} & \vdash \overline{I_{\text{[SESSION]}}} \end{split}$$

Invariant: SYM sessionDev[SESSION*SESSION]

Invariant: $ASY\ sessionDev[SESSION*SESSION]$

session Dev[SESSION*SESSION] is antisymmetric session Dev $^{\backsim}$ \cap session Dev \vdash $I_{\text{[SESSION]}}$

Invariant: $UNI\ sessionDev[SESSION*SESSION]$

 $sessionDev[SESSION*SESSION] is univalent \\ sessionDev; \overline{I_{[SESSION]}}; sessionDev \cupe \vdash \overline{I_{[SESSION]}}$

Invariant: INJ sessionDev[SESSION*SESSION]

$$\begin{split} & sessionDev[SESSION*SESSION] \text{ is injective} \\ & sessionDev \, \bar{'}\,; \overline{I_{[SESSION]}}; sessionDev \, \bar{'}\, \overline{I_{[SESSION]}} \end{split}$$

Invariant: INJ ttTrace[TText*Assignment]

$$\begin{split} & \text{ttTrace}[\text{TText*Assignment}] \text{ is injective} \\ & \text{ttTrace}^{\smile}; \overline{I_{[\text{TText}]}}; \text{ttTrace} \vdash \overline{I_{[\text{Assignment}]}} \end{split}$$

Invariant: UNI ttTrace[TText*Assignment]

 $\begin{aligned} & \text{ttTrace}[\text{TText*Assignment}] \text{ is univalent} \\ & \text{ttTrace}; \overline{I_{[\text{Assignment}]}}; \text{ttTrace} & \vdash \overline{I_{[\text{TText}]}} \end{aligned}$

Invariant: TOT ttTrace/TText*Assignment/

$$\begin{split} & \text{ttTrace}[\text{TText*Assignment}] \text{ is total} \\ & I_{[\text{TText}]} \vdash \text{ttTrace}; \text{ttTrace} \check{\ } \end{split}$$

Invariant: $UNI \ asmVar[Assignment*TText]$

 $\begin{aligned} & \operatorname{asmVar}[\operatorname{Assignment*TText}] \text{ is univalent} \\ & \operatorname{asmVar}; \overline{I_{[\operatorname{TText}]}}; \operatorname{asmVar}^{\smile} \vdash \overline{I_{[\operatorname{Assignment}]}} \end{aligned}$

$Invariant: \ TOT \ asmVar[Assignment*TText]$

 $\begin{aligned} & \operatorname{asmVar}[\operatorname{Assignment*TText}] \text{ is total} \\ & I_{[\operatorname{Assignment}]} \vdash \operatorname{asmVar}; \operatorname{asmVar} \check{} \end{aligned}$

$Invariant: \ UNI \ asmVal[Assignment*TTValue]$

$$\begin{split} & \operatorname{asmVal}[\operatorname{Assignment*TTValue}] \text{ is univalent} \\ & \operatorname{asmVal}; \overline{I_{[\operatorname{TTValue}]}}; \operatorname{asmVal}^{\smile} \vdash \overline{I_{[\operatorname{Assignment}]}} \end{split}$$

Invariant: UNI asmPOT/Assignment*PointOfTime/

asmPOT[Assignment*PointOfTime] is univalent asmPOT; $\overline{I_{[PointOfTime]}}$; asmPOT $^{\sim} \vdash \overline{I_{[Assignment]}}$

Invariant: UNI asmHasPred/Assignment*Assignment]

 $asmHasPred[Assignment*Assignment] is univalent \\ asmHasPred; \overline{I_{[Assignment]}}; asmHasPred \vdash \overline{I_{[Assignment]}}$

Invariant: INJ asmHasPred[Assignment*Assignment]

$$\begin{split} & \operatorname{asmHasPred[Assignment*Assignment] is \ injective} \\ & \operatorname{asmHasPred}^{\smile}; \overline{I_{[\mathrm{Assignment}]}}; \operatorname{asmHasPred} \vdash \overline{I_{[\mathrm{Assignment}]}} \end{split}$$

Invariant: INJ comp Var[Computation*TText]

$$\begin{split} & \operatorname{compVar}[\operatorname{Computation*TText}] \text{ is injective} \\ & \operatorname{compVar}^{\smile}; \overline{I_{[\operatorname{Computation}]}}; \operatorname{compVar} \vdash \overline{I_{[\operatorname{TText}]}} \end{split}$$

Invariant: UNI compVar/Computation*TText/

$$\begin{split} & \operatorname{compVar}[\operatorname{Computation*TText}] \text{ is univalent} \\ & \operatorname{compVar}; \overline{I_{[\operatorname{TText}]}}; \operatorname{compVar} ^{\smile} \vdash \overline{I_{[\operatorname{Computation}]}} \end{split}$$

Invariant: TOT comp Var / Computation * TText /

compVar[Computation*TText] is total $I_{[Computation]} \vdash compVar; compVar$

Invariant: $UNI\ compRes[Computation*TTValue]$

compRes[Computation*TTValue] is univalent compRes; $\overline{I_{[\text{TTValue}]}}$; compRes $\[\vdash \overline{I_{[\text{Computation}]}} \]$

Invariant: SYM sessionLoginAssist/SESSION*SESSION|

 $sessionLoginAssist[SESSION*SESSION] is symmetric \\ sessionLoginAssist = sessionLoginAssist ^{\smile}$

Invariant: ASY sessionLoginAssist[SESSION*SESSION]

sessionLoginAssist[SESSION*SESSION] is antisymmetric sessionLoginAssist $\ \cap$ sessionLoginAssist $\ \vdash I_{[SESSION]}$

Invariant: UNI sessionLoginAssist/SESSION*SESSION/

sessionLoginAssist[SESSION*SESSION] is univalent sessionLoginAssist; $\overline{I_{[SESSION]}}$; sessionLoginAssist $\ \vdash \overline{I_{[SESSION]}}$

Invariant: INJ sessionLoginAssist/SESSION*SESSION|

$$\begin{split} & \operatorname{sessionLoginAssist}[\operatorname{SESSION*SESSION}] \text{ is injective} \\ & \operatorname{sessionLoginAssist}^{\smile}; \overline{I_{\operatorname{[SESSION]}}}; \operatorname{sessionLoginAssist} \vdash \overline{I_{\operatorname{[SESSION]}}} \end{split}$$

Invariant: $SYM \ sessionLogoutAssist[SESSION*SESSION]$

 $sessionLogoutAssist[SESSION*SESSION] \ is \ symmetric \\ sessionLogoutAssist = sessionLogoutAssist ^{\smile}$

Invariant: ASY sessionLogoutAssist[SESSION*SESSION]

 $sessionLogoutAssist[SESSION*SESSION] is antisymmetric \\ sessionLogoutAssist ` \cap sessionLogoutAssist \vdash I_{[SESSION]}$

Invariant: $UNI\ sessionLogoutAssist/SESSION*SESSION$

$$\begin{split} & sessionLogoutAssist[SESSION*SESSION] \text{ is univalent} \\ & sessionLogoutAssist}; \overline{I_{\text{[SESSION]}}}; sessionLogoutAssist} ^{\smile} \vdash \overline{I_{\text{[SESSION]}}} \end{split}$$

Invariant: INJ sessionLogoutAssist/SESSION*SESSION|

$$\begin{split} & \operatorname{sessionLogoutAssist}[\operatorname{SESSION}^*\operatorname{SESSION}] \text{ is injective} \\ & \operatorname{sessionLogoutAssist}^\smile; \overline{I_{[\operatorname{SESSION}]}}; \operatorname{sessionLogoutAssist} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: SYM sessionSRI/SESSION*SESSION

 $sessionSRI[SESSION*SESSION] \ is \ symmetric \\ sessionSRI = sessionSRI \widetilde{\ }$

Invariant: ASY sessionSRI[SESSION*SESSION]

session SRI[SESSION*SESSION] is antisymmetric session SRI $\vdash I_{[\text{SESSION}]}$

Invariant: UNI sessionSRI/SESSION*SESSION

$$\begin{split} & \operatorname{sessionSRI[SESSION*SESSION]} \text{ is univalent} \\ & \operatorname{sessionSRI;} \overline{I_{[\operatorname{SESSION}]}}; \operatorname{sessionSRI}^{\smile} \vdash \overline{I_{[\operatorname{SESSION}]}} \end{split}$$

Invariant: INJ sessionSRI/SESSION*SESSION]

 $sessionSRI[SESSION*SESSION] is injective \\ sessionSRI^{\sim}; \overline{I_{[SESSION]}}; sessionSRI \vdash \overline{I_{[SESSION]}}$

Invariant: SYM sessionSIA [SESSION*SESSION]

session SIA
[SESSION*SESSION] is symmetric session SIA = session SIA
 $\widetilde{}$

Invariant: ASY sessionSIA[SESSION*SESSION]

sessionSIA[SESSION*SESSION] is antisymmetric

sessionSIA $\vdash I_{\text{[SESSION]}}$

Invariant: UNI sessionSIA/SESSION*SESSION

sessionSIA[SESSION*SESSION] is univalent

sessionSIA; $\overline{I_{[SESSION]}}$; sessionSIA $^{\sim} \vdash \overline{I_{[SESSION]}}$

Invariant: INJ sessionSIA/SESSION*SESSION

sessionSIA[SESSION*SESSION] is injective

 ${\rm sessionSIA}^{\smallsmile}; \overline{I_{\rm [SESSION]}}; {\rm sessionSIA} \vdash \overline{I_{\rm [SESSION]}}$

Invariant: UNI scopeID/Scope*ScopeID/

scopeID[Scope*ScopeID] is univalent

scopeID; $\overline{I_{[ScopeID]}}$; scopeID $\overset{\sim}{}$ $\vdash \overline{I_{[Scope]}}$

Invariant: UNI scopeOwner/Scope*Account/

scopeOwner[Scope*Account] is univalent

scopeOwner; $\overline{I_{[Account]}}$; scopeOwner $\ \vdash \overline{I_{[Scope]}}$

Invariant: UNI scopeDescr|Scope*ScopeDescr|

scopeDescr[Scope*ScopeDescr] is univalent

scopeDescr; $\overline{I_{[ScopeDescr]}}$; scopeDescr $\overset{\sim}{\vdash} \overline{I_{[Scope]}}$

Invariant: IRF scopeIII/Scope*Scope

scopeIII[Scope*Scope] is irreflexive

 $scopeIII \vdash \overline{I_{[Scope]}}$

Invariant: ASY scopeIII/Scope*Scope/

scopeIII[Scope*Scope] is antisymmetric

 $scopeIII \subset scopeIII \vdash I_{[Scope]}$

Invariant: SYM scopeIsaCC/Scope*Scope

scopeIsaCC[Scope*Scope] is symmetric

scopeIsaCC = scopeIsaCC

Invariant: ASY scopeIsaCC/Scope*Scope

scopeIsaCC[Scope*Scope] is antisymmetric

 $scopeIsaCC \subset scopeIsaCC \vdash I_{[Scope]}$

Invariant: UNI scopeIsaCC/Scope*Scope

scopeIsaCC[Scope*Scope] is univalent

 $scopeIsaCC; \overline{I_{[Scope]}}; scopeIsaCC^{\smile} \vdash \overline{I_{[Scope]}}$

 $Invariant: \ INJ \ scope Is a CC[Scope*Scope]$

 $scope Isa CC [Scope *Scope] \ is \ injective$

 $scopelsaCC \subset \overline{I_{[Scope]}}; scopelsaCC \vdash \overline{I_{[Scope]}}$

Invariant: UNI ttName/TText*TTName/

 $ttName[TText*TTName] \ is \ univalent$

 $\mathsf{ttName}; \overline{I_{[\mathsf{TTName}]}}; \mathsf{ttName}^{\smile} \vdash \overline{I_{[\mathsf{TText}]}}$

Invariant: $UNI \ ttValue [TText*TTValue]$

ttValue[TText*TTValue] is univalent

 $\mathsf{ttValue}; \overline{I_{[\mathsf{TTValue}]}}; \mathsf{ttValue}^{\smile} \vdash \overline{I_{[\mathsf{TText}]}}$

Invariant: UNI ttTemplate[TText*TTPhrase]

 $ttTemplate[TText*TTPhrase] \ is \ univalent$

 $\texttt{ttTemplate}; \overline{I_{[\texttt{TTPhrase}]}}; \texttt{ttTemplate} \c^{\smile} \vdash \overline{I_{[\texttt{TText}]}}$

Invariant: $UNI\ ttInstance[TText*TTPhrase]$

 $ttInstance[TText*TTPhrase] \ is \ univalent$

 $\texttt{ttInstance}; \overline{I_{[\texttt{TTPhrase}]}}; \texttt{ttInstance} \vdash \overline{I_{[\texttt{TText}]}}$

Invariant: UNI ttICO/TText*TText]

ttICO[TText*TText] is univalent

 $\mathbf{ttICO}; \overline{I_{[\mathrm{TText}]}}; \mathbf{ttICO}^{\smile} \vdash \overline{I_{[\mathrm{TText}]}}$

Invariant: $IRF\ ttICO[TText*TText]$

ttICO[TText*TText] is irreflexive

 $ttICO \vdash \overline{I_{[TText]}}$

Invariant: ASY ttICO/TText*TText/

 ${\it ttICO[TText*TText]} \ is \ antisymmetric$

 $\mathbf{ttICO}^{\smile} \cap \mathbf{ttICO} \vdash I_{[\mathsf{TText}]}$

Invariant: UNI ttICCO[TText*TText]

ttICCO[TText*TText] is univalent

 $\mathsf{ttICCO}; \overline{I_{[\mathsf{TText}]}}; \mathsf{ttICCO}^{\smile} \vdash \overline{I_{[\mathsf{TText}]}}$

Invariant: IRF ttICCO[TText*TText]

ttICCO[TText*TText] is irreflexive

 $\mathsf{ttICCO} \vdash \overline{I_{[\mathsf{TText}]}}$

Invariant: $ASY\ ttICCO[TText*TText]$

 ${\it ttICCO[TText*TText]} \ is \ antisymmetric$

 $ttICCO^{\smile} \cap ttICCO \vdash I_{[TText]}$

Invariant: $UNI\ ttDescr[TText*TTPhrase]$

ttDescr[TText*TTPhrase] is univalent

 $\mathsf{ttDescr}; \overline{I_{[\mathsf{TTPhrase}]}}; \mathsf{ttDescr}^{\smallsmile} \vdash \overline{I_{[\mathsf{TText}]}}$

Invariant: UNI ttOwner/TText*Account/

ttOwner[TText*Account] is univalent

 $\mathbf{ttOwner}; \overline{I_{[\mathbf{Account}]}}; \mathbf{ttOwner}^{\smile} \vdash \overline{I_{[\mathbf{TText}]}}$

Invariant: identity_TTexts

Identity rule, following from identity TTexts

 $\texttt{ttScope}\texttt{;}\,\texttt{ttScope}\texttt{`}\cap\texttt{ttName}\texttt{;}\,\texttt{ttName}\texttt{`}\vdash I_{\texttt{[TText]}}$

Invariant: identity_Scopes

Identity rule, following from identity Scopes

 $scopeOwner; scopeOwner \subset scopeID; scopeID \subset I_{Scope}$

Invariant: identity_UIDs

Identity rule, following from identity UIDs

uidIssuer; uidIssuer $\check{} \cap$ uidUserid; uidUserid $\check{} \vdash I_{[\text{UID}]}$

$Invariant:\ identity_SIAMPersonRefComputations$

Identity rule, following from identity SIAMPersonRefComputations

 ${\tt siamcompFirstName; siamcompFirstName"} \cap {\tt siamcompLastName; siamcompLastName"} \vdash I_{\tt [SIAMPersonRefComputation]}$

Invariant: identity_Computations

Identity rule, following from identity Computations

 $\operatorname{compArg} : \operatorname{compVar} : \operatorname{compVar} : \operatorname{Computation} : I_{[\operatorname{Computation}]}$

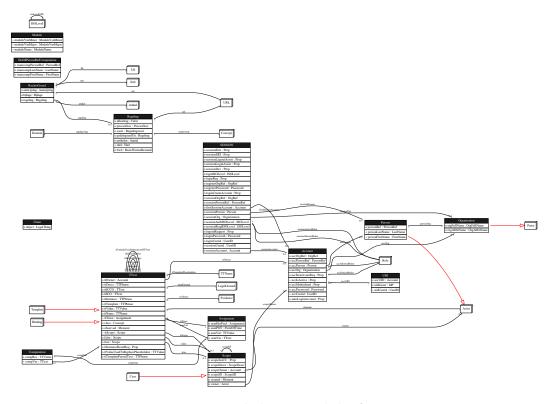


Figure 5.2: Logical data model of MirrorMe

5.3 Logical data model

The functional requirements have been translated into a data model. This model is shown by fig. 5.2.

There are 14 entity types. The details of each entity type are described (in alphabetical order) in the following two tables:

Logical entity types

Concept	Meaning	Type
Account	een verzameling gegevens	
	die (een deel van) de	
	gebruikerscontext van	
	één gebruiker binnen het	
	systeem beschrijft	
	Als iemand inlogt in het	
	systeem moet diens	
	'context' worden	
	geactiveerd, d.w.z. de	
	gegevens over de persoon	
	die het systeem nodig	
	heeft om te kunnen	
	berekenen wat hij/zij wel	
	en niet mag doen, en	
	welke gegevens van het	
	systeem daarbij mogen	
	worden gebruikt. Om	
	zulke computations te	
	kunnen maken wordt een	
	aantal zaken	
	geregistreerd en aan één	
	persoon gekoppeld.	
Assignment	the registration of an event where either a	
	TText was assigned a value, or its value was	
	deleted	
Claim		
Computation	(for a TText T): a set of Assignments for all	
	TTexts T[i] that are mentioned in the	
	ttTemplate of T, for the purpose of computing	g
	a ttValue for T	
Module		
Organization	an organized group of people with a particula	r
	purpose, such as a business or government	
	department	
Person	a human body of flesh and blood	
RechtsGrond		
Regeling		

Concept Meaning Type

SESSION

een verzameling van gegevens die de context beschrijven waarin één persoon het systeem gebruikt Een persoon gebruikt een gegevensruimte (en heet dan 'user') door met een browser (bijv. Chrome of Firefox) het systeem te benaderen dat de gegegensruimte beheert. Als meerdere personen een gegevensruimte delen, moet het systeem de context van elk van hen kunnen onderscheiden, bijvoorbeeld om:

- de interactie 'klein' te houden, d.w.z. alleen gegevens te laten zien die voor hem/haar relevant zijn;
- ervoor te zorgen dat een user niet ziet wat hij niet mag zien;
- te kunnen bijhouden welke persoon, of welk(e) organisatie(onderdeel) verantwoordelijk is voor een zekere transactie;
- automatisch gegevens betreffende de user of zijn context aan transacties toe te kunnen voegen

We gebruiken de term 'SESSION' of 'session' om de verzameling van gegevens betreffende één (actieve) user mee aan te geven. Deze term correspondeert met de gelijknamige term browsers gebruiken, nl. 'een verbinding (door de browser) met een

Concept	Meaning	Type	
SIAMPersonRefC	omputation		
Scope	a set of TTexts that are controlled by a single authority and (together) serve a particular purpose a set of TTexts that are controlled by a single authority and (together) serve a particular		
TText	purpose the collection of one template string, one instance thereof, and one value within some scope If people make claims, they do observations, and they will reason about them. In order to decide whether a claim is true, statements are organised in a structure that represents the argument. The concept "statement" is used to represent claims, observations, and all other utterances that can be considered true or false. Statements and Claims are modeled/implemented as		
UID	TTexts.		
OID			

Other attributes

Concept	Meaning	Type
Aantal		
Aanwijzing		
Actor		
Artikel		
BasisWettenBestand		
Bijlage		
Binding		
Case		
Concept		
Domein		
Evidence		
FirstName		
	73	

Concept	Meaning	Type
ISOLevel	a 'level of assurance' as specified in the ISO/IEC 29115:2013 standard	
IdP		
IfcText		
LastName		
LegalGround		
LegalThing		
Lid		
ModuleName	the name (text) by which the module is known	
ModuleVsnMajor	the major version indicator of the module -	
	incremented when changes are not backwards	
MadulaVan Minan	compatible	
ModuleVsnMinor	the minor version indicator of the module -	
	incremented when changes exist that are	
Moment	backwards compatible	
OrgAbbrName		
OrgFullName		
OrgRef		
Party		
Password	een rij karakters, die	
I diss it of d	geheim gehouden kan	
	worden door een persoon,	
	en door die persoon	
	gebruikt moet worden	
	om toegang te krijgen	
	tot het systeem	
	Om het moeilijk te	
	maken dat het Account	
	van een zeker persoon	
	door een ander wordt	
	gebruikt, registreert het	
	systeem wachtwoorden.	
	Door een wachtwoord	
	geheim te houden	
	ontstaat enige mate van	
	zekerheid dat het	
	systeem gebruikt word	
	door (dan wel met medeweten van) de	
	persoon op wiens naam	
	het wachtwoord is	
	geregistreerd.	
PersonRef	Sorogionicora.	
PointOfTime	a specific moment on a specific date	
Procesflow	a specific memory on a specific date	
Regelingsoort		
0		

Concept	Meaning	Type	
Role	de naam voor het mogen		
	inzien en/of wijzigen van		
	zekere gegevens, die kan		
	worden toegekend aan		
	accounts en in sessions		
	kan worden geactiveerd		
	Niet iedereen mag alle		
	gegevens uit een systeem		
	inzien en/of wijzigen.		
	Om de beheerslast te		
	beperken die ontstaat als		
	dit soort rechten per		
	persoon wordt		
	uitgegeven, gebruiken we		
	het concept 'Role'. Een		
	Role wordt enerzijds		
	toegekend aan Accounts		
	(en daarmee aan		
	Personen) en anderzijds		
	wordt hij gebruikt om		
	gegevens in te zien en/of		
	te wijzigen. Als een user		
	inlogt worden de Rollen		
	die aan hem zijn		
	toegekend in de session		
	geactiveerd		
	(sessionrollen). Interfaces		
	gebruiken deze		
	sessionrollen om al dan		
	niet gegevens te tonen		
_	en/of te editen.		
ScopeDescr	descriptive text or purpose of the Scope		
ScopeID	a short name that helps to refer to the Scope	e	
Sub		`	
TTName	a label used to identify a TText (within a scop	pe)	
TTPhrase	a sequence of words		
TTValue	(the representation of) a value of a TText		
Tekst			
Template			
Titel			
URL			

Concept	Meaning	Type	
UserID	een rij karakters (bijvoorbeeld het e-mailadres van de user), die een account identificeert binnen het systeem Het UserID (gebruikersnaam) van een account identificeert dat account en impliciet daarmee ook diens eigenaar (d.w.z. de persoon die als enige geacht wordt met dit account in te loggen. Een goed gebruik is om hiervoor een e-mailadres te gebruiken waarop de user van het account bereikbaar is.		

5.3.1 Entity type: Account

Als iemand inlogt in het systeem moet diens 'context' worden geactiveerd, d.w.z. de gegevens over de persoon die het systeem nodig heeft om te kunnen berekenen wat hij/zij wel en niet mag doen, en welke gegevens van het systeem daarbij mogen worden gebruikt. Om zulke computations te kunnen maken wordt een aantal zaken geregistreerd en aan één persoon gekoppeld.

This entity type has the following attributes:

Attribute	Type	
Id	Account	Primary key
accOrgRef	OrgRef	Optional
accPersonRef	PersonRef	Optional
accPerson	Person	Optional
accOrg	Organization	Optional
accDeactivateReq	Prop	Optional
accIsActive	Prop	Mandatory
accIsInitialized	Prop	Optional
accPassword	Password	Optional
accUserid	UserID	Mandatory
${\bf auto Login Account}$	Prop	Optional

Account has the following associations:

- 1. sessionAccount (from SESSION to Account).
- 2. accOrg (from Account to Organization).
- 3. accPerson (from Account to Person).
- 4. accAllowedRoles (from Account to Role).

- 5. accDefaultRoles (from Account to Role).
- 6. firstSessionAccount (from SESSION to Account).
- 7. accUID (from Account to UID).
- 8. scopeOwner (from Scope to Account).
- 9. ttOwner (from TText to Account).

5.3.2 Entity type: Assignment

This entity type has the following attributes:

Attribute	Type	
Id	Assignment	Primary key
asmHasPred	Assignment	Optional
asmPOT	PointOfTime	Optional
asmVal	TTValue	Optional
asmVar	TText	Mandatory

Assignment has the following associations:

- 1. ttTrace (from TText to Assignment).
- 2. asmVar (from Assignment to TText).
- 3. compArg (from Computation to Assignment).

5.3.3 Entity type: Claim

This entity type has the following attributes:

Attribute	Type	
Id	Claim	Primary key
object	LegalThing	Optional

Claim has no associations.

5.3.4 Entity type: Computation

This entity type has the following attributes:

Attribute	Type	
Id compRes compVar	Computation TTValue TText	Primary key Optional Mandatory

Computation has the following associations:

- 1. compVar (from Computation to TText).
- 2. compArg (from Computation to Assignment).

5.3.5 Entity type: Module

This entity type has the following attributes:

Attribute	Type	
Id	Module	Primary key
${\it module Vsn Minor}$	ModuleVsnMinor	Mandatory
moduleVsnMajor	ModuleVsnMajor	Mandatory
moduleName	ModuleName	Mandatory

Module has no associations.

5.3.6 Entity type: Organization

This entity type has the following attributes:

Attribute	Type	
Id	Organization	Primary key
orgFullName	OrgFullName	Optional
orgAbbrName	OrgAbbrName	Optional

Organization has the following associations:

- 1. accOrg (from Account to Organization).
- 2. sessionOrg (from SESSION to Organization).
- 3. personOrg (from Person to Organization).

5.3.7 Entity type: Person

This entity type has the following attributes:

Attribute	Type	
Id	Person	Primary key
personRef	PersonRef	Optional
personLastName	LastName	Mandatory
$\operatorname{personFirstName}$	${\bf FirstName}$	Mandatory

Person has the following associations:

- 1. personOrg (from Person to Organization).
- 2. accPerson (from Account to Person).
- 3. sessionPerson (from SESSION to Person).

5.3.8 Entity type: RechtsGrond

This entity type has the following attributes:

Attribute	Type	
Id	RechtsGrond	Primary key
aanwijzing	Aanwijzing	Optional
bijlage	Bijlage	Optional
regeling	Regeling	Optional

RechtsGrond has the following associations:

- 1. url (from RechtsGrond to URL).
- 2. regeling (from RechtsGrond to Regeling).
- 3. artikel (from RechtsGrond to Artikel).
- 4. lid (from RechtsGrond to Lid).
- 5. sub (from RechtsGrond to Sub).

5.3.9 Entity type: Regeling

This entity type has the following attributes:

Attribute	Type	
Id	Regeling	Primary key
afkorting	Tekst	Optional
procesflow	Procesflow	Optional
soort	Regelingsoort	Optional
gedelegeerdUit	Regeling	Optional
artikelen	Aantal	Optional
titel	Titel	Mandatory
bwb	${\bf Basis Wetten Best and}$	Optional

Regeling has the following associations:

- 1. regeling (from RechtsGrond to Regeling).
- 2. url (from Regeling to URL).
- 3. regelgeving (from Domein to Regeling).
- 4. onderwerp (from Regeling to Concept).

5.3.10 Entity type: SESSION

Een persoon gebruikt een gegevensruimte (en heet dan 'user') door met een browser (bijv. Chrome of Firefox) het systeem te benaderen dat de gegegensruimte beheert. Als meerdere personen een gegevensruimte delen, moet het systeem de context van elk van hen kunnen onderscheiden, bijvoorbeeld om:

- de interactie 'klein' te houden, d.w.z. alleen gegevens te laten zien die voor hem/haar relevant zijn;
- ervoor te zorgen dat een user niet ziet wat hij niet mag zien;
- te kunnen bijhouden welke persoon, of welk(e) organisatie(onderdeel) verantwoordelijk is voor een zekere transactie;

• automatisch gegevens betreffende de user of zijn context aan transacties toe te kunnen voegen

We gebruiken de term 'SESSION' of 'session' om de verzameling van gegevens betreffende één (actieve) user mee aan te geven. Deze term correspondeert met de gelijknamige term browsers gebruiken, nl. 'een verbinding (door de browser) met een webservice (die een URL heeft)'. Het systeem houdt één session bij voor elke actieve user, d.w.z. voor elke browser die het systeem benadert. Merk op dat dit in het bijzonder geldt als de user in verschillende tabbladen van dezelfde browser het systeem benadert - er is dan toch maar één session (en één user).

This entity type has the following attributes:

Attribute	Type	
Id	SESSION	Primary key
sessionSIA	Prop	Optional
sessionSRI	Prop	Optional
sessionLogoutAssist	Prop	Optional
sessionLoginAssist	Prop	Optional
sessionDev	Prop	Optional
loginISOLevel	ISOLevel	Optional
loginReq	Prop	Optional
registerOrgRef	OrgRef	Optional
registerPassword	Password	Optional
loginCreateAccount	Prop	Optional
sessionOrgRef	OrgRef	Optional
sessionPersonRef	PersonRef	Optional
firstSessionAccount	Account	Optional
sessionPerson	Person	Optional
sessionOrg	Organization	Optional
${\it session} Auth ISO Level$	ISOLevel	Optional
${\it session} \\ {\it Reqd} \\ {\it ISOLevel}$	ISOLevel	Optional
logoutRequest	Prop	Optional
loginPassword	Password	Optional
loginUserid	UserID	Optional
sessionUserid	UserID	Optional
sessionAccount	Account	Optional

SESSION has the following associations:

- 1. sessionAccount (from SESSION to Account).
- 2. sessionOrg (from SESSION to Organization).
- 3. sessionPerson (from SESSION to Person).
- 4. sessionAllowedRoles (from SESSION to Role).
- 5. sessionActiveRoles (from SESSION to Role).
- 6. firstSessionAccount (from SESSION to Account).

5.3.11 Entity type: SIAMPersonRefComputation

This entity type has the following attributes:

Attribute	Type	
Id	SIAMPersonRefComputation	Primary key
siamcompPersonRef	PersonRef	Optional
siam compLastName	LastName	Optional
siam compFirstName	FirstName	Optional

SIAMPersonRefComputation has no associations.

5.3.12 Entity type: Scope

This entity type has the following attributes:

Attribute	Type	
Id	Scope	Primary key
scopeIsaCC	Prop	Optional
scopeDescr	ScopeDescr	Optional
scopeOwner	Account	Optional
scopeID	ScopeID	Optional
created	Moment	Optional
owner	Actor	Optional

Scope has the following associations:

- 1. owner (from Scope to Actor).
- 2. true (from TText to Scope).
- 3. false (from TText to Scope).
- 4. ttScope (from TText to Scope).
- 5. scopeOwner (from Scope to Account).
- 6. scopeIII (from Scope to Scope).

5.3.13 Entity type: TText

If people make claims, they do observations, and they will reason about them. In order to decide whether a claim is true, statements are organised in a structure that represents the argument. The concept "statement" is used to represent claims, observations, and all other utterances that can be considered true or false. Statements and Claims are modeled/implemented as TTexts.

This entity type has the following attributes:

Attribute	Type	
Id	TText	Primary key
ttOwner	Account	Optional
ttDescr	TTPhrase	Optional
ttICCO	TText	Optional
ttICO	TText	Optional
ttInstance	TTPhrase	Optional
ttTemplate	TTPhrase	Optional
ttValue	TTValue	Optional

Attribute	Туре	
ttName	TTName	Optional
ttTrace	Assignment	Mandatory
class	Concept	Optional
observed	Moment	Optional
ttScope	Scope	Mandatory
false	Scope	Optional
true	Scope	Optional
ttInstanceResetReq	Prop	Optional
tt Value Used To Replace Place holders	TTValue	Optional
tt Template Parsed Text	TTPhrase	Optional

TText has the following associations:

- 1. claimant (from TText to Actor).
- 2. requires (from TText to TText).
- 3. ttTemplatePlaceholders (from TText to TTName).
- 4. ttIsUsedBy (from TText to TText).
- 5. ttIsUsedByCopy (from TText to TText).
- 6. ttIsUsedByStar (from TText to TText).
- 7. ttValueIsUsedInInstanceOfTText (from TText to TText).
- 8. true (from TText to Scope).
- 9. false (from TText to Scope).
- 10. ttScope (from TText to Scope).
- 11. legalGround (from TText to LegalGround).
- 12. evidence (from TText to Evidence).
- 13. ttTrace (from TText to Assignment).
- 14. asmVar (from Assignment to TText).
- 15. compVar (from Computation to TText).
- 16. ttOwner (from TText to Account).

5.3.14 Entity type: *UID*

This entity type has the following attributes:

Attribute	Type	
Id	UID	Primary key
accUID	Account	Optional
uidIssuer	IdP	Optional
uidUserid	UserID	Mandatory

UID has the following associations:

1. accUID (from Account to UID).

5.4 Technical datamodel

The functional requirements have been translated into a technical data model. This model is shown by fig. 5.3.

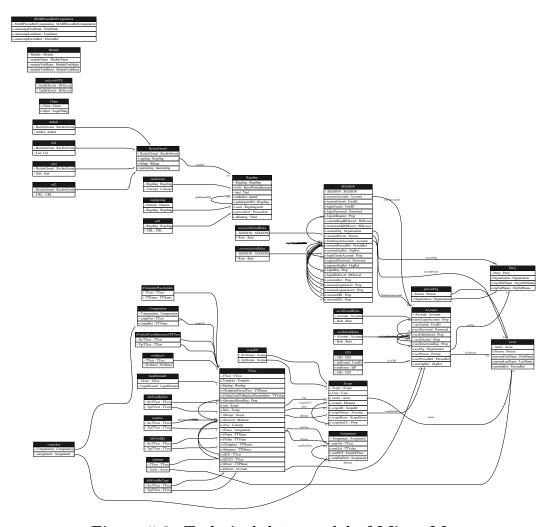


Figure 5.3: Technical data model of MirrorMe $\,$

The technical datamodel consists of the following 78 tables:

5.4.1 Table: Aantal

This table has the following 1 attributes:

Aantal

This attribute implements the identity relation of *Aantal*. INTEGER, Mandatory, Unique.

5.4.2 Table: Aanwijzing

This table has the following 1 attributes:

Aanwijzing

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.3 Table: accAllowedRoles

This is a link-table, implementing the relation $Account \xrightarrow{accAllowedRoles} Role$. It contains the following columns:

Account

This attribute is a foreign key to Account OBJECT, Mandatory.

• Role

This attribute is a foreign key to Role OBJECT, Mandatory.

5.4.4 Table: accDefaultRoles

This is a link-table, implementing the relation $Account \xrightarrow{accDefaultRoles} Role$. It contains the following columns:

Account

This attribute is a foreign key to Account OBJECT, Mandatory.

• Role

This attribute is a foreign key to Role OBJECT, Mandatory.

5.4.5 Table: Account

This table has the following 11 attributes:

Account

This attribute is the primary key. OBJECT, Mandatory, Unique.

• autoLoginAccount

This attribute is a foreign key to Account OBJECT, Optional, Unique.

• accUserid

This attribute is a foreign key to UserID ALPHANUMERIC, Mandatory, Unique.

accPassword

This attribute implements the relation $Account \xrightarrow{accPassword} Password$. PASSWORD, Optional.

• accIsInitialized

This attribute is a foreign key to Account OBJECT, Optional, Unique.

accIsActive

This attribute is a foreign key to Account OBJECT, Mandatory, Unique.

• accDeactivateReq

This attribute is a foreign key to Account OBJECT, Optional, Unique.

accOrg

This attribute is a foreign key to Organization OBJECT, Optional.

accPerson

This attribute is a foreign key to Person OBJECT, Optional.

accPersonRef

This attribute is a foreign key to PersonRef ALPHANUMERIC, Optional.

accOrgRef

This attribute is a foreign key to OrgRef OBJECT, Optional.

5.4.6 Table: Actor

This table has the following 5 attributes:

• Actor

This attribute is the primary key. OBJECT, Mandatory, Unique.

• Person

This attribute implements the identity relation of *Person*. OBJECT, Optional, Unique.

• personFirstName

This attribute is a foreign key to FirstName ALPHANUMERIC, Optional.

• personLastName

This attribute is a foreign key to LastName

ALPHANUMERIC, Optional.

personRef

This attribute is a foreign key to PersonRef ALPHANUMERIC, Optional.

5.4.7 Table: artikel1

This is a link-table, implementing the relation $RechtsGrond \xrightarrow{artikel} Artikel$. It contains the following columns:

• RechtsGrond

This attribute is a foreign key to RechtsGrond OBJECT, Mandatory.

Artikel

This attribute is a foreign key to Artikel OBJECT, Mandatory.

5.4.8 Table: Artikel2

This table has the following 1 attributes:

• Artikel

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.9 Table: Assignment

This table has the following 5 attributes:

Assignment

This attribute is the primary key. OBJECT, Mandatory, Unique.

• asmVar

This attribute is a foreign key to TText OBJECT, Mandatory.

• asmVal

This attribute is a foreign key to TTValue ALPHANUMERIC, Optional.

• asmPOT

This attribute is a foreign key to PointOfTime DATETIME, Optional.

• asmHasPred

This attribute is a foreign key to Assignment OBJECT, Optional, Unique.

5.4.10 Table: BasisWettenBestand

This table has the following 1 attributes:

• BasisWettenBestand

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.11 Table: Bijlage

This table has the following 1 attributes:

• Bijlage

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.12 Table: Claim

This table has the following 2 attributes:

• Claim

This attribute is the primary key. OBJECT, Mandatory, Unique.

• object

This attribute is a foreign key to LegalThing OBJECT, Optional.

5.4.13 Table: claimant

This is a link-table, implementing the relation $TText \xrightarrow{claimant} Actor$. It contains the following columns:

• TText

This attribute is a foreign key to TText OBJECT, Mandatory.

• Actor

This attribute is a foreign key to Actor OBJECT, Mandatory.

5.4.14 Table: compArg

This is a link-table, implementing the relation $Computation \xrightarrow{compArg} Assignment$. It contains the following columns:

Computation

This attribute is a foreign key to Computation OBJECT, Mandatory.

• Assignment

This attribute is a foreign key to Assignment OBJECT, Mandatory.

5.4.15 Table: Computation

This table has the following 3 attributes:

• Computation

This attribute is the primary key. OBJECT, Mandatory, Unique.

• compVar

This attribute is a foreign key to TText OBJECT, Mandatory, Unique.

• compRes

This attribute is a foreign key to TTValue ALPHANUMERIC, Optional.

5.4.16 Table: Concept

This table has the following 1 attributes:

• Concept

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.17 Table: Domein

This table has the following 1 attributes:

Domein

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.18 Table: evidence1

This is a link-table, implementing the relation $TText \xrightarrow{evidence} Evidence$. It contains the following columns:

• TText

This attribute is a foreign key to TText OBJECT, Mandatory.

• Evidence

This attribute is a foreign key to Evidence OBJECT, Mandatory.

5.4.19 Table: Evidence2

This table has the following 1 attributes:

• Evidence

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.20 Table: FirstName

This table has the following 1 attributes:

FirstName

This attribute implements the identity relation of FirstName. ALPHANUMERIC, Mandatory, Unique.

5.4.21 Table: IdP

This table has the following 1 attributes:

IdP

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.22 Table: IfcText

This table has the following 1 attributes:

• IfcText

This attribute implements the identity relation of IfcText. ALPHANUMERIC, Mandatory, Unique.

5.4.23 Table: ISOLevel

This table has the following 1 attributes:

ISOLevel

This attribute implements the identity relation of *ISOLevel*. INTEGER, Mandatory, Unique.

5.4.24 Table: isoLevelGTE

This is a link-table, implementing the relation $ISOLevel \xrightarrow{isoLevelGTE} ISOLevel$. It contains the following columns:

• SrcISOLevel

This attribute is a foreign key to ISOLevel INTEGER, Mandatory.

• TgtISOLevel

This attribute is a foreign key to ISOLevel INTEGER, Mandatory.

5.4.25 Table: LastName

This table has the following 1 attributes:

• LastName

This attribute implements the identity relation of LastName. ALPHANUMERIC, Mandatory, Unique.

5.4.26 Table: legalGround1

This is a link-table, implementing the relation $TText \xrightarrow{legalGround} LegalGround$. It contains the following columns:

• TText

This attribute is a foreign key to TText OBJECT, Mandatory.

• LegalGround

This attribute is a foreign key to LegalGround OBJECT, Mandatory.

5.4.27 Table: LegalGround2

This table has the following 1 attributes:

• LegalGround

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.28 Table: LegalThing

This table has the following 1 attributes:

• LegalThing

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.29 Table: lid1

This is a link-table, implementing the relation $RechtsGrond \xrightarrow{lid} Lid$. It contains the following columns:

• RechtsGrond

This attribute is a foreign key to RechtsGrond OBJECT, Mandatory.

• Lid

This attribute is a foreign key to Lid OBJECT, Mandatory.

5.4.30 Table: Lid2

This table has the following 1 attributes:

• Lid

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.31 Table: Module

This table has the following 4 attributes:

Module

This attribute is the primary key. OBJECT, Mandatory, Unique.

moduleName

This attribute is a foreign key to ModuleName ALPHANUMERIC, Mandatory.

• moduleVsnMajor

This attribute is a foreign key to ModuleVsnMajor INTEGER, Mandatory.

• moduleVsnMinor

This attribute is a foreign key to ModuleVsnMinor INTEGER, Mandatory.

5.4.32 Table: ModuleName

This table has the following 1 attributes:

• ModuleName

This attribute implements the identity relation of *ModuleName*. ALPHANUMERIC, Mandatory, Unique.

5.4.33 Table: ModuleVsnMajor

This table has the following 1 attributes:

• ModuleVsnMajor

This attribute implements the identity relation of ModuleVsnMajor. INTEGER, Mandatory, Unique.

5.4.34 Table: ModuleVsnMinor

This table has the following 1 attributes:

• ModuleVsnMinor

This attribute implements the identity relation of ModuleVsnMinor. INTEGER, Mandatory, Unique.

5.4.35 Table: Moment

This table has the following 1 attributes:

Moment

This attribute implements the identity relation of *Moment*. DATETIME, Mandatory, Unique.

5.4.36 Table: onderwerp

This is a link-table, implementing the relation $Regeling \xrightarrow{onderwerp} Concept$. It contains the following columns:

• Regeling

This attribute is a foreign key to Regeling OBJECT, Mandatory.

• Concept

This attribute is a foreign key to Concept OBJECT, Mandatory.

5.4.37 Table: ONE

This table has the following 1 attributes:

• ONE

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.38 Table: OrgAbbrName

This table has the following 1 attributes:

• OrgAbbrName

This attribute implements the identity relation of OrgAbbrName. ALPHANUMERIC, Mandatory, Unique.

5.4.39 Table: OrgFullName

This table has the following 1 attributes:

• OrgFullName

This attribute implements the identity relation of OrgFullName. ALPHANUMERIC, Mandatory, Unique.

5.4.40 Table: OrgRef

This table has the following 1 attributes:

• OrgRef

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.41 Table: Party

This table has the following 4 attributes:

• Party

This attribute is the primary key. OBJECT, Mandatory, Unique.

Organization

This attribute implements the identity relation of *Organization*. OBJECT, Optional, Unique.

• orgAbbrName

This attribute is a foreign key to OrgAbbrName ALPHANUMERIC, Optional, Unique.

• orgFullName

This attribute is a foreign key to OrgFullName ALPHANUMERIC, Optional, Unique.

5.4.42 Table: Password

This table has the following 1 attributes:

• Password

This attribute implements the identity relation of *Password*. PASSWORD, Mandatory, Unique.

5.4.43 Table: personOrg

This is a link-table, implementing the relation $Person \xrightarrow{personOrg} Organization$. It contains the following columns:

• Person

This attribute is a foreign key to Person OBJECT, Mandatory.

Organization

This attribute is a foreign key to Organization OBJECT, Mandatory.

5.4.44 Table: PersonRef

This table has the following 1 attributes:

PersonRef

This attribute implements the identity relation of PersonRef. ALPHANUMERIC, Mandatory, Unique.

5.4.45 Table: PointOfTime

This table has the following 1 attributes:

• PointOfTime

This attribute implements the identity relation of PointOfTime. DATETIME, Mandatory, Unique.

5.4.46 Table: Processlow

This table has the following 1 attributes:

Procesflow

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.47 Table: RechtsGrond

This table has the following 4 attributes:

• RechtsGrond

This attribute is the primary key. OBJECT, Mandatory, Unique.

regeling

This attribute is a foreign key to Regeling OBJECT, Optional.

• bijlage

This attribute is a foreign key to Bijlage OBJECT, Optional.

aanwijzing

This attribute is a foreign key to Aanwijzing OBJECT, Optional.

5.4.48 Table: regelgeving

This is a link-table, implementing the relation $Domein \xrightarrow{regelgeving} Regeling$. It contains the following columns:

Domein

This attribute is a foreign key to Domein OBJECT, Mandatory.

• Regeling

This attribute is a foreign key to Regeling OBJECT, Mandatory.

5.4.49 Table: Regeling

This table has the following 8 attributes:

Regeling

This attribute is the primary key. OBJECT, Mandatory, Unique.

• bwb

This attribute is a foreign key to BasisWettenBestand OBJECT, Optional.

titel

This attribute is a foreign key to Titel OBJECT, Mandatory.

artikelen

This attribute is a foreign key to Aantal INTEGER, Optional.

• gedelegeerdUit

This attribute is a foreign key to Regeling OBJECT, Optional.

• soort

This attribute is a foreign key to Regelingsoort OBJECT, Optional.

procesflow

This attribute is a foreign key to Processlow OBJECT, Optional.

afkorting

This attribute is a foreign key to Tekst OBJECT, Optional.

5.4.50 Table: Regelingsoort

This table has the following 1 attributes:

• Regelingsoort

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.51 Table: requires

This is a link-table, implementing the relation $TText \xrightarrow{requires} TText$. It contains the following columns:

• SrcTText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TgtTText

This attribute is a foreign key to TText OBJECT, Mandatory.

5.4.52 Table: Role

This table has the following 1 attributes:

• Role

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.53 Table: Scope

This table has the following 8 attributes:

• Scope

This attribute is the primary key. OBJECT, Mandatory, Unique.

• Case

This attribute implements the identity relation of Case. OBJECT, Optional, Unique.

owner

This attribute is a foreign key to Actor OBJECT, Optional.

created

This attribute is a foreign key to Moment DATETIME, Optional.

• scopeID

This attribute is a foreign key to ScopeID ALPHANUMERIC, Optional.

scopeOwner

This attribute is a foreign key to Account OBJECT, Optional.

• scopeDescr

This attribute is a foreign key to ScopeDescr ALPHANUMERIC, Optional.

scopeIsaCC

This attribute is a foreign key to Scope OBJECT, Optional, Unique.

5.4.54 Table: ScopeDescr

This table has the following 1 attributes:

ScopeDescr

This attribute implements the identity relation of ScopeDescr. ALPHANUMERIC, Mandatory, Unique.

5.4.55 Table: ScopeID

This table has the following 1 attributes:

• ScopeID

This attribute implements the identity relation of ScopeID. ALPHANUMERIC, Mandatory, Unique.

5.4.56 Table: scopeIII

This is a link-table, implementing the relation $Scope \xrightarrow{scopeIII} Scope$. It contains the following columns:

• SrcScope

This attribute is a foreign key to Scope OBJECT, Mandatory.

• TgtScope

This attribute is a foreign key to Scope OBJECT, Mandatory.

5.4.57 Table: SESSION

This table has the following 23 attributes:

• SESSION

This attribute is the primary key. OBJECT, Mandatory, Unique.

• sessionAccount

This attribute is a foreign key to Account OBJECT, Optional.

• sessionUserid

This attribute is a foreign key to UserID ALPHANUMERIC, Optional.

• loginUserid

This attribute is a foreign key to UserID ALPHANUMERIC, Optional.

loginPassword

This attribute implements the relation $SESSION \xrightarrow{loginPassword} Password$. PASSWORD, Optional.

• logoutRequest

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

• sessionReqdISOLevel

This attribute is a foreign key to ISOLevel INTEGER, Optional.

• sessionAuthISOLevel

This attribute is a foreign key to ISOLevel INTEGER, Optional.

sessionOrg

This attribute is a foreign key to Organization OBJECT, Optional.

• sessionPerson

This attribute is a foreign key to Person OBJECT, Optional.

• firstSessionAccount

This attribute is a foreign key to Account OBJECT, Optional.

sessionPersonRef

This attribute is a foreign key to PersonRef ALPHANUMERIC, Optional.

• sessionOrgRef

This attribute is a foreign key to OrgRef OBJECT, Optional.

• loginCreateAccount

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

registerPassword

This attribute implements the relation $SESSION \xrightarrow{registerPassword} Password$. PASSWORD, Optional.

• registerOrgRef

This attribute is a foreign key to OrgRef OBJECT, Optional.

loginReq

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

• loginISOLevel

This attribute is a foreign key to ISOLevel INTEGER, Optional.

sessionDev

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

• sessionLoginAssist

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

• sessionLogoutAssist

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

sessionSRI

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

• sessionSIA

This attribute is a foreign key to SESSION OBJECT, Optional, Unique.

5.4.58 Table: sessionActiveRoles

This is a link-table, implementing the relation $SESSION \xrightarrow{sessionActiveRoles} Role$. It contains the following columns:

SESSION

This attribute is a foreign key to SESSION OBJECT, Mandatory.

• Role

This attribute is a foreign key to Role OBJECT, Mandatory.

5.4.59 Table: sessionAllowedRoles

This is a link-table, implementing the relation $SESSION \xrightarrow{sessionAllowedRoles} Role$. It contains the following columns:

SESSION

This attribute is a foreign key to SESSION OBJECT, Mandatory.

• Role

This attribute is a foreign key to Role OBJECT, Mandatory.

5.4.60 Table: SIAMPersonRefComputation

This table has the following 4 attributes:

• SIAMPersonRefComputation

This attribute is the primary key. OBJECT, Mandatory, Unique.

siamcompFirstName

This attribute is a foreign key to FirstName ALPHANUMERIC, Optional.

• siamcompLastName

This attribute is a foreign key to LastName ALPHANUMERIC, Optional.

siamcompPersonRef

This attribute is a foreign key to PersonRef ALPHANUMERIC, Optional.

5.4.61 Table: sub1

This is a link-table, implementing the relation $RechtsGrond \xrightarrow{sub} Sub$. It contains the following columns:

• RechtsGrond

This attribute is a foreign key to RechtsGrond OBJECT, Mandatory.

• Sub

This attribute is a foreign key to Sub OBJECT, Mandatory.

5.4.62 Table: Sub2

This table has the following 1 attributes:

• Sub

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.63 Table: Tekst

This table has the following 1 attributes:

• Tekst

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.64 Table: Titel

This table has the following 1 attributes:

• Titel

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.65 Table: TText

This table has the following 20 attributes:

• TText

This attribute is the primary key. OBJECT, Mandatory, Unique.

• Template

This attribute implements the identity relation of *Template*. OBJECT, Optional, Unique.

Binding

This attribute implements the identity relation of *Binding*. OBJECT, Optional, Unique.

• ttTemplateParsedText

This attribute implements the relation $TText \xrightarrow{ttTemplateParsedText} TTPhrase$. BIGALPHANUMERIC, Optional.

$\bullet \ \ ttValueUsedToReplacePlaceholders$

This attribute is a foreign key to TTValue ALPHANUMERIC, Optional.

$\bullet \ ttInstanceResetReq \\$

This attribute is a foreign key to TText OBJECT, Optional, Unique.

• true

This attribute is a foreign key to Scope OBJECT, Optional.

• false

This attribute is a foreign key to Scope OBJECT, Optional.

• ttScope

This attribute is a foreign key to Scope OBJECT, Mandatory.

observed

This attribute is a foreign key to Moment DATETIME, Optional.

class

This attribute is a foreign key to Concept OBJECT, Optional.

• ttTrace

This attribute is a foreign key to Assignment OBJECT, Mandatory, Unique.

• ttName

This attribute is a foreign key to TTName ALPHANUMERIC, Optional.

• ttValue

This attribute is a foreign key to TTValue ALPHANUMERIC, Optional.

• ttTemplate

This attribute implements the relation $TText \xrightarrow{ttTemplate} TTPhrase$. BIGALPHANUMERIC, Optional.

ttInstance

This attribute implements the relation $TText \xrightarrow{ttInstance} TTPhrase$. BIGALPHANUMERIC, Optional.

• ttICO

This attribute is a foreign key to TText OBJECT, Optional.

• ttICCO

This attribute is a foreign key to TText OBJECT, Optional.

• ttDescr

This attribute implements the relation $TText \xrightarrow{ttDescr} TTPhrase$. BIGALPHANUMERIC, Optional.

• ttOwner

This attribute is a foreign key to Account OBJECT, Optional.

5.4.66 Table: ttIsUsedBy

This is a link-table, implementing the relation $TText \xrightarrow{ttIsUsedBy} TText$. It contains the following columns:

• SrcTText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TgtTText

This attribute is a foreign key to TText OBJECT, Mandatory.

5.4.67 Table: ttIsUsedByCopy

This is a link-table, implementing the relation $TText \xrightarrow{ttIsUsedByCopy} TText$. It contains the following columns:

• SrcTText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TgtTText

This attribute is a foreign key to TText OBJECT, Mandatory.

5.4.68 Table: ttIsUsedByStar

This is a link-table, implementing the relation $TText \xrightarrow{ttIsUsedByStar} TText$. It contains the following columns:

• SrcTText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TgtTText

This attribute is a foreign key to TText OBJECT, Mandatory.

5.4.69 Table: TTName

This table has the following 1 attributes:

• TTName

This attribute implements the identity relation of TTName. ALPHANUMERIC, Mandatory, Unique.

5.4.70 Table: TTPhrase

This table has the following 1 attributes:

• TTPhrase

This attribute implements the identity relation of TTPhrase. BIGALPHANUMERIC, Mandatory, Unique.

5.4.71 Table: ttTemplatePlaceholders

This is a link-table, implementing the relation $TText \xrightarrow{ttTemplatePlaceholders} TTName$. It contains the following columns:

• TText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TTName

This attribute is a foreign key to TTName ALPHANUMERIC, Mandatory.

5.4.72 Table: TTValue

This table has the following 1 attributes:

• TTValue

This attribute implements the identity relation of TTValue. ALPHANUMERIC, Mandatory, Unique.

5.4.73 Table: ttValueIsUsedInInstanceOfTText

This is a link-table, implementing the relation $TText \xrightarrow{ttValueIsUsedInInstanceOfTText} TText$. It contains the following columns:

• SrcTText

This attribute is a foreign key to TText OBJECT, Mandatory.

• TgtTText

This attribute is a foreign key to TText OBJECT, Mandatory.

5.4.74 Table: UID

This table has the following 4 attributes:

• UID

This attribute is the primary key. OBJECT, Mandatory, Unique.

• uidUserid

This attribute is a foreign key to UserID ALPHANUMERIC, Mandatory.

uidIssuer

This attribute is a foreign key to IdP OBJECT, Optional.

accUID

This attribute is a foreign key to Account OBJECT, Optional.

5.4.75 Table: url1

This is a link-table, implementing the relation $Regeling \xrightarrow{url} URL$. It contains the following columns:

• Regeling

This attribute is a foreign key to Regeling OBJECT, Mandatory.

• URL

This attribute is a foreign key to URL OBJECT, Mandatory.

5.4.76 Table: url2

This is a link-table, implementing the relation $RechtsGrond \xrightarrow{url} URL$. It contains the following columns:

• RechtsGrond

This attribute is a foreign key to RechtsGrond OBJECT, Mandatory.

• URL

This attribute is a foreign key to URL OBJECT, Mandatory.

5.4.77 Table: URL3

This table has the following 1 attributes:

• URL

This attribute is the primary key. OBJECT, Mandatory, Unique.

5.4.78 Table: UserID

This table has the following 1 attributes:

• UserID

This attribute implements the identity relation of UserID. ALPHANUMERIC, Mandatory, Unique.

5.5 Logical data model

Concept	С	R	U	D
Case		Case		
Case		Cases		
Artikel		Rechtsgrond		
Template		Case		
ScopeID		Cases		

Concept	С	R	U	D
Scope		Case		
Scope		Cases		
Scope		Garbage		
OrgRef		Login/Register		
OrgRef		ShowAcc		
URL		Rechtsgrond		
RechtsGrond		Rechtsgrond		
ISOLevel		Login/Register		
IfcText		Login/Register		
TTValue		Case		
LegalGround		Case		
LegalGround		Garbage		
Regeling		Rechtsgrond		
Evidence		Garbage		
Lid		Rechtsgrond		
Bijlage		Rechtsgrond		
TTName		Case		
Moment		Cases		
Moment		Garbage		
PersonRef		Login/Register		
PersonRef		ShowAcc		
Role		SessionRoles		
TTPhrase		Case		
TTPhrase		Garbage		
Aanwijzing		Rechtsgrond		
Password		Login/Register		
Sub		Rechtsgrond		