Al Assignment 1

Using Informed and Uninformed Search Algorithms to Solve 8-Puzzle

Names

Name	ID
Arsanuos Essa	17
Amr Mohamed Fathy	46
Muhammed Essam Khamis	66

CONTENTS

CONTENTS	•
A* algorithm	
Code	2
Data Structures	2
Explanation	2
Sample runs	3
DFS	
Code	Ę
Data Structures	5
Explanation	Ę
Sample runs	6
BFS	7
Code	7
Data Structures	7
Explanation	7
Sample runs	8
Assumptions	ç

1) A* algorithm

Code

```
def init (self, heuristic):
   self. heuristic = heuristic
   super(). init ()
def search (self, initial arr):
    .....
    self. initial state = State(initial arr, None, self. heuristic)
   if not self.check solvable (self. initial state):
        return {'steps': [[-1] * 9], 'cost': -1,
                'search depth': -1, 'nodes expanded': -1}
   states heap = []
   heappush (states heap, self. initial state)
    end = None
   while len(states heap):
        # break tie by FIFO criteria
       current explored state = heappop(states heap)
        if current_explored_state.is_goal():
            end = current explored state
           break
        if current explored state not in self.vis:
           self. explored += 1
           self.vis.add(current explored state)
            child states = self.expand(current explored state)
            for child in child states:
               heappush (states heap, child)
   res = {}
   res['steps'] = self.get_steps(end)
   res['cost'] = end.cost
   res['search depth'] = end.cost
   res['nodes expanded'] = self. explored
    return res
```

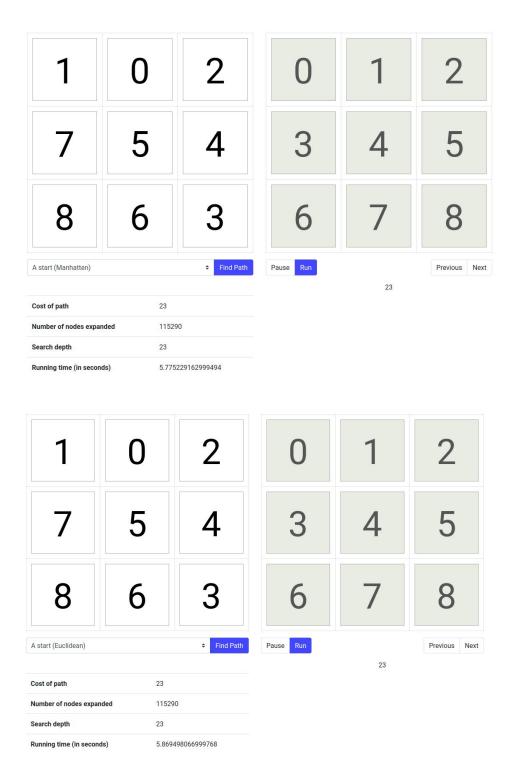
Data Structures

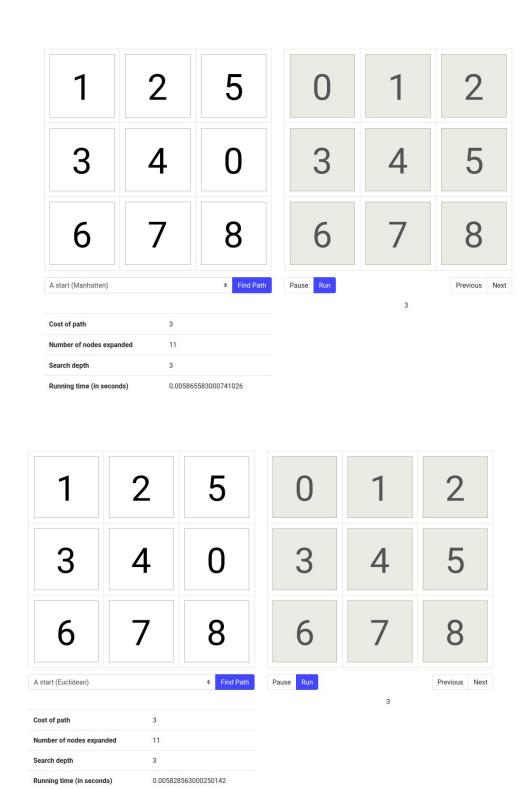
- Heap that represents a priority queue that returns the smallest element in the queue in terms of total cost.
- Set to keep the Visited(explored set) states to avoid repeating states using a hashing search in nearly O(1) for a faster search.

Explanation

- State object that contains a total cost representing f(n)
- We found that searching in frontier list (heap) is costly since it will be linear search O(n)
 So we assumed that we will insert all non visited children and also check if the not not visited the do nothing.
- At the end returning all required numbers to be displayed in UI.

• Sample runs





2) DFS

Code

```
def search(self, initial_state):
   res = {}
    start = State(initial state, None)
    if not self.check_solvable(start):
        return {'steps': [[-1] * 9], 'cost': -1,
                'search depth': -1, 'nodes expanded': -1}
    frontier = [start]
    final state = None
    while frontier:
        state = frontier.pop()
        self.vis.add(state)
      self. explored += 1
       if state.is goal():
            final state = state
           break
       neighbours = self.expand(state)
        for neighbour in neighbours:
            #if neighbour not in frontier:
            frontier.append(neighbour)
   steps = self.get steps(final state)
   res['steps'] = steps
   res['cost'] = final_state.cost
   res['search depth'] = final state.cost
    res['nodes expanded'] = self. explored
    return res
```

Data Structures

o Stack.

Explanation

- Frontier list that represents a stack that returns the first child of the current processed state.
- Visited(explored set) set that using a hashing search in nearly O(1) for a faster search to check for no duplication of the states visited to avoid loops in the search operation.
- We found that searching in stack is costly since it will be linear search O(n) So we
 assumed that we will insert all non visited children and also check if the not not visited the
 do nothing.
- At the end returning all required numbers to be displayed in UI.

• Sample runs

1	2	5	0	1	2
3	4	0	3	4	5
6	7	8	6	7	8
Cost of path Number of nodes exp Search depth	27 anded 28 27		Pause Run	27	Previous Next
1	0	2	0	1	2
7	5	4	3	4	5
78	5	3	3	7	5

3) **BFS**

Code

```
class BFS (Agent):
    def init (self):
        super().__init__()
        self. optimize flag = True
    def search(self, initial_state):
        curr_state = State(initial_state, None)
        if not self.check solvable(curr state):
            return { 'steps' : [[-1] * 9], 'cost' : -1,
                    'searc depth' : -1, 'nodes expanded' : -1}
        frontier = [curr state]
        while frontier:
           curr state = frontier.pop(0)
            self.vis.add(curr state)
         self._explored += 1
            if curr state.is goal():
                break
           children = self.expand(curr state)
            for child in children:
                #if child not in frontier:
                frontier.append(child)
        steps = self.get steps(curr state)
        res = {}
        res['steps'] = steps
        res['cost'] = curr state.cost
        res['search depth'] = curr state.cost
        res['nodes expanded'] = self._explored
        return res
```

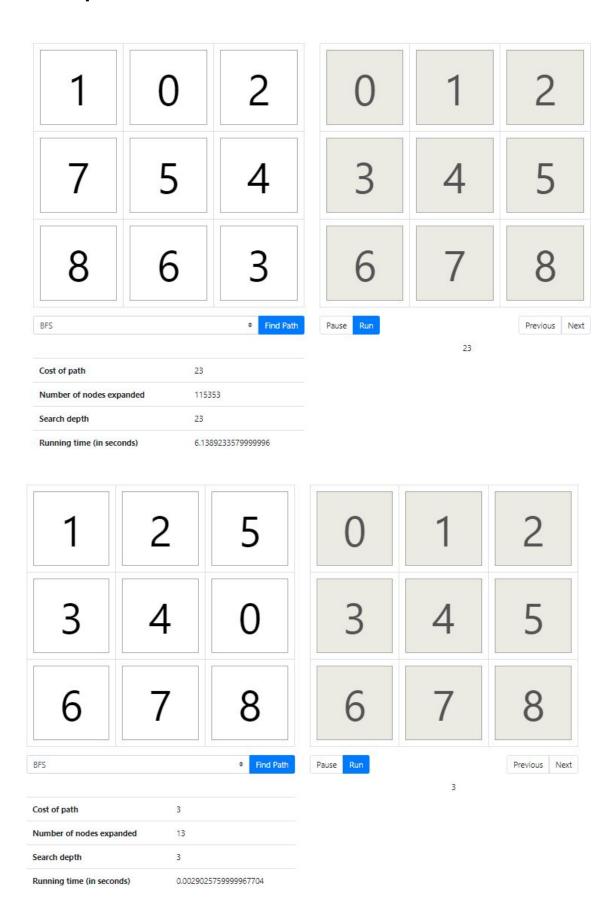
Data Structures

- Queue to represent the frontier list.
- Set to keep the Visited(explored set) states to avoid repeating states using a hashing search in nearly O(1) for a faster search.

Explanation

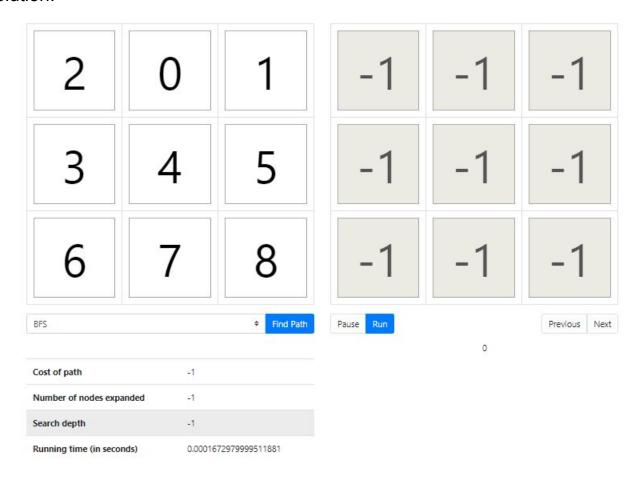
- State object that contains a total cost representing f(n)
- Visited(explored set) set that using a hashing search in nearly O(1) for a faster search.
- We pop the first inserted states (FIFO) then expand it to find all the child states and insert them in the queue to be explored later.
- At the end after we reached the goal state we return all required outputs to be displayed in UI.

• Sample runs



Assumptions

 We check If the puzzle is unsolvable by counting number of inversions in the initial grid, if it is unsolvable we will not run the algorithm as it will be waste of time then we will make alert with an error message in UI and fill the UI output with -1 to indicate an error occured. Otherwise the grid is solvable and we will find the solution.



 We set a timeout with 50 seconds, So if the algorithm can't solve the problem within 50 seconds then we will ignore the request.