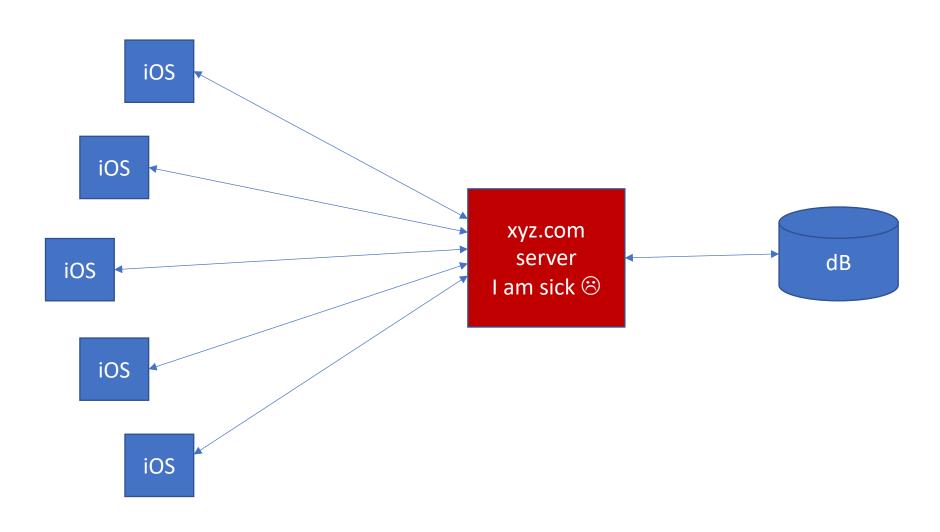


What We will be learning?



1. Benchmarking

1. Benchmarking(Recap)



1. Benchmarking Report

- Benchmarking Report: requests/sec = 100(-c), latency in response Time
 = 10 ms(Round Trip Time)
- 2. What is the latency of 101th, 201th response coming from the server?
- 3. Latency for 101th request = 10 ms(for first 100 requests) + 10 ms = 20 ms
- **4. Latency for 201th request** = 10 ms(for first 100 requests) + 10ms(for next 100 requests) + 10 ms = 30 ms
- 5. Eventually the server will get slower as the number of requests increase

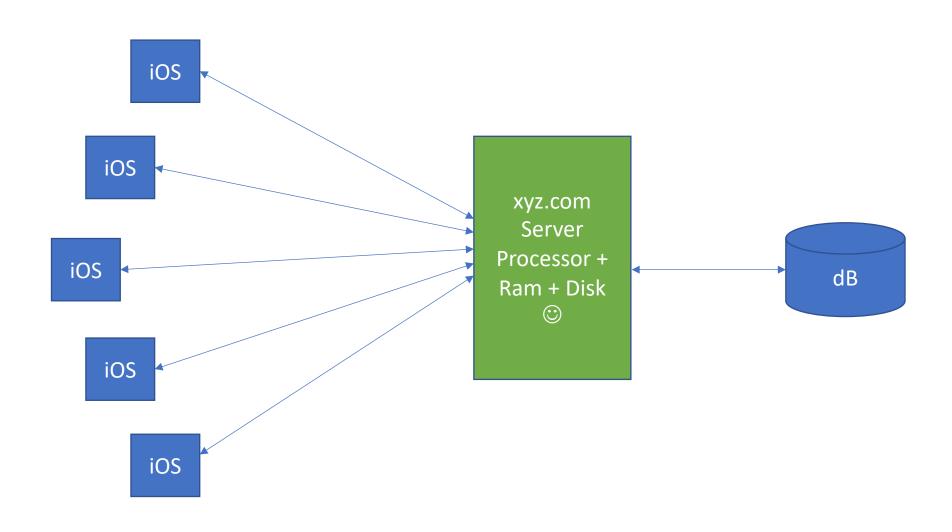
2. Vertical Scaling

2. Vertical Scaling

- 1. We increase the computation power(processer), memory and disk to scale the server. We **throw money to the problem** and buy a **bigger server machine**. This is called vertical scaling.
- 2. BM Report: 1000 req/sec, latency = 5 ms

xyz.com Server Processor + Ram + Disk

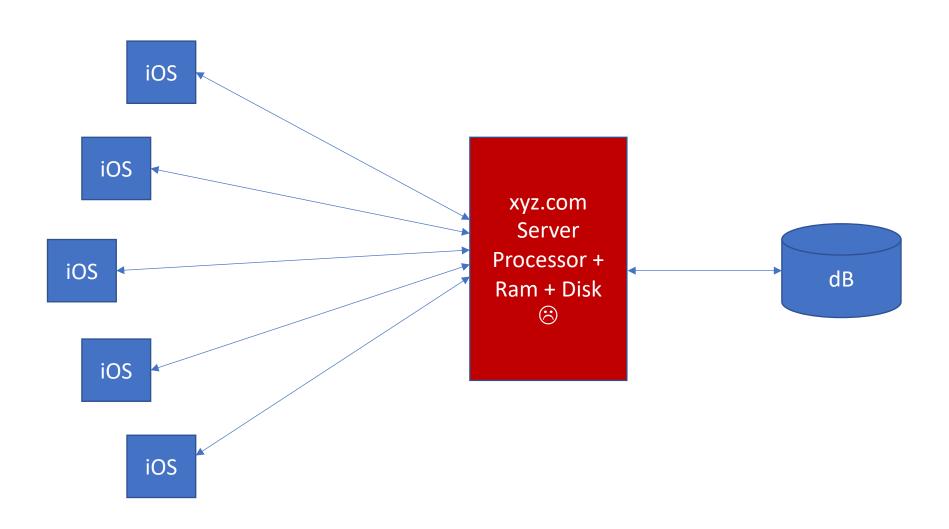
2. Vertical Scaling

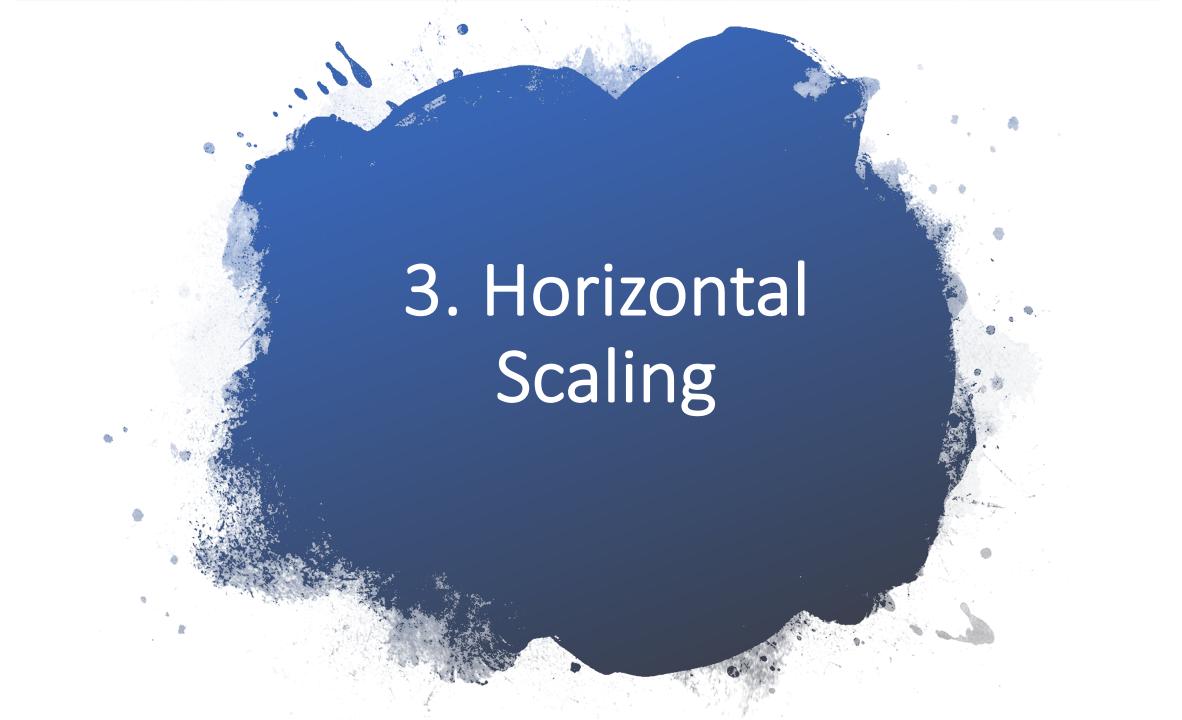


2. Concerns with Vertical Scaling

- 1. What happens if the number of requests/sec increases further?
- 2. There is a **ceiling** to it.!!
- 3. Eventually we are going to **exhaust our financial resources** or the **state of the art technology** we use.
- 4. This has a single point of failure.
- 5. What do we do when the number of requests increases to 10K?

2. Concerns with Vertical Scaling



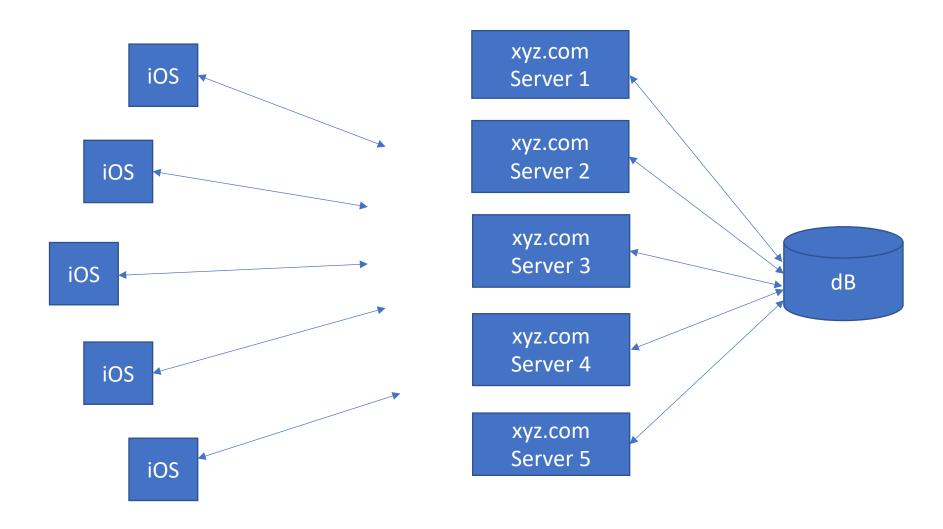


3. Horizontal Scaling

- 1. How should we architect our system to not reach that ceiling?
- 2. Instead of one state of the art costlier server, we get multiple cheaper servers to handle requests. This is called Horizontal Scaling.
- 3. Number of servers varies almost linearly with #requests(10K req/s)



3. Horizontal Scaling



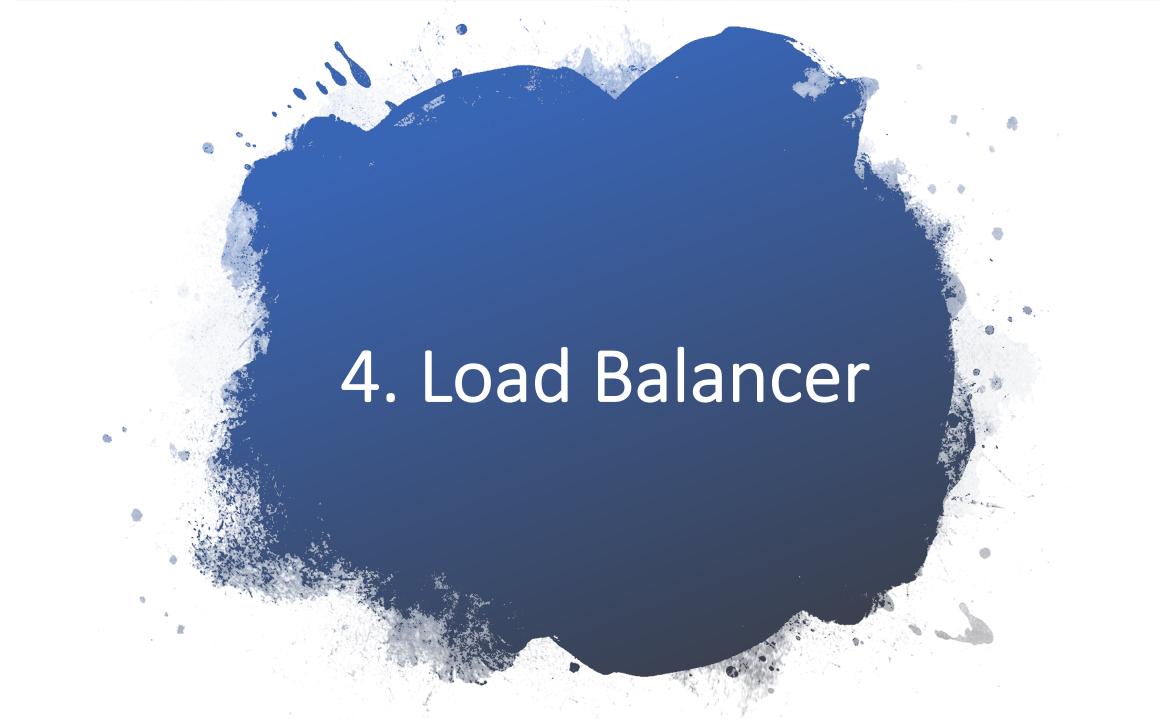
3. Horizontal Scaling

- 1. Number of servers varies almost linearly with #requests(1000 req/s). So we have solved the scalability problem
- 2. What **problems** arise now that we have 5 servers instead of 1?

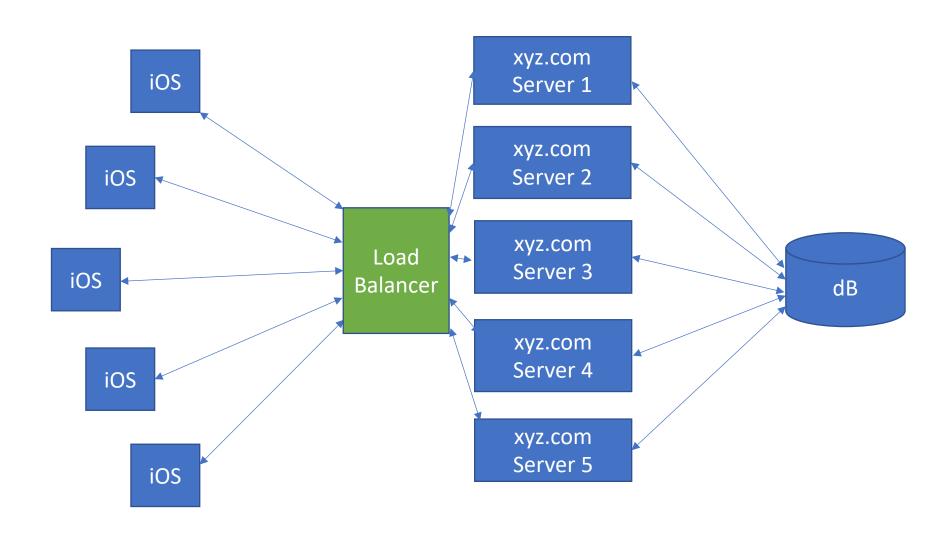
3. Concerns with Horizontal Scaling

1. One potential problem is that we have a single Postgre dB which all 5 servers are trying to access at the same time and concerns might arise. We will understand dB scalability later.

- 2. Second, How to figure out which server to send client requests to?
- This is addressed by another piece of hardware that's sits between client and servers, **Load Balancer**



4. Load Balancer



4. Load Balancer

How does a load balancer decide which server to send client requests to ?

4. Load Balancing Algorithms

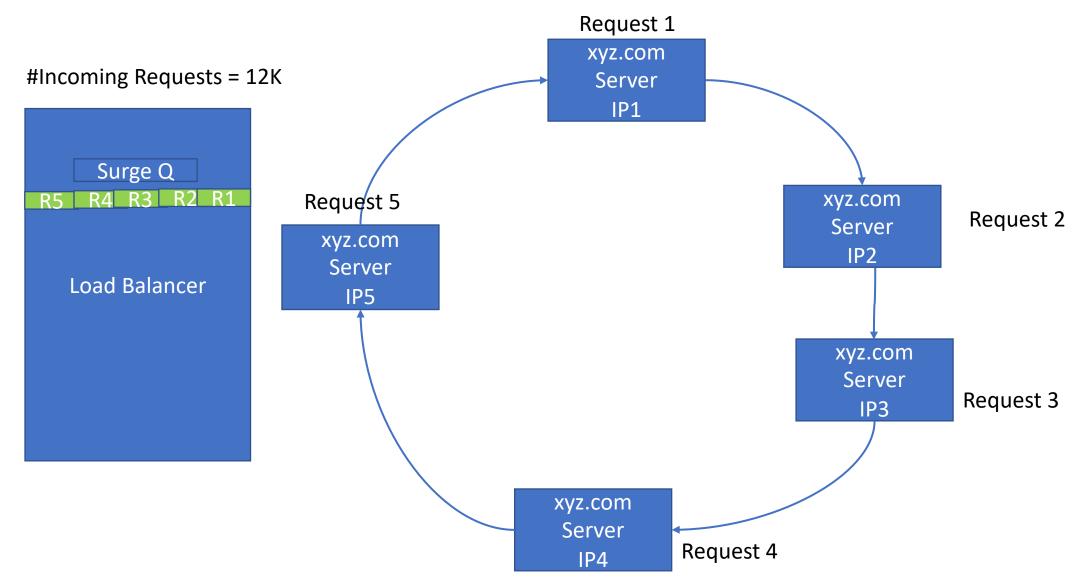
- 1. Random Choice
- 2. Round Robin/ Weighted Round Robin
- 3. Fewest Connections/ Least Loaded
- 4. Layer 4 load Balancing
- 5. Layer 7 load Balancing
- 6. IP Hashing Load Balancing

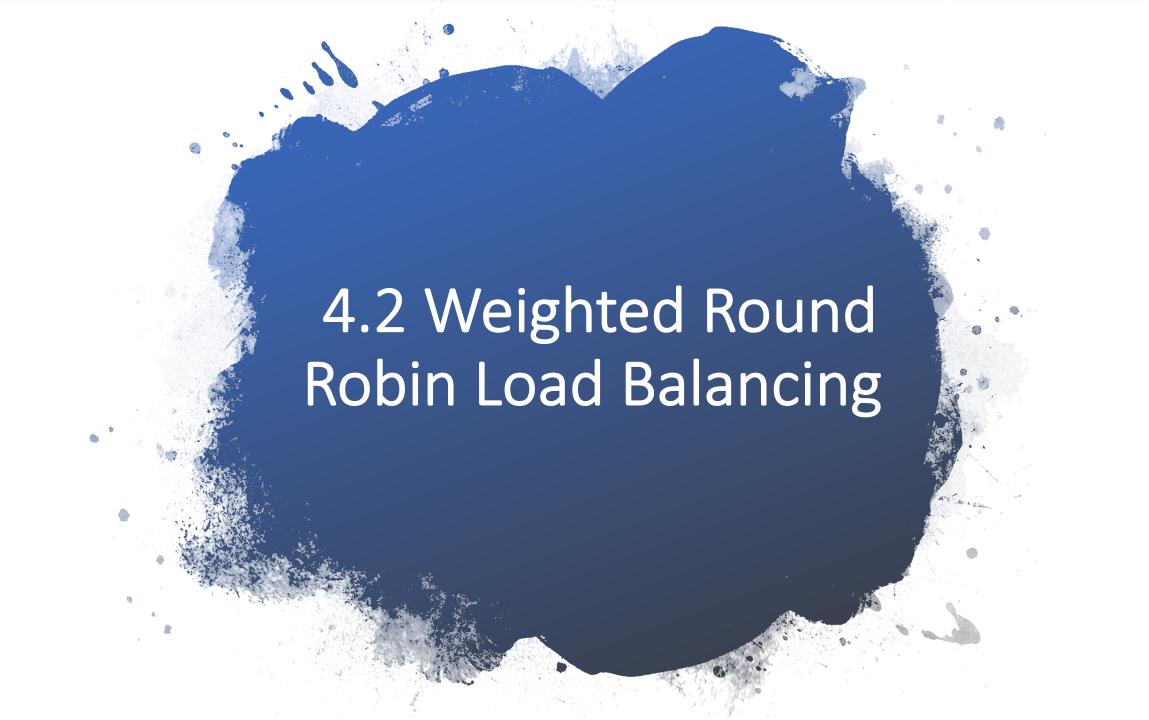
4.2 Round Robin Load Balancing

4.2 Round Robin Load Balancing

- 1. Each server has the domain xyz.com but unique IP address
- 2. Load Balancer has the list of all unique IP addresses for xyz.com
- 3. Incoming requests from Clients enq on a surge Queue inside LB.
- 4. Those requests are assigned in a rotating sequential manner
- 5. Request Cycle 1: Req 1 \rightarrow IP1, Req2 \rightarrow IP2, Req 3 \rightarrow IP3, Req 4 \rightarrow IP4, Req 5 \rightarrow IP5
- 6. Request Cycle 2: Req6 \rightarrow IP1, Req7 \rightarrow IP2...

4.2 Round Robin Load Balancing



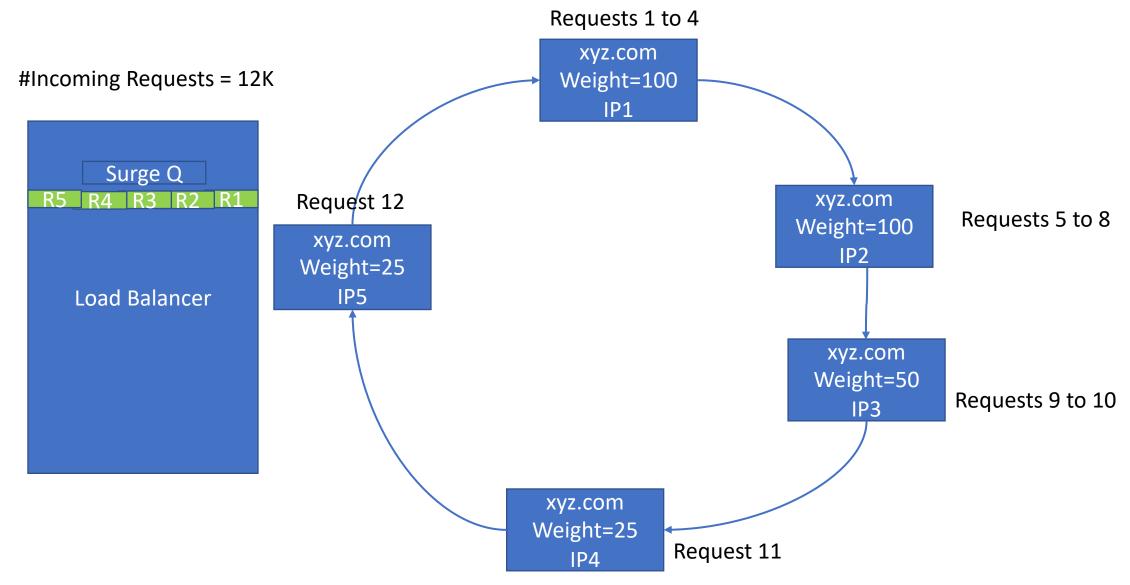


4.2 Weighted Round Robin Load Balancing

- 1. Let's say **Server 1** and **Server 2** are **most powerful** and highly efficient and remaining servers are less efficient.
- 2. So we **assign weights** to each server based on its performance capacity.
- 3. Incoming requests are assigned in a **rotating sequential manner based on server weights** on every request cycle.
- 4. Weights: S1: S2: S3: S4: S5 = 100: 100: 50: 25: 25 = 4: 4: 2: 1: 1

xyz.com Weight=100 IP1 xyz.com Weight=100 IP2 xyz.com Weight=50 IP3 xyz.com Weight=25 IP4 xyz.com Weight=25 IP5

4.2 Weighted Round Robin Load Balancing



Drawbacks of Round Robin?



4.3 Fewest Connections

- 1. Neither Round Robin or Weighted Round Robin take the current server load/connections into consideration when distributing requests
- 2. **Active Connections**: Those HTTP or HTTPS requests that have not yet received a response.
- 3. Current Active Connections:

xyz.com
Active
Conn=3
IP1

xyz.com Active Conn=5 IP2 xyz.com Active Conn=0 IP3

xyz.com Active Conn=10 IP4

xyz.com Active Conn=20 IP5

4.3 Fewest Connections

#Incoming Requests = 12K

Surge Q
R5 R4 R3 R2 R1
LB

xyz.com AC=3 IP1

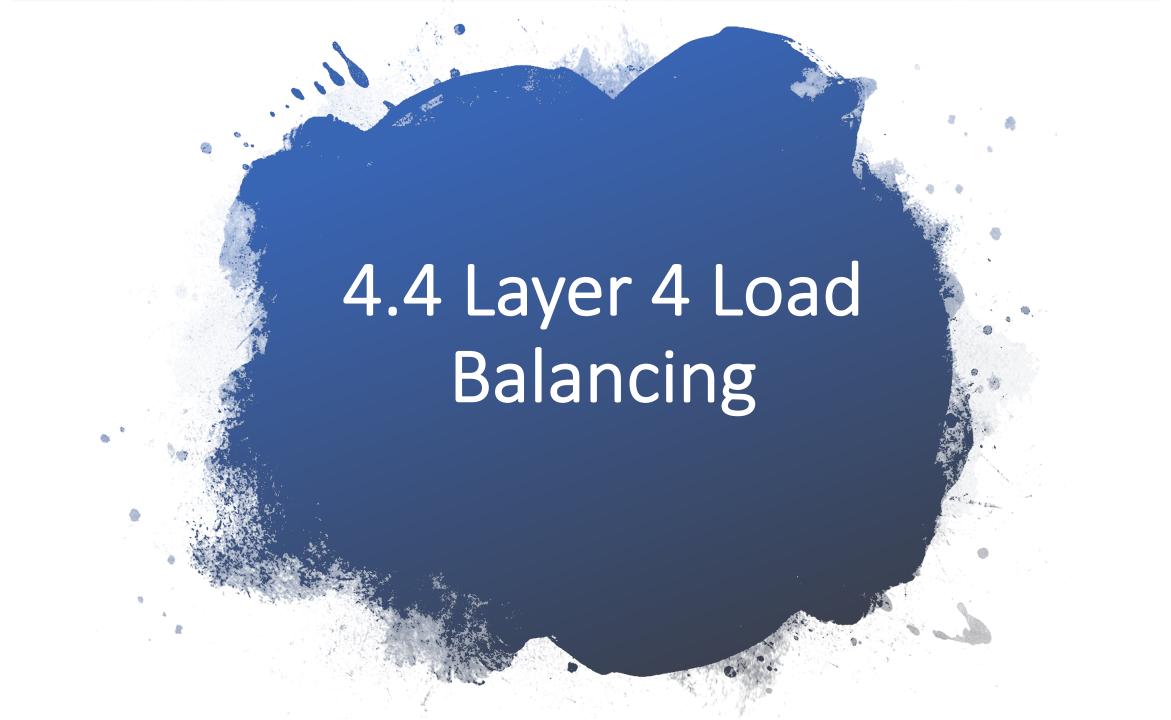
xyz.com AC=20 IP5 xyz.com AC=5 IP2

xyz.com AC=0 IP3

xyz.com AC=10 IP4

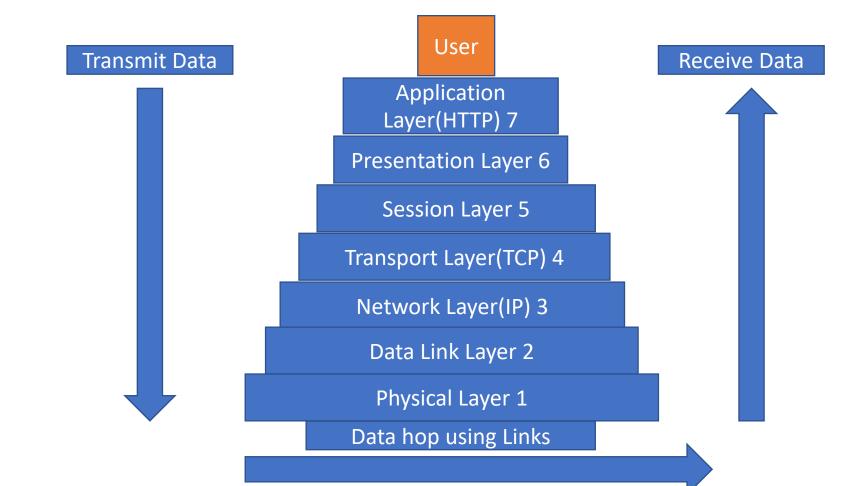
4.3 Fewest Connections

Request #	Server 1(AC= 3)	Server 2(AC=5)	Server 3(AC=0)	Server 4(AC=10)	Server 5(AC=20)
R1	3	5	1	10	20
R2	3	5	2	10	20
R3	3	5	3	10	20
R4	4	5	3	10	20
R5	4	5	4	10	20
R6	5	5	4	10	20
R7	5	5	5	10	20

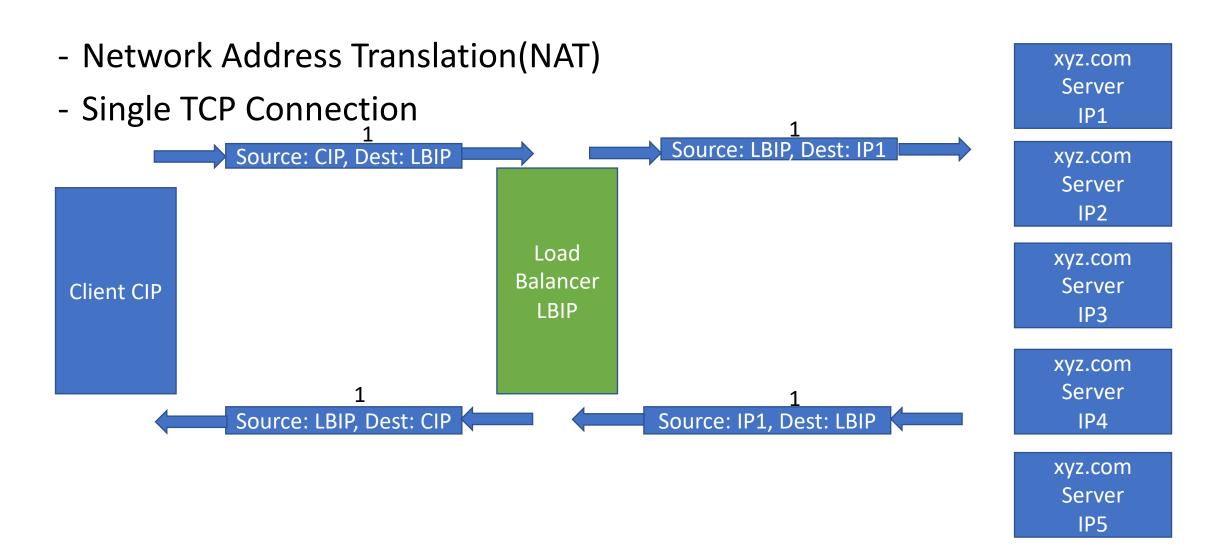


4.4 Layer 4 Load Balancing

7 Layer OSI Model:

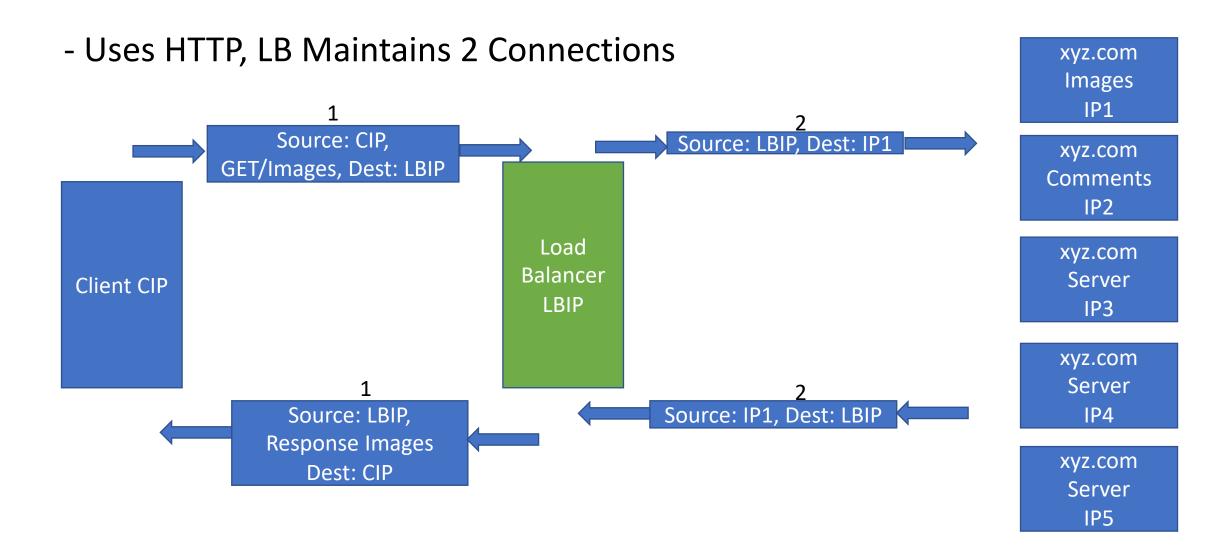


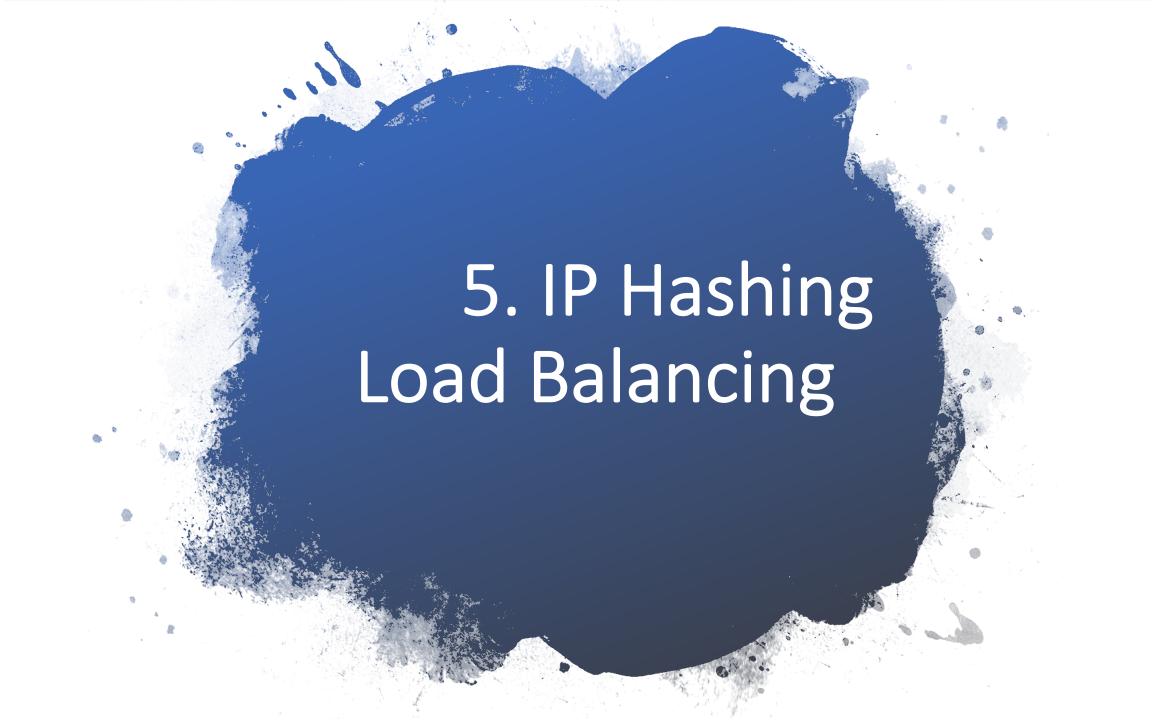
4.4 Layer 4 Load Balancing



4.5 Layer 7 Load Balancing

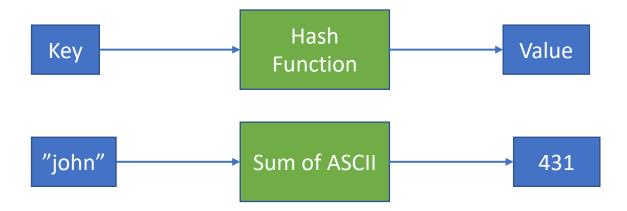
4.5 Layer 7 Load Balancing



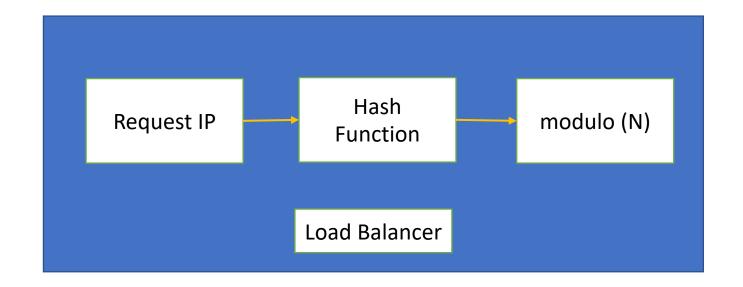


How do we load balance if the **number of servers change dynamically**? What happens if we want to **add/remove servers** based on client load? How do we **manage client sessions** during load balancing?

- 1. What is Hashing?
- Choosing an efficient Hash function to avoid collision (one to one mapping)



N = Number of Servers = 5



xyz.com Server ID - 0

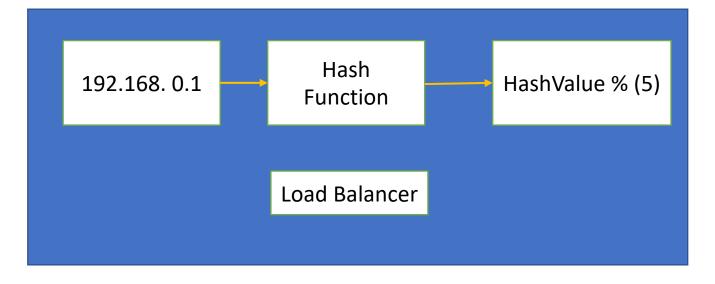
xyz.com Server ID - 1

xyz.com Server ID - 2

xyz.com Server ID - 3

xyz.com Server ID - 4

Hash Value	85	99	62	38	71
Server ID	0	4	2	3	1



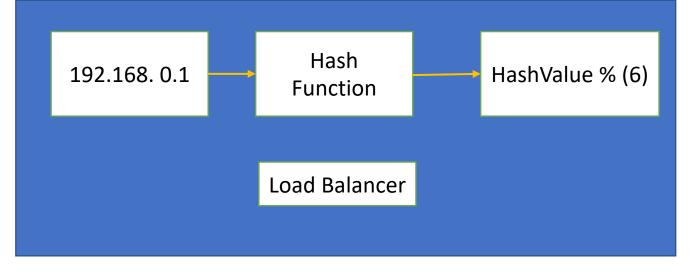
xyz.com Server ID - 0 xyz.com Server ID - 1 xyz.com Server ID - 2 xyz.com Server ID - 3 xyz.com Server ID - 4

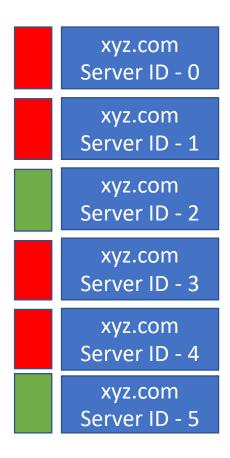
Session Cache

- 1. Incoming **Requests are mapped** to a server based on a hashing function **randomly**.
- 2. **Session Awareness**: The same server ID can be generated for next request in same session because the **request IP and hash function is same**.
- 3. **Server cache** can be reused.

What happens if we want to add a new server to my cluster?

Hash Value	85	99	62	38	71	42
Old SID	0	4	2	3	1	NA
New SID	1	3	2	2	5	0





5. Concerns with IP Hashing Load Balancing

- 1. All server cache gets invalidated because of redirection caused by addition of new server. This is costly if we have >10K servers.
- 2. How can we **reduce the impact on other servers** while adding or removing servers?

- We follow a three step process:

Step 1: Map requests to locations on a ring

Step 2: Map servers to locations on the ring

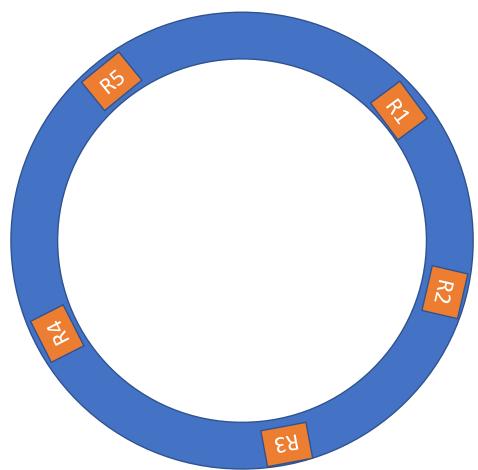
Step 3: Move clockwise and map requests to servers

Step 1: Map requests to locations on a ring using hash function H

- Min Hash Value = 0 degrees, Max Hash Value = 360 degrees

Req ID/IP	1	2	3	4	5
Hash Value	345	269	434	865	1056
Angle	0	56.5	100.3	170.4	220.7

Request Mapping:

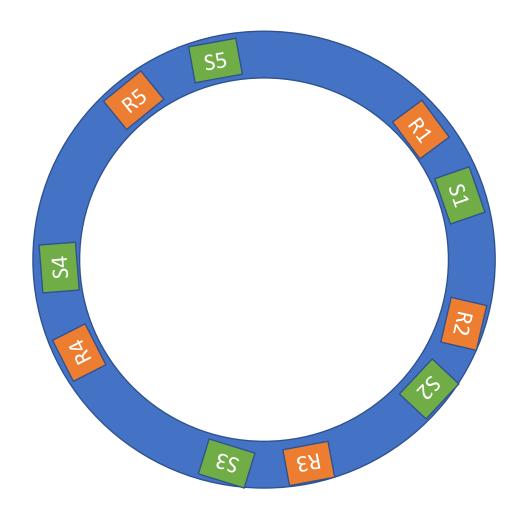


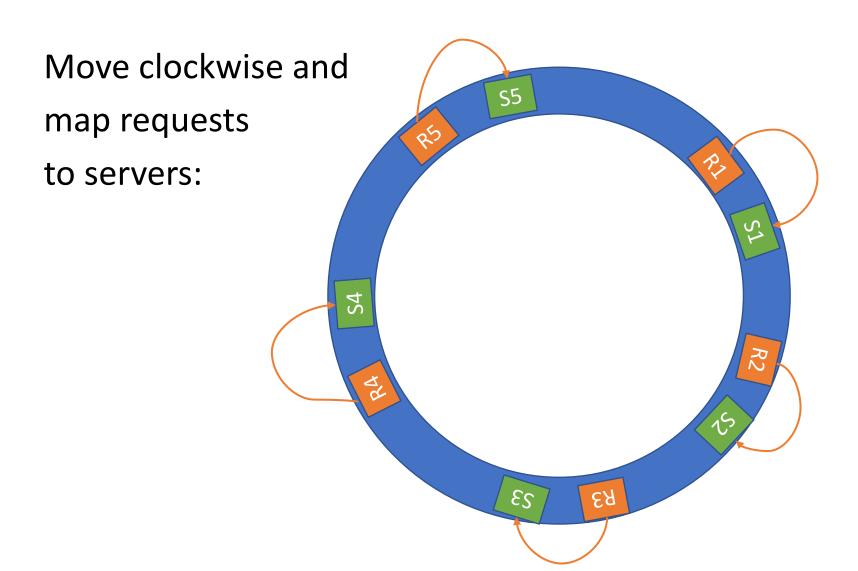
Step 2: Map servers to locations on a ring using same Hash function H

- Min Hash Value = 0 degrees, Max Hash Value = 360 degrees

Server ID/IP	10	12	14	16	18
Hash Value	20	300	543	810	1259
Angle	10.1	70.2	120.9	210.0	340.7

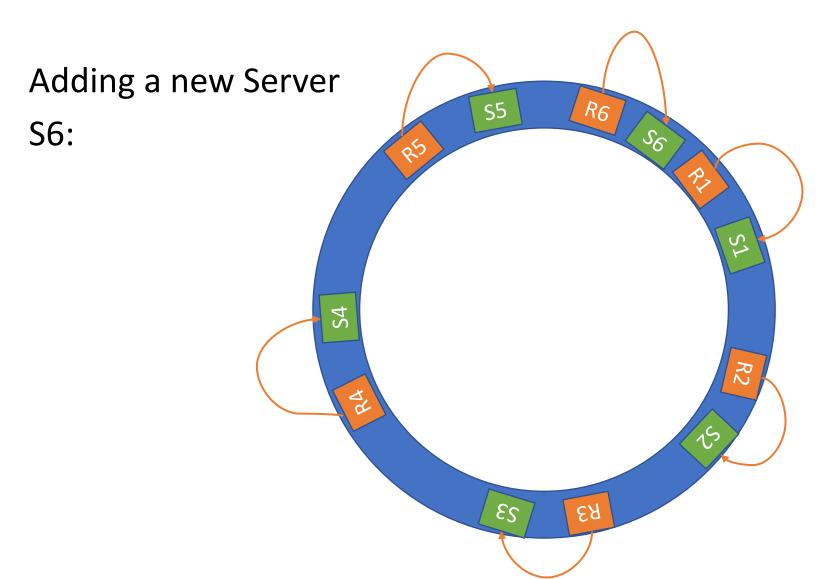
Server Mapping:

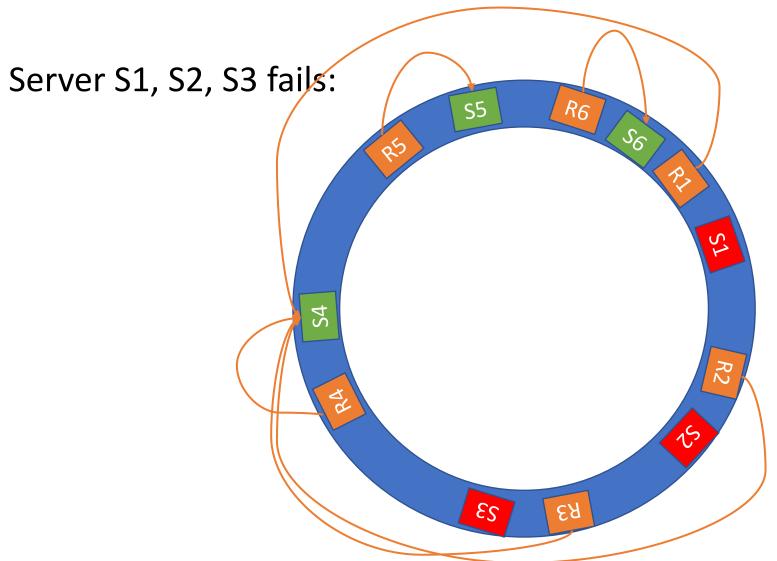




A new Request R6 Comes in:

A new Request R6 Comes in:





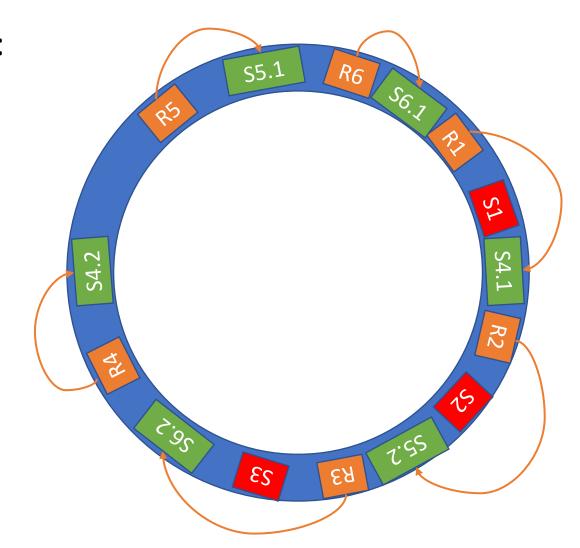
- 1. Virtual Nodes: Use another hash function to map position of servers S4, S5, and S6
- 2. Now Server S4, S5 and S6 have **two virtual positions** each. This helps us in **uniform distribution of load** among remaining healthy servers.

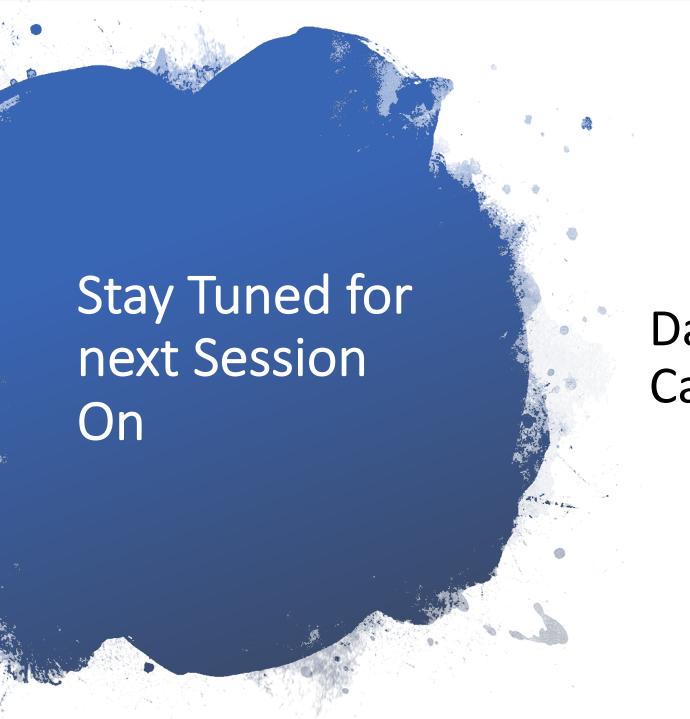
Virtual Nodes:

S4: 1, 2

S5: 1, 2

S6: 1, 2





Database Scaling, Caching



Questions

