

# Introduce the zkVM

The zkVM part for Engineers

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# What is a VM?

In [computing](#), a **virtual machine (VM)** is the [virtualization](#) or [emulation](#) of a [computer system](#). Virtual machines are based on [computer architectures](#) and provide the functionality of a physical computer.

keynotes: Emulate a physical computer

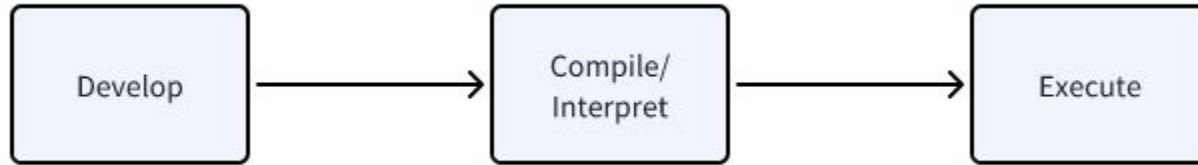
# Goal: running programs like a computer

OS/Application

# What is a VM in ZK?

- Same goal, running programs like a computer.
- Get the execution trace for proving system.

# The program lifecycle



- Expressing ideas using high-level programming languages.
- Lowering these expressions into assembly.
- Executing and verifying the result.

# Engineers only care about the experience

- Modern high-level language: Rust, Go, C++ etc.
- The mature ecosystem: Ide, Package manager, Building, Testing, Debugging etc.
- A robust and well-developed third-party libraries in work field: Web2, Web3 etc.

# ISA: the bridge of program&computer(VM)

In [computer science](#), an **instruction set architecture (ISA)** is an [abstract model](#) that generally defines how [software](#) controls the [CPU](#) in a computer or a family of computers.<sup>[1]</sup> A device or program that executes instructions described by that ISA, such as a central processing unit (CPU), is called an [implementation](#) of that ISA.

# Technology selection: ISA

## ZK10: Analysis of zkVM Designs - Wei Dai & Terry Chung

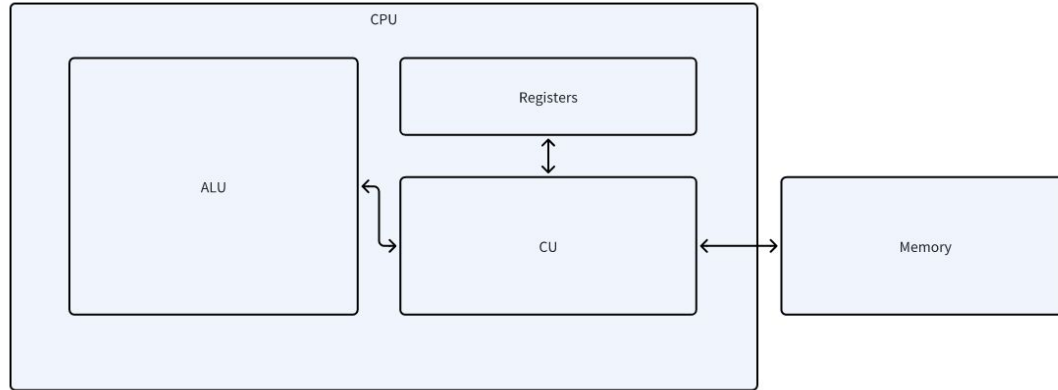
- For engineers:
  - Popular and widely supported ISA.
    - We can use any language we prefer.
- For ZK: **Performance! Performance! Performance!**
  - Hardware acceleration: SIMD and GPU friendly.
    - 32 bit data type for both SIMD and GPU(flops: FP32 >> FP64)
    - Tons of address operations(usize) and memory limitation in zkVM
    - Cache friendly for both CPU and GPU
  - Using small finite field.
    - Less than 32 bit maximum value.
  - The fewer the instructions, always better.

RISCV32/MIPS32 etc.



How to implement?

# The computer simple architecture

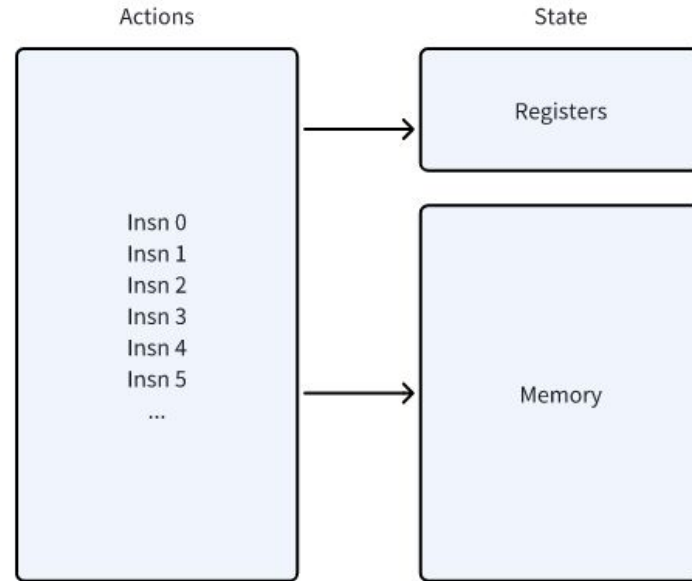


- Control Unit: Fetch/Decode/Dispatch the Instructions
- Arithmetic logic unit: Execute the Instructions
- Register: The data for Instructions
- Memory: The data for Instructions/Programs

# We don't need a “real computer” in the VM

High-level abstraction: Be like a finite-state machine

- State: Memory+Registers
- Actions: Instruction



# SP1 source code analysis: State

```
1 pub struct Executor<'a> {  
2     /// The program.  
3     pub program: Arc<Program>,  
4     /// The state of the execution.  
5     pub state: ExecutionState,  
6     /// The collected records, split by cpu cycles.  
7     pub records: Vec<ExecutionRecord>,  
8 }  
9
```

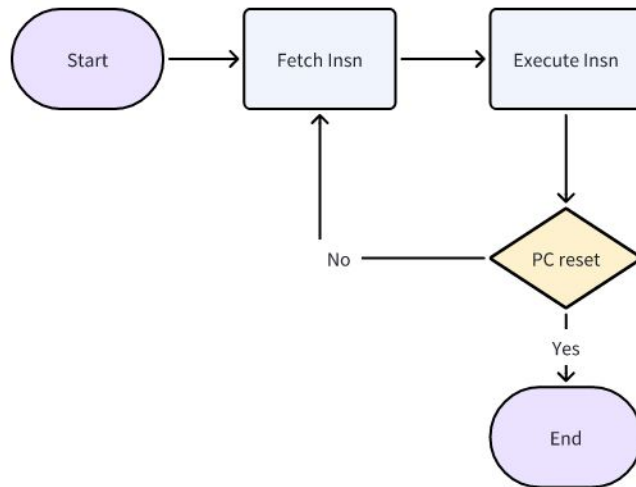
```
10 pub struct ExecutionState {  
11     /// The program counter.  
12     pub pc: u32,  
13     /// The memory which instructions operate over. Values contain the memory value and  
14     /// + timestamp that each memory address was accessed.  
15     pub memory: PagedMemory<MemoryRecord>,  
16     /// The clock increments by 4 (possibly more in syscalls) for each instruction that  
17     /// executed in this shard.  
18     pub clk: u32,  
19     /// A stream of input values (global to the entire program).  
20     pub input_stream: Vec<Vec<u8>>,  
21     /// A ptr to the current position in the input stream incremented by 'HINT_READ' ops  
22     pub input_stream_ptr: usize,  
23     /// A stream of public values from the program (global to entire program).  
24     pub public_values_stream: Vec<u8>,  
25     /// A ptr to the current position in the public values stream, incremented when read  
26     /// 'public_values_stream'.  
27     pub public_values_stream_ptr: usize,  
28 }
```

- The ELF format Program
- PC
- Memory/Register
- Clock
- Witness input
- Public value

# SP1 source code analysis: Actions

```
1  /// Executes the program, returning whether the program  
2  /// has finished.  
3  pub fn execute(&mut self) -> Result<bool, ExecutionError>;  
4
```

- Fetch and Decode Insn
- Execute Insn
- Check PC == 0



# SP1 source code analysis: Actions

```
1 // method execute
2 loop {
3     if self.execute_cycle()? {
4         break;
5     }
6 }
7
```

- Long loop until program exit
- Category for different purposes
- Changing the State

```
9 // method execute_cycle
10 let instruction = self.fetch();
11 self.execute_instruction(&instruction)?;
12
```

```
12
13 // method execute_instruction
14 match instruction.opcode {
15     ... => self.execute_alu(...)
16     ... => self.execute_load(...)
17     ... => self.execute_store(...)
18     ... => self.execute_branch(...)
19     ... => execute_jump
20     ... => execute_syscall
21 }
```

# Same design as ZKM

```
1 pub struct State {  
2     pub memory: Box<Memory>,  
3     pub registers: [u32; 32],  
4     pub pc: u32,  
5     pub input_stream: Vec<Vec<u8>>,  
6     pub input_stream_ptr: usize,  
7     pub public_values_stream: Vec<u8>,  
8     pub public_values_stream_ptr: usize,  
9 }
```

State

```
11 // method split_prog_into_segs  
12 loop {  
13     if instrumented_state.state.exited {  
14         break;  
15     }  
16     let cycles = instrumented_state.step();  
17 }
```

Actions

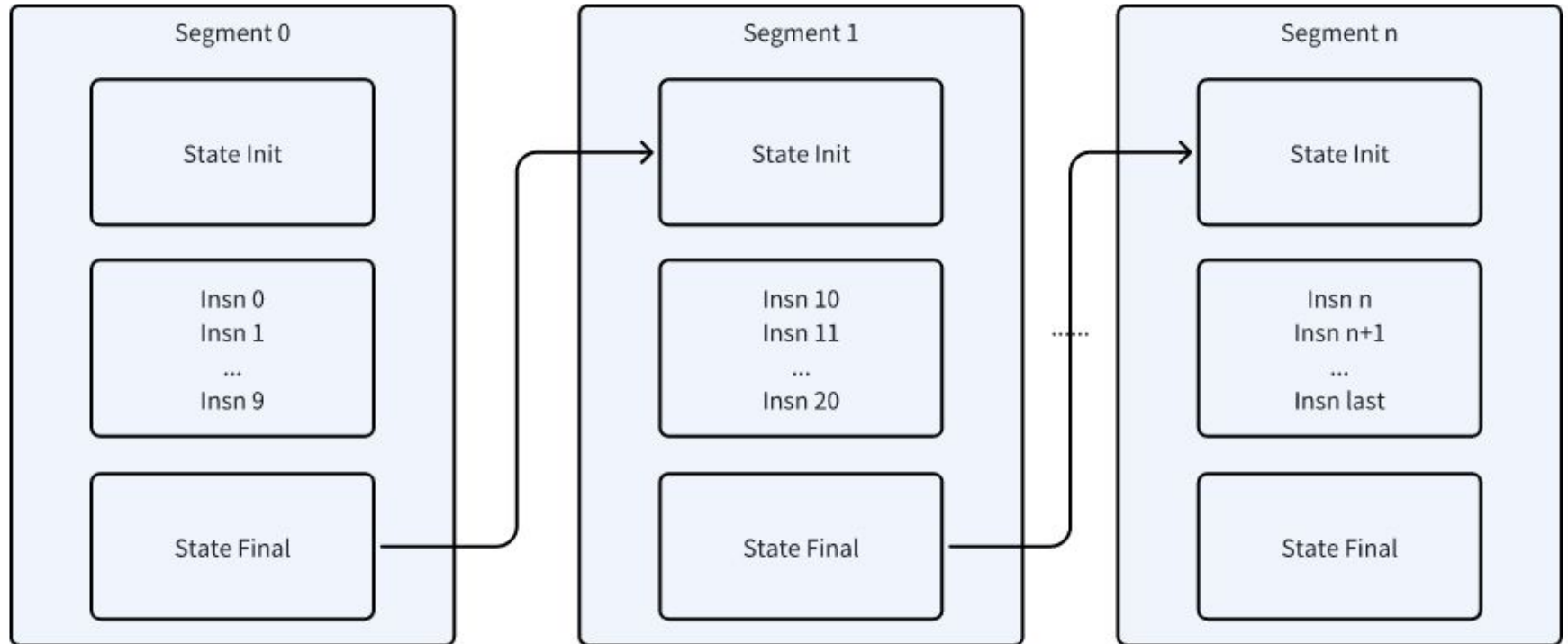
A deeper dive into



# Precompiles: performance boost

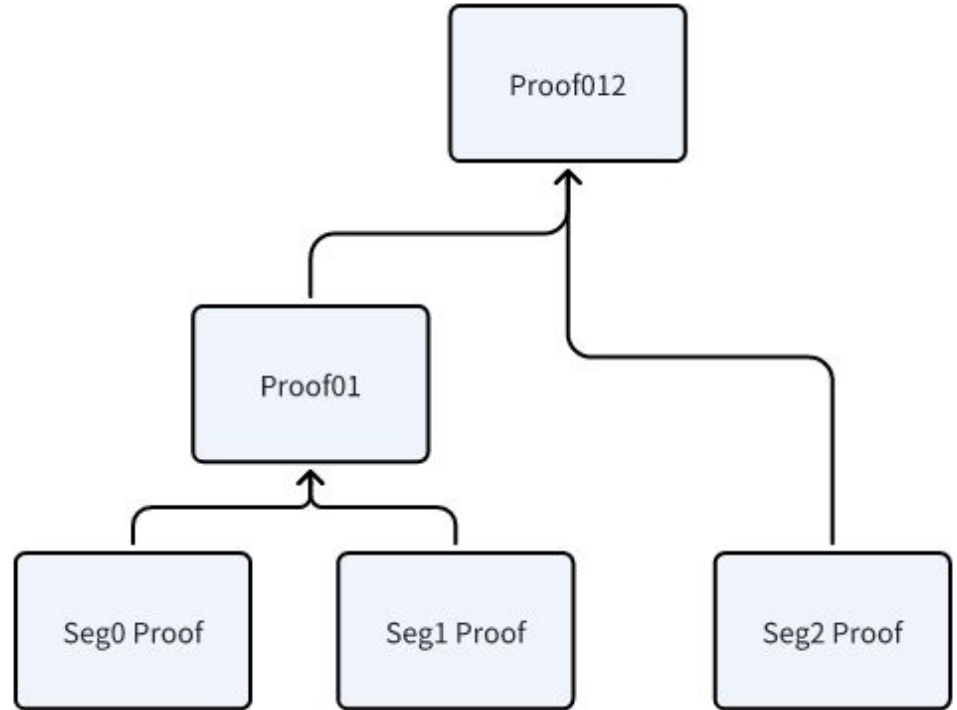
- Precompile universal and stable actions: elliptic curve arithmetic, hashing etc. [\[1\]](#)
- How to?
  - Custom instructions:
    - Pros:
      - Almost language support custom assembly, like ``asm!`` in Rust.
        - `asm!("myinsn", args)`
      - Less patches for language, only for third-party library you wanted.
    - Cons:
      - Unfriendly to third-party VMs and toolchains.
  - Custom syscalls(ABI):
    - Pros:
      - Without break the spec of ISA.
    - Cons:
      - More patches for language
        - New target tier, like ``zkvm`` in Rust.
        - New ABI for new tier

# Continuation/Segment



# Aggregation: Recursive proving

- Verify two proofs recursively each time.



# And more ...

- IO stuff
  - Witness input/Public values.
  - Standard IO: stdin, stdout, stderr.
    - Debug
    - Trace
    - ...
- Memory Pager
  - Improve Merkle proof efficiency.
  - Improve the read/write of VM.

# Summary

“Zero-knowledge” for Users

What's next?

# Looking Ahead

- zkVM based toolchains: custom ABI(IO)
  - Debugger: lldb, gdb [issue](#)
  - Profiler: CPU/Memory profiling
  - ...
- Auto precompiles:
  - “Inline Recursive”
    - Find and verify the proof of instruction sequence in custom program. (what the zkVM do manually)
    - We can even share these general time-consuming algorithm proofs online
  - Circuit Compiler, JIT like
    - Zircen Circuit Compiler
- Standard VM for zkVMs
  - Standardized syscalls(precompiles).
  - Standardized execution trace.
  - We can share the same patches for the upstream.
  - We can share the zkVM based toolchains.