On the Time Complexity of Common All-Reduce Methods

Kun Yuan

(Internal Reference Only)

September 7, 2019

Outline

1. Training framework for distributed deep learning

2. Time complexity for each all-reduce method

 ${\it 3. \,\, Decentralized \,\, algorithms}$

Outline

1. Training framework for distributed deep learning

2. Time complexity for each all-reduce method

3. Decentralized algorithms

Empirical Risk Minimization

• The empirical risk minimization (ERM) problem is

$$\min_{w} F(w) = \frac{1}{N} \sum_{n=1}^{N} Q(w; x_n)$$

where $\{x_n\}_{n=1}^N$ is the dataset; $Q(w;x_n)$ is some cost function

• The above ERM problem can be rewritten as

$$\min_{w} F(w) = \frac{1}{P} \sum_{p=1}^{P} F_p(w), \quad \text{where} \quad F_p(w) = \frac{1}{L} \sum_{n=1}^{L} Q(w; x_{p,n})$$
 (1)

where $\{x_{p,n}\}_{n=1}^{L}$ is the data assigned to node p, and N=LP.

Gradient Descent Method

• The gradient descent method to solve problem (1) is

$$w^{k+1} = w^k - \frac{\alpha}{P} \sum_{p=1}^P \nabla F_p(w^k),$$
 where
$$\nabla F_p(w^k) = \frac{1}{L} \sum_{n=1}^L \nabla Q(w; x_{p,n})$$

where α is the step-size.

- The computation of the true gradient $\nabla F_p(w^k)$ is time-consuming.

Stochastic Gradient Descent

• We compute an approximate gradient for $\nabla F_p(w)$, i.e., $\widehat{\nabla F_p}(w)$:

$$w^{k+1} = w^k - \frac{\alpha}{P} \sum_{p=1}^P \widehat{\nabla F_p}(w^k),$$
 where $\widehat{\nabla F_p}(w^k) = \frac{1}{B} \sum_{n \in \mathcal{R}} \nabla Q(w; x_{p,n})$

where \mathcal{B} is a small batch of index, and $B = |\mathcal{B}|$.

- The batch is usually sampled without replacement.
- SGD is a standard training algorithm for deep learning.
- More advanced training methods: ADAM and its variant

Training framework for deep learning: Parameter Server

$\textbf{Algorithm 1} \ \, \text{Distributed training framework: Parameter Server}$

Setting: P nodes; Partitioned data set $\{\{x_{1,n}\}_{n=1}^{N/K}, \cdots, \{x_{K,n}\}_{n=1}^{N/K}\}$

For $k=1,2,3\cdots$ until meet a criterion:

All (P-1) slave nodes do in parallel:

Pull w^k from the master node P;

Sample a batch of local data $\{x_{p,\mathcal{B}}\}$;

Evaluate a local stochastic gradient $\widehat{\nabla F}_p(w^k; \{x_{p,\mathcal{B}}\});$

Push $\widehat{\nabla F}_p(w^k; \{x_{p,\mathcal{B}}\})$ to the master node P;

The master node P does:

Reduce all received gradients: $\widehat{\nabla F}(w^k) = \frac{1}{P} \sum_{n=1}^{P} \widehat{\nabla F}_p(w^k; \{x_{p,\mathcal{B}}\})$

Update the variable: $w^{k+1} = w^k - \alpha \widehat{\nabla F}(w^k)$

Output: w^k

Training framework for deep learning: Parameter Server

- ullet The all-reduce procedure in PS bases on "reduce ightarrow broadcast"
- Has to distinguish master node from slave nodes
- ullet Four implementations on "reduce ightarrow broadcast"
 - Star implementation
 - Tree implementation
 - Cycle implementation
 - BytePS implementation
- Will discuss each implementation later

Training framework for deep learning: Ring-Allreduce

Algorithm 1 Distributed training framework: Ring-Allreduce

Setting: P nodes; Partitioned data set $\{\{x_{1,n}\}_{n=1}^{N/K},\cdots,\{x_{K,n}\}_{n=1}^{N/K}\}$

For $k = 1, 2, 3 \cdots$ until meet a criterion:

Each node does in parallel:

Sample a batch of local data $\{x_{p,\mathcal{B}}\}$;

Evaluate a local stochastic gradient $\widehat{\nabla F}_p(w^k; \{x_{p,\mathcal{B}}\})$;

Run ring-allreduce and get $\widehat{\nabla F}(w^k) = \frac{1}{P} \sum_{n=1}^{P} \widehat{\nabla F}_p(w^k; \{x_{p,B}\});$

Update the variable: $w^{k+1} = w^k - \alpha \widehat{\nabla} \widehat{F}(w^k)$

Output: w^k

Important note: after ring-allreduce, every node will get $\widehat{\nabla F}(w^k)$

Training framework for deep learning: Ring-Allreduce

- Ring-allreduce framework is fundamentally different from PS
- In ring-allreduce:
 - Every node will get $\widehat{\nabla F}(w^k)$ after ring-allreduce
 - Every node will perform the SGD update
- In parameter server:
 - The master node will get $\widehat{\nabla F}(w^k)$ after reduce
 - The master node will perform the SGD update

Outline

1. Training framework for distributed deep learning

2. Time complexity for each all-reduce method

3. Decentralized algorithms

Star All-reduce

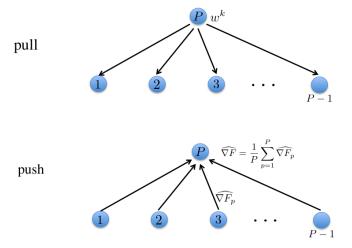


Figure: Illustration of the star all-reduce.

Star All-reduce

- We assume some parameters
 - M: the dimension of the vector
 - B: the bytes used to store each element in the vector
 - t_m : the time used to transmit a unit byte
 - t_c : the time used to establish one connection
- The time complexity for one star-allreduce step:

$$T_{\text{star}} = 2(P-1)(MBt_m + t_c)$$

• Can be implemented in an asynchronous manner; has theoretical guarantee

Tree All-reduce

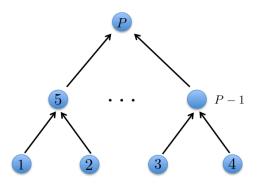


Figure: Illustration of the tree all-reduce.

Tree All-reduce

- The time complexity for one tree-allreduce step:
 - Time spent in each level of the tree: $2(MBt_m+t_c)$
 - There are $\log(P)$ levels in total
 - There are a "push" step and a "pull" step

$$T_{\text{tree}} = 4\log(P)(MBt_m + t_c)$$

Hard to be extended to the asynchronous version

Cycle All-reduce

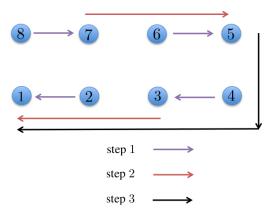


Figure: Illustration of the cycle all-reduce. In the figure, all information are merged to node $1. \ \ \,$

Cycle All-reduce

- The time complexity for one cycle-allreduce step:
 - Time spent in each step: $MBt_m + t_c$
 - There are $\log(P)$ steps in total
 - There are a "reduce" step and a "broadcast" step

$$T_{\text{cycle}} = 2\log(P)(MBt_m + t_c)$$

Hard to be extended to the asynchronous version

BytePS All-reduce

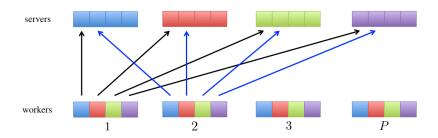


Figure: Illustration of the BytePS all-reduce.

BytePS All-reduce

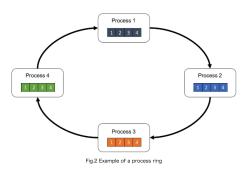
• The time complexity for one BytePS-allreduce step:

$$T_{\text{BytePS}} = 2(MBt_m + P^2t_c)$$

Can be improved to better complexity (but with more complicated ops.)

$$T_{\text{BytePS}} = 2(MBt_m + Pt_c)$$

- Can be extended to the asynchronous version (not quite sure)
- Require more (but cheap) servers



 $\label{eq:Figure: Illustration of the ring-all reduce.} \\$

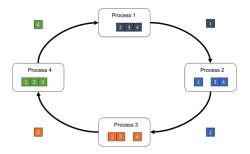


Figure: Illustration of the ring-allreduce.

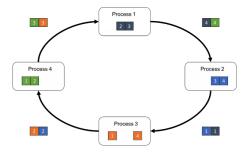


Figure: Illustration of the ring-allreduce.

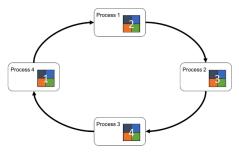


Fig.5 After P-1 steps, each process has a reduced subarray.

Figure: Illustration of the ring-allreduce.

- The time complexity for one Ring-allreduce step:
 - Time spent at each step: $\frac{MB}{P}t_m + t_c$
 - There are 2(P-1) steps in total

$$T_{\rm ring} = 2(P-1)(\frac{MB}{P}t_m + t_c) \approx 2(MBt_m + Pt_c)$$

- Ring-allreduce has more elegant operations compared to BytePS
- Hard to be extended to the asynchronous version

Outline

1. Training framework for distributed deep learning

2. Time complexity for each all-reduce method

 ${\it 3. \,\, Decentralized \,\, algorithms}$

Core Idea

- All the above all-reduce methods traverse through all nodes
- A brand new idea:

Use the local average to approximate global average asympototically

Further lower down the above-mentioned time complexity

Average Consensus

- Calculate the average $\bar{x} = \frac{1}{P} \sum_{p=1}^{N} x_p$ over an arbitrary graph
- Core idea: each agent does local average

$$x_p^{k+1} = \sum_{q \in \mathcal{N}_p} a_{qp} x_q^k, \quad \text{where} \quad \sum_{q \in \mathcal{N}_p} a_{qp} = 1$$

- Convergence guarantee: $\|x_p^k \bar{x}\| = \rho^k$ where $\rho < 1$ (linear convergence)
- This implies $x_p^k \to \bar{x}$ asymptotically for any $p=1,2,\cdots,P$.
- Proof:

$$\|\mathbf{x}^{k} - \bar{\mathbf{x}}^{k}\| = \|A\mathbf{x}^{k-1} - \frac{1}{n}\mathbb{1}\mathbb{1}^{T}\bar{\mathbf{x}}^{k-1}\| \le \|A - \frac{1}{n}\mathbb{1}\mathbb{1}^{T}\|\|\mathbf{x}^{k-1} - \bar{\mathbf{x}}^{k-1}\|$$

$$\le \rho \|\mathbf{x}^{k-1} - \bar{\mathbf{x}}^{k-1}\|$$

where
$$\bar{\mathbf{x}}^k = \bar{\mathbf{x}}^{k-1} = \cdots = \bar{\mathbf{x}}^0$$
.

Average Consensus

$$x_P^{k+1} = \sum_{q \in \mathcal{N}_P} a_{qP} x_P^k$$

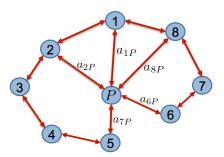


Figure: Illustration of an arbitrary graph and the consensus update.

Decentralized Stochastic Gradient Descent (Diffusion)

Recall the standard SGD:

$$\begin{split} w^{k+1} &= w^k - \frac{\alpha}{P} \sum_{p=1}^P \widehat{\nabla F_p}(w^k), \\ \text{where} \quad \widehat{\nabla F_p}(w^k) &= \frac{1}{B} \sum_{n \in \mathcal{B}} \nabla Q(w^k; x_{p,n}) \end{split}$$

Diffusion:

$$w_p^{k+1} = \sum_{q \in \mathcal{N}_p} a_{qp} \Big(w_p^k - \alpha \widehat{\nabla F_p}(w^k) \Big) \quad \text{(Local average)}$$

Theoretical Guarantee

• When each local cost function $F_p(w)$ is convex

$$\begin{split} & \limsup_{k \to \infty} \frac{1}{P} \sum_{p=1}^P \mathbb{E} \| w_p^k - w^\star \|^2 = O \bigg(\frac{\alpha \sigma^2}{P \nu} \bigg), \quad \text{(SGD)} \\ & \limsup_{k \to \infty} \frac{1}{P} \sum_{p=1}^P \mathbb{E} \| w_p^k - w^\star \|^2 = O \bigg(\frac{\alpha \sigma^2}{P \nu} + \frac{\lambda^2 \alpha^2 \sigma^2}{1 - \lambda} \bigg), \quad \text{(Diffusion)} \end{split}$$

where σ^2 is the gradient noise, λ indicates the connectivity of the network

• When α is sufficiently small or the network is well-connected $(\lambda \to 0)$,

$$SGD \approx Diffusion$$

• We can use local average per iteration without effecting the convergence!!

Local Averaging Time Complexity

- We use the Power-2 graph.
- Star local-averaging: $T = (MBt_m + t_s) \log P$; support asynchrony.
- Ring local-averaging: $T=2(MBt_m+\sqrt{P}\,t_s)$; not support asynchrony.
- Better than their counterparts.

References: