

## CSCI 551 Lab 3 Part 1 Report

**Overview:** For lab 3a we implemented our own version of the transport layer using protocols such as: stop-and-wait, sliding window, piggyback acknowledgement, selective arq, etc. This implementation took place at the user level; however, with the support of `ctcp_sys_internal.c/h`, we are able to interoperate with actual tcp socket connections, which takes place at the kernel level. Our program was able to use http to get webpages.

### Tools & Environment:

- Virtual Box: Custom CSCI 551 image, which supports tcp changes at the user level
  - We use this environment to test our code
- Visual Studio Code:
  - Text editor with various functionalities, providing ease of development
  - Ssh-remote: to avoid using VM's GUI
- GDB:
  - Allows us to debug our code
- Valgrind:
  - Allows us to detect memory leaks within our code

**Program Structure and Design:** To keep things simple, the whole implementation took place within `ctcp.c`. By isolating our changes to this specific file, we didn't have to change other source files or change the Makefile, which affects the way the executable is built.

### Design Goal & Implementations:

1. Avoid overwhelming the receiver
  - a. `Config.send_window` & `Config.recv_window` to know the initial window size of sender and receiver
  - b. Utilize a `current_window_size` variable to track how much buffer space we currently have
2. Reliable transmission
  - a. Utilize `checksum()` function & `segment->cksum` variable to check for packet corruption
  - b. Utilize `config.rt_timeout` to handle packet delay
3. In-order deliver
  - a. Cumulative Ack allows us to know which segments have already been acknowledged and outputted
  - b. Utilize the `curr_ackno` and `prev_ackno` variables to keep track of packet order

## Function implementations:

**ctcp\_init():** Here we initialize the member variables of `ctcp_state`. Linked list member variables are initialized via: “`ll_create()`,” `curr_seqno` & `curr_ackno` are initialized to 1, and everything else is initialized to 0.

**ctcp\_destroy(ctcp\_state \* state):** Here we clear memory for a specific connection state. Specifically, we traverse any `linked_list` members within the state, remove the nodes, `free()` the objects, and finally `free()` the list itself.

**ctcp\_read():** This function keeps using `conn_input`, until either 0 or -1 is detected.

- If `read_result > 0`, we package these input into segments
- If `read_result == -1`, we set our `EOF_FLAG` to true, meaning the host stops packaging & sending segments, then the sender sends a `TH_FIN` segment to the receiver to tell the receiver that it's ready to close down on its part

**ctcp\_receive():** This function processes segments with data as well as segments without data but has flags to control the communication between 2 hosts.

- We start out the function by doing various checks on the incoming segment & do not proceed in acknowledging the segments if:
  - If the `segment->len < length`: this indicates that the segment has been truncated
  - If the checksum doesn't match: the packet is corrupted
  - If the segment isn't in the receiving windows range
- If everything passes we send an acknowledgement back for the incoming segment

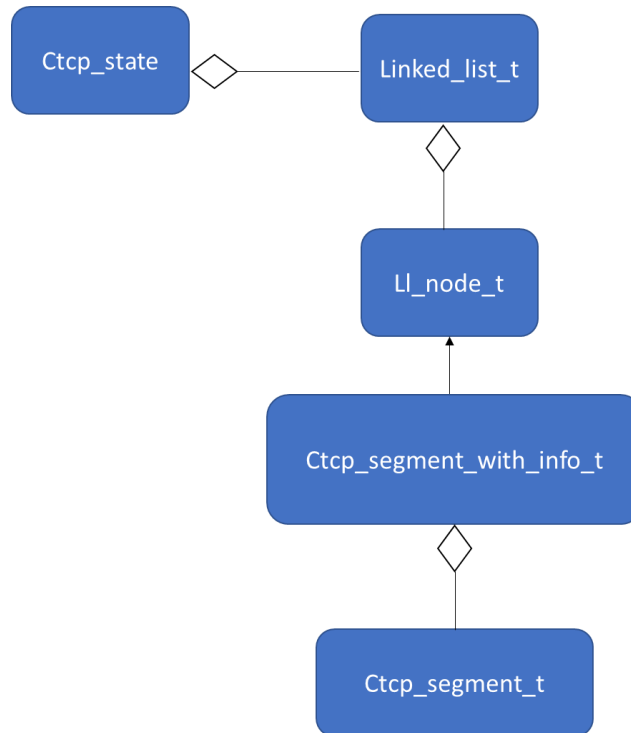
**ctcp\_output():** This function outputs segments that are acknowledged and ready to be outputted; however, it starts by checking if there is enough buffer space to output such a segment. We do not output segments that doesn't have data, such as `ACK` & `FIN` segments.

**ctcp\_timer(ctcp\_state \* state):** The function traverses all the `conn_states` & then traverses all the unacknowledged segments of the `curr_state`. The inner loop checks for timeouts & retransmission numbers, while the outer loop checks the entire `conn_state` if all conditions are met to be `ctcp_destroy(curr_state)`.

**Implementation Challenges:** One of the biggest challenges in implementing `ctcp` was defining the necessary data structures & adding the right fields into `ctcp_state_t` to support the lab. Changing the struct's member variables causes code refactor that propagates to different functions and sections of the code. For example, after defining my “`ctcp_segment_with_info_t`”

struct, I had to propagate the changes to `ctcp_state` class & the different functions that use it. The following picture shows the dependencies of classes.

### Complex Class Diagram & Dependencies:



I introduced a new class called “`ctcp_segment_with_info_t`,” which is a wrapper class over the `ctcp_segment_t` class. This wrapper added extra information to keep track of the number of retransmits & timestamp of each `ctcp_segment_t`. This design was necessary, since none of the variables in the original `ctcp_segment_t` struct keeps track of there information.

### Testing:

1. Basic TCP Functionality: For these test cases, I can get anywhere from 15/17-17/17 test cases correct. It’s unpredictable how the results will come out
2. Dumbbell topology: Unable to receive files over mini-net; however, we’re passing the valgrind test for memory leaks
3. Mini-Internet topology: Unable to cget over the I2 network topology

**Remaining bugs:** I believe my remaining bug is having the client receive file from the server, while the client doesn’t send anything (meaning it has an EOF from the very beginning & shouldn’t shut down)