UNIT 2: Cryptography and Symmetric encryption:

LIST OF TOPIC:

Cryptography, Cryptanalysis and Brute-Force Attack, Symmetric cipher model, Substitution techniques: Caesar cipher, Monoalphabetic ciphers, Polyalphabetic ciphers, One-Time Pad, Hill cipher, Playfair cipher, Transposition techniques: Rail fence, Columnar.

Cryptography

Definition

- -Cryptography is the science of using mathematics to encrypt and decrypt data.
- -Cryptography is the art and science of keeping messages secure.
- -The art and science of concealing the messages to introduce secrecy in information security is recognized as cryptography.

A message is **plaintext** (sometimes called cleartext). The process of disguising a message in such a way as to hide its substance is **encryption**. An encrypted message is ciphertext. The process of turning ciphertext back into plaintext is **decryption**.

A cipher (or cypher) is an algorithm for performing encryption or decryption—a series of well-defined steps that can be followed as a procedure

A cryptosystem is an implementation of cryptographic techniques and their accompanying infrastructure to provide information security services. A cryptosystem is also referred to as a cipher system. The various components of a basic cryptosystem are as follows —

☐ Plaintext ☐ Encryption Algorithm ☐ Ciphertext ☐ Decryption Algorithm ☐ Encryption Key
□ Decryption Key

While **cryptography** is the science of securing data, **cryptanalysis** is the science of analyzing and breaking secure communication. Classical cryptanalysis involves an interesting combination of analytical reasoning, application of mathematical tools, pattern finding, patience, determination, and luck. Cryptanalysts are also called attackers. **Cryptology** embraces both cryptography and cryptanalysis.

Cryptography

☐ An original message is known as the plaintext.
☐ The Coded message is called the ciphertext.
The Process of converting from plaintext to ciphertext is known as enciphering or encryption.
☐ Restoring the plaintext from the ciphetext is deciphering or decryption.

☐ The many schemes used for encryption constitute the area of study known as cryptography.
☐ Techniques used for deciphering a message without any knowledge of the enciphering details
is known as cryptanalysis. It also known as "Breaking the Code".
☐ The areas of cryptography and cryptanalysis together are called cryptology.
☐ A cryptanalyst develops mathematical methods and codes that protect data from computer
hackers. This involves the decryption of a cipher text into plain text in order to transmit a
message over insecure channels.

Symmetric cipher model

- ☐ Symmetric encryption is a form of cryptosystem in which encryption and decryption are performed using the same key.
- \Box It is also known as conventional encryption.
- \Box Symmetric encryption transforms plaintext into cipher text using a secret key and an encryption algorithm.
- \Box Using the same key and a decryption algorithm, the plaintext is recovered from the cipher text.
- ☐ A symmetric encryption scheme has five ingredients
- o **Plaintext:** This is the original intelligible message or data that is fed into the algorithm as input.
- o **Encryption algorithm:** The encryption algorithm performs various substitutions and transformations on the plaintext.
- o **Secret key:** The secret key is also input to the encryption algorithm. The key is a value independent of the plaintext and of the algorithm.
- o **Ciphertext:** This is the scrambled message produced as output. It depends on the plaintext and the secret key. The ciphertext is an apparently random stream of data and, as it stands, is unintelligible.
- o **Decryption algorithm:** This is essentially the encryption algorithm run in reverse. It takes the ciphertext and the secret key and produces the original plaintext.

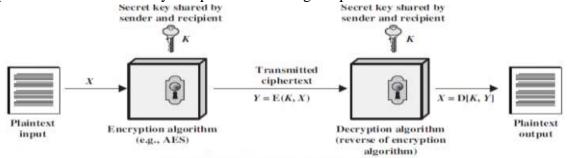
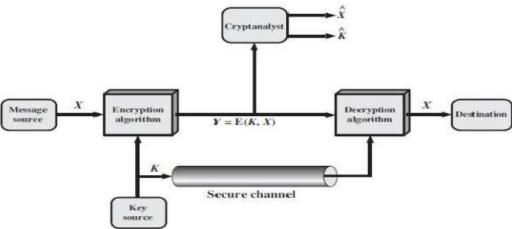


Fig: Simplified Model of Symmetric Encryption

Symmetric Cipher Model



A symmetric cipher model are broadly contains five parts. ☐ Plaintext: This is the original intelligible message. ☐ Encryption algorithm: The encryption algorithm performs various substitutions and transformations on the plaintext. It takes in plaintext and key and gives the cipher text. ☐ Secret key: The key is a value independent of the plaintext and of the algorithm. Different keys will yield different outputs. ☐ Ciphertext: This is the scrambled message produced as output. It depends on the plaintext and the secret key. ☐ Decryption algorithm: Runs on the cipher text and the key to produce the plaintext. This is essentially the encryption algorithm run in reverse. ☐ Two basic requirements of encryption are: 1. Encryption algorithm should be strong. An attacker knowing the algorithm and having any number of cipher text should not be able to decrypt the cipher text or guess the key. 2. The key shared by the sender and the receiver should be secret. \Box Let the plaintext be X = [X1, X2, ..., XM], key be K = [K1, K2, ..., KJ] and the cipher text produced be Y = [Y1, Y2,..., YN]. Then, we can write Y = E(K, X)☐ Here E represents the encryption algorithm and is a function of plaintext X and key K. \Box The receiver at the other ends decrypts the cipher text using the key. X = D(K, Y)☐ Here D represents the decryption algorithm and it inverts the transformations of encryption algorithm. ☐ An opponent not having access to X or K may attempt to recover K or X or both. ☐ It is assumed that the opponent knows the encryption (E) and decryption (D) algorithms. ☐ If the opponent is interested in only this particular message, then the focus of the effort is to recover by generating a plaintext estimate ^X.

☐ If the opponent is interested in being able to read future messages as well then he will attempt

to recover the key by making an estimate K .

☐ Cryptographic systems are characterized along three independent dimensions.
1. The types of operations used for transforming plaintext to ciphertext. All encryption
algorithms are based on two general principles substitution, and transposition. Basic requirement
is that no information be lost. Most systems referred to as product system, involves multiple
stages of substitutions and transpositions.
2. The number of keys used. If both sender and receiver use the same key, the system is referred
to as symmetric, single-key, secret-key, or conventional encryption. If the sender and receiver
use different keys the system is referred to as asymmetric, two-key, or public-key encryption.
3. The way in which the plaintext is processed. A block cipher process a block at a time and
produce an output block for each input block. A stream cipher process the input element
continuously, producing output one element at a time, as it goes along.
Cryptanalysis and Brute-Force Attack
☐ Cryptanalysis: Cryptanalytic attacks rely on the nature of the algorithm plus perhaps some
knowledge of the general characteristics of the plaintext or even some simple plaintext-
ciphertext pairs. This type of attack finds characteristics of the algorithm to find a specific
plaintext or to find key.
☐ Brute-force attack : The attacker tries every possible key on a piece of ciphertext until
plaintext is obtained. On average, half of all possible keys must be tried to achieve success.
$\ \square$ Based on the amount of information known to the cryptanalyst cryptanalytic attacks can be
categorized as:
o Cipher text Only Attack: The attacker knows only cipher text only. It is easiest to defend.
o Known plaintext Attack: In this type of attack, the opponent has some plaintext-cipher text
pairs. Or the analyst may know that certain plaintext patterns will appear in a message. For
example, there may be a standardized header or banner to an electronic funds transfer message
and the attacker can use that for generating plaintext-cipher text pairs.
o Chosen plaintext: If the analyst is able somehow to get the source system to insert into the
system a message chosen by the analyst, then a chosen-plaintext attack is possible. In such a
case, the analyst will pick patterns that can be expected to reveal the structure of the key.
o Chosen Cipher text: In this attack, the analyst has cipher text and some plaintext-cipher text
pairs where cipher text has been chosen by the analyst.
o Chosen Text: Here, the attacker has got cipher text, chosen plaintext-cipher text pairs and
chosen cipher text-plaintext pairs.

 $\hfill\Box$ Chosen cipher text and chosen text attacks are rarely used.

 $\hfill \square$ It is assumed that the attacker knows the encryption and decryption algorithms.

☐ Generally, an encryption algorithm is designed to withstand a known-plaintext attack.

Substitution Techniques

It is one in which the letters of plaintext are replaced by other letters or by numbers or symbols.

Caesar cipher

 \Box The encryption rule is simple; replace each letter of the alphabet with the letter standing 3 places further down the alphabet.

 \Box The alphabet is wrapped around so that Z follows A.

☐ Generally Plain text is in lower case and Cipher text is Upper Case.

☐ Example:

Plaintext: meet me after the party

Ciphertext: PHHW PH DIWHU WKH SDUWB

а	b	c	d	е	f	g	h	1	1	k	1.	m
0	1	2	3	4	5	6	7	8	9	10	11	12
n:	0	р	q	r	S	t	u	V	w	x	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

Here, the key is 3. If different key is used, different substitution will be obtained.

☐ Mathematically, starting from a=0, b=1 and so on, Caesar cipher can be written as:

$$E(p) = (p + k) \bmod (26)$$

 $D(C) = (C - k) \bmod (26)$

Encryption k=3 $E(p) = (p + k) \mod (26)$	Result	Result Cipher Text $D(C) = (C - k) \mod (26)$				
M=E(M)=(12+3)mod26	15=P	D(P)=(15-3)mod26	12=m			
E=E(E)=((4+3)mod26	7=H	D(H)=(7-3)mod26	4=e			
E= E(E)=((4+3)mod26	7=H	D(H)=(7-3)mod26	4=e			
T=E(T)=((19+3)mod26	21=V	D(V)=(21-3)mod26	19=t			

This cipher can be broken

- o If we know one plaintext-cipher text pair since the difference will be same.
- o By applying Brute Force attack as there are only 26 possible keys.

Monoalphabetic Substitution Cipher

VIDEO LINK

https://www.youtube.com/watch?app=desktop&v=ZJWKpviXPCo&ab channel=OnlineTeacher

Instead of shifting alphabets by fixed amount as in Caesar cipher, any random permutation i
assigned to the alphabets. This type of encryption is called monoalphabetic substitution cipher.
$\hfill \Box$ For example, A is replaced by Q, B by D, C by T etc. then it will be comparatively stronger
than Caesar cipher.
☐ The number of alternative keys possible now becomes 26!.
☐ Thus, Brute Force attack is impractical in this case.
☐ However, another attack is possible. Human languages are redundant i.e. certain characters
are used more frequently than others. This fact can be exploited.
☐ In English 'e' is the most common letter followed by 't', 'r', 'n', 'o', 'a' etc. Letters like 'q',
'x', 'j' are less frequently used.
☐ Moreover, digrams like 'th' and trigrams like 'the' are also more frequent.

☐ Tables of frequency of these letters exist. These can be used to guess the plaintext if the
plaintext is in uncompressed English language.
☐ The most common two letter combinations are called as digrams. e.g. th, in, er, re and an.
☐ The most common three letter combinations are called as trigrams. e.g. the, ing, and, and ion.

Playfair Cipher

- \Box In this technique multiple (2) letters are encrypted at a time.
- \Box This technique uses a 5 X 5 matrix which is also called key matrix.

M	0	N	Α	R
С	Н	Υ	В	D
E	F	G	ι/J	К
L	Р	Q	S	Т
U	٧	w	х	Z

The plaintext is encrypted two letters at a time:

- o Break the plaintext into pairs of two consecutive letters.
- o If a pair is a repeated letter, insert a filler like 'X'in the plaintext, eg. "Balloon" is treated as "ba lx lo on".
- o If both letters fall in the same row of the key matrix, replace each with the letter to its right (wrapping back to start from end), eg. "AR" encrypts as "RM".
- o If both letters fall in the same column, replace each with the letter below it (again wrapping to top from bottom), eg. "MU" encrypts to "CM".
- o Otherwise each letter is replaced by the one in its row in the column of the other letter of the pair, eg. "HS" encrypts to "BP", and "EA" to "IM" or "JM" (as desired)
- \Box Security is much improved over monoalphabetic as here two letters are encrypted at a time and hence there are 26 X 26 =676 diagrams and hence it needs a 676 entry frequency table.

However, it can be broken even if a few hundred letters are known as much of plaintext structure is retained in cipher text.

Example 2: PlainText: "instruments" keyword: monarchy

After Split: 'in' 'st' 'ru' 'me' 'nt' 'sz'

cipher text : ga tl mz cl rq tx

For both encryption and decryption, the same key is to be used.

in:	M	0	N	A	R	st:	M	0	N.	A	R	ru:	M	0	N	A	R
	C	H	Y	В	D		C	н	Y	В	D		C	н	Y	В	D
	E	F	G	1	К		E	F	G	1	K.		E	F	G	1	K
	L	P	Q	s	T		L	P	Q	S	T		L	P	Q	S	T
	U	V	W	×	Z		U	V	W	X	Z		U	V	W	×	Z
me:	M	0	N	A	R	nt:	M	0	N	A	R	SZ;	M	0	N	Α	R
	C	H	Y	В	D		C	н	Y	В	D		C	Н	Y	В	D
	E	F	G	1	K		E	F	G	1	K		E	F	G	1	K
	L	P	Q	S	T		L	P	Q	S	T		L	Р	Q	s	T
	U	V	W	X	Z		U	V	W	X	Z		U	V	W	х	Z

video link: https://www.youtube.com/watch?v=quKhvu2tPy8

Strength of playfair cipher Playfair cipher is a great advance over simple mono alphabetic ciphers. Since there are 26 letters, 26x26 = 676 diagrams are possible, so identification of individual diagram is more difficult.

Hill Cipher

source link

https://www.youtube.com/watch?v = sYdq6-kquKc

https://www.educative.io/edpresso/what-is-the-hill-cipher

 \Box This cipher is based on linear algebra.

 \Box Each letter is represented by numbers from 0 to 25 and calculations are done modulo 26.

 \Box This encryption algorithm takes m successive plaintext letters and substitutes them with m cipher text

letters.

 \Box The substitution is determined by m linear equations. For m = 3, the system can be described as:

 $c1 = (k11p1 + k21p2 + k31p3) \mod 26$

 $c2 = (k12p1 + k22p2 + k32p3) \mod 26$

 $c3 = (k13p1 + k23p2 + k33p3) \mod 26$

☐ This can also be expressed in terms of row vectors and matrices.

k11 k12 k13

 $(c1\ c2\ c3) = (p1\ p2\ p3)\ (k21\ k22\ k23)\ mod\ 26$

k31 k32 k33

Where C and P are row vectors of length 3 representing the plaintext and cipher text, and K is a 3 X 3

matrix representing the encryption key

☐ Key is an invertible matrix K modulo 26, of size m. For example:

17 17 5 4 19 15

 $K = (21\ 18\ 21)\ K - 1 = (15\ 17\ 6)$

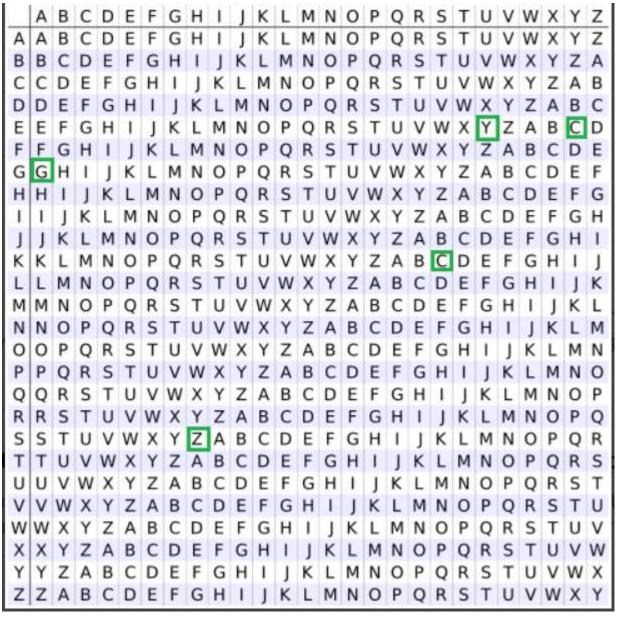
2 2 19 24 0 17

☐ Encryption and decryption can be given by the following formulae:
Encryption: $C = PK \mod 26$
Decryption: $P = CK-1 \mod 26$
The strength of the Hill cipher is that it completely hides single-letter frequencies.
☐ Although the Hill cipher is strong against a cipher text-only attack, it is easily broken with a
known plaintext attack.
o Collect m pair of plaintext-cipher text, where m is the size of the key.
o Write the m plaintexts as the rows of a square matrix P of size m.
o Write the m cipher texts as the rows of a square matrix C of size m.
o We have that C=PK mod 26.
o If P is invertible, then K=P-1C mod 26,
o If P is not invertible, then collect more plaintext-cipher text pairs until an invertible P is obtained.
Polyalphabetic ciphers
The Vigenère cipher
\Box This is a type of polyalphabetic substitution cipher (includes multiple substitutions depending on the
key). In this type of cipher, the key determines which particular substitution is to be used.
☐ To encrypt a message, a key is needed that is as long as the message. Usually, the key is a repeating keyword.
☐ For example, if the keyword is deceptive, the message "we are discovered save yourself" is encrypted as

Key: deceptivedecept Plaintext: wearediscovered

follows:

Ciphertext: ZICVTWQNGRZGVTW



☐ Encryption can be done by looking in the Vigenere Table where ciphertext is the letter key's row and plaintext's column or by the following formula:

 $Ci = (Pi + Ki \mod m) \mod 26$

- \Box Decryption is equally simple. The key letter again identifies the row. The position of the cipher text letter in that row determines the column, and the plaintext letter is at the top of that column.
- \Box The strength of this cipher is that there are multiple ciphertext letters for each plaintext letter, one for each unique letter of the keyword.
- \Box Thus, the letter frequency information is obscured however, not all knowledge of the plaintext structure is lost.

Vernam Cipher

- ☐ This system works on binary data (bits) rather than letters.
- ☐ The technique can be expressed as follows:

 $Ci = Pi \oplus Ki$

Where

Pi = ith binary digit of plaintext.
Ki = ith binary digit of key.
Ci = ith binary digit of ciphertext.
\bigoplus = exclusive-or (XOR) operation
☐ Thus, the ciphertext is generated by performing the bitwise XOR of the plaintext and the key.
☐ Decryption simply involves the same bitwise operation:
$Pi = Ci \oplus Ki$
☐ The essence of this technique is the means of construction of the key.
☐ It was produced by the use of a running loop of tape that eventually repeated the key, so that
in fact the system worked with a very long but repeating keyword.
Although such a scheme has cryptanalytic difficulties, but it can be broken with a very long
ciphertext or known plaintext as the key is repeated.
One-Time Pad
☐ In this scheme, a random key that is as long as the message is used.
☐ The key is used to encrypt and decrypt a single message, and then is discarded. Each new
message requires a new key of the same length as the new message.
☐ This scheme is unbreakable.
☐ It produces random output that bears no statistical relationship to the plaintext.
☐ Because the ciphertext contains no information whatsoever about the plaintext, there is simply
no way to break the code.
☐ For any plaintext of equal length to the ciphertext, there is a key that produces that plaintext.
☐ Therefore, if you did an exhaustive search of all possible keys, you would e plaintexts, with
no way of knowing which the intended plaintext was.
☐ Therefore, the code is unbreakable.
☐ The security of the one-time pad is entirely due to the randomness of the key.
☐ The one-time pad offers complete security but, in practice, has two fundamental difficulties:
o There is the practical problem of making large quantities of random keys. Any heavily used
system might require millions of random characters on a regular basis. Supplying truly random
characters in this volume is a significant task.
o Another problem is that of key distribution and protection. For every message to be sent, a key
of equal length is needed by both sender and receiver.
☐ Because of these difficulties, the one-time pad is used where very high security is required.
☐ The one-time pad is the only cryptosystem that exhibits perfect secrecy.
Example with Alphabetic Characters
Let's use a simple example with alphabetic characters where A=0, B=1,, Z=25:
Plaintext: HELLO
Key: XMCKL (randomly generated) T
To encrypt:
1. Convert plaintext and key to numerical values:
○ H=7. F=4. L=11. L=11. O=14

X=23, M=12, C=2, K=10, L=11Add corresponding characters modulo 26:

 \circ (7 + 23) % 26 = 30 % 26 = 4 (E)

 \circ (14 + 11) % 26 = 25 (Z) **Ciphertext:** EQNVZ

To decrypt:

- 1. Convert ciphertext and key to numerical values:
 - o E=4, Q=16, N=13, V=21, Z=25
 - X=23, M=12, C=2, K=10, L=11
- 2. Subtract corresponding characters modulo 26:
 - \circ (4 23) % 26 = -19 % 26 = 7 (H)
 - \circ (16 12) % 26 = 4 (E)
 - o (13 2) % 26 = 11 (L)
 - o (21 10) % 26 = 11 (L)
 - o (25 11) % 26 = 14 (O)
- **Decrypted Plaintext**: HELLO

Transposition Techniques

☐ A very different kind of mapping is achieved by performing some sort of permutation on the	ne
plaintext letters. This technique is referred to as a transposition cipher.	

 \Box The simplest such cipher is the rail fence technique.

Rail Fence Technique

☐ Encryption invo	olves writing plainter	xt letters diagonally	over a number	of rows, th	en read off
cipher row					

by row.

☐ For example, the text "meet me after the party" can be written (in 2 rows) as:

mematrhpry

etefeteoat

☐ Ciphertext is read from the above row-by-row:

MEMATRHPRYETEFETEAT

 \Box This scheme is very easy to cryptanalyze as no key is involved.

☐ Transposition cipher can be made significantly more secure by performing more than one stage of transposition. The result is a more complex permutation that is not easily reconstructed.

Columnar Transposition

Put plaintext into rows of matrix then read ciphertext out of columns. The simplest transposition cipher is the columnar transposition. This comes in two forms, the complete columnar transposition and the incomplete columnar. In both of these systems, the plain text is written horizontally in a rectangle that is as wide as the length of the key.

Example

suppose matrix is 3 x 4
Plaintext: SEETHELIGHT
Ciphertext: SHGEEHELTTIX

Same effect as Scytale. What is the key? Except the transposition of letters based

on 3 x 4 matrix no key is used.

Keyword Columnar Transposition

In Columnar transposition plain text can be padded based on key either in regular or irregular method. In regular way the plain text has been padded so that it nearly fits the matrix/rectangle. But in the case of irregular transposition leaves these characters blank, though this makes decryption slightly more difficult. The columns are now reordered such that the letters in the key word are ordered alphabetically.

Example

Plaintext: CRYPTOISFUN

Matrix 3 x 4 and keyword MATH Ciphertext: ROUPSXCTFYIN

How can Trudy cryptanalyze this cipher?

• Consider the ciphertext

VOESA IVENE MRTNL EANGE WTNIM HTMLL ADLTR NISHO

- Matrix is n x m for some n and m
- Since 45 letters, $n \square m = 45$
- The ciphertext is

VOESA IVENE MRTNL EANGE WTNIM HTMLL ADLTR NISHO DWOEH

• If encryption matrix was 9 x 5, then

Assignment 2

- 1. Explain the concept of a symmetric cipher model. How does it differ from asymmetric cryptography?
- 2. Drow and explain symmetric cipher model.
- 3. Explain cryptanalysis and brute force attack.
- 4. Explain the Caesar cipher with an example. Discuss its vulnerability to brute-force attacks.
- 5. use a key of 3 to encrypt the message "HELLO WORLD" using Caesar cipher.
- 6. Explain the concept of a monoalphabetic with example
- 7. Explain the concept of a polyalphabetic ciphers with examples.
- 8. Using polyalphabetic ciphers encrypt the message **Plaintext**: ATTACKATDAWN **Keyword**: LEMON
- 9. Using polyalphabetic ciphers encrypt the message **Plaintext**: HIDE THE GOLD **Keyword**: MONARCHY
- 10. Explain the concept of a One-Time Pad with example.
- 11. encrypt the plaintext "HELLO" using a One-Time Pad **Plaintext**: HELLO **Random Key**: XMCKL

- 12. Describe the Hill cipher. How does it use matrix multiplication for encryption?
- 13. Discuss the Playfair cipher. Provide an example of its operation.
- 14. Explain the Rail Fence transposition technique. Provide an example of encrypting a message using a rail fence with 3 rows.
- 15. encrypt the message "MEET ME AFTER THE PARTY" using 3 rails.
- 16. Describe the Columnar transposition technique. How does it reorder plaintext letters for encryption?
- 17. encrypt the message "MEET ME AFTER THE PARTY" using the keyword "SECRET" Columnar transposition technique.
- 18. Define cryptography and cryptanalysis. How do they differ in their goals?
- 19. What is a brute-force attack in the context of cryptography?
- 20. Briefly explain how a monoalphabetic cipher encrypts plaintext.
- 21. What makes the One-Time Pad cipher unique in terms of its security properties?
- 22. Explain columnar transposition technic.
- 23. How does the Playfair cipher handle repeated letters in plaintext?
- 24. Give an example of a transposition cipher technique. What is its basic principle?
- 25. Define following term: Cryptography, plain text, Encryption, Decryption, ciphertext, Cipher, Cryptosystem, Cryptanalysis,