

PHASE RETRIEVAL IN HIGH DIMENSIONS : STATISTICAL AND COMPUTATIONAL PHASE TRANSITIONS

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PHASE RETRIEVAL AS A GLM

Generalized Linear Model (GLM)

Observations $Y_\mu \in \mathbb{R}$

$$Y_\mu \sim P_{\text{out}} \left(\cdot \mid \frac{1}{\sqrt{n}} \sum_{i=1}^n \Phi_{\mu i} X_i^* \right) \quad \mu \in \{1, \dots, m\}$$

(Probabilistic) channel

Sensing matrix (real/complex)

Signal (real/complex), n-dimensional

In **phase retrieval**, one only measures the modulus: $P_{\text{out}}(\cdot | z) = P_{\text{out}}(\cdot | |z|)$

Classical problem, non-trivial even in the noiseless case $Y_\mu = |(\Phi \mathbf{X}^*)_\mu|^2 / n$, many algorithms:

- SDP relaxations [Candès&al '11, Candès&al '12, Waldspurger&al '12, Goldstein&al '16, ...]
- Non-convex optimization procedures [Netrapalli&al '13, Candès&al '14, Gerchberg 1972, ...]
- Spectral methods [Mondelli&al '18, Luo&al '18, Dudeja&al '19, ...]

Goal: Fundamental limits of phase retrieval with **random** sensing matrices and **random** signal in the **typical** case and in high dimensions.



Different from the injectivity studies of the "worst-case" [Bandeira&al '13]

$$Y_\mu \sim P_{\text{out}} \left(\cdot \middle| \frac{1}{\sqrt{n}} \sum_{i=1}^n \Phi_{\mu i} X_i^* \right) \quad \mu \in \{1, \dots, m\}$$

In the limit $m, n \rightarrow \infty$ with $\alpha = m/n = \Theta(1)$, what is the smallest α needed to recover \mathbf{X}^* ...

- Better than a random guess ?
- Perfectly ? (up to the possible rank deficiency of Φ)
- With which (polynomial-time) algorithm ?

Our model: i) \mathbf{X}^* and Φ can be **real** ($\beta = 1$) or **complex** ($\beta = 2$)

ii) The signal \mathbf{X}^* is generated using a (known) i.i.d. prior distribution P_0 and $\text{Var}_{P_0}(X^*) = \rho > 0$

iii) The matrix Φ is **right-orthogonally (unitarily) invariant**: $\forall \mathbf{U}, \Phi \stackrel{d}{=} \Phi \mathbf{U}$

iv) The empirical spectral distribution of $\Phi^\dagger \Phi / n$ converges: $\nu_n \equiv \frac{1}{n} \sum_{i=1}^n \delta_{\lambda_i \left(\frac{\Phi^\dagger \Phi}{n} \right)} \xrightarrow{n \rightarrow \infty} \nu \in \mathcal{M}_1^+(\mathbb{R}_+).$

Encompasses many models: **Gaussians, product of Gaussians, random column-orthogonal/unitary, any $\Phi \equiv \mathbf{U} \mathbf{S} \mathbf{V}^\dagger$ with $S_i^2 \stackrel{\text{i.i.d.}}{\sim} \nu$**

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OPTIMAL ERROR IN GLMS

We consider any channel (not necessarily phase retrieval)

Conjecture: Consider the following scalar optimization problem $f = \sup_{q_x, q_z} [\underbrace{I_0^{(\beta)}(q_x)}_{P_0} + \underbrace{I_{\text{out}}^{(\beta)}(q_z)}_{P_{\text{out}}} + \underbrace{\beta I_{\text{int}}(q_x, q_z)}_{\nu}]$

Then the *information-theoretic Minimal Mean Squared Error* is: $\text{MMSE} = \rho - q_x$

(The functions involved in the optimization problem are fully explicit)

Theorem: If either

a) $\Phi_{\mu i} \stackrel{\text{i.i.d.}}{\sim} \mathcal{N}_{\beta}(0, 1)$ (standard Gaussian distribution)

b) P_0 is Gaussian and $\Phi = \text{WB}$

Gaussian matrix

Any matrix

} , the conjecture above stands.

- Conjecture obtained with the replica method of statistical physics [Parisi&al 1987, Takahashi&al '20]
- Proven using probabilistic interpolation methods [Guerra '03, Talagrand '07, Barbier&al '18, Barbier&al '19]

2'

OPTIMAL ERROR IN GLMs

We consider any channel (not necessarily phase retrieval)

$$f = \sup_{q_x, q_z} [I_0^{(\beta)}(q_x) + I_{\text{out}}^{(\beta)}(q_z) + \beta I_{\text{int}}(q_x, q_z)]$$

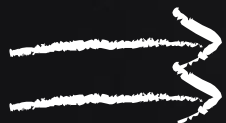
“Replica-symmetric” potential $f(q_x, q_z)$

Strong conjecture: For GLMs, the optimal polynomial-time algorithm is an explicit iterative algorithm:

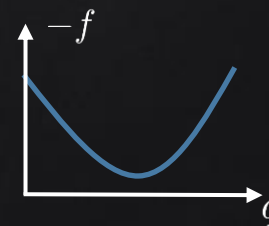
Approximate Message Passing. Called here **G-VAMP** (*Generalized Vector Approximate Message Passing*).

[Mézard '89, Donoho&al '09, Montanari&al '10, Krzakala&al '11, Rangan&al '16, Schniter&al '16, ...]

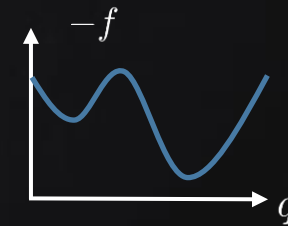
Important result [Schniter&al '16]: The MSE of G-VAMP in the large n limit is given by running *gradient ascent* on the Replica-symmetric potential starting from $q_x = q_z = 0$ (random initialization).



We can investigate “computational-to-statistical” gaps by studying the landscape of $f(q_x, q_z)$!



No gap



Gap : “Hard” phase


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(ALGORITHMIC) WEAK RECOVERY

We consider phase retrieval: $P_{\text{out}}(\cdot|z) = P_{\text{out}}(\cdot||z|)$

What is the minimal number of measurements $\alpha = m/n$ necessary to beat a random guess in polynomial time?

This threshold $\alpha_{\text{WR,Algo}}$ is a solution of:

- This is an implicit equation 
- Derived by a stability analysis of the replica-symmetric potential.

$$\alpha = \frac{\langle \lambda \rangle_\nu^2}{\langle \lambda^2 \rangle_\nu} \left[1 + \left\{ \int_{\mathbb{R}} dy \frac{\left(\int_{\mathbb{K}} \mathcal{D}_\beta z (|z|^2 - 1) P_{\text{out}} \left[y \left| \sqrt{\frac{\rho \langle \lambda \rangle_\nu}{\alpha}} z \right| \right] \right)^2}{\int_{\mathbb{K}} \mathcal{D}_\beta z P_{\text{out}} \left[y \left| \sqrt{\frac{\rho \langle \lambda \rangle_\nu}{\alpha}} z \right| \right]} \right\}^{-1} \right]$$

For any phase retrieval channel and prior, the highest weak recovery threshold is reached by random column-orthogonal/unitary matrices (up to a scaling).

For noiseless phase retrieval:

$$\alpha = \left(1 + \frac{\beta}{2} \right) \frac{\langle \lambda \rangle_\nu^2}{\langle \lambda^2 \rangle_\nu}$$

- Gaussian matrices: $\alpha_{\text{WR,Algo}} = \frac{\beta}{2}$ [Barbier&al '18, Mondelli &al '18]
- Random column-orthogonal/unitary matrices: $\alpha_{\text{WR,Algo}} = 1 + \frac{\beta}{2}$ [Dudeja&al '19] for $\beta = 2$

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STRONG (FULL) RECOVERY

We consider noiseless phase retrieval: $P_{\text{out}}(y|z) = \delta(y - |z|^2)$ and a Gaussian prior $P_0 = \mathcal{N}_\beta(0, 1)$

How many measurements are necessary to be able to information-theoretically achieve the best possible recovery?

If $\frac{1}{n} \text{rk} \left(\frac{\Phi^\dagger \Phi}{n} \right) \rightarrow r \in [0, 1] :$ $\alpha_{\text{FR,IT}} = \beta r$ Does not depend on the precise statistics of Φ

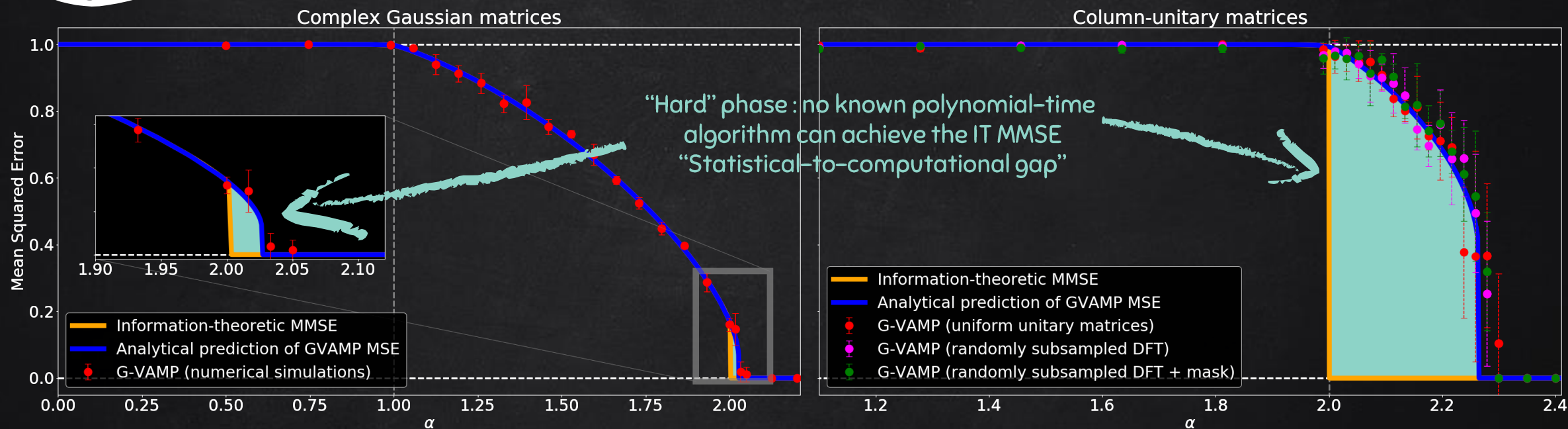
For $\alpha \geq \alpha_{\text{FR,IT}}$, the MMSE reaches a plateau with $\text{MMSE}_{\mathbf{x}} = 1 - r$; $\text{MMSE}_{\Phi \mathbf{x}} = 0$

- The real case $\alpha_{\text{FR,IT}} = r$ can be derived by a counting argument [Candès&al, '05]
- The complex case $\alpha_{\text{FR,IT}} = 2r$ can (as far as we know) only be derived by the replica-symmetric potential!

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APPLICATIONS

We consider noiseless phase retrieval: $P_{\text{out}}(y|z) = \delta(y - |z|^2)$ and a Gaussian prior $P_0 = \mathcal{N}_\beta(0, 1)$



Very good agreement of G-VAMP with the analytical predictions.

Some (funny) remarks:

- For column-unitary matrices $\alpha_{\text{FR,IT}} = \alpha_{\text{WR,Algo}} = 2$: “all-or-nothing” IT transition.
- For all other full-rank complex matrices $\alpha_{\text{WR,Algo}} < \alpha_{\text{FR,IT}}$
- For real matrices, there can be a large gap! Ex: column-orthogonal matrices $\alpha_{\text{FR,IT}} = 1 < \alpha_{\text{WR,Algo}} = 3/2$
- Matrices with **controlled structure** (e.g. randomly subsampled DFT) still perform very well with G-VAMP!

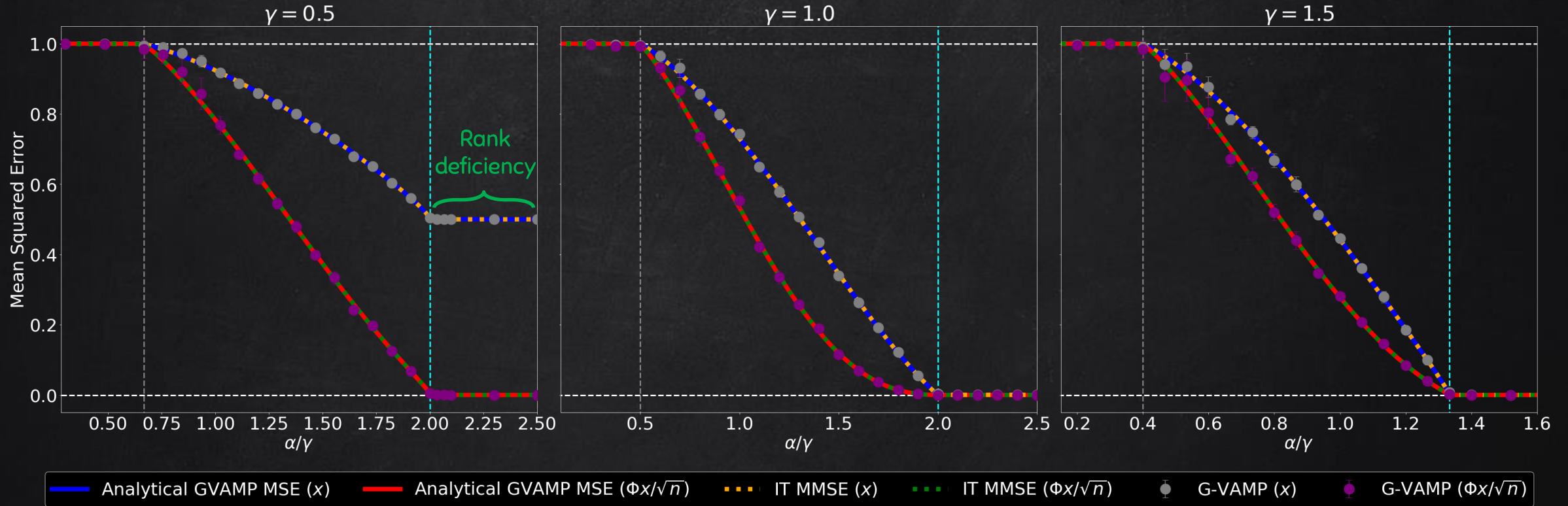
CONCLUSION / SUMMARY (NEW RESULTS IN RED)

	Matrix ensemble and value of β	$\alpha_{\text{WR, Algo}}$	$\alpha_{\text{FR, IT}}$	$\alpha_{\text{FR, Algo}}$
Noiseless phase retrieval with Gaussian prior	Real Gaussian Φ ($\beta = 1$)	0.5	1	$\simeq 1.12$
	Complex Gaussian Φ ($\beta = 2$)	1	2	$\simeq 2.027$
	Real column-orthogonal Φ ($\beta = 1$)	1.5	1	$\simeq 1.584$
	Complex column-unitary Φ ($\beta = 2$)	2	2	$\simeq 2.265$
	$\Phi = \mathbf{W}_1 \mathbf{W}_2$ ($\beta = 1$, aspect ratio γ)	$\gamma/(2(1 + \gamma))$	$\min(1, \gamma)$	Theorem
	$\Phi = \mathbf{W}_1 \mathbf{W}_2$ ($\beta = 2$, aspect ratio γ)	$\gamma/(1 + \gamma)$	$\min(2, 2\gamma)$	Theorem
Generic phase retrieval with any prior	Φ , $\beta \in \{1, 2\}$, $\text{rk}[\Phi^\dagger \Phi]/n = r$	Analytical expression	βr	Conjecture
	Gauss. Φ , $\beta \in \{1, 2\}$, symm. P_0, P_{out}	Analytical expression	Theorem	Theorem
	$\Phi = \mathbf{W} \mathbf{B}$, $\beta \in \{1, 2\}$, Gauss. P_0 , symm. P_{out}	Analytical expression	Theorem	Theorem
	Φ , $\beta \in \{1, 2\}$, symm. P_0, P_{out}	Analytical expression	Conjecture	Conjecture

THANK YOU !

Many numerical simulations were performed using the open-source TrAMP package [\[Baker&al, '20\]](#)

Product of two complex Gaussian matrices $\Phi = W_1 W_2$, with $W_1 \in \mathbb{C}^{m \times p}$, $W_2 \in \mathbb{C}^{p \times n}$ and $\gamma = p/n$



- Very good agreement of G-VAMP with the analytical predictions.
- We recover the two thresholds $\alpha_{\text{WR, Algo}} = \gamma/(1 + \gamma)$ and $\alpha_{\text{FR, IT}} = \min(2, 2\gamma)$
- (Very small) computational-to-statistical gap $\alpha_{\text{FR, Algo}} > \alpha_{\text{FR, IT}}$ for $\gamma \neq 1$

OPTIMAL ALGORITHMS (REAL CASE)

