# Composite Proximal Bundle Method

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#### Abstract

We consider minimization of nonsmooth functions which can be represented as the composition of a positively homogeneous convex function and a smooth mapping. This is a sufficiently rich class that includes max-functions, largest eigenvalue functions, and norm-1 regularized functions. The bundle method uses an oracle that is able to compute separately the function and subgradient information for the convex function, and the function and derivatives for the smooth mapping. With this information, it is possible to solve approximately certain proximal linearized subproblems in which the smooth mapping is replaced by its Taylor-series linearization around the current serious step. Our numerical results show the good performance of the Composite Bundle method for a large class of problems.

## 1 Introduction and motivation

For some years already, nonsmooth optimization research has focused on exploiting structure in the objective function as a way to speed up numerical methods. Indeed, for convex optimization, complexity results establish that oracle based methods have a linear rate of convergence; [NY83]. The  $\mathcal{U}$ -Lagrangian theory [LS97], [LOS00], and the  $\mathcal{V}\mathcal{U}$ -space decomposition [MS00], can be seen as tools to "extract" smooth structure from general convex functions. The approach was extended to a class of nonconvex functions in [MS04]. Similar ideas for more general nonconvex functions, having a smooth representative on a certain manifold corresponding to the  $\mathcal{U}$ -subspace, are explored in [Lew02], [Har03]. Another line of work concentrates efforts on identifying classes of functions structured enough to have some type of second-order developments, such as the primal-dual gradient structured functions from [MS00], [MS03], or the composite functions in [Sha03]. Composite functions, that were first considered in [Fle87, Ch. 14], and studied in [BF95], [LW02], have been more recently revisited in [Nes07] and [LW08], in the convex and nonconvex settings, respectively.

Most of the work above is conceptual, in the sense that if some algorithmic framework is considered, often it is not implementable. Among the exceptions, we find the convex optimization algorithms in [Ous00], [MS05], [Nes07], and the  $\mathcal{V}$ -space identification method in [DSS09]. In this paper, we develop an implementable bundle method that exploits structure for certain nonconvex composite optimization problems. More precisely, given a smooth mapping  $c: \Re^n \to \Re^m$  and a positively homogeneous (of degree 1) convex function  $h: \Re^m \to \Re$ , we consider how to solve the unconstrained problem

$$\min_{x \in \Re^n} (h \circ c)(x) \,. \tag{1}$$

In this composite objective, we refer to c as the *inner* mapping and to h as the *outer* function. Note in particular that the outer function is real-valued on  $\Re^m$ , an assumption that excludes indicator functions, but still covers a rich enough family of interesting problems; see § 2.

For solving the possibly nonconvex problem (1) we develop a specialized bundle method that takes full advantage of the composite structure of the objective function. Specifically, rather than using composite data:

 $(bb_{hoc})$ : Given  $x \in \Re^n$ , a composite black-box computes  $(h \circ c)(x)$  and  $\gamma \in \partial(h \circ c)(x)$ ,

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we suppose the oracle has the ability to make separate computations. More precisely,

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 \begin{cases} \text{ } (\mathtt{bb_c}): & \text{Given } x \in \Re^n \\ \text{ } (\mathtt{bb_h}): & \text{Given } C \in \Re^m \end{cases} \text{ an inner black-box computes } c(x) \text{ and its Jacobian } Dc(x)
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While  $(bb_{hoc})$  would build a classical cutting-planes model for the composite objective, having the richer oracle  $(bb_c) - (bb_h)$  at hand, we can make a better approximation and build a cutting-planes model for the so-called *conceptual model*. This model, introduced in [LW08] for general composite functions, and given by (4) below, defines certain proximal linearized subproblems in which the smooth mapping is replaced by its Taylor-series linearization. The corresponding algorithmic framework, named ProxDescent in [LW08], is rather broad (outer functions can be extended-valued and prox-regular), but relies heavily on the exact computation of proximal points at each iteration. In this work, we provide a fully implementable variant of such algorithmic framework for a less general, but still large enough, class of outer functions (real-valued, positively homogeneous, and convex).

In order to deal with approximate proximal computations, our fully implementable variant, given by Algorithm 1 below, incorporates typical features from bundle methods, such as *serious* and *null* steps iterations. But this is not the only modification of ProxDescent: there is an additional level of approximation, introduced by the need to estimate the conceptual model, which is handled in the implementable algorithm by backtracking. For these reasons, our convergence analysis deviates significantly from that of [LW08]. In addition, instead of a standard proximal term as in ProxDescent, in our development we consider a variable prox-metric that opens the way to exploiting second-order information of the smooth mapping, when available.

An alternative line of work, based on somewhat different richer oracles (that can compute more than one subgradient at any given point), refers to the proximity control bundle method [NPR08]. These methods, devised for nonconvex optimization problems appearing in eigenvalue optimization and automatic control, use certain "first-order local models" that contain the conceptual model from [LW08] as a special case.

Our paper is organized as follows. In Sec. 2 we give several classes of functions with composite structure. The conceptual model and the proposed cutting-planes function, that makes the conceptual model computationally tractable, are described in Sec. 3. Sections 4 and 5 contain, respectively, the Composite Bundle method and its convergence results. In Sec. 6 we report our numerical experience, on many of the functions described in Sec. 2. Our results show the good performance of the composite algorithm, when compared to several solvers, on a large number of problems of various dimensions. In the final section with concluding remarks, we compare the main features of our composite method with the proximity control bundle method from [NPR08].

# 2 Composite functions in the family

Our class of composite functions includes many examples from [LW08], [Nes07], [Sha03]. We mention in passing that the wording "composite" means slightly different things for different authors. Specifically, while we refer to a composition that can be nonconvex (as in [LW08], [Sha03]), composite functions in [Nes07] are the convex sum of two terms, one convex and another one smooth.

We suppose the outer function is positively homogeneous and convex, so it is also sublinear. In [HUL01, Lesson C] there are many examples of real-valued sublinear functions, such as norms, quadratic seminorms, gauges of closed convex sets containing 0 in their interior, infimal convolutions of sublinear functions, and support functions of bounded sets. The composition of any of such (outer) functions with any smooth (inner) mapping defines a function in our family. We review below some interesting special cases.

**Example 1 (max-functions)** The function given by the pointwise maximum of a finite collection of  $C^2$ -functions in  $\Re^n$ ,  $\{c_1(\cdot), \ldots, c_m(\cdot)\}$ , can be defined by composing the outer function  $h(C) = \max(C_1, \ldots, C_m)$  with the inner mapping with components  $c_j(\cdot)$ , for  $j = 1, \ldots, m$ .

For any  $x \in \Re^n$ , the Jacobian mapping Dc(x) is the  $m \times n$  matrix with rows given by the transposed gradients  $\nabla c_j(x)^{\top}$ . As for the second-order derivative,  $D^2c(x)$ , for any  $d, \tilde{d} \in \Re^n$ ,  $[D^2c(x)\tilde{d}]d$  is a vector in  $\Re^m$  with  $j^{th}$  component given by  $\tilde{d}^{\top}\nabla^2c_j(x)d$ .

We now give a more involved function, appearing in semidefinite programming.

Example 2 (eigenvalue optimization) Let  $\mathcal{S}^p$  denote the linear space of symmetric matrices of order p, and consider the function  $\lambda_{\max}(X)$ , given by the maximum eigenvalue of a symmetric matrix  $X \in \mathcal{S}^p$ , whose elements are  $C^2$ -functions of a vector  $x \in \Re^n$ . Suppose that for a fixed X, r denotes the multiplicity of the maximum eigenvalue. Then from the analysis in [BS00, Ex. 3.98], the orthogonal projection onto the eigenspace corresponding to the maximum eigenvalue is an analytic function near X. By applying Gram-Schmidt orthonormalization to the columns of such projection, it is possible to define a "diagonalizing" analytic mapping  $\mathcal{C}: \mathcal{S}^p \to \mathcal{S}^r$ , such that:  $\mathcal{C}(X) = \lambda_{\max}(X) \mathbf{I}_r$  with  $\mathbf{I}_r$  the identity matrix of order r; the Jacobian  $\mathcal{DC}(X)$  is surjective; and the eigenvalues  $\lambda(\mathcal{C}(X))$  coincide with  $\lambda_{\max}(X)$ . The inner smooth mapping corresponds to  $\mathcal{C}$ , while the outer function is the eigenvalue function  $h(\mathcal{C}) = \lambda(\mathcal{C})$ , a positively homogeneous convex function.

The next function is typically used in applications such as wavelet-based image restoration, sparse representations, sparse regression, and compressive sensing problems.

Example 3 (Regularized minimization maps) Many signal or image reconstruction problems, as well as approximation problems, minimize a function of the form

$$\frac{1}{2}|Ax - y|^2 + \tau |x|,$$

for  $x \in \Re^n$ ,  $y \in \Re^k$ , A a matrix  $k \times n$ ,  $\tau$  a positive parameter, and for given appropriate norms. In the expression above, the first *data fidelity* term is smooth, and the second term corresponds to some regularization or penalty, for example ensuring sparsity of the solution. To fit our composite framework, it suffices to take m = n + 1,

$$c_j(x) = x_j \text{ for } j = 1, \dots, n, \qquad c_{n+1}(x) = \frac{1}{2} |Ax - y|^2;$$

and define the outer function  $h(C_1, \ldots, C_{n+1}) = C_{n+1} + \tau | (C_1, \ldots, C_n) |$ .

For any  $x \in \Re^n$ , the Jacobian mapping Dc(x) is the  $(n+1) \times n$  matrix made by appending to the identity matrix of order n a row with the transpose of  $A^{\top}(Ax-y)$ . As for the second-order derivative,  $D^2c(x)$ , for any  $d, \tilde{d} \in \Re^n$ ,  $[D^2c(x)\tilde{d}]d$  is a vector with its first n components equal to 0, its last component equal to  $\tilde{d}^{\top}A^{\top}Ad$ , a constant with respect to x.

We now give a possibly nonconvex function, appearing in approximation problems.

**Example 4 (sum of Euclidean norms)** Given a collection of smooth vector functions,  $\{\phi_1, \dots, \phi_J\}$  with  $\phi_j: \Re^n \to \Re^{m_j}$  for  $j=1,\dots,J$ , the function  $\sum_{j=1}^J |\phi_j(x)|$  is the composition of the following smooth mapping with  $m=\sum_{j=1}^J m_j$  components, and outer function:

$$c(x) = \left(\phi_1(x), \dots, \phi_J(x)\right) \quad \text{and} \quad h(C_1, \dots, C_{m_1}, C_{m_1+1}, \dots, C_m) = \sum_{j=1}^J \left| \left(C_{\sum_{k=1}^{j-1} m_k + 1}, \dots, C_{\sum_{k=1}^{j} m_k}\right) \right|.$$

When  $n = J = m_1 = 1$ , letting  $\phi_1(x) = 0.5a_2x^2 + a_1x + a_0$  for  $a_{2,1,0}$  given scalars, it is easy to see that the function f is convex if  $a_2 > 0$  with  $a_1^2 < 2a_0a_2$ , and concave otherwise.

The next function shows how to cast  $\ell_1$ -penalizations of nonlinear programming problems into the composite framework (provided no indicator function is used for polyhedral sets, to preserve finite-valuedness of the outer function).

Example 5 (Uryasev's exact penalty function) This is an optimal control problem over a discrete horizon with n time steps. The control variable  $x \in \mathbb{R}^n$  is constrained to the box  $[-0.2, 0.2]^n$ . Given initial values  $(\xi_0, \psi_0)$ , the state variables  $(\xi, \psi) \in \mathbb{R}^2$  follow a trajectory given by a recursive state equation of the form

$$\begin{split} \xi_1(x_1) &= \xi_0 + 0.2 \psi_0 \,, \\ \xi_{i+1}(x_{1:i+1}) &= \xi_i(x_{1:i}) + 0.2 \psi_i(x_{1:i}) \,, \end{split} \quad \psi_1(x_1) = -0.004 + 0.2 x_1, \text{ and } \\ \psi_1(x_1) &= -0.004 + 0.2 x_1, \text{ and } \\ \psi_{i+1}(x_{1:i+1}) &= \psi_i(x_{1:i}) - \text{cubic} \psi_i(x_{1:i})^2 - 0.004 \xi_i(x_{1:i}) + 0.2 x_{i+1} \,, \end{split}$$

for  $i=1,\ldots,n-1$  and where  $\mathtt{cubic} \geq 0$  is a given parameter. In the expressions above, we use the notation  $x_{1:i}$  to refer to the dependence of the  $i^{th}$ -state variables on the first  $i^{th}$  components of the control

variable x. The control problem to be solved is:

$$\begin{cases} \min_{x} & \frac{1}{2} \sum_{i=1}^{n} \xi_{i}(x_{1:i})^{2} \\ x_{i} \in [-0.2, 0.2] & \text{for } i = 1, \dots, n \\ \psi_{i}(x_{1:i}) \geq -1 & \text{for } i = 1, \dots, n-1 \\ \psi_{n}(x_{1:n}) = 0. \end{cases}$$

The corresponding  $\ell_1^+$ -penalty function uses parameters  $c_1$ ,  $c_2$ ,  $c_3 > 0$ , and can be written in composite form by letting m = 3n + 1,

$$c_{j}(x) = \begin{cases} -1 - \psi_{j}(x_{1:j}) & j = 1, \dots, n-1 \\ \psi_{n}(x_{1:n}) & j = n \\ \frac{1}{2} \sum_{i=1}^{n} \xi_{i}(x_{1:i})^{2} & j = n+1 \\ x_{j-n-1} - 0.2 & j = n+2, \dots, 2n+1 \\ -x_{j-2n-1} - 0.2 & j = 2n+2, \dots, 3n+1 \end{cases},$$

and taking the outer function

$$h(C) = C_{n+1} + c_1 \left( |C_{n+2:2n+1}|^+ + |C_{2n+2:3n+1}|^+ \right) + c_2 |C_n| + c_3 |C_{1:n-1}|^+.$$

The Jacobian mapping Dc(x) can be computed using the adjoint state equations, to obtain the trajectory derivatives  $\nabla \xi_i$  and  $\nabla \psi_i$ . When the parameter cubic is null, the control problem is linear-quadratic and the composite function  $h \circ c$  is convex. In this case, the second-order derivative,  $[D^2c(x)\tilde{d}]d$  is a vector with all of its components equal to 0, except for the  $(n+1)^{th}$  component, equal to  $\tilde{d}^{\top} \sum_{i=1}^{n} \nabla \xi_i \nabla \xi_i^{\top} d$  for any  $d, \tilde{d} \in \Re^n$ . When cubic > 0, the function is nonconvex, and second-order derivatives are difficult to compute for higher dimensions. For n=2, only  $c_2(x)=\psi_2$  and  $c_3(x)=\frac{1}{2}(\xi_1^2+\xi_2^2)$  have a nonzero Hessian, constantly equal to a scalar factor of the identity matrix of order n (the corresponding factors are -0.08 cubic and 0.0016, respectively).

Our final function, modifying [Lew02, Sec. 7] as in [MS03] to have a smooth inner mapping, is an example of a composite function having a nonconvex positively homogeneous outer function.

**Example 6 (Modified Lewis function)** For  $x = (x_1, x_2)^{\top}$  consider the following function, defined on a partition of  $\Re^2$ :

$$f(x) := \begin{cases} x_1^2 - x_2 & \text{on} \quad \{(x_1, x_2) \in \Re^2 : x_2 \le 0\} \\ x_1^2 + x_2 & \text{on} \quad \{(x_1, x_2) \in \Re^2 : 0 < x_2 \le x_1^2\} \\ 3x_1^2 - x_2 & \text{on} \quad \{(x_1, x_2) \in \Re^2 : 0 < x_1^2 \le x_2 \le 4x_1^2\} \\ -5x_1^2 + x_2 & \text{on} \quad \{(x_1, x_2) \in \Re^2 : 4x_1^2 < x_2\}. \end{cases}$$

This nonconvex function has the origin as unique critical point. The composite form,  $f = h \circ c$ , has smooth inner mapping  $c: \Re^2 \to \Re^2$  defined by  $c_1(x) = x_1^2$  and  $c_2(x) = x_2$ ; and outer function  $h: \Re^2 \to \Re$  given by

$$h(C) := \begin{cases} & |C_1|^+ - C_2 & \text{on} \quad C \in \{(C_1, C_2) \in \Re^2 : C_2 \le 0\} \\ & |C_1|^+ + C_2 & \text{on} \quad C \in \{(C_1, C_2) \in \Re^2 : 0 < C_2 \le C_1\} \\ & 3|C_1|^+ - C_2 & \text{on} \quad C \in \{(C_1, C_2) \in \Re^2 : 0 < C_1 \le C_2 \le C_1\} \\ & -5|C_1|^+ + C_2 & \text{on} \quad C \in \{(C_1, C_2) \in \Re^2 : 4C_1 < C_2\} \,. \end{cases}$$

This outer function is posivitely homogeneous, but not convex, so  $h \circ c$  is not included in our family.

# 3 Outer function and conceptual model

Even though the composite function  $h \circ c$  may be nonconvex, it is always locally Lipschitz continuous and directionally differentiable. Furthermore, as shown in [Sha03, Prop.3.2], the composite function is semismooth [Mif77], and strongly semismooth if the inner mapping is twice continuously differentiable.

Since the outer function is assumed to be positively homogeneous of degree 1, its directional derivative satisfies  $h'(0;\cdot) = h(\cdot)$ , so  $(h \circ c)'(x;d) = h(Dc(x)d)$  for any  $x,d \in \Re^n$ . As a result, the function  $(h \circ c)'(x;\cdot)$  is convex and its convex subdifferential at 0 coincides with the Clarke subdifferential  $\partial(h \circ c)(x)$ , in turn equal to the regular subdifferential from [RW98, Def. 8.3].

The optimality condition for  $\bar{x}$  to be a critical point of (1) is  $0 \in \partial(h \circ c)(\bar{x})$ . Being real-valued and convex, the outer function is continuous and, using a chain rule [BS00, Ch.3.4.1],

$$\forall x \in \Re^n \quad \partial(h \circ c)(x) = Dc(x)^\top \partial h(C) \quad \text{for } C = c(x).$$
 (2)

In particular, critical points satisfy the inclusion

$$0 \in Dc(\bar{x})^{\top} \partial h(\bar{c}) \quad \text{ for } \bar{c} = c(\bar{x}).$$

Another important consequence of positive homogeneity is the generalized Euler formula, stating that

$$\forall C \in \Re^m \text{ and for any } G \in \partial h(C) \text{ it holds that } h(C) = G^\top C.$$
 (3)

(the result follows from the subgradient inequality  $h((\alpha+1)C) \ge h(C) + G^{\top}((\alpha+1)C - C)$ , using that  $h((\alpha+1)C) = (\alpha+1)h(C)$  for any  $\alpha \ge -1$ ). This particular structural property will be exploited in the bundle method to define a cutting-planes approximation of the conceptual model.

The conceptual model from [LW08] composes the outer function with the linearization of the inner mapping around a point  $x^k$ . Specifically, if  $D_k = Dc(x^k)$  denotes the Jacobian mapping at  $x^k$ , then for any  $d \in \Re^n$ , letting

$$c_k(d) := c(x^k) + D_k d,$$

the conceptual model is given by the function

$$h(c_k(d)). (4)$$

To make tractable this conceptual model, we use a cutting-planes approximation  $\check{h}$  for the outer function and, for each pair  $(H^i := h(C^i), G^i \in \partial h(C^i))$ , define the plane

$$H^{i} + G^{i \top}(c_{k}(d) - C^{i}) = h(c(x^{k})) - \Delta_{i}^{k} + G^{i \top}c_{k}(d),$$

with

$$\Delta_i^k := h(c(x^k)) - H^i + G^{i} \, {}^{\mathsf{T}}C^i \ge G^{i} \, {}^{\mathsf{T}}c(x^k) \,, \tag{5}$$

by the subgradient inequality. The cutting-planes model for the conceptual model has the form

$$\check{h}(c_k(d)) = h(c(x^k)) + \max_{i \in \mathcal{B}} \left\{ -\Delta_i^k + G^{i \top} c_k(d) \right\}, \tag{6}$$

where the index set  $\mathcal{B}$  represents the bundle of information, varying along iterations. By convexity of h the cutting-planes model is always a lower approximation for the conceptual model, as shown in Figure 1 below.

When h is positively homogeneous, the equality (3) implies in (5) that  $\Delta_i^k := h(c(x^k))$ , so each plane has the form  $G^{i \top} c_k(d)$ , and the bundle only needs to keep the subgradient information  $(G^i)$ . Accordingly, the composite cutting-planes model can be written as

$$\check{h}(c_k(d)) = \max_{i \in \mathcal{B}} \left\{ G^{i \top} c_k(d) \right\}. \tag{7}$$

The example in Figure 1 shows that the composite model may not be a lower approximation for  $h \circ c$ , however, since the conceptual model is a good estimate for  $h \circ c$ , eventually the composite model is a good approximation for the objective function. This feature is particularly important for nonconvex bundle methods, which need to handle very carefully cutting-planes models (in some cases a tangent line of a nonconvex function can cut-off a section of the graph of the function, leaving out a critical point, for example), [HS10].

Typical situations arising when dealing with nonconvex functions are displayed in Figure 1. Specifically, setting n=1 in Example 4, the function in the figure is defined by adding 4 scalar terms of the form  $|\phi_k(x)| = 0.5a_kx^2 + b_kx + c_k|$ , with  $a_1 > 0$  and  $a_{2,3,4} < 0$ . Both cutting-planes models in the drawing use the same base points  $x^i$  in their respective bundles of information. The classical model is defined using the composite data  $\{x^i, (h \circ c)(x^i), \gamma^i \in \partial(h \circ c)(x^i)\}$ , whereas the composite model uses the outer data  $\{G^i \in \partial h(C^i), \text{ for } C^i = c(x^i)\}$ . Finally, the point  $x^*$  in the x-axis is a critical point found by the composite bundle method given in the next section.

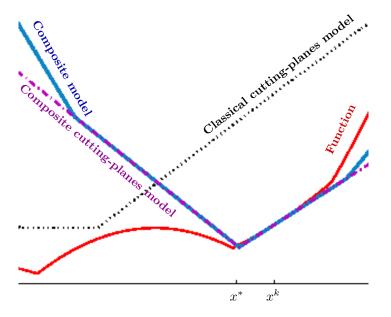


Figure 1: Various models for a nonconvex composite function.

# 4 Composite bundle method

At iteration number  $\ell$ , the composite model  $\check{h}_{\ell}(c_k(d))$ , given by (7) written with bundle set  $\mathcal{B} = \mathcal{B}_{\ell}$ , is used to generate a direction  $d^{\ell}$ . The inner mapping is linearized at a point  $x^k$  with  $k = k(\ell)$ . It is possible for k to be fixed for several consecutive values of  $\ell$ . The quadratic programming problem (QP) solved by  $d^{\ell}$  is nothing but the computation of the proximal point of the composite model, with prox-center  $x^k$  and a variable metric depending on matrices  $M_k^{\ell}$  detailed below.

The bundle of information  $\mathcal{B}_{\ell}$  is formed from (some of the) past information, corresponding to past directions  $d^i$  with  $i < \ell$ , generated using a prox-center  $x^{k(i)}$  and for which a call to the outer black-box ( $bb_h$ ) with  $C^i = c_{k(i)}(d^i)$  gave the value  $H^i = h(C^i)$  and the subgradient  $G^i \in \partial h(C^i)$ .

The variable prox-metric is induced by a positive definite matrix  $M_k^\ell$  of order n having the form

 $M_k^\ell = M_k + \mu_\ell I$  with  $M_k$  a symmetric  $n \times n$  matrix,  $\mu_\ell \ge 0$ , and I the identity matrix of order n.

The corresponding primal and dual metrics in  $\Re^n$  are

$$\forall x, \gamma \in \Re^n \quad |x|_{k,\ell}^2 := x^\top M_k^\ell x \quad \text{and} \quad \|\gamma\|_{k,\ell}^2 := \gamma^\top (M_k^\ell)^{-1} \gamma,$$

respectively. The eigenvalues of  $M_k^{\ell}$  are obtained by adding  $\mu_{\ell}$  to those of  $M_k$ , while the eigenvectors of  $M_k^{\ell}$  coincide with those of  $M_k$ . As a result, if for a matrix M the symbol  $\lambda(M)$  (respectively  $\lambda_{\min/\max}(M)$ ) denotes an eigenvalue (resp., minimum/maximum eigenvalues), then

$$\lambda(M_k^{\ell}) = \lambda(M_k) + \mu_{\ell}, \quad \lambda_{\min}(M_k^{\ell}) = \lambda_{\min}(M_k) + \mu_{\ell} \quad \text{and } \lambda_{\min}((M_k^{\ell})^{-1}) = \frac{1}{\lambda_{\max}(M_k) + \mu_{\ell}}.$$

In particular, letting  $|\cdot|$  denote the Euclidean norm in  $\Re^n$ ,

$$|x|_{k,\ell}^2 \ge \left(\lambda_{\min}(M_k) + \mu_{\ell}\right)|x|^2 \quad \text{and} \quad \|\gamma\|_{k,\ell}^2 \ge \frac{1}{\lambda_{\max}(M_k) + \mu_{\ell}}|\gamma|^2.$$
 (8)

Having the elements above, the next direction  $d^{\ell} \in \Re^n$  solves the quadratic programming problem

$$\min_{d \in \Re^n} \check{h}_{\ell}(c_k(d)) + \frac{1}{2} |d|_{k,\ell}^2. \tag{9}$$

Letting  $C^{\ell} = c_{k(\ell)}(d^{\ell})$ , the corresponding optimality condition with  $k = k(\ell)$  is

$$\exists \hat{G}^{\ell} \in \partial \check{h}_{\ell}(C^{\ell}) = conv\{G^{i} : i \in \mathcal{B}_{\ell}\} : \begin{cases} M_{k}^{\ell} d^{\ell} + D_{k}^{\top} \hat{G}^{\ell} = 0 \\ \check{h}_{\ell}(C^{\ell}) = \hat{G}^{\ell} {}^{\top} C^{\ell} = \hat{G}^{\ell} {}^{\top} c(x^{k}) - \|D_{k}^{\top} \hat{G}^{\ell}\|_{k,\ell}^{2}. \end{cases}$$
(10)

The optimal gradient  $\hat{G}^{\ell}$  is a convex combination only at active  $G^{i}$ 's such that  $\check{h}_{\ell}(C^{\ell}) = G^{i^{\top}}C^{\ell} = \hat{G}^{\ell^{\top}}c(x^{k})$ .

Without loss of clarity, and depending on the context, we sometimes use the short notations k and  $C^{\ell}$ , instead of the longer ones  $k(\ell)$  and  $c_{k(\ell)}(d^{\ell})$ .

### Algorithm 1 (COMPOSITE BUNDLE)

Inner and outer black-boxes,  $(bb_c)$  and  $(bb_h)$  respectively, are available.

#### Step 0 (Input and Initialization)

Select a stopping tolerance tol<sub>stop</sub>  $\geq 0$ , two Armijo and Wolfe-like parameters  $m_1, m_2 > 0$ , and a minimum positive threshold  $0 < \mu_{\min} < +\infty$ .

Initialize the iteration and serious step counters to  $\ell = 0$  and  $k = k(\ell) = 0$ , respectively.

For a starting point  $x^0 \in \mathbb{R}^n$ , call  $(bb_c)$  to compute the inner oracle values  $C^0 = c(x^0)$  and  $D_0 = Dc(x^0)$ . Call  $(bb_h)$  to compute the outer oracle values  $H^0 = (h \circ c)(x^0)$ ,  $G^0 \in \partial h(C^0)$ . Set  $\mathcal{B}_0 := \{G^0\}$ .

For a symmetric matrix  $M_0$  of order n, choose a prox-parameter  $\mu_0 \ge 0$  ensuring that the matrix  $M_0^0 = M_0 + \mu_0 I$  is positive definite.

Step 1 (Model Generation and QP Subproblem) Having the current serious step iterate  $x^k$ ,  $c(x^k)$ , its associated Jacobian  $D_k$ , and the bundle  $\left\{G^i\right\}_{i\in\mathcal{B}_\ell}$ , define the composite cutting-planes model function  $\check{h}_\ell(c_k(\cdot))$  from (7). A matrix  $M_k$  and a scalar  $\mu_\ell \geq 0$  such that  $M_k^\ell = M_k + \mu_\ell I$  is positive definite are also available.

Compute  $d^{\ell}$  by solving the quadratic program (9). These calculations include finding optimal simplicial multipliers  $\alpha^{\ell}$  such that in (10)

$$\hat{G}^{\ell} = \sum_{i \in \mathcal{B}_{\ell}} \alpha_i^{\ell} G^i.$$

Define the aggregate error

$$\hat{e}_{\ell} = (h \circ c)(x^k) - \hat{G}^{\ell \top} c(x^k),$$

(note that  $\hat{e}_{\ell} \geq 0$  by (5) and (3)), and compute the predicted decrease

$$\delta_{\ell} = (h \circ c)(x^{k}) - \check{h}_{\ell}(c_{k}(d^{\ell})) - \frac{1}{2} |d^{\ell}|_{k,\ell}^{2} = \hat{e}_{\ell} + \frac{1}{2} ||D_{k}^{\top} \hat{G}^{\ell}||_{k,\ell}^{2},$$
(11)

where the last equality follows from (10).

#### Step 2 (Stopping test)

If  $\max(\hat{e}_{\ell}, \|D_k^{\top} \hat{G}^{\ell}\|_{k,\ell}^2) \leq \mathsf{tol}_{\mathsf{stop}}$ , stop.

Otherwise, call (bb<sub>h</sub>) to obtain  $h(c_k(d^{\ell}))$  and  $G^{\ell} \in \partial h(c_k(d^{\ell}))$ .

## Step 3 (Serious/backtrack/null step test)

Check the descent condition

$$h(c_k(d^{\ell})) \le (h \circ c)(x^k) - m_1 \delta_{\ell}. \tag{12}$$

If (12) does not hold, declare a null step: take  $\mu_{\ell+1} \ge \mu_{\ell}$ ,  $M_k^{\ell+1} = M_k + \mu_{\ell+1} I$ , and go to Step 4. If (12) is true, call the inner oracle ( $bb_c$ ) to compute  $c(x^k + d^{\ell})$ . Call the outer oracle ( $bb_h$ ) to compute  $\Gamma^{\ell} \in \partial h(c(x^k + d^{\ell}))$ , and check the condition below:

$$\Gamma^{\ell \top}[c_k(d^{\ell}) - c(x^k + d^{\ell})] \ge -m_2 \delta_{\ell}. \tag{13}$$

If (13) holds, declare a serious step, and go to Step 4.

Otherwise, if (13) does not hold, declare a backtracking step: take  $\mu_{\ell+1} \ge \mu_{\min} + \mu_{\ell}$ ,  $M_k^{\ell+1} = M_k + \mu_{\ell+1}$ I,  $\mathcal{B}_{\ell+1} = \mathcal{B}_{\ell}$ , set  $k(\ell+1) = k(\ell)$ , increase  $\ell$  by 1, and loop to Step 1.

#### Step 4 (Bundle update and management)

If needed, compress  $\mathcal{B}_{\ell}$  either by keeping all strongly active elements  $(i \in \mathcal{B}_{\ell} : \alpha_i^{\ell} > 0)$ , or by replacing some strongly active elements by the aggregate gradient  $\hat{G}^{\ell}$ . Define  $\mathcal{B}_{\ell+1}$  by adding to the possibly

compressed set the new outer gradient  $G^{\ell}$ .

If the step was declared null, set  $k(\ell+1)=k(\ell)$ , increase  $\ell$  by 1, and loop to Step 1.

If the step was declared serious, set  $k(\ell+1) = k+1$ ,  $x^{k+1} = x^k + d^\ell$ , call (bb<sub>c</sub>) to compute  $D_{k+1} = Dc(x^{k+1})$ , compute a new symmetric matrix  $M_{k+1}$  and a prox-parameter  $\mu_{\ell+1} \geq 0$  such that  $M_{k+1}^{\ell+1} = M_{k+1} + \mu_{\ell+1} I$  is positive definite. Increase k and  $\ell$  by 1, and loop to Step 1.

In Step 3, the decision to declare a null-steps is different from standard bundle methods, and exploits the structure knowledge in the composite case. Namely, while condition (12) checks if there is descent for the outer function, (13) checks adequacy between the inner mapping and its conceptual model.

The prox-parameter increases at backtracking steps, and can be increased at null steps or decreased at serious steps. In this context, the wording "backtrack" may sound odd, since it corresponds to choosing a larger prox-parameter. The explanation comes from noticing that, due to (10), the value  $1/\mu_{\ell}$  can be seen as a stepsize. In this sense, when the algorithm detects the need of backtracking, increasing the prox-parameter results indeed in a smaller stepsize. Since if the matrix  $M_k$  is positive definite, we may choose a null prox-parameter, in the backtracking step we force positivity by means of the positive threshold  $\mu_{\min}$ , set at the Initialization step.

In order to ensure that matrices  $M_k^{\ell}$  are always positive definite, the prox-parameter update can be done as follows:

$$\mu_{\ell} = \begin{cases} 0 & \text{if } \lambda_{\min}(M_{k(\ell)}) \ge \mu_{\min}, \\ -\lambda_{\min}(M_{k(\ell)}) + \mu_{\min} & \text{otherwise.} \end{cases}$$
 (14)

We mention that there is a difference between the Composite Algorithm 1 and usual bundle methods, in terms of oracle calls. A standard bundle method makes one call to both the inner and outer oracles per iteration, independent of getting a serious or a null step. Instead, Step 2 in Algorithm 1 makes one call to the outer oracle ( $bb_h$ ) at all iterations, needing additional calls to both oracles ( $bb_{h/c}$ ) at Step 3, for deciding between serious or backtracking steps. In this case, the additional subgradient  $\Gamma^{\ell} \in \partial h(c(x^k + d^{\ell}))$  may enter the bundle.

In Step 4 we see that the bundle is formed only by outer gradients, corresponding either to  $G^i \in \partial h(C^i)$  or to some past aggregate gradient  $\hat{G}^i$ . By convexity of h, the cutting-planes model is always a lower approximation to h:

for all 
$$\ell \ge 1$$
 and  $C \in \Re^m$   $\check{h}_{\ell}(C) \le h(C)$ . (15)

Moreover, since  $\hat{G}^{\ell} \in \partial h(C^{\ell})$  by (10), after some algebra involving  $\delta_{\ell}$  and  $\hat{e}_{\ell}$ , we see that

$$h(c(x^k)) - \check{h}_{\ell}(C^{\ell}) - \|D_k^{\top} \hat{G}^{\ell}\|_{k,\ell}^2 = \delta_{\ell} - \frac{1}{2} \|D_k^{\top} \hat{G}^{\ell}\|_{k,\ell}^2 = \hat{e}_{\ell},$$

and, hence,

$$\forall \ell \text{ and } C \in \Re^m \qquad h(C) \ge (h \circ c)(x^k) + \hat{G}^{\ell \top}(C - c(x^k)) - \hat{e}_{\ell}. \tag{16}$$

Remark 1 (The trivial composite case) To see how a classical bundle method compares to Algorithm 1, consider a convex nonsmooth function f, and the, somewhat artificial, outer function  $h \equiv f$  and inner mapping c(x) = x. With respect to our composite structure, h is no longer positive definite, Dc = I,  $c_k(d) = x^k + d = c(x^k + d)$ , and, instead of (7), the cutting-planes model has the form (6). Since the inner mapping is affine, there is no second order to exploit, and the matrices  $M_k \equiv 0$  in Algorithm 1. The QP optimality condition in (10) becomes

$$\exists \hat{G}^{\ell} \in \partial \check{h}_{\ell}(x^{\ell}) = conv\{G^{i} : i \in \mathcal{B}_{\ell}\} : \left\{ \begin{array}{l} \mu_{\ell}d^{\ell} + \hat{G}^{\ell} = 0 \\ \check{h}_{\ell}(x^{\ell}) = h(x^{k}) - \hat{\Delta}_{\ell}^{k} + \hat{G}^{\ell \top} x^{\ell} \end{array}, \right.$$

for  $\hat{\Delta}_{\ell}^{k} := \sum_{i \in \mathcal{B}_{\ell}} \alpha_{i}^{\ell} \Delta_{i}^{k}$ . The aggregate error is defined by  $\hat{e}_{\ell} := \hat{\Delta}_{\ell}^{k} - \hat{G}^{\ell} \, ^{\mathsf{T}} x^{k}$ , consistent with the error definition in Step 1, because when h is positively homogenous,  $\Delta_{i}^{k} = h(x^{k})$ , by (3). As a result, the expression for the predicted decrease in (11) remains true. Finally, in Step 3, because  $c_{k}(d^{\ell}) = x^{k} + d^{\ell} = c(x^{k} + d^{\ell})$ , there is no backtracking step, because the inequality in (13) trivially holds:

$$0 = \Gamma^{\ell \top} (c_k(d^{\ell}) - c(x^k + d^{\ell})) < -m_2 \delta_{\ell}.$$

From the above it follows that Algorithm 1 with  $M_k \equiv 0$  is applicable for the trivial composite case, becoming a classical bundle method for convex functions without any particular structure.

# 5 Convergence results

Since matrices  $M_k^{\ell}$  are always positive definite, they induce a norm and (8) holds. Together with (11) and (10), this means that the relations

$$\delta_{\ell} \ge \begin{cases} \max\left(\frac{1}{2} \|D_{k}^{\top} \hat{G}^{\ell}\|_{k,\ell}^{2}, \hat{e}_{\ell}\right) \ge \frac{1}{2(\lambda_{\max}(M_{k}) + \mu_{\ell})} |D_{k}^{\top} \hat{G}^{\ell}|^{2} & \text{(a)} \\ \max\left(\frac{1}{2} |d^{\ell}|_{k,\ell}^{2}, \hat{e}_{\ell}\right) \ge \frac{1}{2} \left(\lambda_{\min}(M_{k}) + \mu_{\ell}\right) |d^{\ell}|^{2} & \text{(b)} \end{cases}$$

are always satisfied.

We start our convergence analysis showing that there can only be a finite number of consecutive backtracking steps.

**Proposition 1 (Finite backtracking loop)** Suppose the inner mapping is  $C^2$ . If after some iteration  $\hat{\ell}$ , Algorithm 1 makes a last serious step and thereafter generates only null and backtracking steps, the following holds:

- (i) If the sequence  $\{d^{\ell}\}_{\ell \geq \hat{\ell}}$  is bounded, then there exists  $\mu_{noBT} > 0$  such that (13) holds for any  $\mu_{\ell} > \mu_{noBT}$ .
- (ii) There cannot be infinitely many consecutive backtracking steps.

**Proof:** For convenience, let  $\hat{k} = k(\hat{\ell})$ ,  $\hat{x} = x^{\hat{k}} (= x^{k(\ell)})$  for all  $\ell \geq \hat{\ell}$ , and  $\hat{M} = M_{\hat{k}}$  denote, respectively, the k-iteration index, prox-center and matrix corresponding to the last serious step.

Since c is a  $C^2$ -mapping, a mean-value theorem applies to each component  $c_j$ , j = 1, ..., m:

$$c_j(\hat{x} + d^{\ell}) - c_j(\hat{x}) - \nabla c_j(\hat{x})^{\top} d^{\ell} = \frac{1}{2} \nabla^2 c_j(\xi_j) (d^{\ell}, d^{\ell}) \text{ for some } \xi_j \in [\hat{x}, \hat{x} + d^{\ell}].$$

Recall that  $|\cdot|$  is the Euclidean matrix norm. By assumption,  $\{\hat{x} + d^{\ell}\}$  is bounded so  $|\Gamma^{\ell}| \leq L$  for all  $\ell$  and some constant L. Similarly, boundedness of  $\{d^{\ell}\}$  implies that  $|\nabla^2 c_j(\xi_j)| \leq D$  for all  $j = 1, \ldots, m$ , for some constant D. By Cauchy-Schwarz inequality,

$$\Gamma^{\ell \top}[c(\hat{x} + d^{\ell}) - c(\hat{x}) - D_{\hat{k}}d^{\ell}] \le \frac{\sqrt{m}}{2}LD|d^{\ell}|^{2}$$
.

Using the rightmost inequality in (17)(b), we obtain

$$\delta_{\ell} \ge \frac{1}{2} |d^{\ell}|_{\hat{k},\ell}^2 \ge \frac{1}{2} (\lambda_{\min}(\hat{M}) + \mu_{\ell}) |d^{\ell}|^2$$
.

In the notation used in this proof,  $c(x^k + d^\ell) = c(\hat{x} + d^\ell)$  and  $c_k(d^\ell) = c(\hat{x}) + D_{\hat{k}}d^\ell$ . Suppose that (13) is not satisfied and multiply by -1 the corresponding inequality:

$$m_2 \delta_\ell < -\Gamma^{\ell \ \top} [c_k(d^\ell) - c(x^k + d^\ell)] = \Gamma^{\ell \ \top} [c(\hat{x} + d^\ell) - c(\hat{x}) - D_{\hat{k}} d^\ell] \,.$$

Using the bounds above we see that

$$\frac{1}{2} m_2(\lambda_{\min}(\hat{M}) + \mu_{\ell}) |d^{\ell}|^2 \le m_2 \delta_{\ell} < \Gamma^{\ell \top} [c(\hat{x} + d^{\ell}) - c(\hat{x}) - D_{\hat{k}} d^{\ell}] \le \frac{\sqrt{m}}{2} LD |d^{\ell}|^2.$$

Therefore, nonsatisfaction of (13) implies  $\mu_{\ell} < \sqrt{m}LD/m_2 - \lambda_{\min}(\hat{M})$ . Item (i) follows, by taking  $\mu_{\text{noBT}} \ge \sqrt{m}LD/m_2 - \lambda_{\min}(M_k)$ .

The claim in item (ii) is shown by contradiction, assuming that (13) does not hold, with  $\ell \to \infty$  due to backtracking. An infinite backtracking loop drives  $\mu_{\ell}$  to infinity, as well as the minimum eigenvalue of the matrix  $M_{\hat{k}}^{\ell}$ , because  $M_{\hat{k}}^{\ell} = \hat{M} + \mu_{\ell} I$ . Since there are only backtracking steps, the bundle does not change, and the composite model  $\check{h}_{\ell}(c_{\hat{k}}(\cdot))$  is a fixed function, say  $\varphi$ . Hence, by [HUL93, Prop. XV.4.1.5], the minimand in (9) converges to  $\varphi(d)$  for all  $d \in \Re^n$  and  $d^{\ell} \to 0$  as  $\ell \to \infty$ . Therefore,  $\hat{x} + d^{\ell} \to \hat{x}$  with  $\varphi(d^{\ell}) + \frac{1}{2} |d^{\ell}|_{\hat{k},\ell}^2 \to \varphi(0)$ . In particular, this means that the sequence  $\{d^{\ell}\}$  is bounded, and the contradiction follows from item (i).

As a consequence, if the algorithm loops forever, it generates either an infinite sequence of serious steps, or a finite one, followed by infinitely many null steps (possibly nonconsecutive, due to backtracking). Both situations are considered below, adapting classical developments for bundle methods [HUL93, Ch.XV.3, XV.4] to the composite setting.

We start with the case of infinitely many serious steps.

**Lemma 1 (Infinitely many serious steps)** Suppose Algorithm 1 generates an infinite sequence  $\{x^k\}$  of serious steps. Denote by  $\ell_k$  an iteration index yielding a serious step:  $x^{k+1} = x^k + d^{\ell_k}$ . The following holds:

- (i) If  $0 < m_2 < m_1$  then, as  $k \to \infty$ , either  $(h \circ c)(x^k) \to -\infty$  or  $\delta_{\ell_k} \to 0$ .
- (ii) Suppose  $\delta_{\ell_k} \to 0$ . Then  $\lim \hat{e}_{\ell_k} = 0$  and,

if the series 
$$\sum_{k} \frac{1}{\mu_{\ell_k} + \lambda_{\max}(M_k)}$$
 is divergent, (18)

then  $\liminf \|D_k^{\top} \hat{G}^{\ell_k}\|_{k,\ell_k} = 0.$ 

- (iii) If, in addition, the sequence  $\{x^k\}$  is bounded, then it has at least one accumulation point that is critical for (1).
- (iv) If, instead of (18), the stronger condition

the sequence 
$$\{\mu_{\ell_k} + \lambda_{\max}(M_k)\}\$$
 is bounded above, (18')

then all accumulation points are critical for (1).

**Proof:** By (13), since  $\Gamma^{\ell_k} \in \partial h(c(x^k + d^{\ell_k}))$ , we have that

$$h(c(x^k) + D_k d^{\ell_k}) \ge h(c(x^k + d^{\ell_k})) - m_2 \delta_{\ell_k} = (h \circ c)(x^{k+1}) - m_2 \delta_{\ell_k}. \tag{19}$$

Together with satisfaction of (12), this means that  $(h \circ c)(x^{k+1}) \leq (h \circ c)(x^k) - (m_1 - m_2)\delta_{\ell_k}$ . The telescopic sum yields that either  $(h \circ c)(x^k) \setminus -\infty$  or  $0 \leq \delta_{\ell_k} \to 0$ , because  $m_1 - m_2 > 0$ .

The first assertion in item (ii) follows from (11). The rightmost inequality in (17)(a), together with our assumption (18) gives the second result.

To show item (iii), consider a k-subsequence such that  $\|D_k^\top \hat{G}^{\ell_k}\|_{k,\ell_k} \to 0$  and extract a further subsequence of serious steps with accumulation point  $x^{acc}$ . By boundedness of  $\{x^k\}$  and local Lipschitzianity of h, there is an associated subsequence  $\{\hat{G}^\ell\}$  with accumulation point  $\hat{G}^{acc}$ . Passing to the limit in (16) and using that  $\hat{e}_{\ell_k} \to 0$ , we obtain in the limit that  $\hat{G}^{acc} \in \partial h(C^{acc})$  for  $C^{acc} = c(x^{acc})$ . By smoothness of c,  $D_k \to D_{acc} = Dc(x^{acc})$ , and by (2),  $D_{acc}^\top \hat{G}^{acc} \in \partial (h \circ c)(x^{acc})$ . The result follows from item (ii). Item (iv) is similar to item (iii), noticing that if L > 0 denotes an upper bound for  $\{\mu_{\ell_k} + \lambda_{\max}(M_k)\}$ , then the relation  $\frac{1}{2(\mu_{\ell_k} + \lambda_{\max}(M_k))} \ge 1/2L$  in (17)(b) gives that  $|D_k^\top \hat{G}^{\ell_k}| \to 0$  as  $k \to \infty$  (for the whole sequence).

The remaining case refers to an infinite null-step loop and makes use of the aggregate linearization

$$\check{h}_{\ell}^{\mathrm{agg}}(c_k(d)) = \check{h}_{\ell}(c_k(d^{\ell})) + \hat{G}^{\ell \top} D_k(d - d^{\ell}),$$

the associated strongly convex function

$$\mathcal{H}_{\ell}(d) = \check{h}_{\ell}^{\text{agg}}(c_k(d)) + \frac{1}{2}|d|_{k,\ell}^2,$$
 (20)

and the result [HUL93, Lem. XV.4.3.3]:

$$\mathcal{H}_{\ell-1}(d) = \mathcal{H}_{\ell-1}(d^{\ell-1}) + \frac{1}{2}|d - d^{\ell-1}|_{k,\ell-1}^2.$$
(21)

Finally, and similar to [CL93, Sec. 4], note that Step 4 in Algorithm 1 updates the bundle of information in a way ensuring that not only

if iteration 
$$\ell - 1$$
 was declared a null step, then  $\check{h}_{\ell-1}^{\text{agg}}(c_k(d)) \leq \check{h}_{\ell}(c_k(d))$ , (22)

but also

if iteration 
$$\ell - 1$$
 was declared a null step, then  $\check{h}_{\ell}(c_k(d)) \ge G^{\ell-1} \, {}^{\top}c_k(d)$  (23)

for all  $d \in \Re^n$ .

**Lemma 2 (Finitely many serious steps)** Suppose that, after some iteration  $\hat{\ell}$ , Algorithm 1 makes a last serious step  $\hat{x} = \hat{x}^{k(\ell)}$  and thereafter generates an infinite number of null steps, possibly nonconsecutive, due to intermediate backtracking steps. The following holds:

- (i) The sequence  $\{d^{\ell}\}_{\ell>\hat{\ell}}$  is bounded and there is an iteration  $\ell'>\hat{\ell}$  such that only null steps are done for all  $\ell\geq\ell'$ .
- (ii) If  $m_1 < 1$ , then  $\delta_{\ell} \to 0$  and  $\hat{e}_{\ell} \to 0$ .
- (iii) If, in addition, for the (fixed) matrix  $\hat{M} = M_{k(\hat{\ell})}$ .

the series 
$$\sum_{\ell > \ell'} \frac{\mu_{\ell-1} + \lambda_{\min}(\hat{M})}{(\mu_{\ell} + \lambda_{\max}(\hat{M}))^2}$$
 is divergent, (24)

then  $\liminf |Dc(\hat{x})^{\top} \hat{G}^{\ell}| = 0$ ,  $\hat{x} + d^{\ell} \to \hat{x}$  for some  $\ell$ -subsequence, and  $\hat{x}$  is critical for (1).

**Proof:** For convenience, let  $\hat{k} = k(\hat{\ell})$ ,  $\hat{x} = x^{\hat{k}} (= x^{k(\ell)})$  for all  $\ell \geq \hat{\ell}$ , and  $\hat{M} = M_{\hat{k}}$  denote, respectively, the k-iteration index, prox-center and matrix corresponding to the last serious step.

Consider  $\ell > \hat{\ell}$  and recall that, since  $M_k^{\ell} = \hat{M} + \mu_{\ell} I$  and  $\{\mu_{\ell}\}$  is nondecreasing at null and backtracking steps,

$$\frac{1}{2}|d|_{\hat{k},\ell-1}^2 \le \frac{1}{2}|d|_{\hat{k},\ell}^2.$$

The sum of this inequality and (22), together with (20) written with  $\ell$  replaced  $\ell-1$ , results in the relation

$$\mathcal{H}_{\ell-1}(d) \leq \check{h}_{\ell}(c_{\hat{k}}(d)) + \frac{1}{2}|d|_{\hat{k},\ell}^2$$
.

In particular, for  $d = d^{\ell}$ , we obtain that  $\mathcal{H}_{\ell-1}(d^{\ell}) \leq \mathcal{H}_{\ell}(d^{\ell})$  from (20), because  $\check{h}_{\ell}(c_{\hat{k}}(d^{\ell})) = \check{h}_{\ell}^{\text{agg}}(c_{\hat{k}}(d^{\ell}))$ . Together with (21), written at  $d = d^{\ell}$ , we see that

$$\mathcal{H}_{\ell-1}(d^{\ell-1}) \le \mathcal{H}_{\ell-1}(d^{\ell-1}) + \frac{1}{2}|d^{\ell} - d^{\ell-1}|_{\hat{k},\ell-1}^2 = \mathcal{H}_{\ell-1}(d^{\ell}) \le \mathcal{H}_{\ell}(d^{\ell}). \tag{25}$$

By the definitions of  $\check{h}_{\ell}^{\text{agg}}$  and  $\mathcal{H}_{\ell}$ , the optimal value in (9) equals  $\mathcal{H}_{\ell}(d^{\ell})$ . Hence, (25) implies that the sequence of optimal values in (9) is strictly increasing, with  $\mathcal{H}_{\ell}(d^{\ell}) \leq \check{h}_{\ell}(c_k(0)) = \check{h}_{\ell}(c(\hat{x}))$ . Since, by (15),  $\check{h}_{\ell} \leq h$ , then

$$\mathcal{H}_{\ell}(d^{\ell}) \le (h \circ c)(\hat{x}). \tag{26}$$

So  $\{\mathcal{H}_{\ell}(d^{\ell})\} \uparrow \mathcal{H}_{\infty}$  for some  $\mathcal{H}_{\infty} \leq (h \circ c)(\hat{x})$ , with  $|d^{\ell} - d^{\ell-1}|_{\hat{k}, \ell-1}^2 \to 0$  as  $\ell \to \infty$ , by (25). But  $\mu_{\ell} \geq \mu_{\hat{\ell}}$  (at null and backtracking steps prox-parameters are nondecreasing), so the left relation in (8) implies that

$$d^{\ell} - d^{\ell-1} \to 0. \tag{27}$$

Using (20) with d=0, we see that  $\mathcal{H}_{\ell}(0)=\check{h}^{\mathrm{agg}}_{\ell}(c_{\hat{k}}(0))$ . Together with the definition of  $\check{h}^{\mathrm{agg}}_{\ell}$ , the left inequality in the second line in (10), and the definition of  $c_k(d^{\ell})$ , this implies that that  $\mathcal{H}_{\ell}(0)=\check{h}_{\ell}(c_{\hat{k}}(d^{\ell}))-\hat{G}^{\ell} \ ^{\mathsf{T}}D_{\hat{k}}d^{\ell}=\hat{G}^{\ell} \ ^{\mathsf{T}}c(\hat{x})$ . Since  $\hat{G}^{\ell}\in conv\{G^i:i\in\mathcal{B}_{\ell}\}$ , using (5) and (3) we obtain that  $\mathcal{H}_{\ell}(0)\leq (h\circ c)(\hat{x})$ . Therefore, writing (21) with  $\ell-1$  replaced by  $\ell$  at d=0 yields the relations

$$\frac{1}{2} |d^{\ell}|_{\hat{k},\ell}^2 = \mathcal{H}_{\ell}(0) - \mathcal{H}_{\ell}(d^{\ell}) \le (h \circ c)(\hat{x}) - \mathcal{H}_{\hat{\ell}+1}(d^{\hat{\ell}+1}),$$

because the sequence  $\{\mathcal{H}_{\ell}(d^{\ell})\}$  is increasing, by (25), since  $\ell > \hat{\ell}$ . Using once more the left relation in (8) we conclude that the sequence  $\{d^{\ell}\}$  is bounded, and item (i) follows from Proposition 1(i). To show item (ii), consider iteration indices  $\ell - 1$ ,  $\ell \geq \ell'$ , giving two consecutive null steps, and set  $C^{\ell} = c(\hat{x}) + Dc(\hat{x})d^{\ell}$  and  $C^{\ell-1} = c(\hat{x}) + Dc(\hat{x})d^{\ell-1}$ . Since  $\delta_{\ell} \geq 0$ , we substract the inequality  $\delta_{\ell} \leq (h \circ c)(\hat{x}) - \check{h}_{\ell}(C^{\ell})$ , obtained from (11), from nonsatisfaction of (12), both with  $x^{k} = \hat{x}$ , to see that

$$0 \le (1 - m_1)\delta_{\ell} \le h(C^{\ell}) - \check{h}_{\ell}(C^{\ell}). \tag{28}$$

By (3),

$$h(C^{\ell-1}) = G^{\ell-1} \, {}^{\mathsf{T}} C^{\ell-1} \,,$$

and by (23),

$$\check{h}_{\ell}(c_{\hat{k}}(d^{\ell})) = \check{h}_{\ell}(C^{\ell}) \ge G^{\ell-1} \, {}^{\mathsf{T}}C^{\ell}.$$

Since  $\{d^{\ell}\}$  is bounded, any Lipschitz constant L for h gives an upper bound for  $\{|G^{\ell}|\}$ ; as a result,

$$h(C^{\ell}) - \check{h}_{\ell}(C^{\ell}) = h(C^{\ell}) - h(C^{\ell-1}) + h(C^{\ell-1}) - \check{h}_{\ell}(C^{\ell})$$

$$\leq L|C^{\ell} - C^{\ell-1}| + G^{\ell-1} \top (C^{\ell-1} - C^{\ell})$$

$$\leq 2L|Dc(\hat{x})||d^{\ell} - d^{\ell-1}|. \tag{29}$$

From (28) and (27) it follows that  $\delta_{\ell} \to 0$ , and by the right hand side expression in (11),  $\hat{e}_{\ell} \to 0$ , as stated.

Finally, to see item (iii), the left hand side expression in (11) of  $\delta_{\ell}$  and the definitions of  $\check{h}_{\ell}^{\text{agg}}$  and  $\mathcal{H}_{\ell}$  give the identity  $\delta_{\ell} = (h \circ c)(\hat{x}) - \mathcal{H}_{\ell}(d^{\ell})$ . Therefore, by the right inequality in (25), (22) with  $d = d^{\ell}$ , and the left relation in (8),

$$\begin{array}{lcl} \delta_{\ell-1} & \geq & \delta_{\ell} + \frac{1}{2} |d^{\ell} - d^{\ell-1}|_{\hat{k}, \ell-1}^2 \\ & \geq & \delta_{\ell} + \frac{\lambda_{\min}(\hat{M}) + \mu_{\ell-1}}{2} |d^{\ell} - d^{\ell-1}|^2 \,. \end{array}$$

From (29) and (28), we obtain that  $\delta_{\ell} \leq \frac{2L|Dc(\hat{x})|}{1-m_1}|d^{\ell}-d^{\ell-1}|$ , so

$$\delta_{\ell-1} - \delta_{\ell} \ge \frac{\lambda_{\min}(\hat{M}) + \mu_{\ell-1}}{2} \left(\frac{1 - m_1}{2L|Dc(\hat{x})|}\right)^2 \delta_{\ell}^2.$$

Letting  $K := (1 - m_1)^2 / \Big( 8 |Dc(\hat{x})|^2 L^2 \Big)$ , and summing over  $\ell > \hat{\ell}$ ,

$$K\sum_{\ell>\hat{\ell}} (\lambda_{\min}(\hat{M}) + \mu_{\ell-1})\delta_{\ell}^2 \le \delta_{\hat{\ell}} < +\infty.$$

Furthermore, using (17)(a), the series

$$\sum_{\ell \sim \hat{\ell}} \frac{\mu_{\ell-1} + \lambda_{\min}(\hat{M})}{(\mu_{\ell} + \lambda_{\max}(\hat{M}))^2} |Dc(\hat{x})^{\top} \hat{G}^{\ell}|^4$$

converges too. With our assumption (24), this implies that  $\liminf |Dc(\hat{x})^{\top} \hat{G}^{\ell}|^4 = 0$ . Consider indices  $\ell$  in a corresponding convergent subsequence of  $\{d^{\ell}\}$ . Since from the expression for  $d^{\ell}$  in (10),  $|Dc(\hat{x})^{\top} \hat{G}^{\ell}|^2 = |M_k^{\ell} d^{\ell}|^2 \geq (\lambda_{\min}(\hat{M}) + \mu_{\hat{\ell}})^2 |d^{\ell}|^2$ , we see that the subsequence fo  $\{d^{\ell}\}$  converges to zero and, hence,  $\hat{x} + d^{\ell} \to \hat{x}$  on this subsequence. Finally, by boundedness of  $\{C^{\ell} = c(\hat{x}) + Dc(\hat{x})d^{\ell}\}$ , all outer subgradients are bounded and, hence, the subsequence  $\{\hat{G}^{\ell}\}$  has some accumulation point  $\hat{G}^{acc}$  such that  $Dc(\hat{x})^{\top}\hat{G}^{acc} = 0$ . The result follows from passing to the limit in (16) to give  $\hat{G}^{acc} \in \partial h(c(\hat{x}))$ , and using the chain rule (2).

Remark 2 (The trivial composite case, suite and end.) In the setting of Remark 1, when h is not positively homogeneous but merely convex, and c is the identity, the cutting-planes model has the form (6) and (23) states that

if iteration  $\ell - 1$  was declared a null step, then  $\check{h}_{\ell}(c_k(d)) \geq h(x^k) - \Delta_{\ell-1}^k + G^{\ell-1} \, {}^{\top}c_k(d)$ ,

which is consistent with the fact that  $h(x^k) = \Delta_{\ell-1}^k$  when h is positively homogenous, by (5) and (3). In fact, when  $M_k \equiv 0$  for all k, as considered in the trivial structure, Lemmas 1 and 2 boil down to [HUL93, pp.309 and 311, vol.II, Thms.XV.3.2.2 and XV.3.2.4].

Putting together Proposition 1 and Lemmas 1 and 2, we can show convergence for objective functions that are *inf-compact*, i.e., functions having some level set that is nonempty and compact (sometimes also referred to as *level-bounded* functions). In this case, by lower semicontinuity,  $h \circ c$  always attains its minimum and the sequence of serious steps is bounded.

**Theorem 1 (Synthesis)** Consider solving problem (1) with Algorithm 1 and suppose that in (1) the objective function  $h \circ c$  is inf-compact. If  $0 < m_2 < m_1 < 1$ ,  $\mathsf{tol}_{\mathsf{stop}} = 0$ , with both (18) and (24) being satisfied, the following holds.

- (i) Either the sequence of serious steps is infinite and bounded, and at least one of its accumulation points is critical for (1).
- (ii) Or the last serious step  $\hat{x}$  is critical for (1), with  $\{\hat{x} + d^{\ell}\} \to \hat{x}$  for some  $\ell$ -subsequence. If, instead of (18), the stronger condition (18') holds, item (i) can be replaced by
- (i') Either the sequence of serious steps is infinite and bounded, and all its accumulation points are critical for (1). □

Our conditions (18) and (24), on the variable prox-metric, are fairly general and not difficult to enforce. For (18) to hold, it is enough to take matrices  $M_k^{\ell}$  that are uniformly bounded from above at serious steps (when  $\ell = \ell_k$ ). As for null steps, choosing  $\mu_{\ell+1} \in [\mu_{\ell}, \mu_{\text{max}}]$  for some finite bound  $\mu_{\text{max}} > -\lambda_{\min}(\hat{M})$ , rule (14) ensures satisfaction of (24). Depending on the particular problem, the more general condition (24) may help in preventing bad choices (too small) for the bound  $\mu_{\text{max}}$ .

We finish our analysis by supposing the stopping tolerance is positive. In this case, if (18') and (24) hold, by Lemmas 1 and 2, the nominal decrease goes to 0. Then, by (11), both  $\hat{e}_{\ell}$  and  $\hat{G}^{\ell}$  go to 0, and eventually the stopping test will be triggered. For such an iteration index, say  $\ell$ best, (16) gives the following approximate optimality condition for the last serious step  $x^{\text{best}} = x^{k(\ell \text{best})}$ :

$$\forall x \in \Re^n \qquad (h \circ c)(x) \quad \geq \quad (h \circ c)(x^{\mathrm{best}}) + \hat{G}^{\ell \mathrm{best}} \, {}^{\top} \Big( c(x) - c(x^{\mathrm{best}}) \Big) - \hat{e}_{\ell \mathrm{best}} \\ \qquad \geq \quad (h \circ c)(x^{\mathrm{best}}) - \hat{G}^{\ell \mathrm{best}} \, {}^{\top} \Big( Dc(x^{\mathrm{best}})(x - x^{\mathrm{best}}) + o(|x - x^{\mathrm{best}}|) \Big) - \hat{e}_{\ell \mathrm{best}} \\ \qquad \geq \quad (h \circ c)(x^{\mathrm{best}}) - \mathsf{tol}_{\mathsf{stop}} |x - x^{\mathrm{best}}| - \mathsf{tol}_{\mathsf{stop}} + o(|x - x^{\mathrm{best}}|) \, .$$

# 6 Numerical experience

In order to assess from a practical point of view the Composite Bundle method, we coded Algorithm 1 in MATLAB and ran it on several collections of functions described in Sec. 2, using a computer with one 3GHz processor and 1.49GB RAM.

## 6.1 Solvers in the benchmark

We compared the performance of the Composite Bundle method with the MATLAB HANSO package, implementing a "Hybrid Algorithm for NSO", and downloadable from http://cs.nyu.edu/overton/software/index.html. As explained in [LO08], for nonsmooth optimization problems, BFGS may fail theoretically. However, the HANSO package can provide good benchmarks for the purpose of comparison, helping to shed some light on important NSO issues, as shown in our results below.

The package is organized in two phases:

- A first phase runs a BFGS method for smooth unconstrained optimization, with a linesearch capable of handling kinks explained in [LO08], and with multiple starting points. If the termination test is not satisfied at the best point found by BFGS, HANSO continues to the next phase.
- The second phase executes up to three runs of the Gradient Sampling method in [BLO02], starting from the lowest point found by BFGS, and with decreasing sampling radii. As initial information, the Gradient Sampling uses BFGS final bundle, corresponding to the last min[100, 2n, n + 10] generated points and their (bb) information. For locally Lipschitz functions, the method converges to Clarke critical points in a probabilistic sense; [BLO02].

We also created another hybrid algorithm, the *Hybrid Composite Bundle*, that after the BFGS phase switches to the Composite Bundle method, starting like HANSO from BFGS's final bundle.

Therefore, the benchmark considers the four solvers below:

- CBun, the Composite Bundle method,
- BFGS, the first phase in the HANSO package,
- Hanso, the hybrid variant combining BFGS and Gradient Sampling methods; and
- HYCB, the hybrid variant combining BFGS and CBUN.

### 6.2 Parameters for the different solvers

Letting n be the problem dimension, the maximum number of iterations and calls to the oracle were set to  $\mathtt{maxit} = 150 \min(n, 20)$  and  $\mathtt{maxsim} = 300 \min(n, 20)$ , respectively. We set the stopping tolerance  $\mathtt{tol}_{\mathtt{stop}} = 10^{-5}$  if n < 50 and multiply it by  $\sqrt{n}$  when  $n \ge 50$ .

### 6.2.1 Parameters for CBUN and HYCB

The two Armijo and Wolfe-like parameters are  $m_1 = 0.9$  and  $m_2 = 0.55$ , respectively. The minimum and maximum positive thresholds are  $\mu_{\min} = 10^{-6}$  and  $\mu_{\max} = 10^{8}$ . In all of our runs, only strongly active bundle elements are kept at each iteration, but, when there are more than  $|\mathcal{B}| = 50$  strongly active elements, the bundle is compressed.

If no second-order information is available for the inner mapping, in the variable prox-metric we take  $M_k = 0$ . Otherwise, we use Newton-type matrices

$$M_k = \sum_{j=1}^{m} \hat{G}_j^{\ell_k} \nabla^2 c_j(x^k) , \qquad (30)$$

exploiting sparsity patterns, if they exist. To update the prox-parameter, we use an estimate for  $\lambda_{\min}(M_k)$ , obtained by a modified Cholesky factorization.

The prox-parameter was started at

$$\mu_0 = \frac{5|\gamma(x^0)|^2 10^{-\text{fact}}}{1 + |(h \circ c)(x^0)|}$$
.

The parameter fact is 0 if the function is convex, and 3 otherwise (a large value for this parameter constrains the search of the next iterate in a region close to  $x^0$ , which makes sense for a nonconvex function). At later iterations, every time a serious step is declared, the update is done as follows:

$$\mu_{\ell+1} = \min(\mu, \mu_{\text{max}}) \text{ for } \mu = \begin{cases} \max(\mu_{\text{min}}, \mu_{\ell}, -1.01\lambda_{\text{min}}(M_k)) & \text{if } \lambda_{\text{min}}(M_k) < 0 \\ \max(\mu_{\text{min}}, \mu_{\ell}^{\text{qN}}) & \text{if } \lambda_{\text{min}}(M_k) = 0 \\ \max(0, \lambda_{\text{min}}(M_k) - \mu_{\ell}^{\text{nlcv2}}) & \text{if } \lambda_{\text{min}}(M_k) > 0. \end{cases}$$

In these relations,  $\mu_{\ell}^{qN}$  is computed using the reversal quasi-Newton scalar update in [BGLS03, § 9.3.3]:

$$\mu_{\ell}^{\text{qN}} := \min \Bigl\{ \frac{|\gamma^k - \gamma^{k-1}|^2}{(\gamma^k - \gamma^{k-1})^\top (x^k - x^{k-1})}, \text{ for } \begin{array}{l} \gamma^k \in \{Dc_k^\top \hat{G}^{\ell_k}, Dc_k^\top G^{\ell_k}\} \\ \gamma^{k-1} \in \{Dc_k^\top \hat{G}^{\ell_{k-1}}, Dc_k^\top G^{\ell_{k-1}}\} \end{array} \Bigr\},$$

recalling that  $\ell_k$  and  $\ell_{k-1}$  denote the last two iterations declaring a serious step ( $x^k$  and  $x^{k-1}$ , respectively).

Backtracking steps multiply the current prox-parameter by a factor of two. At consecutive null steps, the prox-parameter is defined as  $\mu_{\ell+1} = \max(\mu_{\ell}, \sqrt{(\mu_{\ell} + \lambda_{\min}(M_k)) \text{nul}})$ , for nul a counter of null steps. This update satisfies the convergence condition (24). In all cases, if  $\lambda_{\min}(M_k) < 0$ ,  $\mu_{\ell+1} \ge -1.1\lambda_{\min}(M_k)$ , and  $\mu_{\ell+1} \in [\mu_{\min}, \mu_{\max}]$ .

#### 6.2.2 Parameters for BFGS and HANSO

Both tolerances normtol and evaldist are set to  $\sqrt{\text{tol}_{\text{stop}}}$ . In the linesearch, Armijo's and Wolfe's parameters were taken equal to 0.01 and 0.5, respectively. The option for a Strong Wolfe's criterion is not activated, and the method quits when the linesearch fails. As for the quasi-Newton updates, they are of the full memory type, with scaling of the initial Hessian. Finally, the *final bundle* keeps  $\min[100, 2n, n+10]$  past gradients.

#### 6.3 Benchmark rules

In an effort to make comparisons fair, we adopted the following rules:

- All solvers use the same black-boxes.

- Each solver has a different computational effort per iteration, which depends not only on the solver, but also on how many times the black-box is called per iteration. The number of iterations is not a meaningful measure for comparison, and the number of (bb) calls is different for each solver. While each iteration of Hanso calls (bbc) and (bbh), the Composite Bundle calls (bbh) at Step 2, and can call (bbc) and (bbh) at Step 3. Moreover, Hanso always uses gradient values, but the Composite Bundle method requires additional first order information for the inner mapping only at a serious step, to define a new inner model. Striving for fairness, we defined the counters below:
  - -For Hanso, one call to both  $(bb_c)$  and  $(bb_h)$  counted as 1.
  - -For the Composite Bundle method, we kept separate counters for the number of c, Dc, and h/G evaluations: nc, nDc, nf, respectively. To each counter corresponds a number of scalars needed for the corresponding calculation: m, mn, and 1 + m, respectively. This gives the number of scalars required for a single evaluation of  $h \circ c$  and a subgradient (as in HANSO):

$$OneEval = m(1+n) + 1 + m$$
.

Accordingly, the number of (bb) calls in the Composite Bundle method was defined as

$$\frac{\mathtt{nc} m + \mathtt{nDc} m n + \mathtt{nf}(1+m)}{\mathtt{OneEval}} \,.$$

At this point a potential advantage of CBun becomes clear: when for a given function the method makes many consecutive null steps, this is not expensive in terms of (bb) calls (only (bb<sub> $\hbar$ </sub>) is needed).

- We also computed total CPU time for each solver to reach its stopping test. This information should mostly be taken as a complement to the counter of (bb) calls. The reason is that CPU times can be misleading, because for almost all the tested functions the (bb) calls take a negligible time, a feature that is rare in real-life problems. For example, for the nuclear generation planning problems in [ES10, Sec. 7], 95% of the total CPU time is spent in producing the black-box information. By contrast, in all of ours runs and for all but one instance, (bb) calls represent less than 10% of the total CPU times (taking up to 40% for one test-function, written in Fortran and requiring a mex-interface).
- All solvers use the same quadratic programming packages, for which there are two possibilities: a Fortran library with a mex-interface, or a special MATLAB solver. Quadratic programs like (9) can be solved by QUADPROG, the built-in MATLAB QP solver, but we prefer the method in [Kiw86], and made a mex-interface for the Fortran code developed by the author. This method is specially tailored for quadratic minimization over a simplex, as it is the case for the problem dual to (9) and, hence, often outperforms QUADPROG, which is a general solver. Hanso also has a special QP solver, written in MATLAB. Since Hanso QP problems amount to setting  $M_k^{\ell} \equiv 0$  in

$$0 \in conv\{G^i : i \in \mathcal{B}_\ell\} + M_k^\ell d^\ell,$$

which is the optimality condition for (9), we could modify Hanso QP solver to handle (9). Reciprocally, it is possible to use the Fortran QP solver in [Kiw86] to solve Hanso QP problems.

- For large-scale instances, Hanso offers a limited memory BFGS method explained in [Ska10]. However, since for the prox-variable metric the calculation of  $M_k$  corresponds to a Newton method, the limited-memory option was not activated in the comparisons.
- Each solver has specific stopping tests, and since BFGS uses a smooth method, the triggers terminating the runs are only heuristical. For each solver we declare a run a success when:
  - -For BFGS and HANSO, when the tolerance on the smallest vector in the convex hull of certain subgradients is met (for BFGS, these are the subgradients in its *final bundle*).
  - -For CBun and HyCB, when the stopping test in Step 2 is reached.
- All non-successful runs are declared failures, with a special counter for when a solver reached the maximum number of iterations or calls to the black-box (maxit or maxsim, denoted by max in the tables). Possible reasons for failures are detection of unboundedness in x or in the function values, errors in the QP solver, and, for BFGS, a nondescent direction, or a problem in the linesearch.
- The hybrid variants are not initiated if BFGS detected unboundedness. They are started when BFGS succeeds, reaches maxit, cannot descend from the generated direction, or if the linesearch or the QP solver failed.

- Both BFGS and the Composite Bundle use the same starting points, each function was run for 10 different starting points. Unless otherwise specified, the starting points have all components randomly drawn in [-1, 1].
- To measure the accuracy reached by each solver, we only considered test-functions with known optimal value  $\bar{f}$ , or such that all the solvers converged to the same final value within a tolerance of  $10^{-5}$ . In this case, the optimal value  $\bar{f}$  is given by the smallest value function found by all the solvers. Letting  $f^{\text{best}}$  denote the function value of the analyzed case, then

$$RA := -\log_{10} \max \left( 10^{-16}, \frac{f^{\text{best}} - \bar{f}}{1 + |\bar{f}|} \right)$$

measures the number of digits of accuracy achieved by the solver.

- We exclude from the tables with results those nonconvex cases for which different solvers found different critical points, see Table 10 in the appendix for details.
- Since full tables are large, for the reader's convenience in this section we only reproduce the overall results of each full table in the Appendix.

## 6.4 Battery of convex test problems

We first consider typical functions for convex NSO benchmarking, described in Table 1.

Name	n	$\overline{f}$	Reference
Maxquad	10	-0.84140833459641	[BGLS03, p. 131]
BadGuy	10	-2048	[HUL93, p. 277, vol.II Ex.XV.1.1.2]
TR48	48	-638565	[HUL93, p. 21, vol.II, Ex.IX.2.2.6]
TSP proble	ems in [	HUL93, p. 22, vol.II, Ex.	IX.2.2.7], data from TSPLIB95
TSP	29	-9015	bayg29
TSP	442	-50500	pcb442
TSP	1173	-56349	pcb1173
TSP	3038	-136601	pcb3038
Convex Ur	y proble	ems, low dimension	Ex. 5 with cubic=0
Ury-cvx	10	500	
Ury-cvx	20	911.833349450300716	
Ury-cvx	30	1118.219919518173128	

Table 1: Convex problems in Group 1

The composite structure  $h \circ c$  for functions Maxquad, Ury is the one described in the respective examples. Both TR48 and TSP are the piecewise maximum of affine functions, so the inner mapping is affine. However, our Fortran (bb) code for TSP was too involved to identify the c-components (corresponding to the *minimal 1-trees* in the underlying graph), so we just took the trivial composite case for TSP, as in Remark 1. Similarly for BadGuy, because it does not have a positively homogeneous outer mapping.

Table 2 summarizes the results from Table 8 in the Appendix, obtained for Group 1. In Table 8, each column corresponds to one of the four solvers, CBun, BFGS, Hanso, HyCB, run with the Fortran or Matlab special QP solver (denoted by f and m, respectively). Since for all of our runs we observed that the Matlab QP special solver is systematically less efficient and/or less reliable than the Fortran one, the shorter Table 2 contains the indicators for the Fortran variants only: CBunf, BFGSf, Hansof, HyCBf.

To each function and each case in Table 1 corresponds a row in Table 8, reporting the figures obtained with each solver, averaged over ten runs, with random starting points. All solvers used the same starting points, with components in [-1,1] except for BadGuy, taken in [-512,512]. Results are displayed in three columns, with the accuracy RA, the mean CPU time in seconds  $\overline{sec}$ , and the average number of black-box calls  $\overline{(bb)}$ , respectively. At the end of each function there are two lines. A first line with the mean values mean averaged for all the considered cases (in this line the second column reports between parenthesis

the considered number of runs). The second line gives the number of failures and how many of these failures corresponded to having reached the maximum allowed for iterations or simulations (max/fails). Finally, the bottom in Table 8 contains the same indicators, averaged over all the problems in the group, as well of the total number of instances considered for the test. These bottom lines and the average of each function are reproduced in Table 2 as a summary.

(bb)	Group 1		CBunf			BFGSf			Hansof			HyCBf	
	(mean over)	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(\mathtt{bb})}$	RA	sec	$\overline{(\mathtt{bb})}$
BadGuy	(10)	8	1.70	135	9	1.92	289	9	2.53	377	9	2.00	301
$_{ m MQ}$	(10)	11	0.16	40	8	0.20	432	8	0.32	683	12	0.24	443
TR48	(10)	12	3.51	48	13	32.57	2738	13	290.53	29439	15	33.16	2746
TSP	(40)	8	102.05	1045	13	1946.97	8188	13	1986.19	18050	14	2014.92	8686
Ury	(30)	11	1.82	45	9	2.64	1469	11	20.03	17762	15	2.72	1475
MEAN	(100)	10	22	262	10	397	2623	11	460	13262	13	411	2730
max/fa	ails		1/1			3/38			0/38			0/0	

Table 2: Group 1 overall results: problems in Table 1.

For Group 1, we observe that all methods are very accurate. When compared to CBun, BFGS exhibits a significant increase both in CPU times and number of (bb) calls; failing to trigger its termination criteria 38% of the times (only thrice the reason was max). For the TSP family, BFGS is 60% more accurate than CBun, but for getting 5 more digits, BFGS spends 8 (20) times more CPU ((bb) calls) than CBun does. We conjecture that endowing TSP with a nontrivial composite structure can improve CBun's average figures (we observed a significant change for TR48, when comparing CBun performance on TR48 black-boxes with and without composite structure). Hanso is the slowest solver, and uses the most (bb) calls, but it is not the most accurate method: this hybrid variant did not seem to be adequate for this set of problems, probably because they are all convex. By contrast, HyCB eliminated all of BFGS's failures, with a relatively low additional computational effort: HyCB extra CPU times and (bb) calls represent less than 5% of BFGS totals. With respect to CBun, HyCB gain of 30% in accuracy is obtained at the stake of an increase of almost 2000% and 1000% in CPU times and (bb) calls. For this group of problems, CBun performs better than all the other solvers.

### 6.5 Convex and nonconvex problems

The two other groups include a mix of convex and nonconvex problems given in Table 3. Group 2 gathers problems with low dimensions ( $n \le 50$ ), while Group 3 contains high dimensional problems ( $n \in \{100, 500\}$ ). For each instance in Groups 2 and 3, we give the optimal values when known, or the lowest function value found by all solvers, in Table 10 in the Appendix.

For the MQ, EucSum, TiltedNorm, and CPS-collections, matrices and vectors were generated randomly. All A-matrices are symmetric positive semidefinite, with condition number equal to  $range A^2$ . The B-matrices in CPS are symmetric positive semidefinite and condition number equal to  $n^2$ . To make calculations possible in our computer, for all the sparse matrices the density was set to 0.1, 0.01, 0.01, 0.001 for n = 10, 50, 100, 500, respectively.

Tables 4 and 5 summarize the results for Group 2 and 3, respectively. In the Appendix, Tables 11 and 12 report the respective full details.

For Group 2 the instances excluded because different solvers found different critical points correspond to two variants of the functions NK and LV, namely F8 and T3 in [Ska10]; and GenModRos with starting point in [0,2] and [-2,2].

For this group, the overall results show again that CBun performs better in average that the other three solvers. However, BFGS did better for some instances of LV, as well as for GenModRos and ModRos. CBun had difficulties solving the second instance of LV, corresponding to T3 in [Ska10] with n=10. For the excluded instances of GenModRos, the optimal value (5.337) was found often by BFGS while CBun found only a critical point (with value 9.3283). For the function NesChebRos, nine out of the thirty starting points very difficult to handle by BFGS, but not by CBun, explaining the huge

Name	Parameters	Reference
CPS	$n \in \{10, 50, 100, 500\}$	[Ska10, Sec.4.2.2]
CVX	$range A \in \{0.2, 0.8\}n$	$f(x) = \sqrt{x^{\top} A x} + x^{\top} B x$
MQ	$n \in \{10, 50, 100, 500\}$	Ex. 1,with
CVX	$range A \in \{0.2, 0.8\}n$	$c_j(x) = \frac{1}{2}x^{\top}A_jx + b_j^{\top}x$
	m = range A + 3	$\{A_j\} \geq 0$ and $\{b_j\}_{j=1}^{range\ A}$ LI
EucSum	same than MQ, but $n \in \{4, 10, 50, 100\}$	Ex. 4 with the $\ell_1$ -norm, $m_j = 1$ , $J = m$ , and
NCV		$\phi_j = c_j \text{ from MQ}$
TiltedNorm	$n \in \{10, 50, 100, 500\}$	[Ska10, Sec.4.2.1]
CVX	w = 4	$f(x) = w Ax  + (w-1)e_1^{\top}Ax$
GenModRos	n = 12	[Ska10, Sec.4.2.4]
NCV	U = 1, V = 10	$f(x) = \sum_{i=1}^{n-1} \left( V_{\frac{i}{n}}   x_{i+1} - x_i^2 / n   + U_{\frac{i}{n}} (1 - x_i)^2 \right)$
	$x^{0} \in \{[-0.1, -0.1], [-1, 1], [-2, 0], [0, 2], [-2, 2]\}$	,
ModRos	n=2	[LO08, Sec. 5.7]
NCV	$w \in \{1, 2, 4, 8\}$	$f(x) = w x_2 - x_1^2  + (1 - x_2)^2$
NesChebRos	$n \in \{5, 10, 50, 100\}$	[LO08, The nonsmooth variation in Sec. 5.8]
NCV	$x^0 \in [0, 2]$	$f(x) = \sum_{i=1}^{n-1}  x_{i+1} - 2x_i^2 + 1  + 0.25(x_1 - 1)^2$
Ferrier	$n \in \{10, 50, 100\}$	[HS10]
	- (4, 0)	$\sum_{i=1}^{n}  ix_i^2 - 2x_i + \sum_{i=1}^{n} x_i $ if $case = 1$
NCV	$case \in \{1, 3\}$	$f(x) = \begin{cases} \sum_{i=1}^{n}  ix_i^2 - 2x_i + \sum_{j=1}^{n} x_j  & \text{if } case = 1\\ \max_{i=1}^{n}  ix_i^2 - 2x_i + \sum_{j=1}^{n} x_j  & \text{if } case = 3 \end{cases}$
NK	$n \in \{10, 50, 100, 500\}$	Problems Fcase in [Ska10, Sec.5.4.2]
CVX/NCV	$case \in \{1, 3, 4, 5, 8, 9\}$	see also [HMM04, Sec. 3]
LV	$n \in \{10, 50, 100, 500\}$	Problems Tcase in [Ska10, Sec.5.4.3]
NCV	$case \in \{3, 4, 5, 6\}$	see also [LV00]
Ury	$n = \in \{10, 20, 30, 100\}$	Ex. 5
CVX/NCV	cubic $= \in \{0, 0.01\}$ if $n = 100$ , cubic=0.01 otherwise	

Table 3: Convex and nonconvex problems in Groups 2 and 3

difference in accuracy obtained by such solvers. For these problems, we observed that BFGS got stuck at a nonoptimal kink and exited having triggered the heuristical stopping test (the projection of zero on its final bundle was smaller than the tolerance). By contrast, the Rosenbrock modifications GenModRos and ModRos put CBun into trouble: these are the only problems for which CBun systematically makes more (bb) evaluations than BFGS. For these functions, CBun finds a very precise minimizer, after taking many serious steps (very short ones); since for each new serious step the mapping Jacobian  $Dc(\hat{x}^k)$  is computed, this significantly increases the total (bb) counter. Finally, and as observed for Group 1, HyCB seems to be a better hybrid variant than Hanso.

(bb)	Group 2		CBun:	f		BFGS	f		HANSO	f		НуСВ	f
	(mean over)	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(\mathtt{bb})}$
CPS	(60)	10	0.06	10	6	0.22	348	6	0.31	470	8	0.24	352
EucSum	(50)	9	0.72	17	6	0.80	574	6	2.17	3708	14	0.85	578
Ferrier	(40)	6	0.46	21	4	0.55	777	4	0.93	1775	7	0.60	783
$\operatorname{GenModRos}$	(30)	10	1.52	839	8	0.23	472	8	0.99	1821	14	0.31	514
LV	(60)	8	1.42	334	11	1.91	1292	11	8.47	6355	13	1.94	1297
MQ	(40)	9	0.14	11	7	3.79	2212	7	20.00	9611	9	3.90	2217
ModRos	(40)	6	0.12	130	6	0.03	104	6	0.04	124	9	0.17	260
NK	(100)	11	0.51	28	8	0.73	577	8	6.08	6195	14	0.86	589
NesChebRos	(30)	16	0.20	71	1	0.70	775	2	8.77	10335	2	1.15	1021
TiltedNorm	(30)	11	0.99	62	7	0.12	434	7	0.35	1132	10	0.14	439
Ury	(30)	12	1.95	58	11	2.67	1444	11	19.78	19460	14	2.99	1456
MEAN	(510)	10	0.73	143	7	1.07	819	7	6.17	5544	10	1.20	864
max/fails			29/29			9/76			0/60			23/23	

Table 4: Group 2 overall results: problems in Table 3,  $n \le 50$ .

For Group 3 the instances excluded because different solvers found different critical points correspond

to four variants of the functions NK and LV, namely F8 and T3, T5, and T6 in [Ska10]; and EucSum with n = 100 and  $range A_i = 400$ .

(bb)	Group 3		CBunf			$\operatorname{BFGSf}$			Hansof			HyCBf	
	(mean over)	RA	sec	$\overline{(\mathtt{bb})}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(\mathtt{bb})}$	RA	sec	$\overline{(\mathtt{bb})}$
CPS	(60)	6	0.70	7	5	83.31	1945	5	106.62	2925	7	83.64	1948
EucSum	(10)	16	0.27	8	5	0.78	433	5	1.08	734	6	0.84	435
Ferrier	(20)	4	6.54	64	10	5.83	3203	10	9.41	10437	10	5.89	3207
LV	(40)	4	3.97	65	13	101.84	8856	16	232.08	29816	13	102.08	8860
MQ	(40)	4	1398.41	40	3	3951.10	6646	9	20483.63	33343	8	5350.40	6657
NK	(100)	9	6.49	19	10	85.94	5189	10	195.28	21102	14	87.44	5193
NesChebl	Ros (10)	16	0.84	19	1	29.63	5247	1	112.88	44548	1	29.73	5251
TiltedNo	rm (10)	3	7.84	105	7	1.59	1372	7	2.04	2776	7	1.62	1374
Ury	(20)	0	48.02	92	12	61.73	5770	12	273.85	45070	16	91.74	5836
MEAN	(310)	7	163.67	47	7	480.19	4295	8	2379.65	21195	9	639.26	4307
${\tt max}/{\rm fails}$			29/29			21/53			0/47			15/15	

Table 5: Group 3 overall results: problems in Table 3, n = 100 and n = 500.

As expected, this higher dimensional group is more difficult for all the solvers. The low mean  $(\overline{bb})$  for CBun indicates that the methods often stalled making many null/backtracking steps, rather than serious steps. However, the second instance of Ferrier functions (corresponding to outer function  $h(\cdot) = \max(\cdot)$  and n = 100) was difficult for CBun, which made many short serious steps, expensive in terms of (bb) calls. Problem MQ with n = 500 and  $range\ A_j = 400$  was very difficult to solve by all methods. The 100 runs of the NK family did not seem difficult for any solver. CBun exited problem TiltedNorm having reached the maximum number of iterations, while BFGS found a good point and triggered its heuristical stopping test. For Ury, both CBun and BFGS ended by having reached the maximum number of iterations, but BFGS final point is much better than the one found by CBun. For NesChebRos, BFGS fails in the linesearch, stuck at a nonoptimal kink, and none of the hybrid variants succeeds in getting out of the ridge. For this group, and especially for the nonconvex functions, we see a more erratic behaviour of CBun, even though it still has the best indicators in mean.

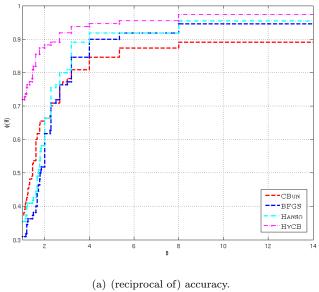
#### 6.6 Performance Profiles

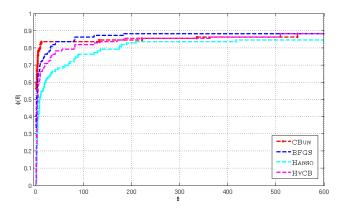
Figures 2(a), 2(b) and 2(c) contain performance profiles over all the 920 runs, excluding cases converging to different critical points, but including failures, like in the tables of results. This choice was done not to handicap BFGS, whose heuristical stopping test may fail to be triggered sometimes.

Each line in a performance profile can be interpreted as a cumulative probability distribution of a resource of interest: accuracy, CPU time, (bb) calls. Usually, "smaller" values mean "better performance" of the consider resource. Therefore, for both accuracy and CPU time we plotted the reciprocal of the figures obtained by each solver. In this manner in all the profiles below, the solver with the highest line is the best one for the given indicator of performance. In each profile, we look at the following points:

- The leftmost abscissa values, indicating for each solver the percentage of runs for which each solver had the best performance for the given indicator. The highest value corresponds to the best solver.
- The abscissa of the intersection between two lines gives the factor that makes the respective solvers comparable.
- For an abscissa value  $\theta$  with ordinate  $\phi(\theta)$ , the value  $1 \phi(\theta)$  corresponds to the fraction of problems that a solver cannot solve within a factor  $\theta$  of the best one.

The first profile, in Figure 2(a), shows the performance in terms of accuracy. Looking at the highest value for the leftmost abscissa, we conclude that the hybrid variant HyCB is the most precise solver in 72% of the runs. CBun, Hanso, and BFGS are the most accurate solvers in 37%, 36%, 31% runs, respectively. The abscissa values of 2 show that CBun, BFGS and Hanso failed in 34%, 48% and 34% of the runs to achieve half of HyCB's precision (the respective ordinate values are  $\phi(2) = 0.66, 0.52, 0.66$ ).





(b) (reciprocal of) CPU time.

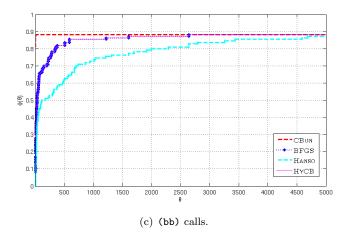


Figure 2: Performance Profiles.

Profile 2(b) measures the performance in terms of CPU time in seconds, and shows that CBuN is the fastest solver in 56% of the runs, followed by BFGS, which was fastest in 33% of the runs. At  $\theta = 40$ , the solvers ordinates are  $\phi(40) = 84, 84, 77, 66$ . This means that BFGS is as fast as CBuN, within a factor of 40, for 84% of the runs, and takes more than 40 times CBuN's CPU time in 16% of the runs. HANSO and HYCB fail to take less than 40 times CBuN's CPU time in 23% and 34% of the runs, respectively.

The final profile, in Figure 2(c), measures the performance of the different solvers in terms of (bb) calls, and shows a clear superiority of CBun, which appears as the most economic solver in 80% of the cases. Hanso makes extensive use of (bb) calls, so it should mostly be used for unstructured nonconvex functions that are not too difficult to evaluate (possibly like the matrix problems in [Ska10, Sec. 5.3]). The lines of BFGS and HyCB practically coincide, making both methods indistinguishable in terms of (bb) calls. Since BFGS is faster and HyCB is more precise, the choice between these two solvers should be driven the user's preference (speed or accuracy), keeping in mind that HyCB is more reliable in terms of stopping test.

## 6.7 Determining V-dimension

Many composite functions are partly smooth [Lew02], a notion that generalizes to the nonconvex setting the VU-space decomposition for convex functions in [LOS00] and [MS00].

Identification of the  $\mathcal{V}\mathcal{U}$  subspaces can be used to determine directions along which the function behaves smoothly and, hence, identify a region where a Newton-like method is likely to succeed. Such smooth directions lie in the  $\mathcal{U}$ -subspace; its orthogonal complement, the  $\mathcal{V}$ -subspace, concentrates all the relevant nonsmoothness of the function, at least locally. Near a critical point  $\bar{x}$ , the  $\mathcal{V}$ -subspace is spanned by the subdifferential  $Dc(\bar{x})^{\top}\partial h(\bar{C})$ , with  $\bar{C}=c(\bar{x})$ , and the  $\mathcal{U}$ -subspace is the orthogonal complement of  $\mathcal{V}$ . Alternatively, in the wording of [Lew02], the  $\mathcal{U}$ -subspace is the subspace tanget to the *smooth or activity manifold* at  $\bar{x}$ , and  $\mathcal{V}=\mathcal{U}^{\perp}$ .

In [LO08] and [Ska10] it is observed that BFGS can retrieve VU-information by analyzing the eigenvalues of the inverse Hessian used to define a new iterate. For comparison purposes, we consider CBuN and BFGS only and estimate the dimension of the respective generated V-subspaces as follows:

- For CBun we compute the range of the subspace spanned by the final strongly active gradients in the bundle:

$$\dim \mathcal{V}_{\text{CBun}} := range \left\{ Dc(x^k)^\top (G^i - \hat{G}^\ell \,) : i \in \mathcal{B}_\ell \text{ with } \alpha_i^\ell > 0 \right\} \,,$$

for  $x^k$  the last generated serious step and  $\ell$  the iteration triggering the stopping test; recalling that in Step 4 of Algorithm 1 the bundle sizes are kept controlled by a parameter  $|\mathcal{B} \max|$ ,

- For BFGS we count how many eigenvalues of the final inverse Hessian H cluster near 0:

$$\dim \mathcal{V}_{\text{BFGS}} := \operatorname{card} \left\{ i \leq n : \frac{\lambda_i(H)}{\lambda_{\text{max}}(H)} \leq \epsilon \right\} \,,$$

for  $\epsilon$  a given tolerance.

	BadGuy		Euc	Sum		Maxquad		M	Q		TR48
n	10	10	10	10	10	10	10	10	10	10	48
$dim \ \mathcal{V}$	10	8	6	4	2	3	8	6	4	2	??
CBun	10	10	8	6	4	4	8	7	4	3	47
BFGS	2.3	8	3.8	3.4	2.6	3.2	8	6.1	5	7	47

Table 6: V dimensions for BadGuy, EucSum, Maxquad, MQ, and TR48.

Table 6 reports the obtained results for some of the problems in Group 1 and 2, with low dimension and  $\mathcal{V}$ -dimensionality depending on the case. The parameters setting were  $|\mathcal{B} \max| = 50$  and  $\epsilon \in \{0.1, 0.01\}$  (which gave identical results for this group of runs). Each problem was run 10 times with random starting points

Even though the rules adopted for determining the V-dimensions are rather rough, both CBuN and BFGS estimations are reasonable, with a few exceptions. For both solvers the worst results are those obtained for the nonconvex EucSum functions. The extremely low V-dimension estimated by BFGS

for BadGuy comes from the fact that it is hard to automatically separate almost null from nonnull, yet small, eigenvalues. An aposteriori (visual) examination of the eigenvalues obtained for each starting point shows a rather erratic behaviour of BFGS for this function over the different starting points, even though BFGS's heuristic stopping test was always triggered. Such oscillation could be explained by a lack of stability of the Hessian with respect to small perturbations, a common phenomenon for a nonsmooth function near a kink.

We made a second group of runs, to determine the impact of smaller or larger  $\mathcal{V}$ -dimension, with respect to the dimension of the full space. We considered the CPS function, with dimension  $n \in \{10, 50, 100\}$  and varying  $\mathcal{V}$ -dimension. For this example, the  $\mathcal{V}$ -dimension coincides with the range of the matrix A (taken with sparse density equal to 0.1 for all cases).

Table 7 reports the  $\mathcal{V}$ -dimensions estimated by CBun and BFGS for different parameters. For CBun, the maximum bundle size was set to 50 and 100: we expect results to be worse if  $|\mathcal{B} \max| < \dim \mathcal{V}$  and the bundle needs to be compressed to an insufficient number of elements. For BFGS, we took two values of  $\epsilon$ , as before. In the table, the parameters appears between parentheses after the name of each solver.

										CP	S								
n	10	10	10	10	50	50	50	50	50	100	100	100	100	100	100	100	100	100	100
$dim  \mathcal{V}$	8	6	4	2	40	30	20	10	2	90	80	70	60	50	40	30	20	10	2
CBun(50)	8	6	4	2	40	28	20	10	2	8	46	36	44	44	30	29	20	10	2
CBun (100)	8	6	4	2	40	30	20	10	2	77	78	70	56	50	40	30	20	10	2
$\mathrm{BFGS10^{-2}}$	8	6.2	4.4	2.6	44	39.5	35.5	35.5	36	95.5	91	86.5	88	86	85	86	85.5	86.5	86
$\mathrm{BFGS10^{-3}}$	8	6	4	2	43	37	31	24	19	93	86	81	80	75	71	67	61	63	64

Table 7:  $\mathcal{V}$  dimensions for CPS.

We observe that problems with larger V-subspace are more difficult for both solvers. In general, CBun(100) seems to give a reasonable estimate, but this is not always true, especially when n = 100.

We conclude our analysis with Figure 3, with the real and estimated  $\mathcal{V}$ -dimensions for the 30 different functions, showing that in general BFGS overestimates the size of the  $\mathcal{V}$ -space. We emphasize that this set of tests determining  $\mathcal{V}$ -dimensionality is only preliminary, and rather crude. For this reason, the conclusions above should not be taken as an indication of goodness or badness of a solver. The subject of determining  $\mathcal{V}\mathcal{U}$  subspaces is still rather unexplored, with a few exceptions in [LO08], [Ska10], and the MQ functions considered in [DSS09].

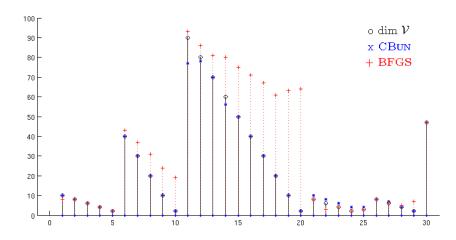


Figure 3: True and estimated V-dimensions for the 30 functions.

# Concluding Remarks

The composite bundle method presented in this work makes tractable the algorithm ProxDescent in [LW08] for a large class of composite functions, with real-valued, positively homogeneous, and convex outer functions. In particular, the method can be applied to minimize some nonconvex nonsmooth functions, a challenging issue for bundle methods. Our composite cutting-planes model, approximating the conceptual model, avoids typical pitfalls in nonconvex bundle methods (recall Figure 1).

The numerical experience reported in this work shows the good performance of the Composite Bundle method for problems of moderate size. For large dimensions, the use of variable prox-metrics may increase the solution times too much, even if there are sparse patterns to exploit. The impact of such increase is problem dependent: for some functions (like CPS and TSP) there is a clear advantage in applying a bundle method ( $n \le 500$  in CPS, and  $n \le 3038$  in TSP, but  $M_k \equiv 0$ ). The advantage is less clear for other functions, especially some of the nonconvex ones in Group 3. Since we sometimes also observed that too many short serious steps made CBUN stall, we conjecture that a linesearch (replacing or complementing the curved search modifying  $\mu_\ell$ ) can improve the performance of Algorithm 1 for nonconvex functions, but this is a subject of future research.

Although BFGS is accurate and fast (at least for our examples, with blackboxes computationally light), neither BFGS nor HANSO appeared as the best alternatives for many classes of functions considered in our runs. However, conclusions can be different for a different set of test-functions. Also, the usefulness of a solver depends on the specific purpose sought by the user: since BFGS descends fast from a starting point, it could be an interesting alternative if not much accuracy is required, or if the user seeks for a "better" point, without caring if it is the best one. For some problems, we observed that BFGS got stuck at a nonoptimal kink and exited having triggered the heuristical stopping test (the projection of zero on its final bundle was smaller than the tolerance). If reliability is a concern, the output of BFGS can be plugged into a bundle method, as in HyCB, to satisfy a theoretical stopping test.

However, if too much accuracy is desired, the hybrid variant is likely to increase the computational bulk of BFGS too much (at least when compared to applying directly CBun). As for Hanso, since it makes extensive use of (bb) calls, we think it should mostly be used for unstructured nonconvex functions that are not too difficult to evaluate (possibly like the matrix problems in [Ska10, Sec. 5.3]).

Another important issue to consider for a heavy duty application is that, even in the presence of a composite structure, the resulting smooth mapping may be large, or have no special second order sparse patterns to exploit. In this case, it can be sound to use null or diagonal matrices  $M_k$  in Algorithm 1, or apply the limited memory variants in [HMM04] and [Ska10].

We mention the work [KBM10], comparing several NSO general purpose solvers for different type and size of problems, as well as for different (bb) available information. Table 3 therein, analizing the efficiency and reliability of the considered solvers, can be useful as a complement of information to the conclusions drawn in our numerical results, keeping in mind that solvers are different and that the test-functions are not exactly the same, although there is some intersection.

Comparison with [NPR08]. The proximity control (nonconvex) bundle algorithm [NPR08] considers models for functions such that several Clarke subgradients at one point can be computed at reasonable cost. The proposed scheme is fairly general and bears some resemblance to our composite approach, which we explain next.

Instead of assuming that the objective function enjoys some particular structure, in [NPR08] the authors suppose there is available certain local model,  $\phi(\cdot, x^k)$ , for the objective function f at the current iterate  $x^k$ . In our notation,  $f = h \circ c$ , and the local model is  $\phi(x^k + \cdot, x^k) = h(c_k(\cdot))$ . As explained in [NPR08, Rems. 2.9 and 6.3], such composite model is both a strong and strict first-order model for f. The local model is approximated by a working model,  $\phi_k(\cdot, x)$ , which would correspond to our composite cutting-planes model,  $\check{h}_{\ell}(c_k(\cdot))$ . However, in [NPR08, Def. 3.3], the first-order working model satisfies the conditions  $\phi_k(x,x) = (h \circ c)(x)$  and  $\partial_1 \phi_k(x,x) \subset \partial_1 \phi(x,x)$ . For our composite model, this would mean to require that

$$\check{h}_{\ell}(c(x^k)) = (h \circ c)(x^k)$$
 and  $conv\{G^i \text{ for all } i \in \mathcal{B}_{\ell} : \check{h}_{\ell}(c(x^k)) = G^{i \top}c(x^k)\} \subset \partial(h \circ c)(x^k)$ ,

which only holds in our case if the outer subgradient information for  $C = c(x^k)$  was kept in the bundle.

Another related important difference is that when building the working model, in addition to (15), (22), and (23), the method in [NPR08] needs to incorporate certain *exact* cutting planes. In our setting, an exact cutting plane would correspond to requiring that

$$\forall \ell \geq 1$$
, given some  $\gamma \in \partial(h \circ c)(x^k)$ ,  $(h \circ c)(x^k) + \gamma^\top \leq \check{h}_{\ell}(c_k(\cdot))$ .

Instead, in (13) we use the outer subgradient  $\Gamma^{\ell} \in \partial h(C^{\ell})$  for  $C^{\ell} = c(x^k + d^{\ell})$  to detect if the linearization of the inner mapping is not good enough and trigger the backtracking process. But if  $x^k + d^{\ell}$  is declared a serious step, the corresponding subgradient  $\Gamma^{\ell}$  does not enter the bundle (nothing prevents the bundle management step to incorporate this data, though).

Like ours, the quadratic programming subproblem in [NPR08] includes a second-order term with a possibly nonpositive definite matrix  $M_k$ , augmented by a (positive enough) matrix  $\mu_{\ell}I$ . But the acceptance test corresponding to (12) uses a predicted decrease that is different from ours. Namely, instead of  $\delta_{\ell}$  in (11), the larger amount  $\delta'_{\ell} := \delta_{\ell} + \frac{1}{2}\mu_{\ell}d^{\ell} \,^{\top}d^{\ell}$  is used. As for null steps, the decision on whether or not to increase the parameter  $\mu_{\ell}$  is done by checking if, for some parameter  $m_3 \in (m_1, 1)$ ,

$$h(c(x^k) + D_k d^{\ell}) \le (h \circ c)(x^k) - m_3 \delta'_{\ell}$$

(the prox-parameter is left unchanged if the inequality above does not hold).

Convergence results for the proximity control bundle method with strong first-order models are similar to ours. The method keeps matrices  $M_k$  bounded below and above by  $\pm q I$  for some  $0 < q < +\infty$ , so (18) always holds. The case of infinite null steps is treated in [NPR08, Lem. 4.1], where it is shown that a subsequence of the prox-parameter diverges and, hence, checking satisfaction of our condition (24) is not straightforward.

Section 7 in [NPR08], devoted to applications, contains several cases showing the good numerical behaviour of the algorithm. An interesting subject of future research would be to compare both algorithms performance on composite objective functions.

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Appendix with full tables

(bb) $\#$ - $n$ CBUNf		CBunf		CBUNm	п	Щ	BFGSf		B	BFGSm		1	HANSOf		HANSOm	_	HY(	HyCBf		HYCBm	m s
	RA,	sec,	(bb) RA,		sec, (bb) RA,	Α,	sec,	(bb) RA,	Α,	sec,	(bb) RA	ιA,	sec,	(bb) RA,	sec,	(bb) R.A.		sec, (1	(bb) RA,	sec,	(qq)
BadGuy 1-10 8	×,	1.70,	135	70, 135   1, 24.32,1501   9,	501	9,	1.92,	586	9	1.94,	586	9	2.53,	377 9,	2.52,	378	9, 2.00		301 9,	22.12, 1793	1793
mean, (10)	×,	1.70,	135	.70, 135  1, 24.32,150	501	9,	1.92,	586	9,	1.94,	289	9,	2.53,	377 9,	2.52,	378	9, 2.00		301 9,	22.12, 1793	1793
max/fails		0/0	-	10/10			0/0			0/0			0/0		0/0		0	0/0		10/10	
MQ 1-10	11,	0.16, 4	40 1	40 11, 1.20, 62  8	62	œ,	0.20,	432	∞ ∞	0.77,	441	×,	0.32,	683 8,	1.66,	1451 12	_	0.24, 44	443 12	0.84	452
mean, (10)	11,	0.16,	40 111,		62 8	œή	0.20,	432	× ×	0.77,	441	ò	0.32,	683 8,	1.66,	1451 12		0.24, 44	443 12,	0.84	452
max/fails		0/0		0/0			0/1			0/1			0/0		0/1		0,	0/0		0/0	
TR48 1-48	12,	3.51,	48	48 3,133.67, 196 13	196[13		32.57, 2	2738 13		46.02,	2738 13		290.53,29439   13,	439 13	232.43,20537 15	0537 11	5, 33.16,		2746 13	115.59,	2816
mean, (10)	12,	3.51,	48	48 3,133.67, 196 13	1961		32.57, 2	2738   13		46.02,	2738 13,	3, 2	290.53,2943913,	43913,	$232.43,20537 \boxed{15}$	0537 1	5, 33.16,		274613,	115.59,	2816
max/fails		0/0		10/10			0/10			0/10			0/10		0/10		0	0/0		4/4	
1-29 16,		0.11,	20	50 6, 9.68, 608 9	3 809	6	0.12,	173 9		0.09,	173 9	9,	0.20, 319 9,	319 9,	0.19, 313		9, 0.1	0.12, 18	180   9,	0.10,	180
2-442	7	13.38, 9	966	7, 13.38, 996 7, 20.30,1000 16	.000		60.39,	5145 16		8.11,	5145 1	6, 2	17.17,44	446 16	$68.11,\ 5145 16,\ 217.17,44446 16,\ 241.89,44446 16,$	4446 10		$66.89,\ 5660 16$	3016,	77.13,5664	5664
TSP 3-1173	,	91.09,20	044	6, 91.09, 2044  6, 83.85, 1363 16, 352.03, 5657 16, 360.70, 5657 16, 352.03, 5657 16, 360.70, 5657 16, 374.85, 6202 16	363/10	6, 3!	52.03, 1	5657 1	6, 36	0.70	5657 1	6, 3	52.03, 5	65716,	360.70	5657 1(	3, 374.8	85, 620	)2 16,	, 391.58, 6216	6216
4-3038	4.3	03.63,10	060	$_{4\cdot 303.} \ \ 4,303.63,1090 \ \ 4,330.90,1092 \ _{10,7375.32,21777} \ _{10,7352.68,21777} \ _{10,7375.33,21777} \ _{10,7352.68,21777} \ _{16,7617.81,22701} \ _{11,7586.84,22592} $	.092	0,73%	75.32,2	1777 1	0.735	2.68.2	1777	0.73	75.33,21	777 10.7	7352.68,2	1777 10	3,7617.8	81,2270	)1111,	7586.84,	22592
mean, (40)		02.05,10	045	$8,102.05,1045 \boxed{6,111.18,1016} \boxed{13,1946.97,\ 8188} \boxed{13,1945.40,\ 8188} \boxed{13,1986.19,18050} \boxed{13,1988.87,18048} \boxed{14,2014.92,\ 8686} \boxed{13,2013.91,\ 8663} \boxed{13,1988.87,18048} \boxed{14,2014.92,\ 8686} \boxed{13,2013.91,\ 8663} \boxed{14,2014.92,\ 8686} \boxed{13,2013.91,\ 8663} \boxed{14,2014.92,\ 8686} $	016	3,194	46.97, 8	8188 1.	3,194	5.40,	8188	3,19	86.19,18	050 13,1	1988.87,18	$8048 1_{-}$	4,2014.	92, 868	3613,	2013.91,	8663
max/fails		1/1		0/0			2/6			2/6			9/0		9/0		0,	0/0		0/0	
1-10	11,	0.18,	18	0, 66.20,	384 8	ος	0.51,	578 14		3.40,	823 14	4.	7.39, 9855   14	855 14	6.60	$6.60,\ 3932 16$		0.53, 58	584 14	13.20,	961
Ury = 3-20	12,	1.08,	39	39 0,175.27, 415 13,	415 1	<del>ن</del> ى	2.12,	1395 13	, ,	0.13, 1395 13	395 1	က	20.29,19696   13	696 13,	25.20, 8770   14	$8770 1_{-}$		2.18, 140	1401   16	14.59,	1430
5-30	10,	4.20,	- 22	77 -1,199.46, 285 7,	285	7,	5.29, 2	2433 7		41.38,	2503   7		32.42,23734	734 7,	82.82,1854316	8543   16		5.44, 244	244012	84.60,	2635
mean, (30)	11,	1.82,	45	45 -0,146.98, 362	362 9	9,	2.64,	1469   11		18.30,	1574   11	•	20.03,17762 11	762 11,	38.20,10415 15	0415 1		2.72, 147	1475 14,	37.46,	1675
$\max/{\text{fails}}$		0/0		30/30			1/21			1/29			0/22		0/29		0,	0/0		9/9	
MEAN, $(100)$ $10$ ,	10,	22, 262	262	4, 83,	83, 627   10,	0,	397,	2623 11	1,	402,	402, 2646   11	1,	460,13262   11	262 11,	453,10	453,10166   13,		411, 273	2730 12	438,	3080
max/fails		1/1	$\dashv$	50/50			3/38			3/46			0/38		0/46		0,	0/0		20/20	

Table 8: Results for Group 1: problems in Table 1.

Name	Parameters	$h \circ c$
CPS	$n \in \{10, 50, 100, 500\}$	0 (opt)
	$n \in \{10, 50, 100, 500\}$ $range A \in \{0.2, 0.8\}n$	o (opt)
CVX		0 (1)
MQ	$n \in \{10, 50, 100, 500\}$	0 (opt)
CVX	$range A \in \{0.2, 0.8\}n$	
- D C	m = range A + 3	
EucSum	4 4	0.020529450442540 (1)
NCV	n = 4, range A = 2	0.930538450443740 (best)
NCV	n = 10, range A = 8	0.465587005455171 (best)
NCV	n = 10, range A = 2	0.666424291184390 (best)
NCV	n = 50, range A = 40	0.399571129728750 (best)
NCV	n = 50, range A = 10	0.002143493387185 (best)
NCV	n = 100, range A = 80	excluded
NCV	n = 100, range A = 20	0.333869560359649 (best)
TiltedNorm	$n \in \{10, 50, 100, 500\}$	0 (opt)
CVX	w = 4 $n = 12$	
GenModRos		5.377690121369670 (best)
NCV	U = 1, V = 10 $n = 2$	2 ( )
ModRos		0 (opt)
NCV	$w \in \{1, 2, 4, 8\}$ $n \in \{5, 10, 50, 100\}$	
NesChebRos	$n \in \{5, 10, 50, 100\}$	0 (opt)
NCV	$x^0 \in [0, 2]$	
Ferrier	$n \in \{10, 50, 100\}$	0 (opt)
NCV	$case \in \{1, 3\}$	
NK		
NCV	case = 8	excluded
CVX/NCV	$case \in \{1, 3, 4, 5, 9\}$	$\{0, -\sqrt{2}(n-1), 2(n-1), 2(n-1), 0\}$ (opt)
LV		
NCV	$case = 3, n \in \{10, 50, 100, 500\}$	excluded
NCV	case = 4, n = 10	106.059118520625645 (best)
NCV	case = 4, n = 50	587.997761620671213 (best)
NCV	case = 4, n = 100	1190.421065495747825 (best)
NCV	case = 4, n = 500	6009.807496496680869 (best)
NCV	$case \in \{5, 6\}, n \le 100$	0 (best)
NCV	$case \in \{5, 6\}, n = 500$	(excluded)
Ury		
NCV	n = 10, cubic=0.01	500 (best)
NCV	n = 20, cubic=0.01	909.889558838787480 (best)
NCV	n = 30, cubic=0.01	1114.734712066170232 (best)
CVX	n = 100, cubic=0	1159.869805021747879 (best)
NCV	n = 100, cubic = 0.01	1162.455887489049701 (best)

Table 10: Optimal (opt) or best (best) function values, for problems in Groups 2 and 3.

(bb)	#-n		CBunf			BFGSf			Hanso	f		HyCBf	
		RA	sec	(bb)	RA	sec	(bb)	RA	sec	(bb)	RA	sec	(bb)
	1-10	12	0.07	14	7	0.07	184	7	0.11	245	8	0.09	187
	2-10	12	0.03	14	7	0.05	131	7	0.09	196	9	0.06	134
	3-10	11	0.02	8	6	0.04	94	6	0.07	156	9	0.05	98
CPS	4-50	6	0.07	9	5	0.51	772	5	0.65	953	5	0.54	774
		6		9	6	0.31	560	6		741	10	0.34 $0.42$	563
	5-50	1	0.08						0.53				
	6-50	10	0.07	8	6	0.27	349	6	0.41	530	9	0.29	352
mean,	(60)	10	0.06	10	6	0.22	348	6	0.31	470	8	0.24	352
max/fails			0/0			0/0			0/0			0/0	
	1-4	10	0.34	9	6	0.08	40	6	0.12	65	16	0.10	43
	2-10	7	0.09	27	7	0.21	590	7	6.36	15622	16	0.25	601
EucSum	3-10	16	0.01	6	5	0.02	75	5	0.06	143	5	0.03	77
	4-50	5	3.02	30	6	3.35	1411	6	3.79	1626	16	3.49	1416
	5-50	6	0.13	11	7	0.34	752	7	0.50	1085	16	0.36	754
mean,	(50)	9	0.72	17	6	0.80	574	6	2.17	3708	14	0.85	578
	(00)			11	0		014			3100	14		010
max/fails			1/1	0		0/0	1.40		0/0	054	-	0/0	1.40
	1-10	8	0.04	8	4	0.11	142	4	0.16	254	7	0.12	148
Ferrier	2-10	7	0.02	9	4	0.03	141	5	0.11	348	9	0.05	147
I CITICI	3-50	5	0.26	13	2	1.65	1865	2	2.76	4820	6	1.76	1874
	4-50	5	1.51	53	4	0.39	961	5	0.69	1678	5	0.46	965
mean,	(40)	6	0.46	21	4	0.55	777	4	0.93	1775	7	0.60	783
max/fails	()	`	0/0		1	0/5		1	0/0		'	0/0	.00
	1.10	9	1.93	1050	7	0.24	489	7	1.10	2044	16	0.35	547
	1-12												
G 14 15	2-12	6	2.19	1229	8	0.21	447	8	0.83	1568	16	0.32	503
GenModRos	3-12	16	0.43	237	9	0.24	480	9	1.03	1850	10	0.26	492
	4-12		ationary										
	5-12	$\neq$ st	ationary	points									
mean,	(30)	10	1.52	839	8	0.23	472	8	0.99	1821	14	0.31	514
max/fails	` ′		2/2			0/0			0/0			0/0	
,	1-10	≠ st	ationary	points		-,-			-,-			-, -	
	2-10	4	4.70	1740	8	0.17	347	8	1.52	2639	16	0.18	351
	3-10	16	0.02	. 8	5	0.09	206	5	0.29	479	6	0.11	209
LV	4-10	16	0.03	15	4	0.15	305	5	0.24	453	7	0.16	311
L V	5-50	$\neq$ st	ationary	points									
	6-50	3	1.89	149	16	5.16	2534	16	42.32	29835	16	5.22	2539
	7-50	3	1.21	55	16	2.07	1582	16	2.33	1764	16	2.11	1586
	8-50	5	0.65	34	16	3.84	2778	16	4.10	2960	16	3.88	2783
mean,	(60)	8	1.42	334	11	1.91	1292	11	8.47	6355	13	1.94	1297
	(00)	0		334	11		1232	111		0333	15		1231
max/fails		10	10/10	1.0	_	0/10	000		0/10	0755	1.0	0/0	00.
	2-10	10	0.06	13	8	0.30	680	9	1.66	3755	10	0.32	687
MQ	3-10	10	0.02	10	7	0.04	144	7	0.07	205	10	0.05	147
11100	4-50	7	0.44	11	4	13.74	5769	4	76.94	31841	8	14.11	5775
	5-50	9	0.06	8	8	1.08	2256	8	1.31	2642	9	1.13	2259
mean,	(40)	9	0.14	11	7	3.79	2212	7	20.00	9611	9	3.90	2217
max/fails	( - /		0/0			7/20			0/10			0/0	
max/ rans	1-2	9	0.07	38	6	0.02	37	6	0.03	68	10	0.04	67
		1											
ModRos	2-2	7	0.04	53	6	0.02	56	6	0.03	80	10	0.06	100
	3-2	4	0.13	150	6	0.03	109	6	0.03	122	9	0.27	382
	4-2	2	0.24	278	5	0.06	214	5	0.07	228	6	0.30	492
mean,	(40)	6	0.12	130	6	0.03	104	6	0.04	124	9	0.17	260
max/fails			13/13			0/0			0/0			17/17	
,	1-10	9	0.10	32	6	0.03	63	6	0.07	126	16	0.08	81
	10-50	16	0.55	23	8	0.06	75	8	0.20	298	11	0.17	79
	11-50		ationary			0.00	10	"	0.20	200		0.11	13
			-		_	0.04	70	-	0.15	0.00	1.0	0.10	H.O.
	12-50	16	0.11	8	5	0.04	72	5	0.15	269	10	0.13	76
	2-10	16	0.30	82	6	0.07	191	7	0.17	408	9	0.32	237
NK	3-10	13	0.05	12	6	0.07	199	6	0.34	741	16	0.09	205
1111	4-10	14	0.09	31	7	0.02	70	7	0.06	140	16	0.05	77
	5-10	$\neq$ st	ationary	points									
	6-10	11	0.03	12	5	0.02	76	5	0.06	141	16	0.06	88
	7-50	8	0.44	22	5	0.09	238	5	0.18	431	16	0.43	251
		5	3.01	44	16	3.14	2351	16	23.55	29652	16	3.32	2357
	8-50										1		
	9-50	6	0.45	11	16	3.79	2438	16	36.04	29739	16	3.98	2443
mean,	(100)	11	0.51	28	8	0.73	577	8	6.08	6195	14	0.86	589
max/fails			0/0			0/19			0/19			1/1	
•	1-5	16	0.36	188	1	0.16	415	2	0.58	1886	2	0.45	614
NesChebRos	2-10	16	0.02	9	2	0.08	184	2	0.15	398	2	1.09	716
1.05011051105	3-50	16	0.02 $0.21$	14	1	1.87	1727	1	25.59	28720	1	1.91	1732
mean,	(30)	16	0.20	71	1	0.70	775	2	8.77	10335	2	1.15	1021
$\mathtt{max}/\mathrm{fails}$			3/3			2/3			0/2			5/5	
	1-10	14	0.10	30	6	0.05	180	6	0.13	407	15	0.07	186
TiltedMax	2-20	13	0.25	46	7	0.09	358	7	0.30	1043	9	0.12	365
	3-50	6	2.63	109	7	0.22	763	7	0.60	1946	7	0.24	766
	0.00		00	100	' '	9.22	100	' '	5.50	1010		J.27	.00

			Tabl	le 11 –	conti	nued fro	m prev	vious j	page				
(bb)	#-n		CBunf			BFGSf			HANSO:	f		HyCBf	
		RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$
mean,	(30)	11	0.99	62	7	0.12	434	7	0.35	1132	10	0.14	439
max/fails			0/0			0/0			0/0			0/0	
	2-10	16	0.18	20	7	0.29	399	7	10.16	14847	10	0.30	405
Ury	4-20	11	1.01	42	11	1.90	1305	11	18.19	19606	16	2.06	1316
	6-30	10	4.66	110	16	5.81	2627	16	30.98	23928	16	6.61	2649
mean,	(30)	12	1.95	58	11	2.67	1444	11	19.78	19460	14	2.99	1456
max/fails	` ′		0/0			0/19			0/19			0/0	
MEAN,	(510)	10	0.73	143	7	1.07	819	7	6.17	5544	10	1.20	864
max/fails			29/29			9/76			0/60			23/23	

Table 11: Results for Group 2: problems in Table 3,  $n \leq 50$ .

(bb)	#-n		CBunf			BFGSf			Hansof			HyCBf	
		RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	(bb)	RA	sec	$\overline{(bb)}$
	10-500	6	1.31	7	3	106.27	2409	3	227.11	6780	3	106.60	2410
	11-500	5	1.26	7	5	229.58	4214	5	238.07	4515	8	230.25	4217
CPS	12-500	5	1.25	7	5	152.19	2011	5	160.79	2312	9	152.93	2014
CFS	7-100	6	0.13	7	5	5.27	1464	5	5.95	1765	5	5.33	1466
	8-100	6	0.14	8	5	4.26	1005	5	4.85	1306	9	4.34	1008
	9-100	6	0.14	8	6	2.29	569	6	2.95	870	10	2.36	572
mean,	(60)	6	0.70	7	5	83.31	1945	5	106.62	2925	7	83.64	1948
max/fails	( )		0/0			0/0			0/0			0/0	
*	6-100	≠ st	ationary po	ints		,			,			,	
EucSum	7-100	16	0.27	8	5	0.78	433	5	1.08	734	6	0.84	435
mean,	(10)	16	0.27	8	5	0.78	433	5	1.08	734	6	0.84	435
max/fails	( - /		0/0			0/0			0/0			0/0	
,	5-100	5	0.88	14	14	9.07	4657	14	15.64	17958	14	9.16	4661
Ferrier	6-100	4	12.20	114	5	2.58	1748	5	3.18	2916	5	2.62	1752
mean,	(20)	4	6.54	64	10	5.83	3203	10	9.41	10437	10	5.89	3207
max/fails	(20)	1	1/1	01	10	0/2	0200	10	0/0	10101	10	0/0	0201
max/ rans	10-100	3	0.54	13	16	22.09	4676	16	112.32	43976	16	22.17	4679
	11-100	3	5.44	113	16	17.69	3676	16	18.38	3978	16	17.81	3680
	12-100	4	5.65	121	4	29.24	5882	16	40.52	10822	5	29.49	5887
	13-500	11	ationary po		4	23.24	3002	10	40.52	10022		29.49	3661
LV	14-500	4	4.25	13	16	338.32	21189	16	757.08	60490	16	338.82	21191
	15-500	11	ationary po		10	336.32	21109	10	131.06	00490	10	336.62	21191
	16-500		ationary po										
	9-100		ationary po		1.0	101.04	0050	1.0	020.00	00016	1.0	100.00	0000
mean,	(40)	4	3.97	65	13	101.84	8856	16	232.08	29816	13	102.08	8860
max/fails		_	0/0	405	١.	1/5	0500		0/4	10010		0/0	0-11
	6-100	2	90.94	105	4	165.89	6536	16	804.07	42648	5	172.30	6544
MQ	7-100	16	0.15	9	5	12.16	5122	5	13.23	6178	6	12.27	5125
	8-500	-0	5023.56	22	0	14634.68	8352	0	75735.20	42086	16	20200.74	8384
	9-500	0	478.97	25	3	991.68	6572	16	5382.02	42461	3	1016.28	6575
mean,	(40)	4	1398.41	40	3	3951.10	6646	9	20483.63	33343	8	5350.40	6657
$\mathtt{max}/\mathrm{fails}$			6/6			6/8			0/5			2/2	
	13-100	16	1.35	23	5	0.63	410	5	0.80	711	5	0.80	413
	14-100	4	7.81	36	16	13.09	4443	16	58.45	43744	16	13.41	4447
	15-100	5	0.91	9	16	16.35	4752	16	79.97	44052	16	16.86	4756
	16-100	12	0.88	21	8	0.13	83	8	0.39	384	16	0.26	86
	17-100	$\neq$ st	ationary po	$_{ m ints}$									
NK	18-100	10	0.21	8	4	0.09	69	4	0.29	370	16	0.31	75
IVIX	19-500	5	19.78	39	4	30.79	902	4	32.57	1204	16	39.20	920
	20-500	2	11.11	7	16	342.25	19620	16	763.23	58920	16	343.85	19621
	21-500	6	16.54	18	16	451.70	21468	16	1006.16	60769	16	453.87	21470
	22-500	16	5.44	22	8	2.25	68	8	6.47	489	11	3.20	72
	23-500	$\neq$ st	ationary po	ints									
	24-500	16	0.83	7	5	2.16	72	5	4.41	373	8	2.64	74
mean,	(100)	9	6.49	19	10	85.94	5189	10	195.28	21102	14	87.44	5193
max/fails	,		0/0			0/8			0/8			0/0	
NesChebRos	4-100	16	0.84	19	1	29.63	5247	1	112.88	44548	1	29.73	5251
mean,	(10)	16	0.84	19	1	29.63	5247	1	112.88	44548	1	29.73	5251
max/fails	()	-	0/0		-	0/10		1	0/10		-	0/0	
TiltedMax	4-100	3	7.84	105	7	1.59	1372	7	2.04	2776	7	1.62	1374
mean,	(10)	3	7.84	105	7	1.59	1372 $1372$	7	$\frac{2.04}{2.04}$	2776	7	1.62	1374 $1374$
	(10)	3		100	'		1312	'		4110	'		13/4
max/fails	<b>=</b> 400	1	2/2	94	7	0/0	EOE 4	7	0/0 260.76	45155	16	$\frac{0}{0}$ 95.00	E 907
Ury	7-100	1	70.85		1	61.95	5854		269.76	45155	I		5897
•	8-100	0	25.19	91	16	61.52	5685	16	277.93	44986	16	. 88.49	5774
											Cont	inued on ne	xt page

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Table 12 – continued from previous page													
(bb)	#-n	CBunf			BFGSf			Hansof			HyCBf		
		RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$	RA	sec	$\overline{(bb)}$
mean,	(20)	0	48.02	92	12	61.73	5770	12	273.85	45070	16	91.74	5836
max/fails			20/20			14/20			0/20			13/13	
MEAN,	(310)	7	163.67	47	7	480.19	4295	8	2379.65	21195	9	639.26	4307
max/fails			29/29			21/53			0/47			15/15	

Table 12: Results for Group 3: problems in Table 3, n=100 and n=500.