# SYSTEM SOFTWARE ASSIGNMENT 1

TOPICS: - 1. CONTROL SECTIONS & ADVANTAGE

OF USING THEM

2. MASM ASSEMBLER

REFERENCES: - Leland L. Beck, System Softwase: An Introduction to Systems Pologonamoning, 3/E

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#### CONTROL SECTIONS

A control section is a past of the program that maintains its identity after assembly; each control section can be loaded & subcrated independently of the others. Different control sections are most often used Jose subscoutines on other logical subdivisions of a program. The programmer can assemble, load & manipulate each of these control sections separately. The resulting flexibility is a major benefit of using control sections.

When contoud sections Joseph logically orelated paorts of a porgocam, it is necessary to porovide some awars Jose linking thom togethere. For example, instructions in 1 contoud section might need to refer to instructions on data located in another section. Because contoud sections are independently loaded & relocated, the assembler is unable to process these references in the usual may. The assembler has no idea where any other contribut section will be located at execution time. Buth references between contribut sections are called exterinal references. The assembler generates information for each exterinal references that will allow the loader to perform the suguirored linking.

The assembless establishes a sepassate location counters (beginning at c) for each control section, just as it closs for program blocks. Control sections differ from program blocks in that they are handled separately by the assembless. Symbols that are defined in 1 corrbool section may not be used directly by another control section; they must be identified as external references for the loader to handle. 2 assembless directives used to identify such references

EXTREF statement in a contout section names symbols, called external symbols, that are defined in this contout section 2 may be used by other sections. Contout section mames don't need to be marmed in an extrement because they are automatically considered to be external symbols. The extrest statement mames symbols that are used in this contout section 2 are defined elevation. The order in which symbols are listed in the extrest 2 EXTRES statements is not significant:

The assembless must also include information in the object psiograms that will cause the boardess to insent the psiopess values whose they are suguished. We need 2 new succosed types in the object psions and a change in a previously defined succosed type. The 2 new succosed types are Define 2 Refers. A Define succosed gives information about external symbols that are defined in this control section - that is, symbols named by EXTDEF, A Refers succosed lists symbols that are used as external suggestions by the control section - that is, symbols marked by EXTREF. The Josimals of these succoseds are as follows.

## Define second:

Col. 1 D

Col. 2-7 Name of exterinal symbol defined in this control section

Col. 8-13 Relative address of symbol within this control section

Choxadocimal)

Col. 14-73 Repeat information in Col. 2-13 Jose other exterinal

symbols.

### Refer recond:

Col. 1 F

Cal. 2-7 Name of extensal symbol referred to in this

contorol section.

Cal. 8-73 Names of other external reference symbols

The other information needed for pregnam linking is added to the Modification record type. The new Josephan is as follows.

#### Modification record (revised):

Col. 1 M

Cal. 2-7 Stasifing address of the field to be modified,

orelative to the beginning of the contour section

(hexadecimal)

Col. 8-9 Length of the field to be amodified, in half-bytes

(hexadecimal)

Col. 10 Modification flag (+ 091 -)

Cal. 11-16 Exterinal symbol whose value is to be added to con

subtoracted Josom the indicated field.

The 2 new items specify the modification to be performed: adding on subtracting the value of some external symbol. The symbol used for another one.

There is a separate set of object program records for each control section. The records Jose each control section are exactly the same as they would be if the sections were assembled separately.

The Define & Refer successed from each contound section include the symbols named in the EXTDET & EXTREF statements. In the case of Define, the successed also indicates the overlative addicess of each external symbol within the contound section. For EXTREF symbols, no addicess information is available. These symbols are simply named in the Refer successed.

The sievised Modification second may still be used to perform proposed medication. In the case of sielocation, the amedification sequiposed is adding the beginning address of the control section to cestain fields in the object program. The gyambol used as the name of the control section mame two as its value the sequiposed address. Since the control section name is automatically an external symbol, it is available from use in Modification succourds. In this way, exactly the same mechanism can be used for program succeptant finking.

When an exposession involves exterinal references, the assembless cannot in general determine whether on not the exposession is legal. The pairing of relative terms to test legality cannot be done without knowing which of the terms occur in the same control sections, I this is unknown at assembly time. In such a case, the assembless evaluates all of the terms it can, I combines these to Joseph an initial exposession value. It also generates Modification records so the loaders can finish the evaluation, The loaders can then check the expression for essences.

#### MASM Assembles

An MASM assembles language pologoram is usuitten as a collection of segments. Each segment is defined as belonging to a pasiticular class, corresponding to its contents. Commonly used classes are CODE, DATA, CONST & STACK.

Dusing peroperam execution, segments are addressed via the ×86 segment sugisters. In most cases, code segments are addressed using sugister 3s. These segment sugisters are automatically set by the system loader when a peroperam is loaded for execution. Register C8 is set to indicate the segment that contains the starting label specified in the END statement of the peroperam. Register 6s is set to indicate the last stack segment perocessed by the loader.

Data segments asie mosimally addressed using DS, ES, FS on GIS. The segment sugisted to be used can be specified explicitly by the pologolammes of the pologolammes closs not specify a segment sugistes, 1 is selected by the assembles.

By default, the assembless assumes that all sufferences to data segments use suggister DS. This assumption can be changed by the assembless dissective ASSUME. For example, the dissective

ASSUME ES: DATASEG2

tells the assembles to assume that sugisteer ES indicates the segment DATASEGIZ. Thus, any sufferences to labels that asse defined in DATASEGIZ will be assembled using sugisteer ES. It is also possible to collect several segments into algorithm I use ASSUME to associate a segment sugisteer with the group.

Registeres DS, ES, FS & G13 must be loaded by the perogram before they can be used to address data segments.

Jump instructions are assembled in 2 different ways, depending on whether the target of the jump is in the same code segment as the jump instruction. In near jump is a jump to a target in the same code segment; a far jump is a jump to a target in a different code segment. A near jump is assembled using the current code segment realister CS. A far jump must be assembled using a different segment register, which is specified in an instruction prefix. The assembled machine instruction for a mear jump occupies 2 or 3 bytes. The assembled instruction for a far jump requires 5 bytes.

Positioned references to labels in the source program can cause problems. For example, consider a jump instruction like TMP TARGIET

If the definition of the label TARGIET occusis in the psignam before the IMP instruction, the assembles can tell whether this is a near jump on a far jump. However, if this is a Jornard reference to TARGIET, the assembles closs not know how many bytes to reserve Jose the instruction.

By default, MASM assumes that a footklasted jump is a meast jump. If the tasget of the jump is in another code segment, the psiogolammes must make the assemblest by usilling IMP FAR PTR TARGIET

If the jump address is within 128 bytes of the current instruction, the peragrammer can specify the sharter (2-byte) near jump by writing JMP SHORT TARGIET

If the IMP to TARGIET is a Jasi Jump, & the pologonommen does not specify the TTR, a poloblem occupies. During Pass 1, the assembles occupied is bytes in the actual assembled instruction suguious 5 bytes. In the earliest versions of MASM, this caused an assembly enripe. In later versions of MASM, the assemblest can supert Pass 1 to generate the connect location counters values.

Segments in an MABM source propriam can be written in more than I part. If a segment directive specifies the came name as a proviously defined regement, it is considered to be a continuation of that regement. All of the parts of a regement are gathered together by the assembly process. Thus, regements can perform a similar function to the program blocks in SIC/XE.

References between segments that are assembled together are automatically handled by the assembler. External deferences between separately assembled modules must be handled by the linker. The MASM directive PUBLIC has approximately the same function as the SIC/XE directive EXTDEF. The MASM directive EXTRES. The MASM directive EXTRES.

The object pousgoiam foot MASM may be in several different footemats, to allow easy & efficient execution of the pousgoiam in a vacility of operating envisionments. MASM can also poudure an imstruction timing listing that shows the number of clock cycles or equived to execute each machine instruction. This allows the pougonomment to execute a great deal of control in optimizing timing-conitical sections of code.