

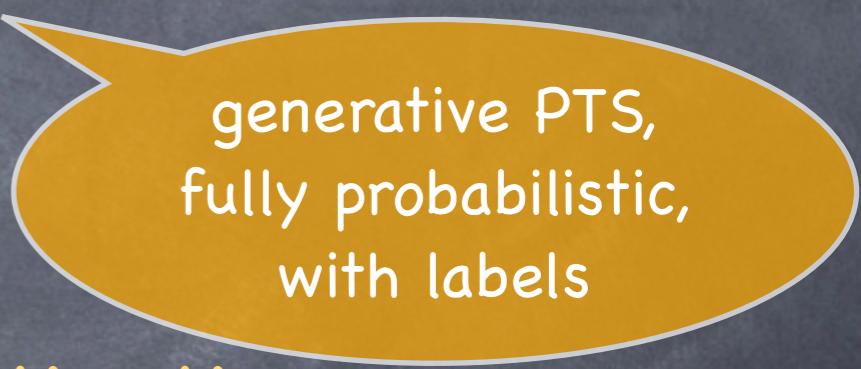
# Sound and Complete Axiomatization of Trace Semantics for Probabilistic Transition Systems

Alexandra Silva and Ana Sokolova  
CWI and University of Salzburg

MFPS 2011, Pittsburgh, 27.5.2011

# We will discuss

- ⦿ history
- ⦿ probabilistic transition systems
- ⦿ (finite) trace semantics
- ⦿ the sound and complete axiomatization
- ⦿ in a coalgebraic setting



generative PTS,  
fully probabilistic,  
with labels

# History

- ⦿ For LTS Milner '84, JCSS
- ⦿ expressions for LTS
- ⦿ Kleene style theorem
- ⦿ axiomatization
- ⦿ sound and complete for bisimilarity

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- axiomatization  $P + Q \equiv Q + P, P + 0 \equiv P, \mu x.P \equiv P[\mu x.P/x], \dots$
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$$P \equiv Q \iff P \sim Q$$

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- ⌚ Milner's result was extended by Rabinovich  
'93 MFPS for trace semantics
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$$P \equiv Q \iff \text{tr}(P) = \text{tr}(Q)$$

# Now we do it for PTS

- Expressions/axioms for PTS come in many flavors  
(mainly for bisimilarity)  
we build on Silva, Bonchi, Bonsangue, Rutten '09/'10
- Trace semantics for PTS also exists in variants  
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same expressions, one more axiom,  
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# Probabilistic transition systems

PTS here are generative, labelled, with explicit termination

$$X \rightarrow \mathcal{D}_\omega(1 + A \times X)$$

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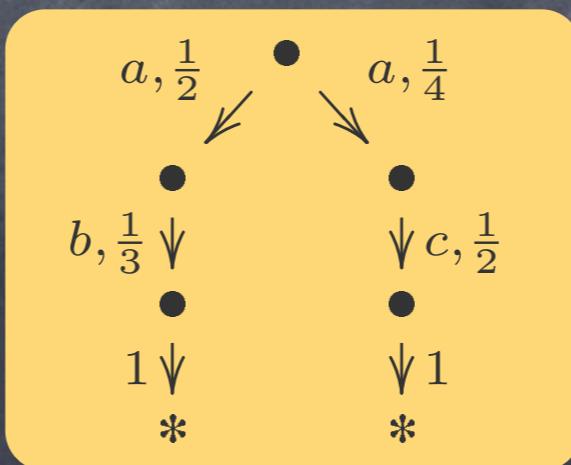
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Example:



# Coalgebra basics

Category  $C$ , functor  $F$ , category of coalgebras:

$\text{Coalg}_F$

Objects:

$$X \xrightarrow{c} FX$$

Arrows:

$$\begin{array}{ccc} X & \xrightarrow{h} & Y \\ c \Downarrow & & \Downarrow d \\ FX & \xrightarrow{Fh} & FY \end{array}$$

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Final coalgebra semantics:

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bisimilarity in Sets  
(for wpp functors)  
trace semantics in  $\mathcal{Kl}(T)$   
(for  $TF$ -coalgebras)

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It instantiates to finite trace distribution:

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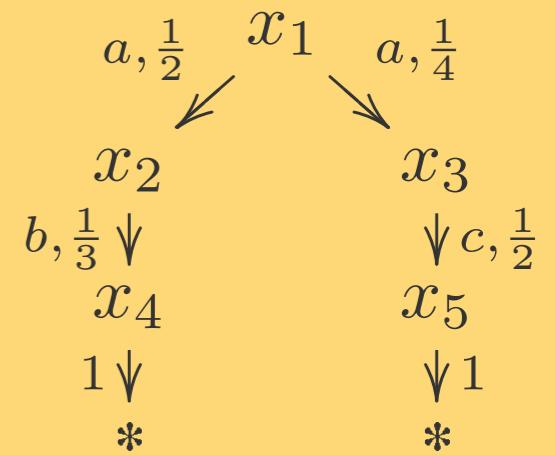
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$$\text{tr}(x_1)(ab) = \frac{1}{6}$$

$$\text{tr}(x_1)(ac) = \frac{1}{8}$$

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Are the same ones as for bisimilarity

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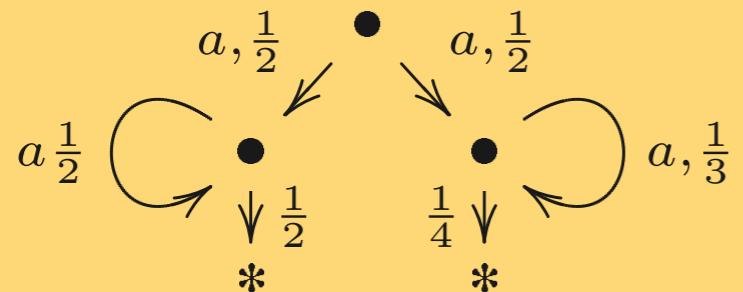
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$$\begin{aligned} &\frac{1}{2} \cdot a \cdot \mu x. \left( \frac{1}{2} \cdot a \cdot x \oplus \frac{1}{2} \cdot * \right) \\ &\oplus \frac{1}{4} \cdot a \cdot \mu y. \left( \frac{1}{3} \cdot a \cdot y \oplus \frac{1}{2} \cdot * \right) \end{aligned}$$

# Expressions for PTS

expressions behave! (Kleene-style theorem)

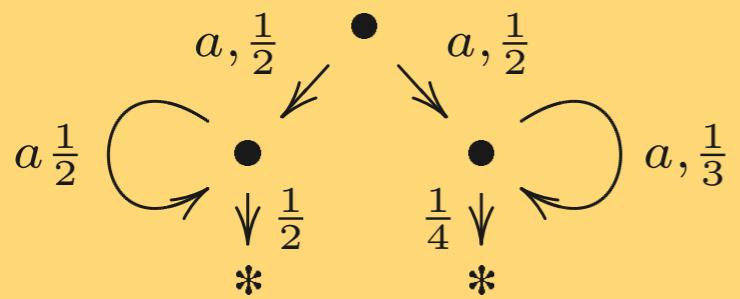
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# Axioms

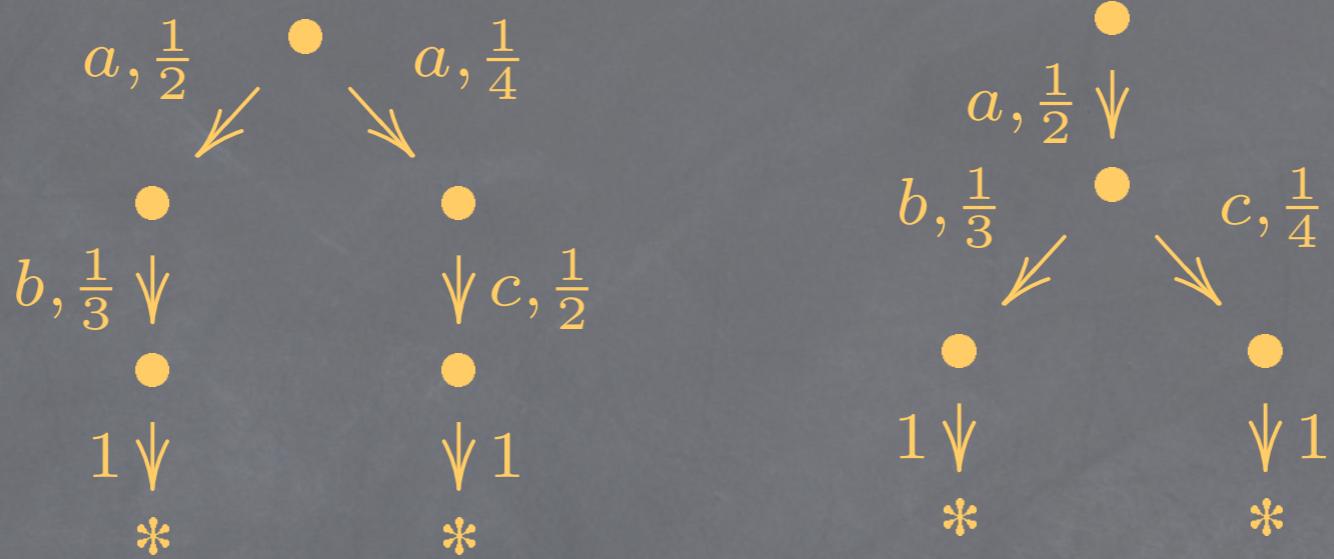
bisimilarity

$E_1 \oplus (E_2 \oplus E_3)$	$\equiv$	$(E_1 \oplus E_2) \oplus E_3$	$(A)$
$E_1 \oplus E_2$	$\equiv$	$E_2 \oplus E_1$	$(C)$
$E \oplus \emptyset$	$\equiv$	$E$	$(E)$
$\mu x. E$	$\equiv$	$E[\mu x. E/x]$	$(FP)$
$\gamma[E/x] \equiv E$	$\Rightarrow$	$\mu x. \gamma \equiv E$	$(UFP)$
$\mu x. E$	$\equiv$	$\mu y. E[y/x]$ if $y$ is not free in $E$	$(\alpha - equiv)$
$E_1 \equiv E_2$	$\Rightarrow$	$E[E_1/x] \equiv E[E_2/x]$	$(Cong)$
$0 \cdot E$	$\equiv$	$\emptyset$	$(Z)$
$p \cdot E \oplus p' \cdot E$	$\equiv$	$(p + p') \cdot E$	$(S)$

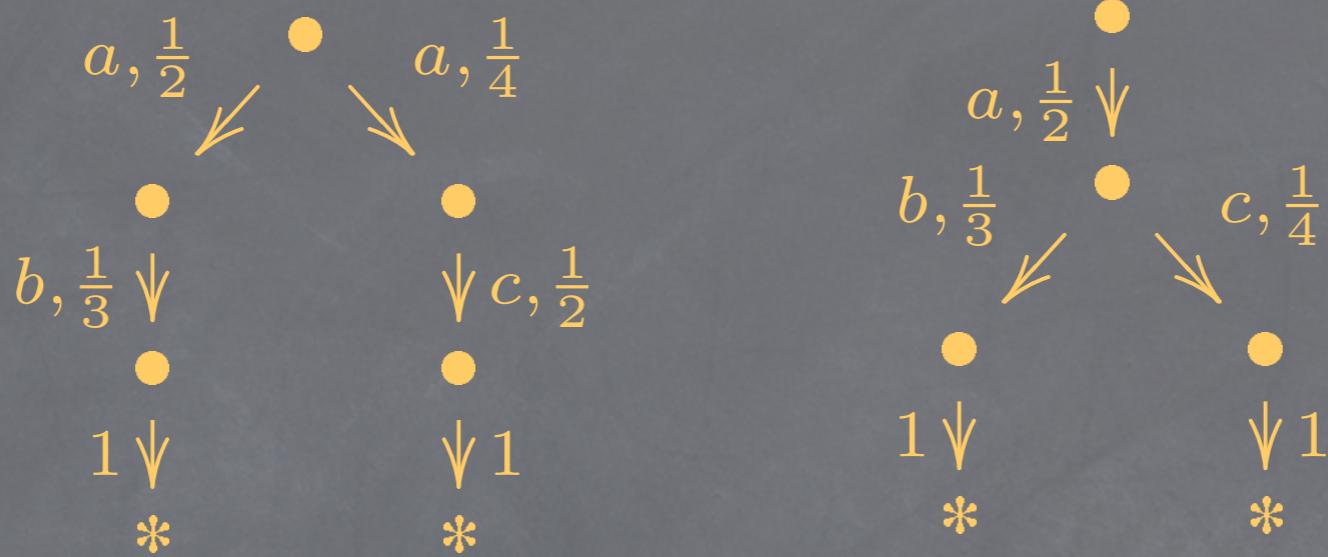
$$p \cdot a \cdot (p_1 E_1 \oplus p_2 E_2) \equiv p_1 \cdot a \cdot p E_1 \oplus p_2 \cdot a \cdot p E_2 \quad (D)$$

trace

# Example

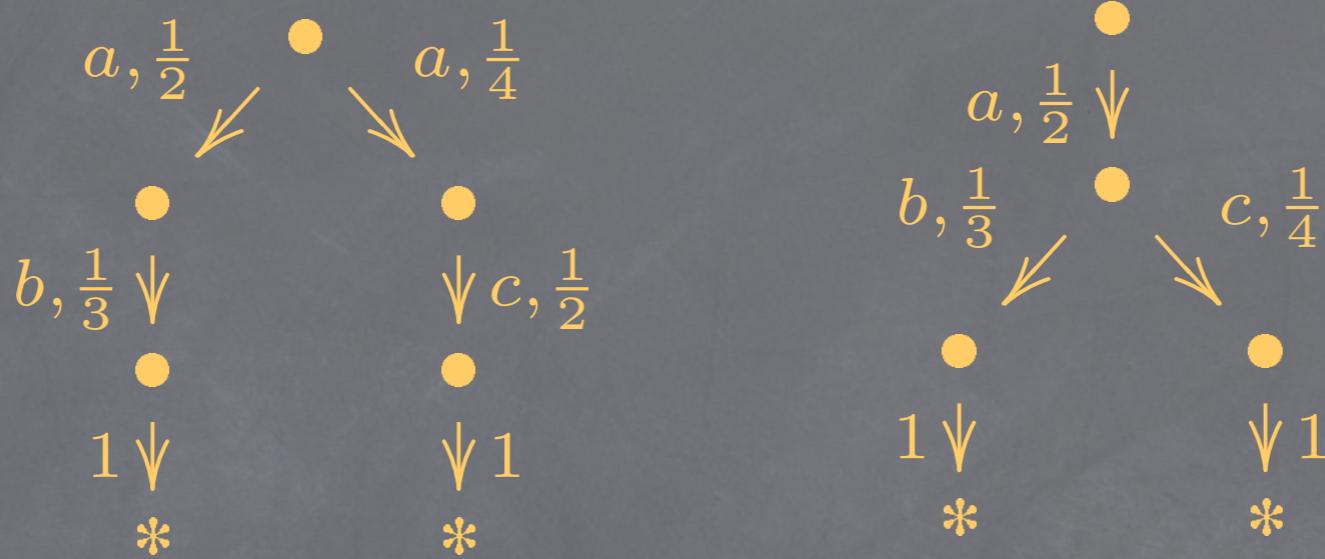


# Example



$$\begin{aligned} \left( \frac{1}{2} \cdot a \cdot \frac{1}{3} \cdot b \cdot 1 \cdot * \right) \oplus \left( \frac{1}{4} \cdot a \cdot \frac{1}{2} \cdot c \cdot 1 \cdot * \right) &\stackrel{(D)}{\equiv} \frac{1}{2} \cdot a \cdot \left( \frac{1}{2} \left( \frac{2}{3} \cdot b \cdot 1 \cdot * \right) \oplus \frac{1}{4} (1 \cdot c \cdot 1 \cdot *) \right) \\ &= \frac{1}{2} \cdot a \cdot \left( \frac{1}{3} \cdot b \cdot 1 \cdot * \oplus \frac{1}{4} \cdot c \cdot 1 \cdot * \right) \end{aligned}$$

# Example



$$\frac{1}{3} \cdot b \cdot 1 \cdot * = \frac{1}{2} \left( \frac{2}{3} \cdot b \cdot 1 \cdot * \right), \quad \frac{1}{2} \cdot c \cdot 1 \cdot * = \frac{1}{2} (1 \cdot c \cdot 1 \cdot *)$$

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 &= \frac{1}{2} \cdot a \cdot \left( \frac{1}{3} \cdot b \cdot 1 \cdot * \oplus \frac{1}{4} \cdot c \cdot 1 \cdot * \right)
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# Soundness and Completeness

Find an injective map  $out_{\equiv}$  with  $\text{tr} = out_{\equiv} \circ [-]$

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## Soundness

$$\begin{array}{lcl} E_1 \equiv E_2 & & \\ \Leftrightarrow & [E_1] = [E_2] & \\ \stackrel{(*)}{\Rightarrow} & out_{\equiv}([E_1]) = out_{\equiv}([E_2]) & \\ \stackrel{(\triangle)}{\Leftrightarrow} & tr(E_1) = tr(E_2) & \end{array}$$

## Completeness

$$\begin{array}{lcl} & tr(E_1) = tr(E_2) & \\ \stackrel{(\triangle)}{\Leftrightarrow} & out_{\equiv}([E_1]) = out_{\equiv}([E_2]) & \\ \stackrel{(\heartsuit)}{\Rightarrow} & [E_1] = [E_2] & \\ \Leftrightarrow & E_1 \equiv E_2 & \end{array}$$

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(\*) - existence of  $out_{\equiv}$     ( $\triangle$ ) -  $tr = out_{\equiv} \circ [-]$     ( $\heartsuit$ ) - injectivity

# How to get out?

Bisimilarity case, F-coalgebras

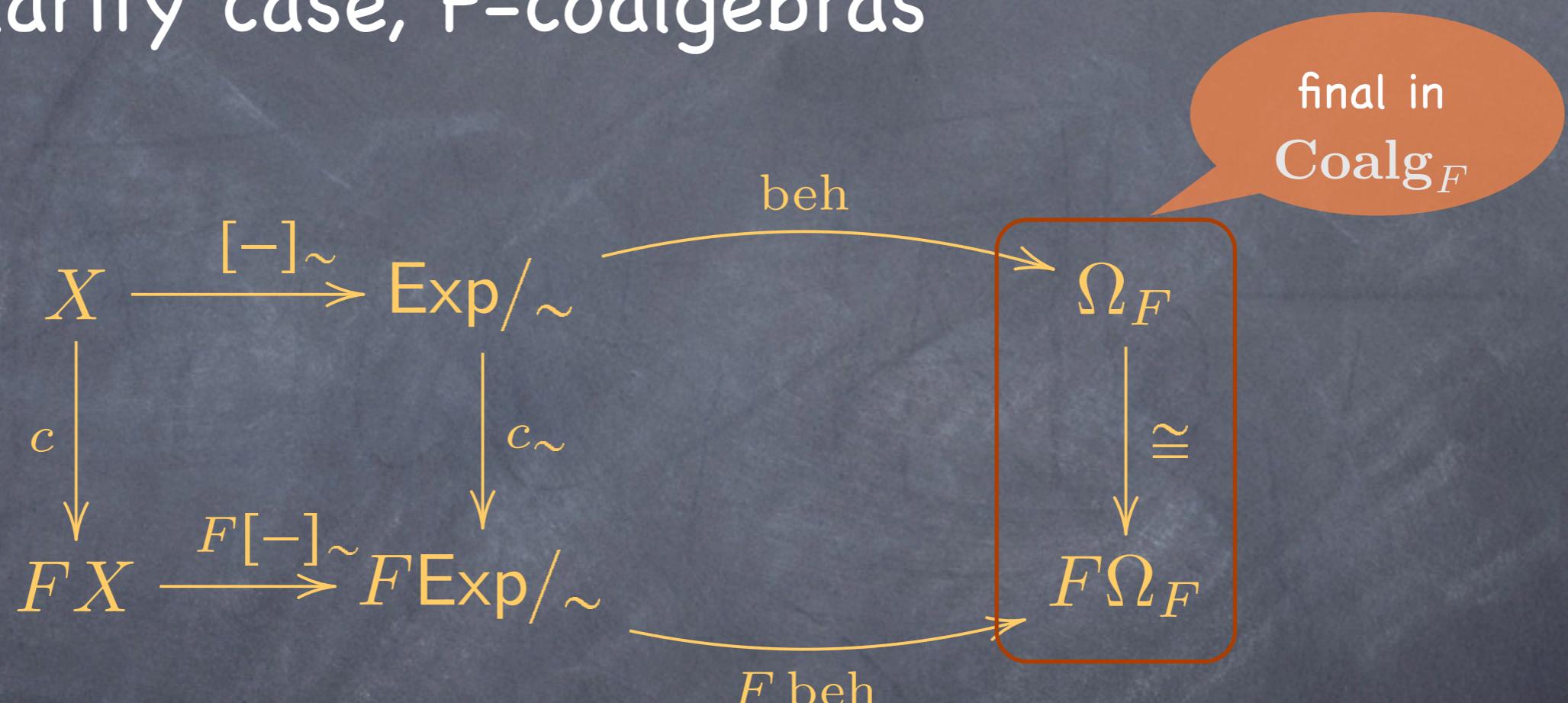
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$$\begin{array}{ccccc} X & \xrightarrow{[-]_\sim} & \text{Exp}/_\sim & \xrightarrow{\text{beh}} & \Omega_F \\ c \downarrow & & \downarrow c_\sim & & \downarrow \simeq \\ FX & \xrightarrow{F[-]_\sim} & F\text{Exp}/_\sim & \xrightarrow{F \text{ beh}} & F\Omega_F \end{array}$$

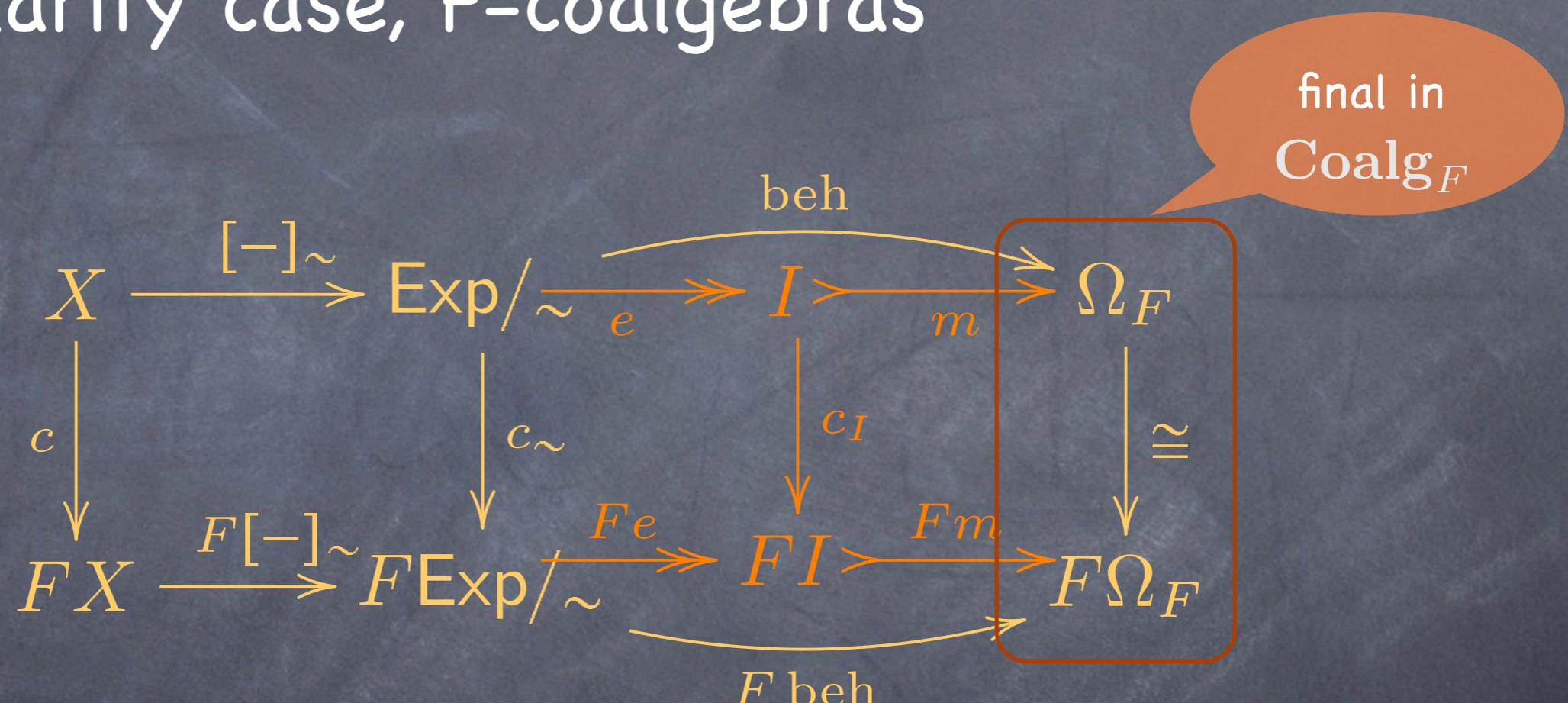
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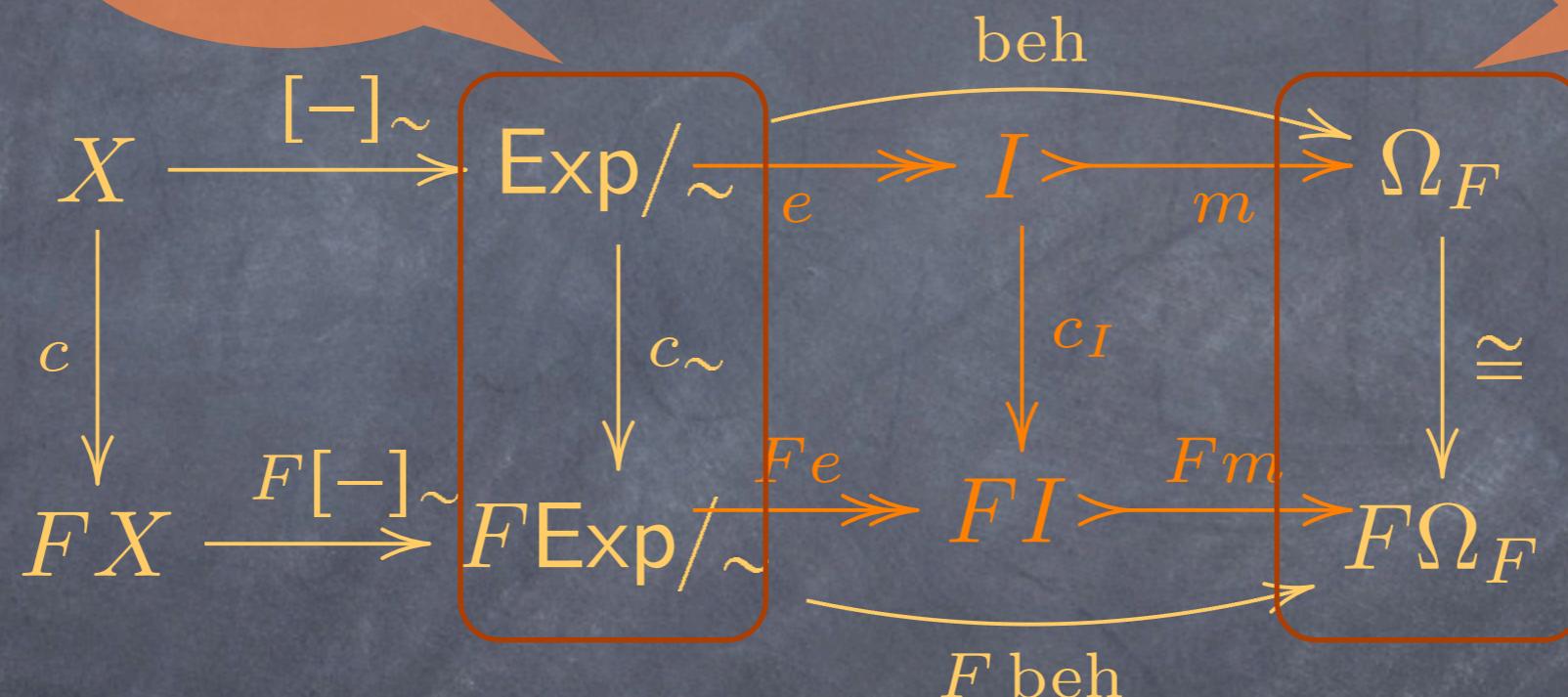


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Bisimilarity case,  $F$ -coalgebras

final  
elsewhere

final in  
 $\text{Coalg}_F$



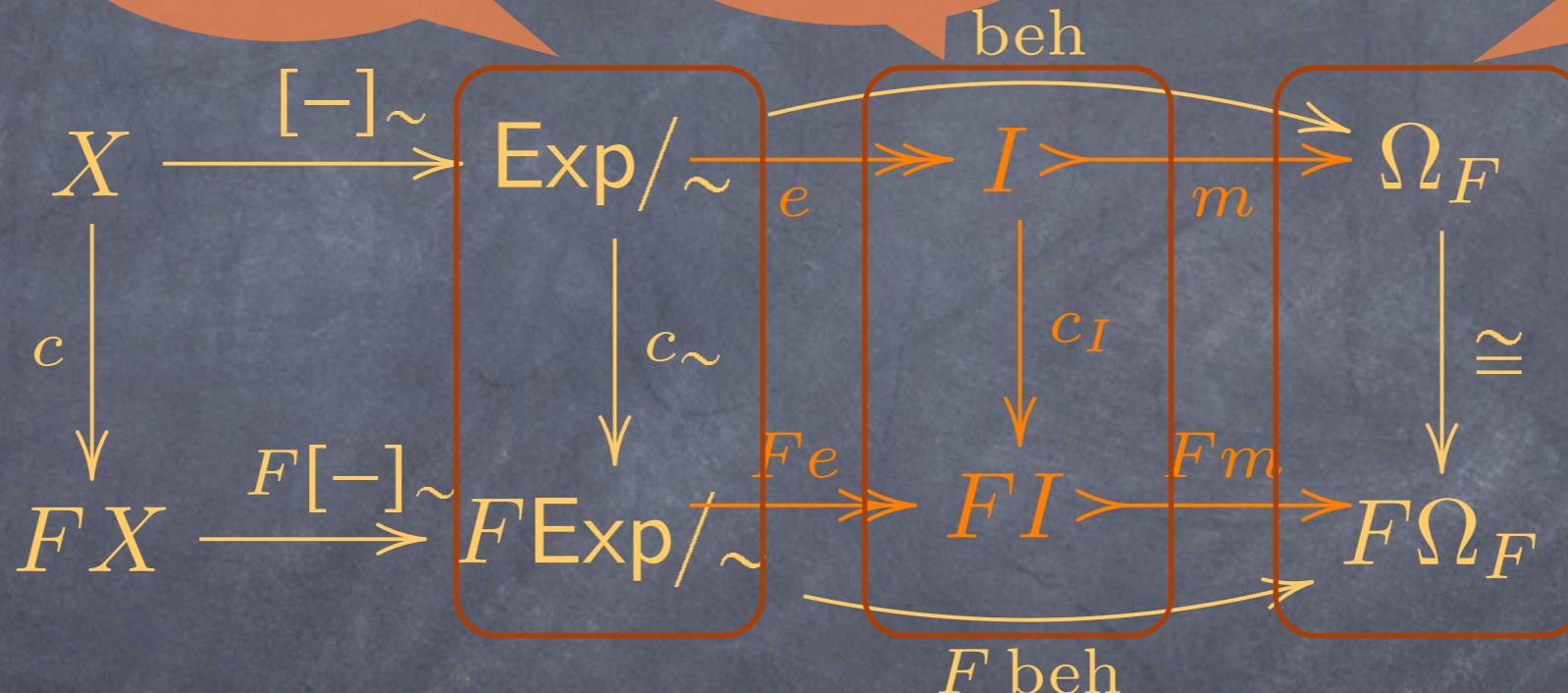
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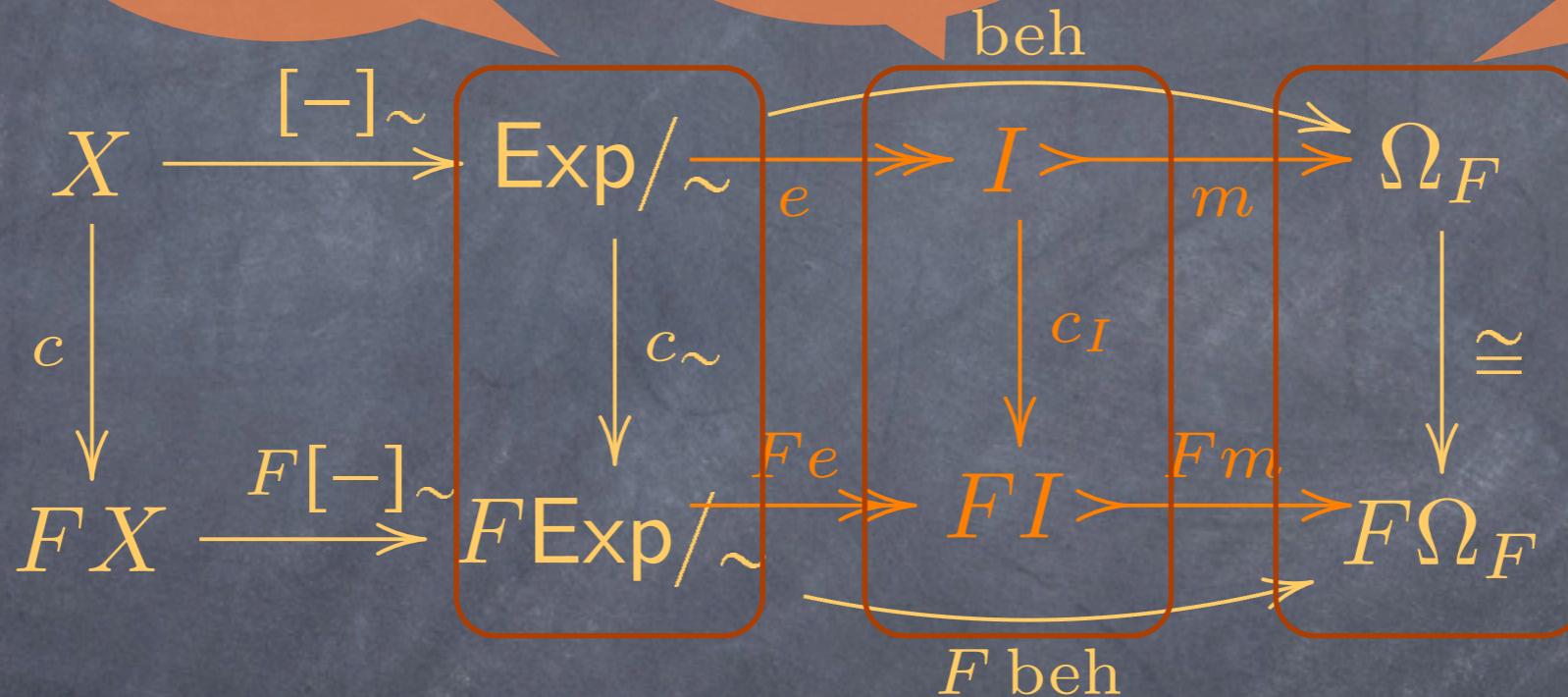
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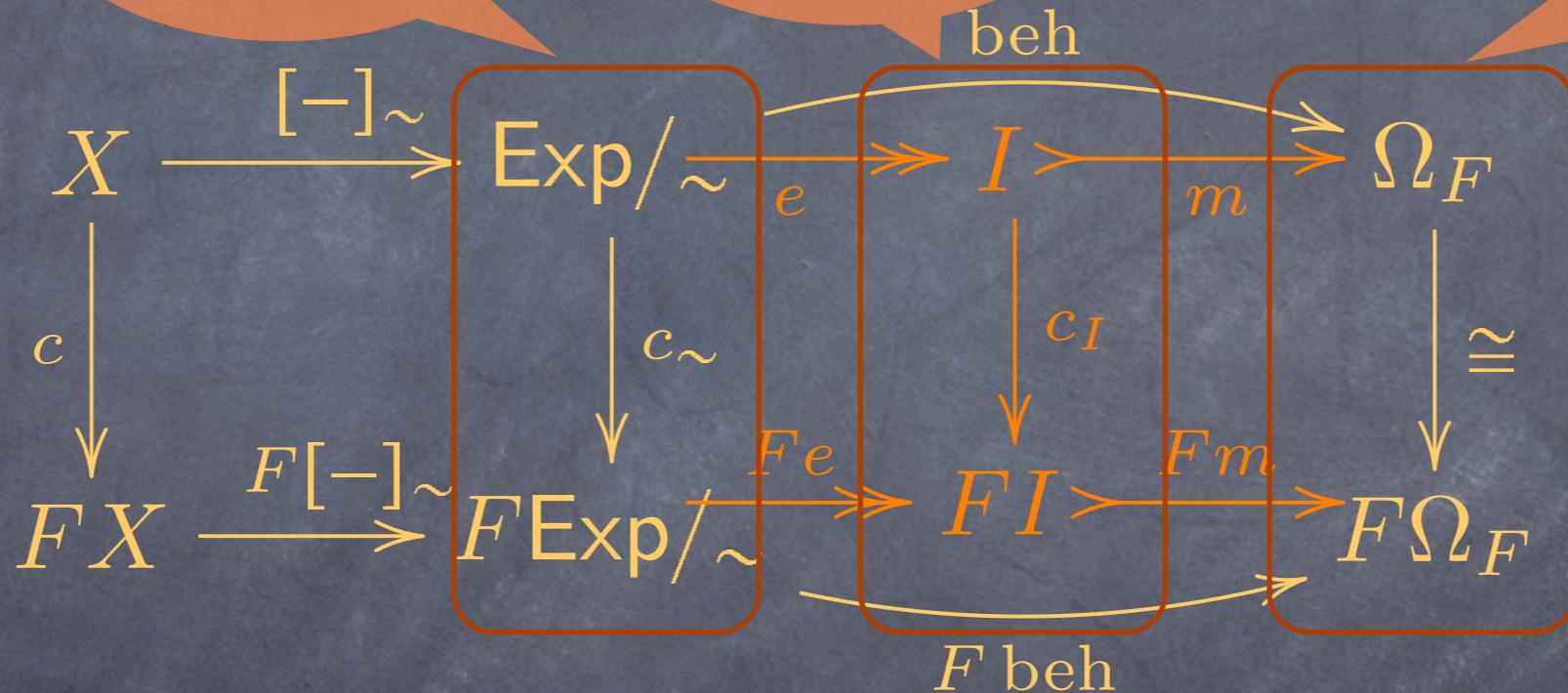
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Silva et al. '08/'09/'10, Jacobs'06

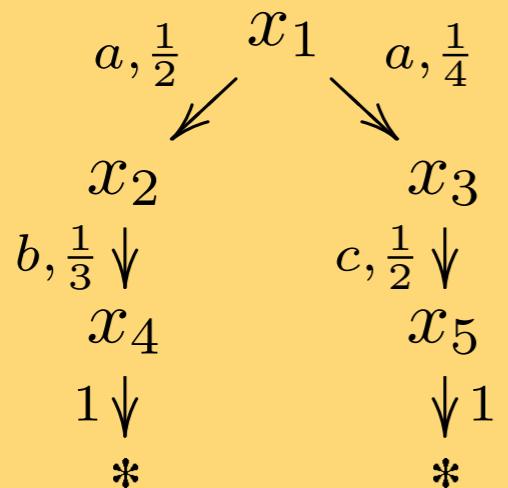
there are also  
algebras around

# A way out for traces?

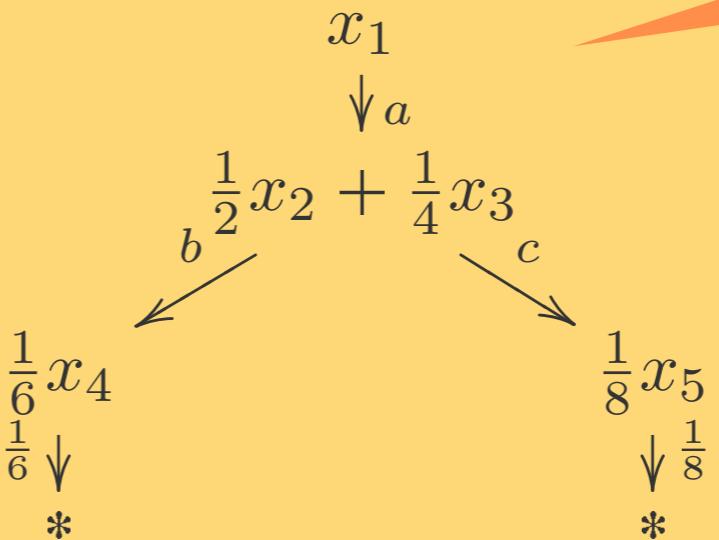
- ⦿ It is tough to work in Kleisli categories
- ⦿ Factorization ?
- ⦿ So we find a way to stay in Sets  
or rather in  $Sets^{\mathcal{D}_\omega}$
- ⦿ A way out - determinization

# Determinization of PTS

PTS example



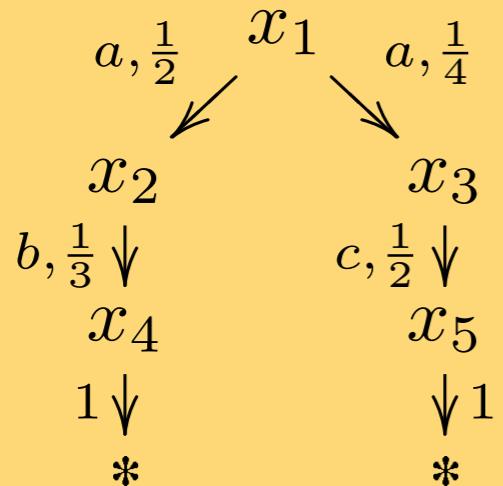
Its determinization



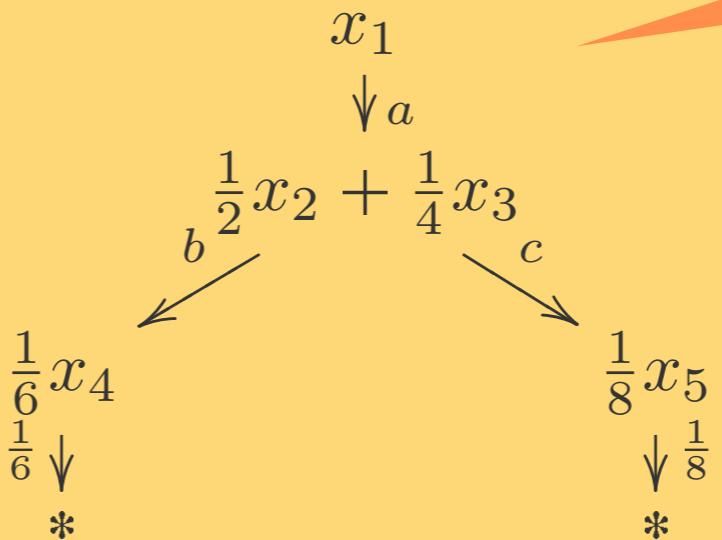
a G-coalgebra

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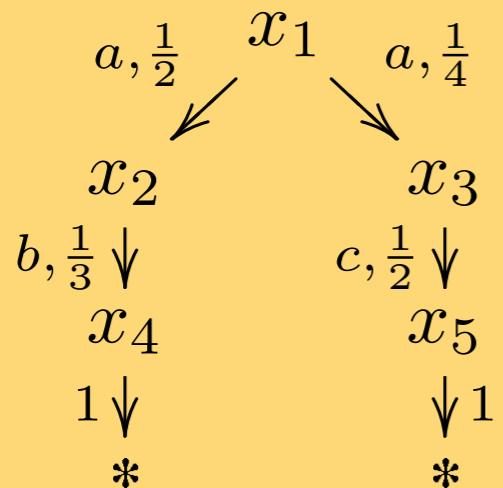
a G-coalgebra

$$GX = [0, 1] \times X^A$$

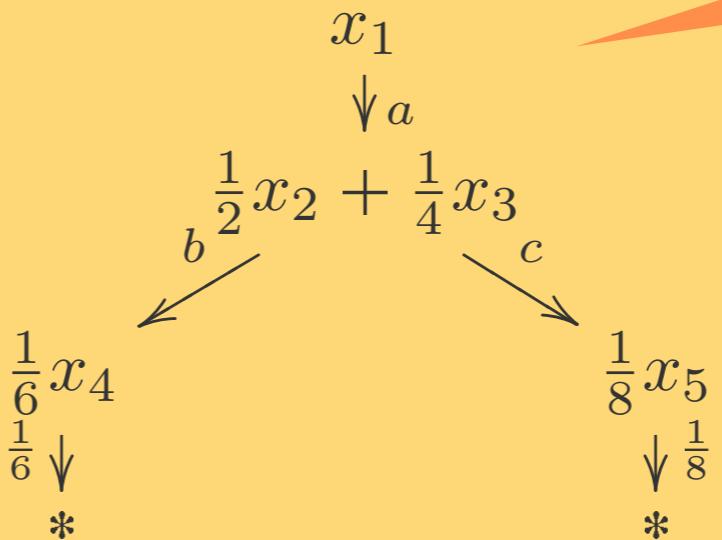
$$\begin{array}{ccc} X & \xrightarrow{\text{out}} & [0, 1]^{A^*} \\ c \downarrow & & \downarrow \cong \\ GX & \xrightarrow{G\text{out}} & G([0, 1]^{A^*}) \end{array}$$

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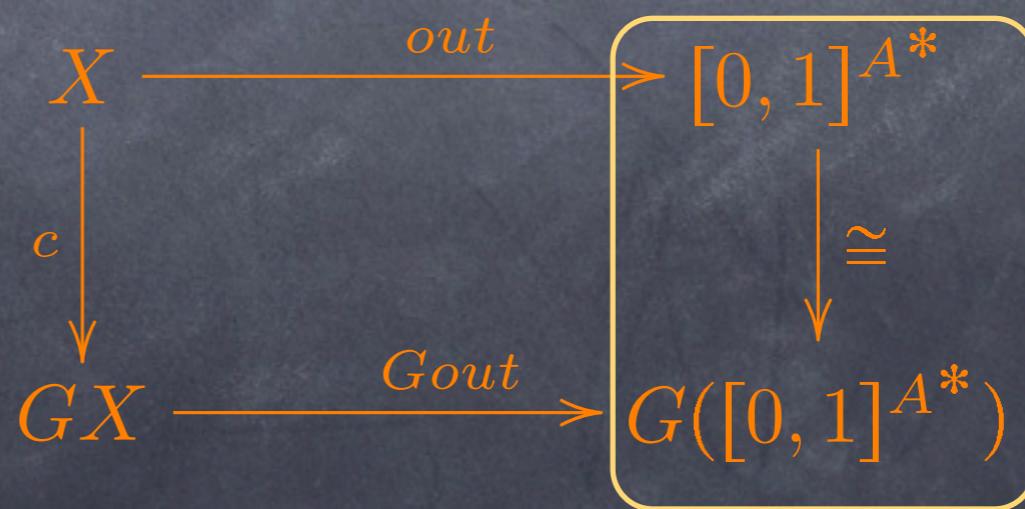


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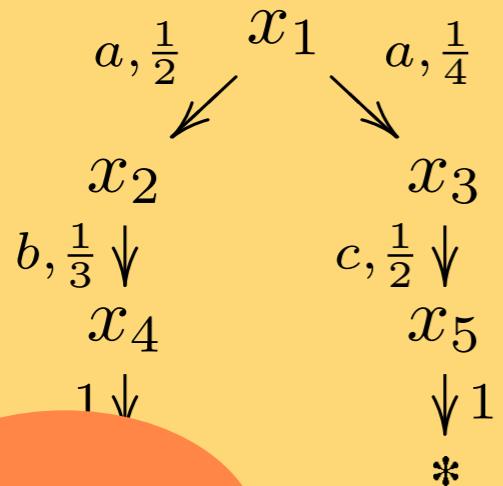
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final  
G-coalgebra

# Determinization of PTS

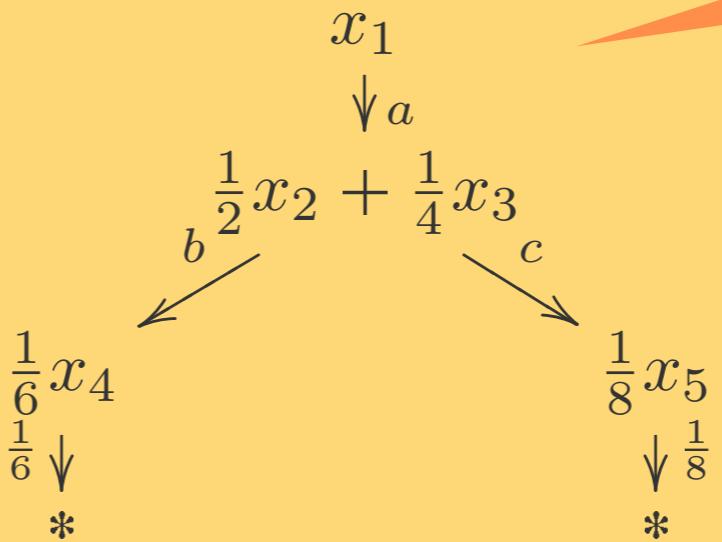
## PTS example



Sets  $\mathcal{D}_\omega$

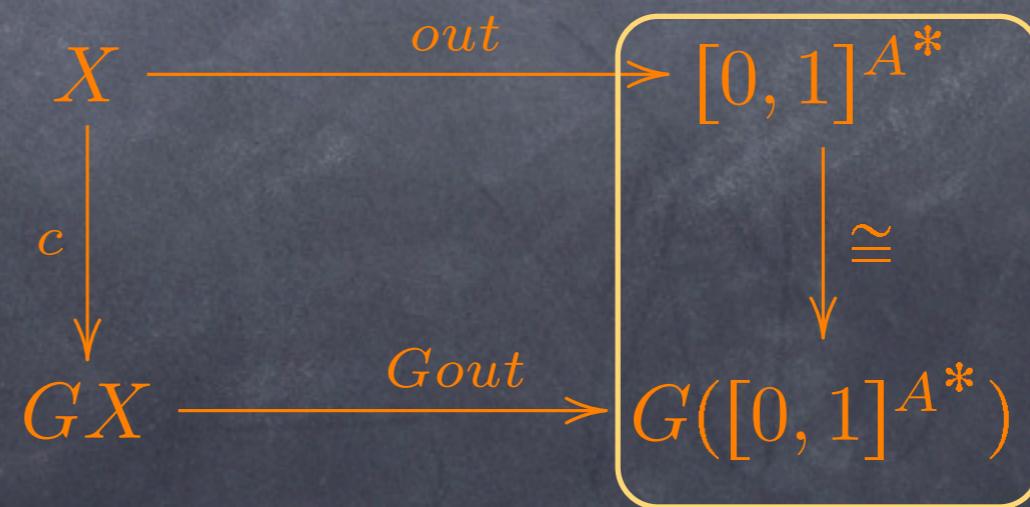
$$GX = [0, 1] \times X^A$$

## Its determinization



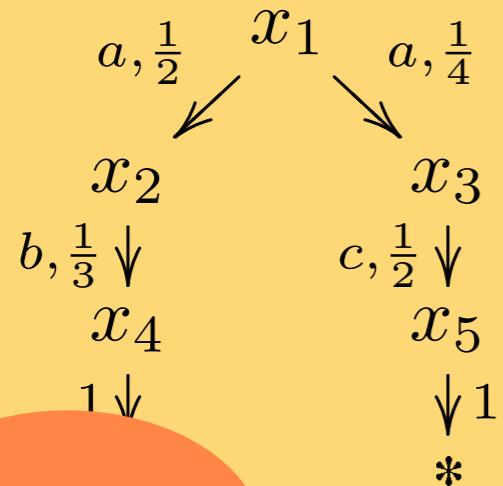
a G-coalgebra

final  
G-coalgebra



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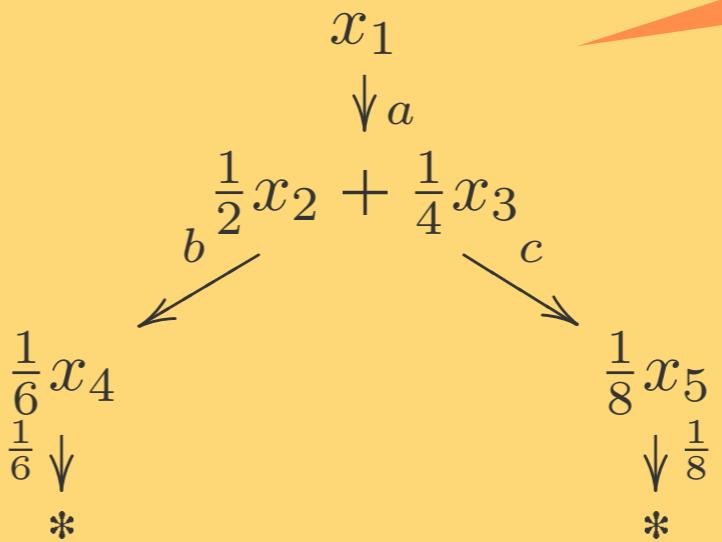
## PTS example



Sets  $\mathcal{D}_\omega$

$$\begin{array}{c}
 GX = [0, 1] \times X^A \\
 X \xrightarrow{\eta} \mathcal{D}_\omega(X) \\
 c \downarrow \qquad \qquad \qquad (\delta \circ c)^\# \downarrow \\
 \mathcal{D}_\omega(1 + A \times X) \xrightarrow{\delta} G\mathcal{D}_\omega(X)
 \end{array}$$

## Its determinization



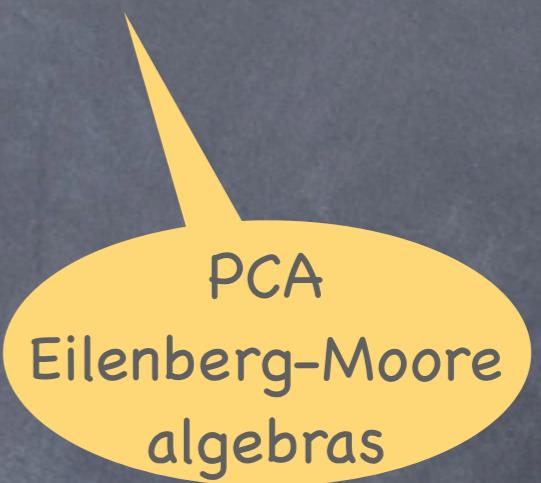
a  $G$ -coalgebra

final  
 $G$ -coalgebra

$$\begin{array}{ccc}
 X & \xrightarrow{out} & [0, 1]^{A^*} \\
 c \downarrow & & \downarrow \cong \\
 GX & \xrightarrow{Gout} & G([0, 1]^{A^*})
 \end{array}$$

# A way out for traces

Trace case, almost G-coalgebras on  $\text{Sets}^{\mathcal{D}_\omega}$



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Trace case, almost G-coalgebras on  $\text{Sets}^{\mathcal{D}_\omega}$

$$\begin{array}{ccc} \text{Exp}/\equiv & \xrightarrow{\quad out \quad} & \Omega_{\hat{G}} \\ \downarrow d & & \downarrow \simeq \\ \hat{G}\text{Exp}/\equiv & \xrightarrow{\quad \hat{G}out \quad} & \hat{G}\Omega_{\hat{G}} \end{array}$$

PCA  
Eilenberg-Moore  
algebras

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Trace case, almost G-coalgebras on  $\text{Sets}^{\mathcal{D}_\omega}$

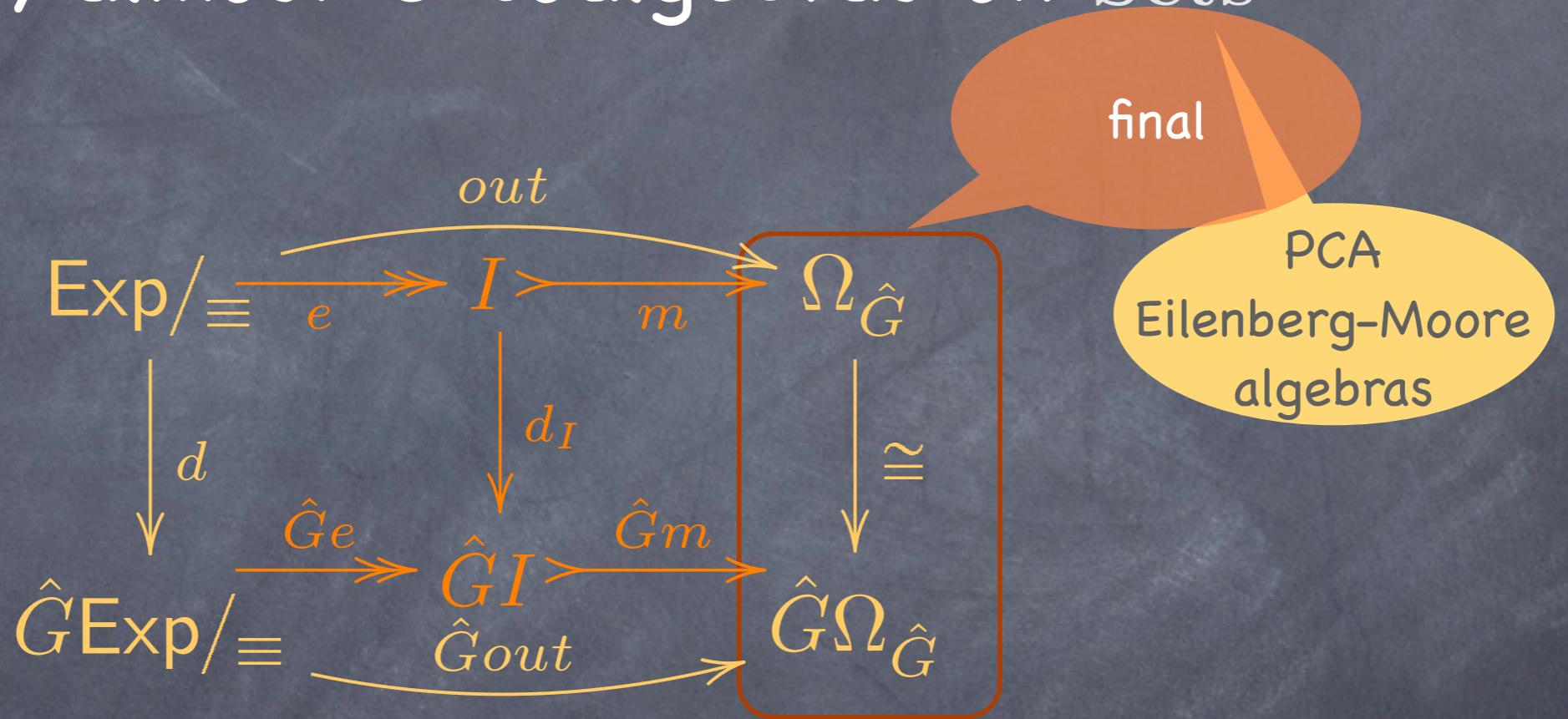
$$\begin{array}{ccccc} \text{Exp}/\equiv & \xrightarrow{e} & I & \xrightarrow{m} & \Omega_{\hat{G}} \\ \downarrow d & & \downarrow d_I & & \downarrow \simeq \\ \hat{G}\text{Exp}/\equiv & \xrightarrow{\hat{G}e} & \hat{G}I & \xrightarrow{\hat{G}m} & \hat{G}\Omega_{\hat{G}} \end{array}$$

*out*

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# A way out for traces

Trace case, almost G-coalgebras on Sets $^{\mathcal{D}_\omega}$



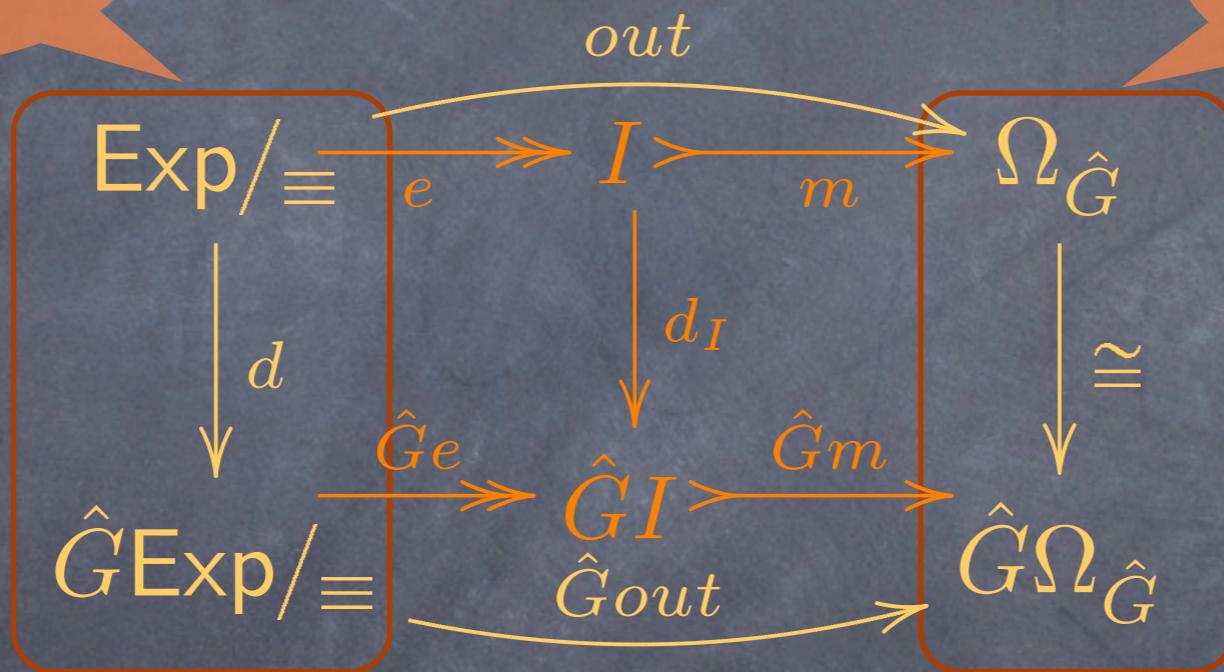
# A way out for traces

Trace case, almost G-coalgebras on  $\text{Sets}^{\mathcal{D}_\omega}$

final  
elsewhere

final

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# A way out for traces

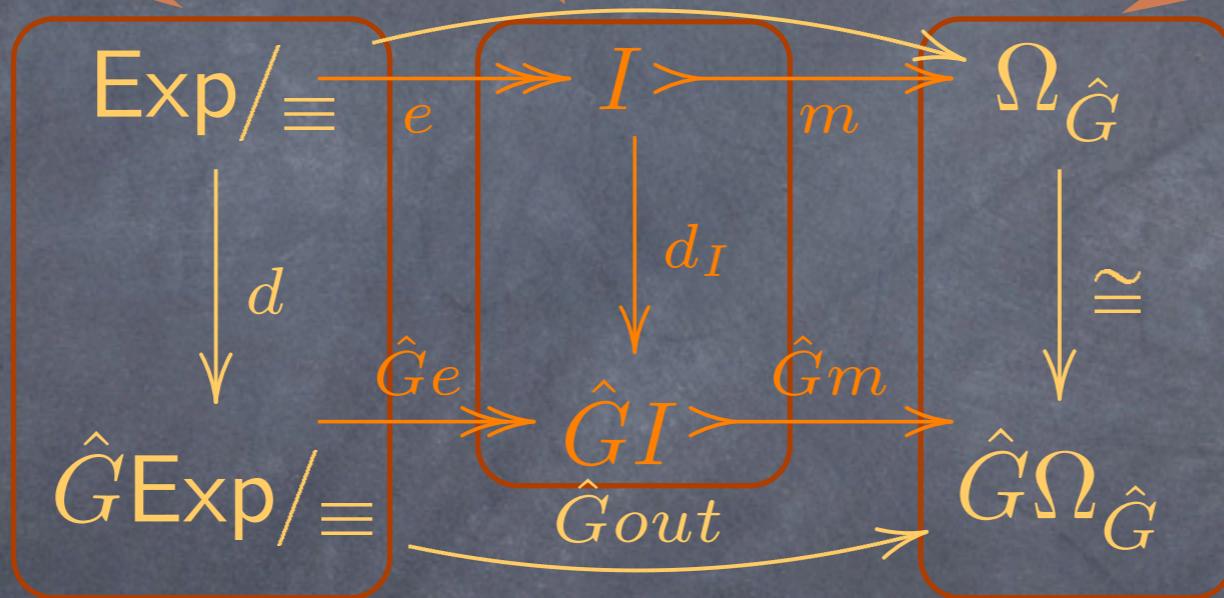
Trace case, almost  $G$ -coalgebras on  $\text{Sets}^{\mathcal{D}_\omega}$

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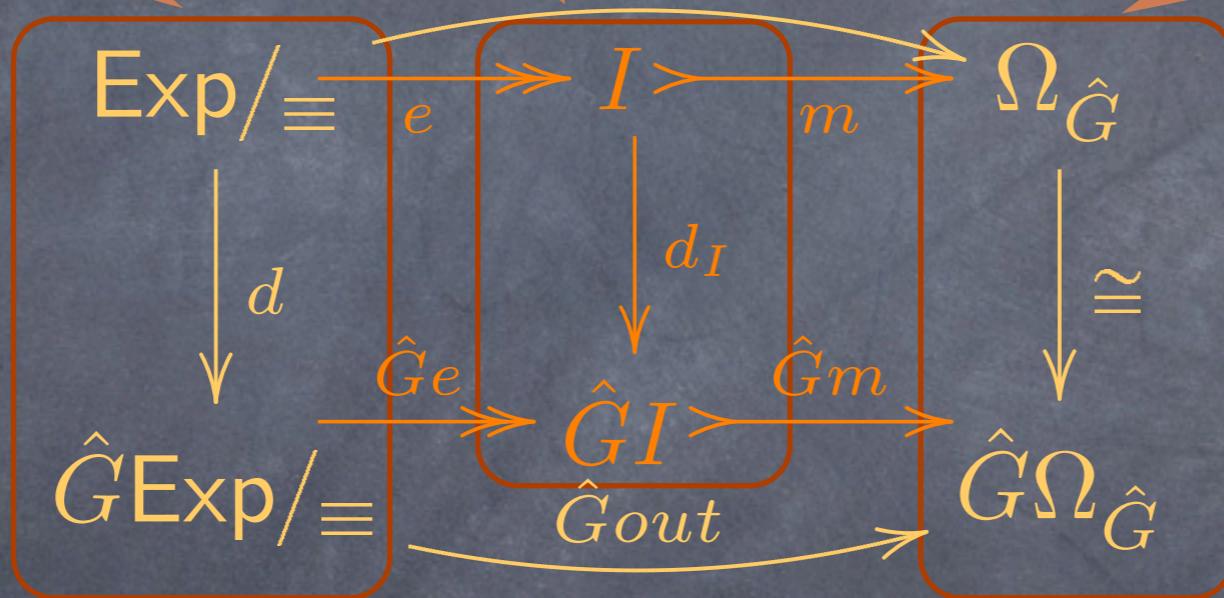
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Hence,  $e$  is iso, and  $\text{out}$  is injective

# A way out for traces

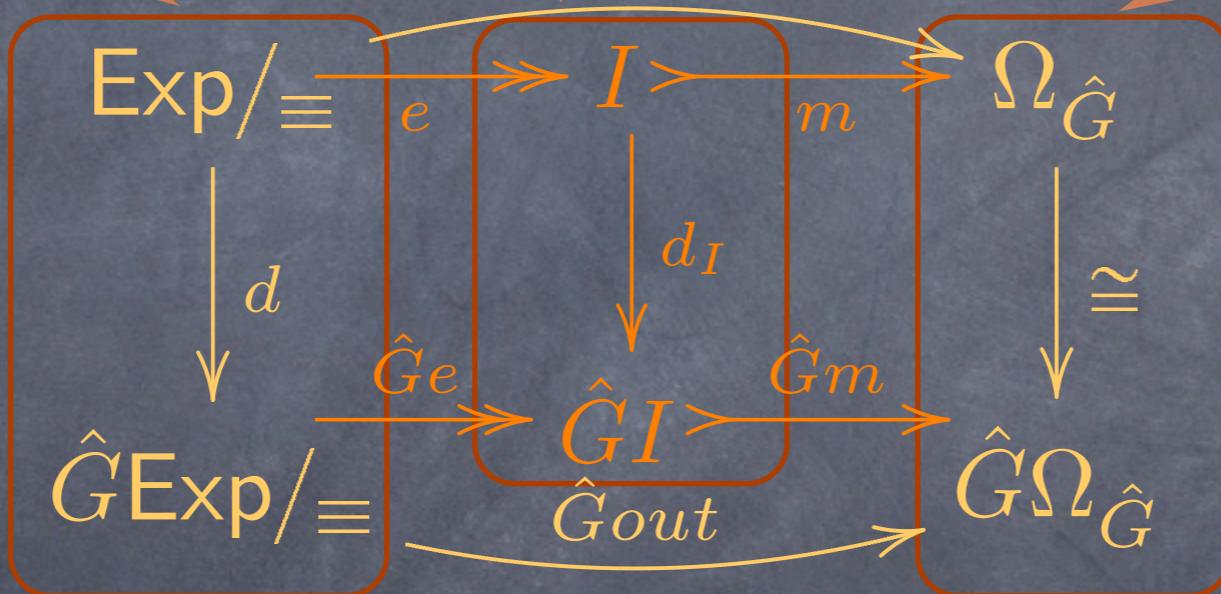
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Hence,  $e$  is iso, and  $out$  is injective

Moreover:  $tr = out \circ [-]$

# Conclusions

- we present a solution to a concrete problem

sound and complete axiomatization  
of traces for PTS

bisimilarity expressions and axioms plus one new axiom

- in a coalgebraic setting
- it opens many generalization questions...
- all about algebra and coalgebra

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Thank you !