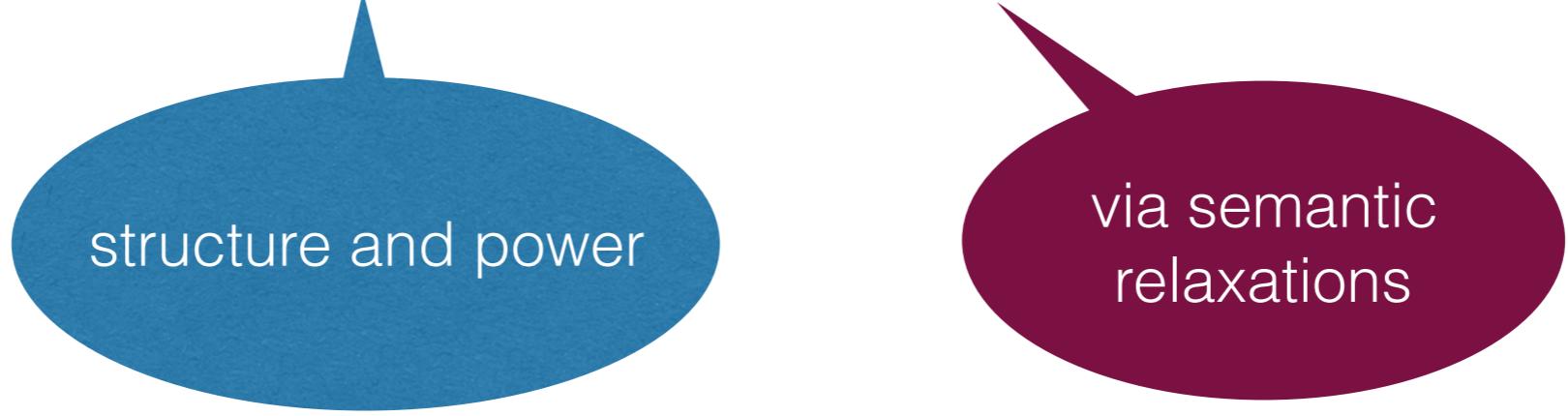


Semantics for Concurrency

Ana Sokolova  UNIVERSITY
of SALZBURG

Introduction - 5.3.19

- Intro: Concurrent data structures
correctness and performance



structure and power

via semantic
relaxations

Concurrent Data Structures Correctness and Relaxations



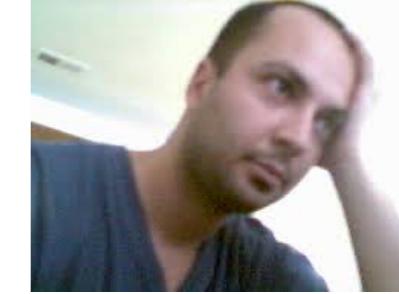
Hannes Payer
Google



Tom Henzinger
IST AUSTRIA



Christoph Kirsch
UNIVERSITY of SALZBURG



Ali Sezgin
UNIVERSITY OF CAMBRIDGE



Andreas Haas **Google**



Michael Lippautz



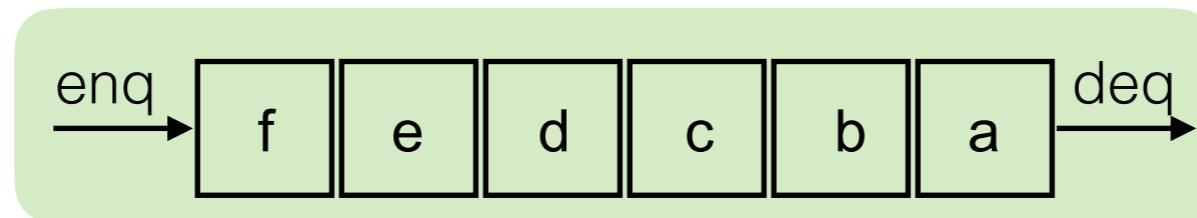
Andreas Holzer
Google



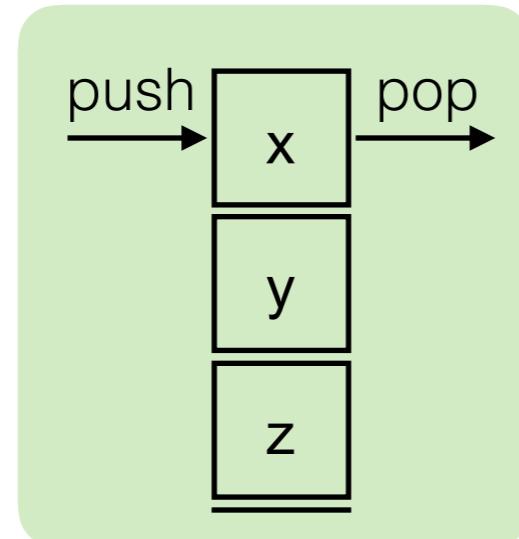
Helmut Veith
TU WIEN

Data structures

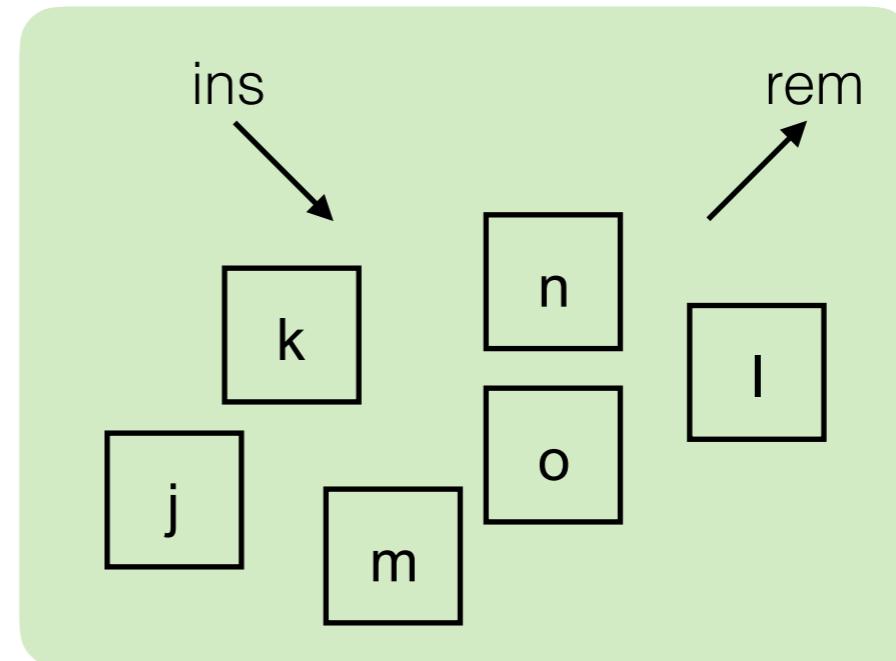
- Queue FIFO



- Stack LIFO

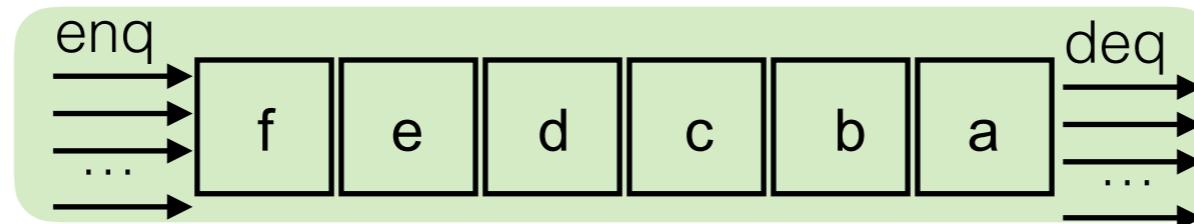


- Pool unordered

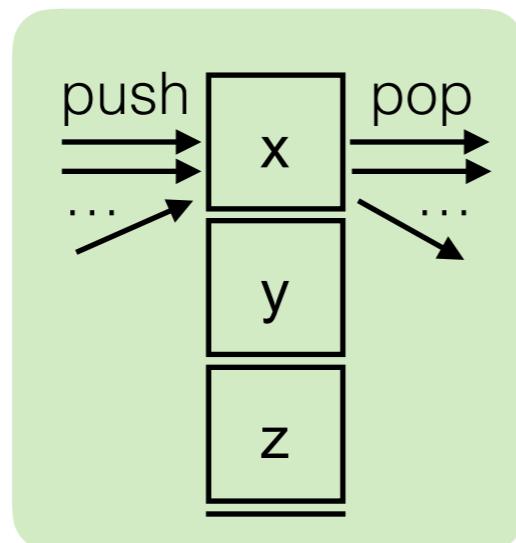


Concurrent data structures

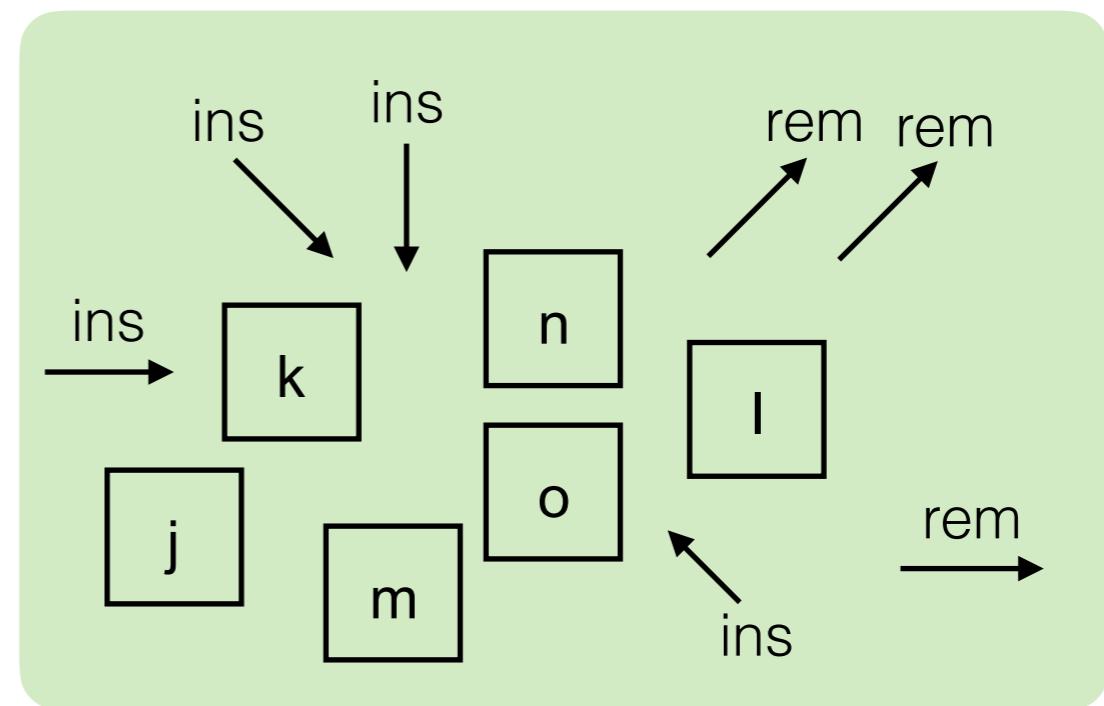
- Queue FIFO



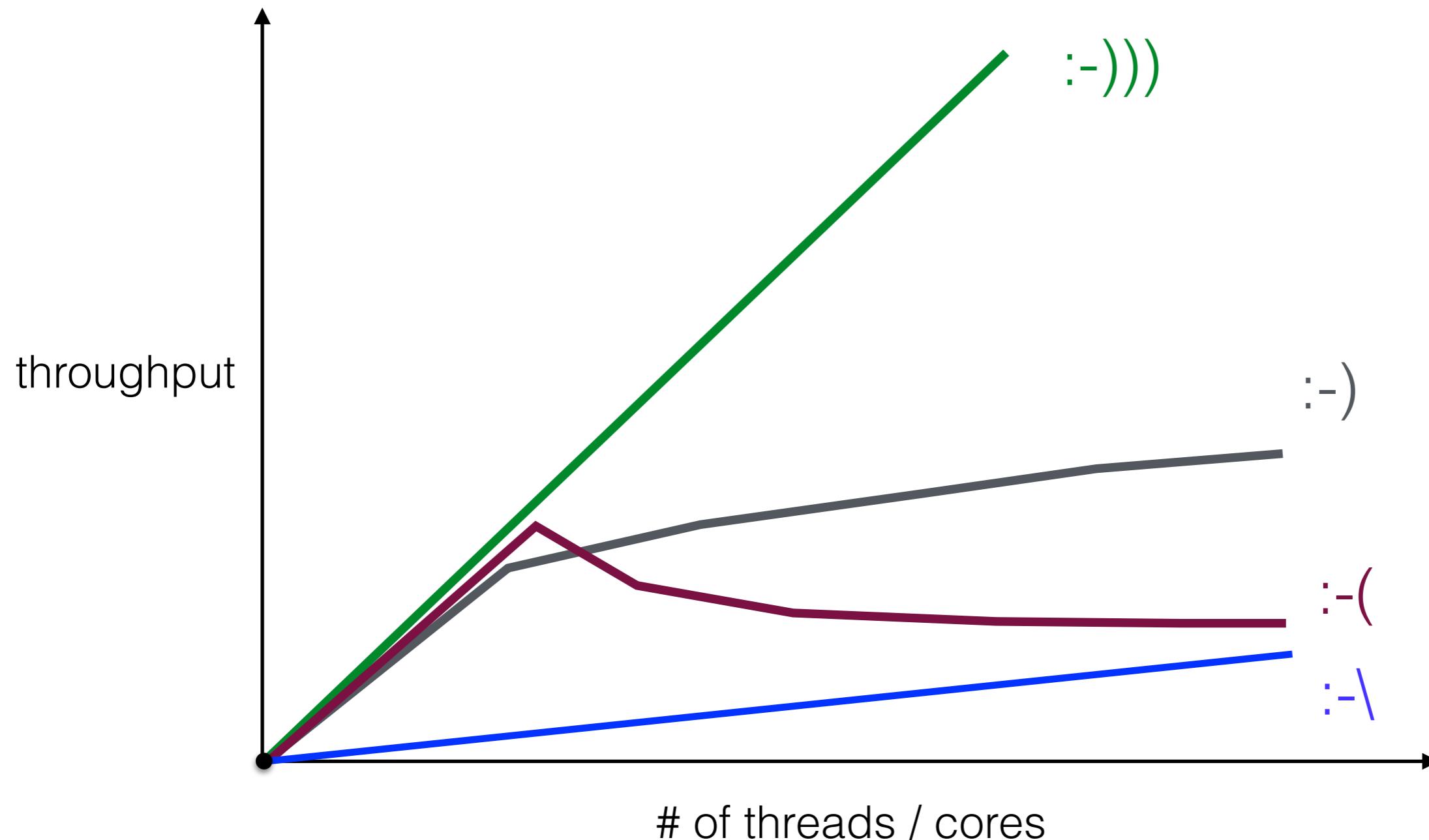
- Stack LIFO



- Pool unordered



Performance and scalability



Semantics of concurrent data structures

t1: enq(2) deq(1)
t2: enq(1) deq(2)

e.g. queues

- Sequential specification = set of legal sequences

e.g. queue legal sequence
enq(1)enq(2)deq(1)deq(2)

- Consistency condition = e.g. linearizability / sequential consistency

e.g. the concurrent history above is a linearizable queue concurrent history

Consistency conditions

there exists a legal sequence that preserves precedence order

consistency is about extending partial orders to total orders

there exists a legal sequence that preserves per-thread precedence (program order)

A history is ... wrt a sequential specification iff

Linearizability [Herlihy,Wing '90]



Sequential Consistency [Lamport'79]

