



Application Programming vs System Programming

Application Programming

- High level functions
- Using standard C library
- Man page level 3
- Libraries maintains buffers to read or write data from system calls. So libraries also called Buffered I/O.
- Libraries doesn't create performance penalty.

System Programming

- low level function
- Using System calls.
- Man page level 2
- System calls doesn't have buffers.
- System calls create performance penalty.

File System

- The file system
 - manages files,
 - allocating file space,
 - administrating free space,
 - controlling access to files
 - and retrieving data for user.
- The internal representation of file is given an ***inode table***.

OS vs File System

DOS: FAT32 (File Allocation Table)

Windows: FAT, NTFS (New Technology File System)

LINUX: ext2, ext3 and ext4 (Extended file system), JFS (Journaling file system), btrfs (b-tree File System)

File System Layout

Boot block	Super block	Inode list	Data Block
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Boot Block: contain bootstrap code that is read into machine to boot, or initialize, the operating system.

Super Block: Describes the state of a file system – How large it is, how many files it can store, where to find free space on the file system.

Inode (index node)list: inode represents the type of the file.

Data Block: contain file data and administrative data.

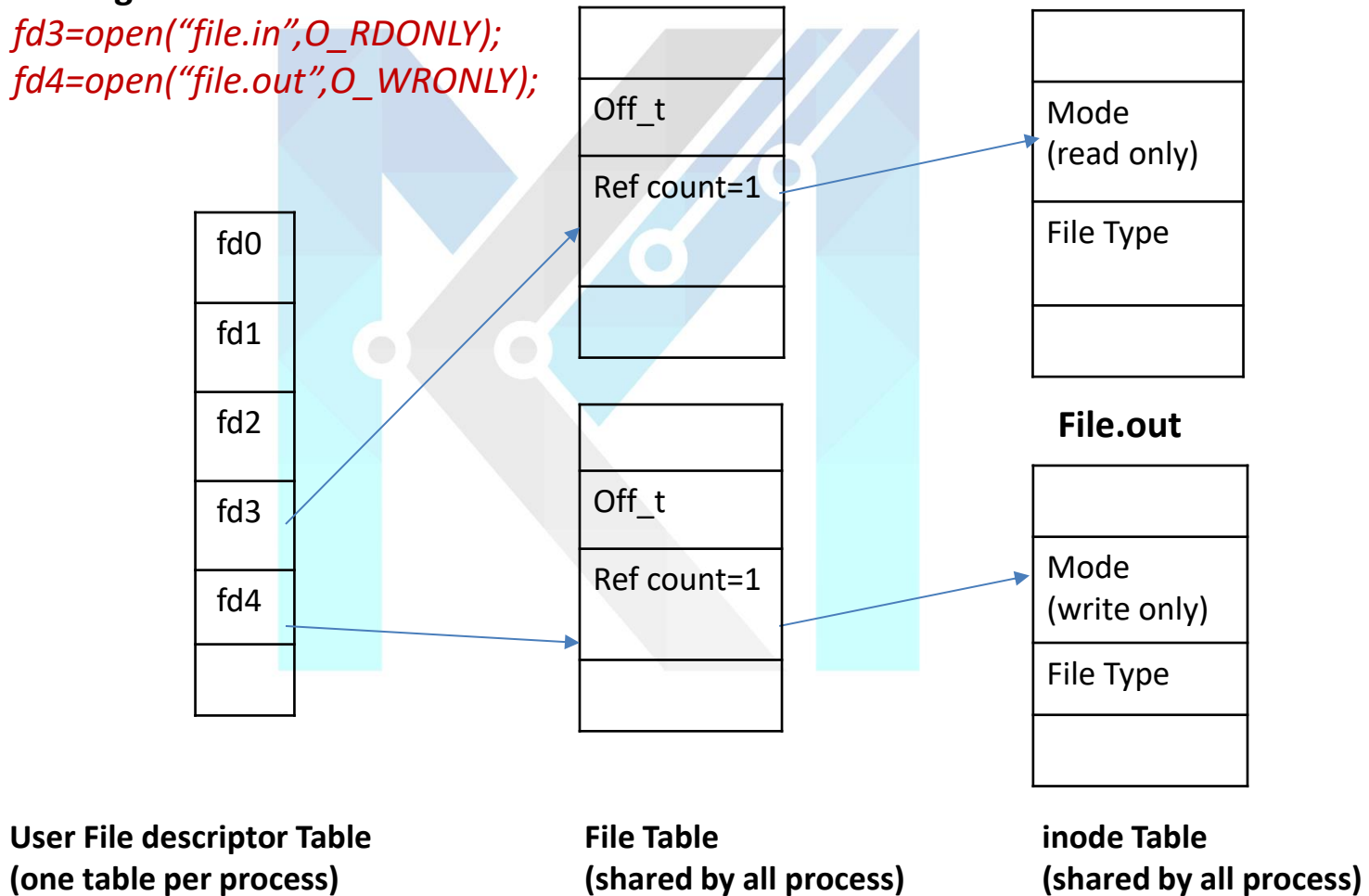
File System - Kernel Data Structures

- ***User File Descriptor table***: Each process has its own separate *descriptor table* whose entries are indexed by the process's open file descriptors. Each open descriptor entry points to an entry in the *file table*.
- ***File table***: Each file table entry consists of (for our purposes) the current file position, a *reference count* of the number of descriptor entries that currently point to it, and a pointer to an entry in the *inode table*.
- ***Inode table***: Each entry contains most of the information in the stat structure, including the st_mode and st_size members.
- ***File table*** is global kernel structure where as ***user file descriptor table*** is allocated per process.

File System - Kernel Data Structures

In this example, two descriptors reference distinct files. **There is no sharing.**

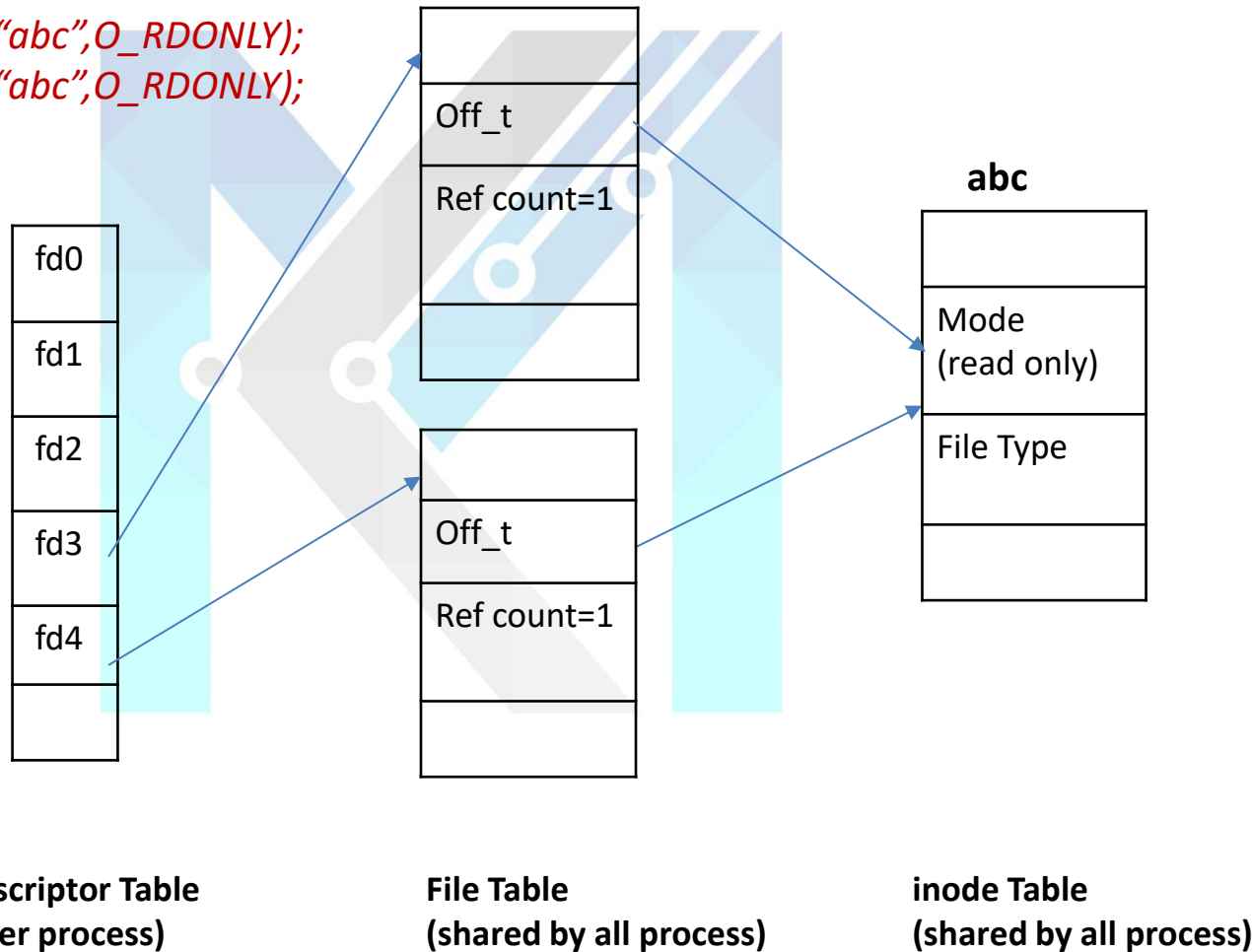
```
fd3=open("file.in",O_RDONLY);  
fd4=open("file.out",O_WRONLY);
```



File System - Kernel Data Structures

This example shows two descriptors **sharing the same disk file** through two open file table entries

```
fd3=open("abc",O_RDONLY);  
fd4=open("abc",O_RDONLY);
```



File System - Kernel Data Structures

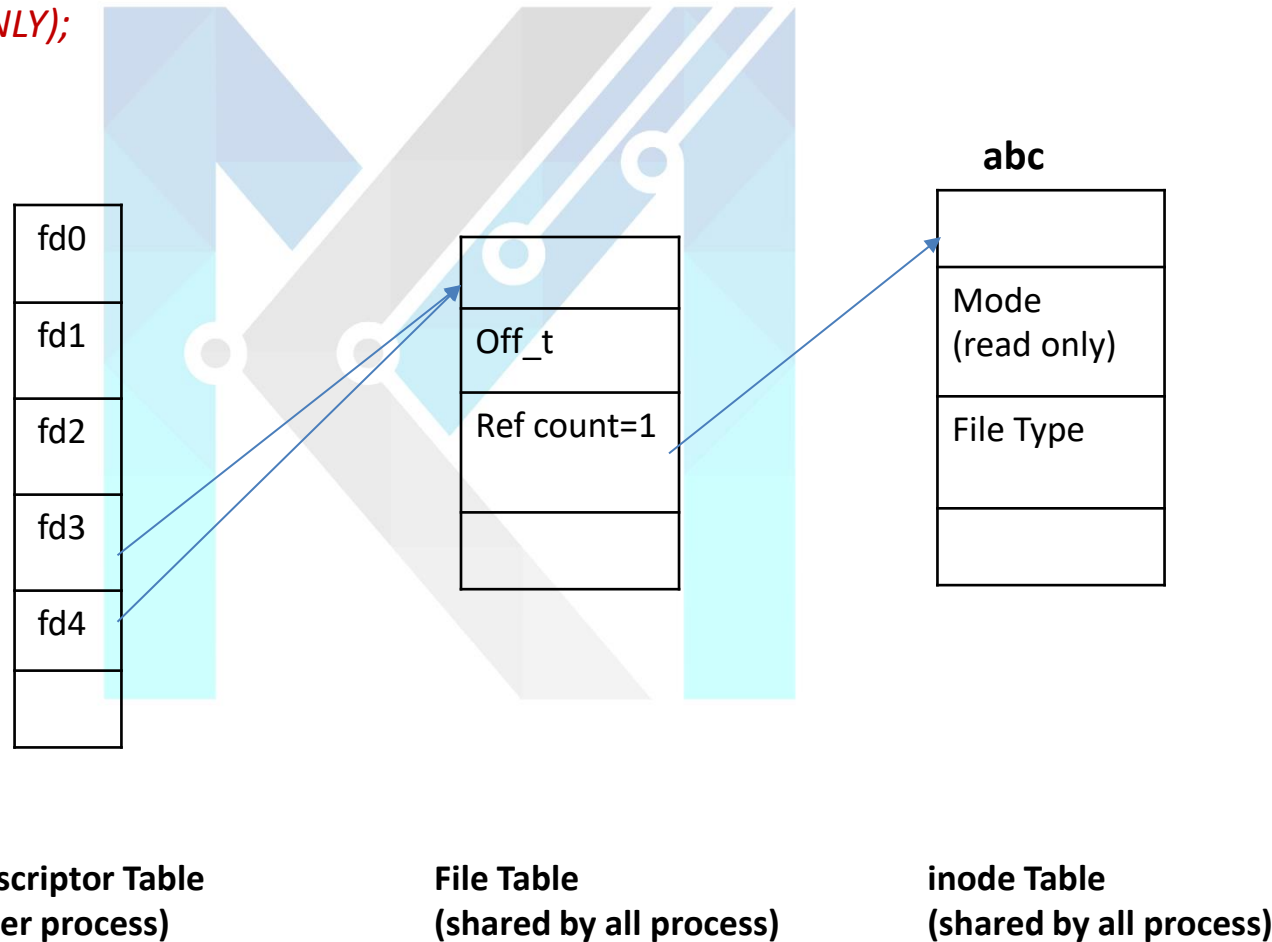
dup() Example:

`dup(int oldfd)`

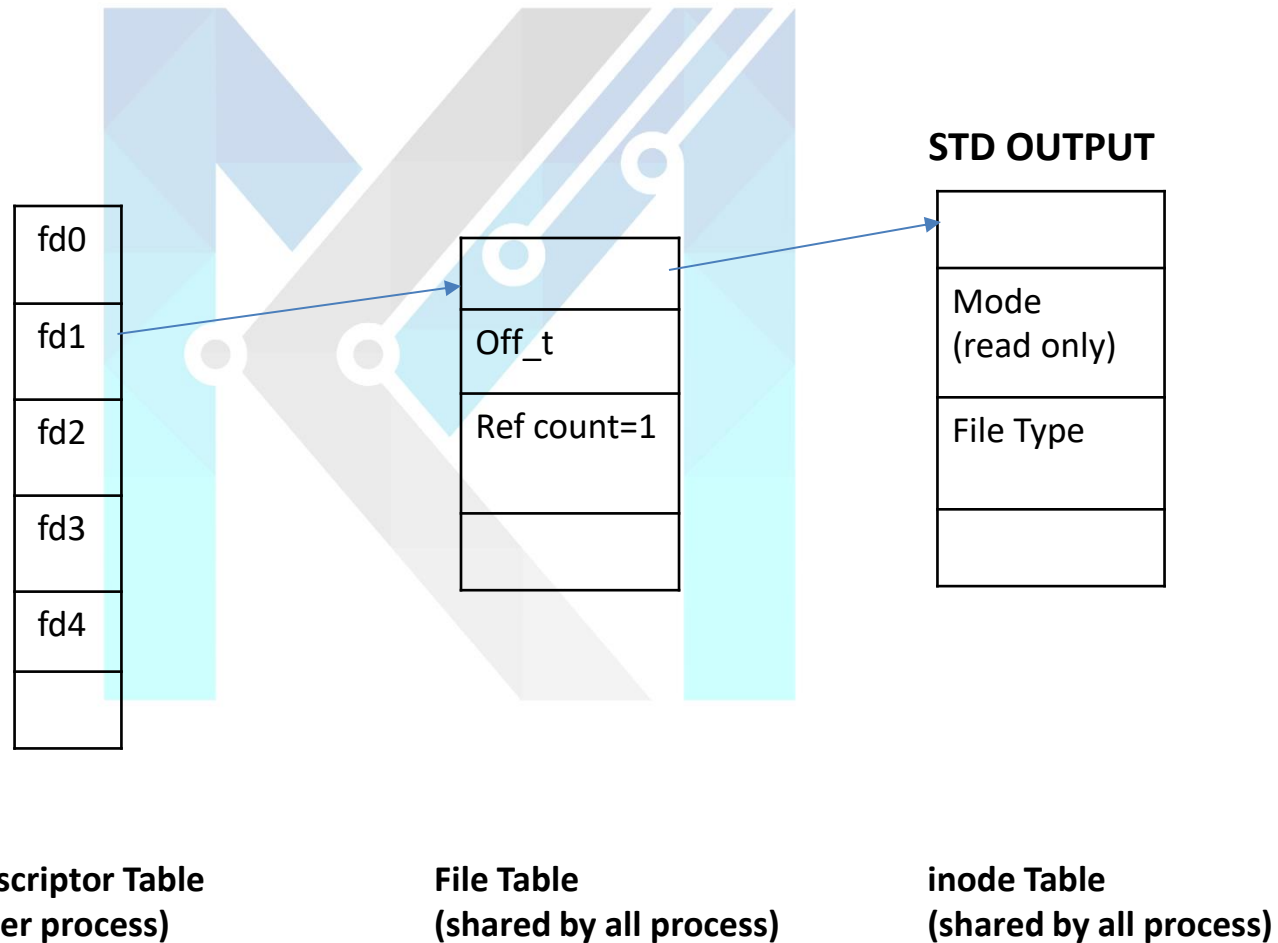
`fd=open("abc",O_RDONLY);`

`Close(fd);`

`dup_fd=dup(fd);`



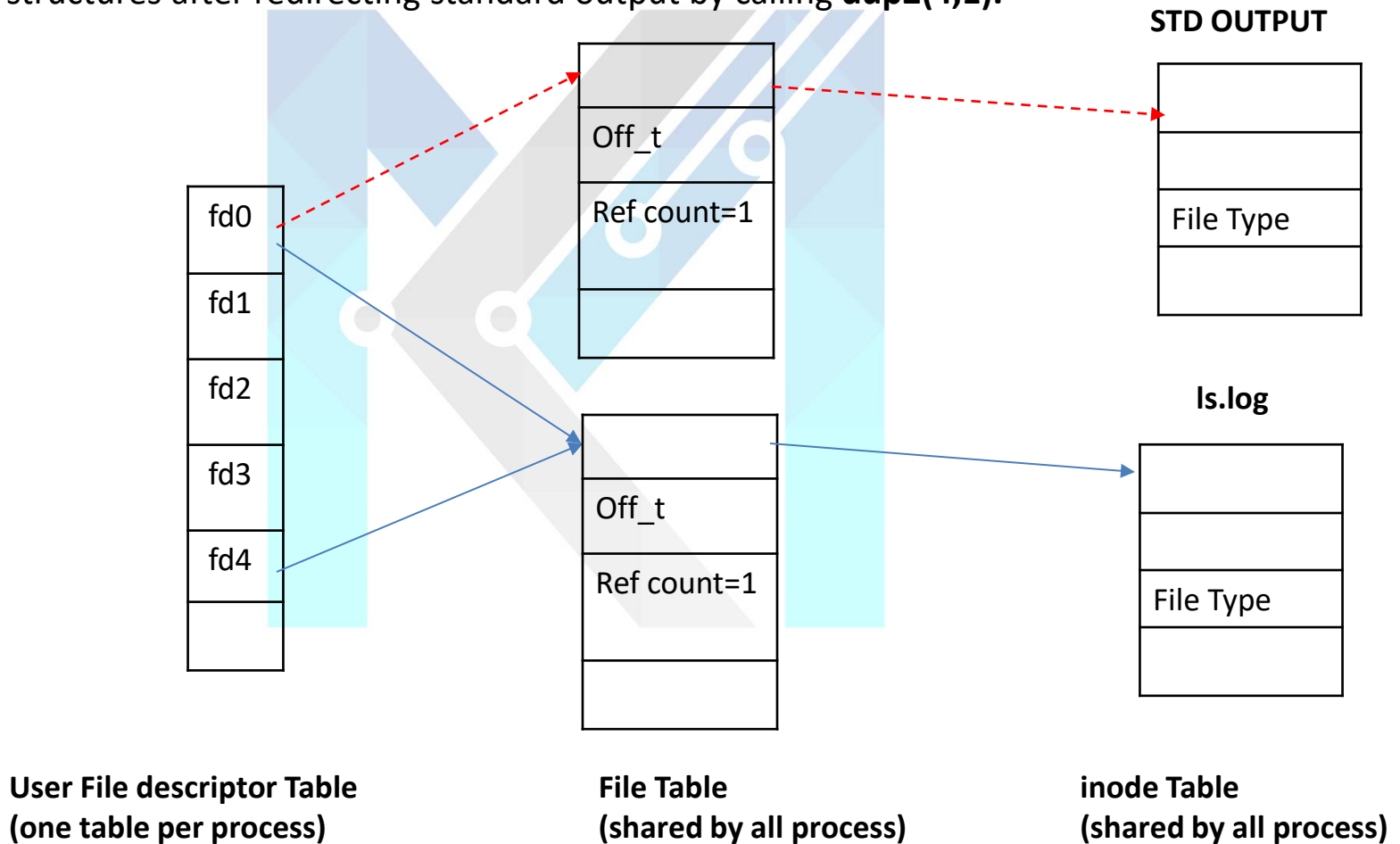
ls command o/p: STD OUT



I/O Redirection – **ls > ls.log** ?

dup2(int oldfd, int newfd) Example:

Kernel data structures after redirecting standard output by calling **dup2(4,1)**.



Monitoring File System Events



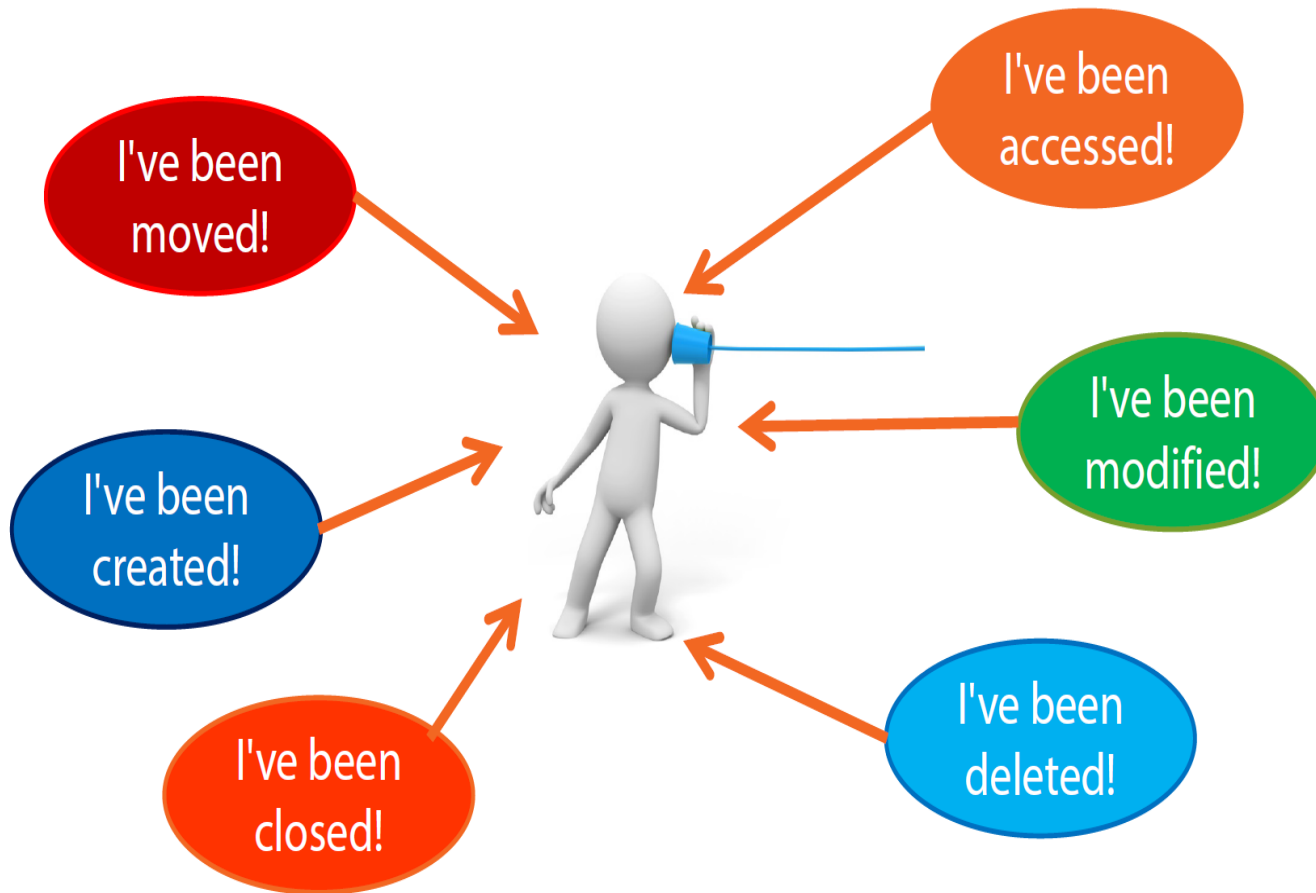
```
$ tail -f <file_name>
```

-follow: output appended data as the file grows.

Monitoring File System Events

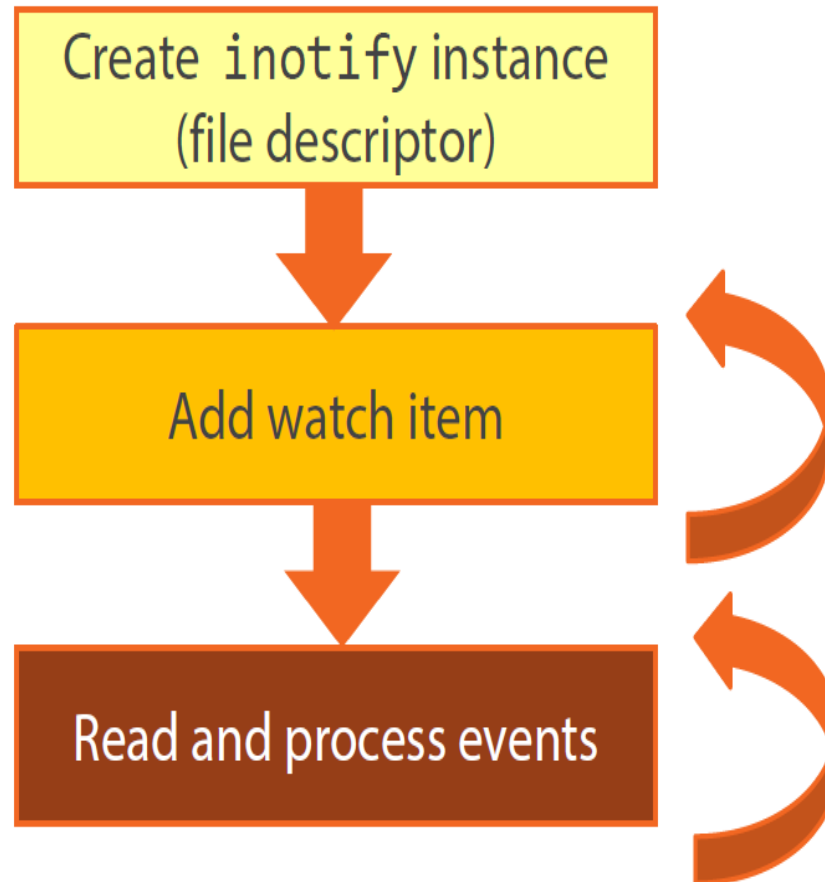
- The inotifyAPI
- Creating an inotifyinstance
- Adding to the watch list
- Reading events

Monitoring File System Events



Individual files or whole directories can be watched

Three Steps



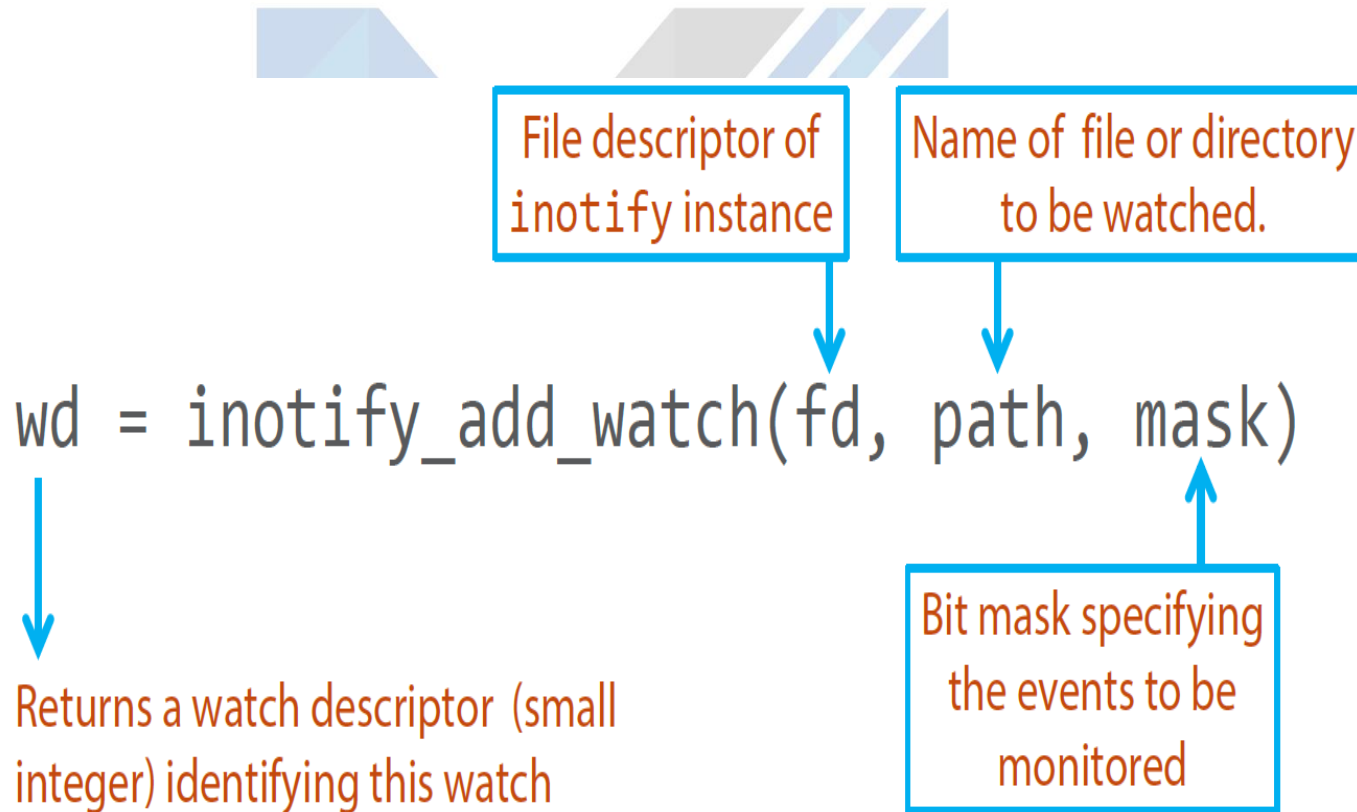
Creating an inotify Instance

```
fd = inotify_init()
```



↓
Returns a file descriptor on which we can later read() the events.

Adding a watch Item

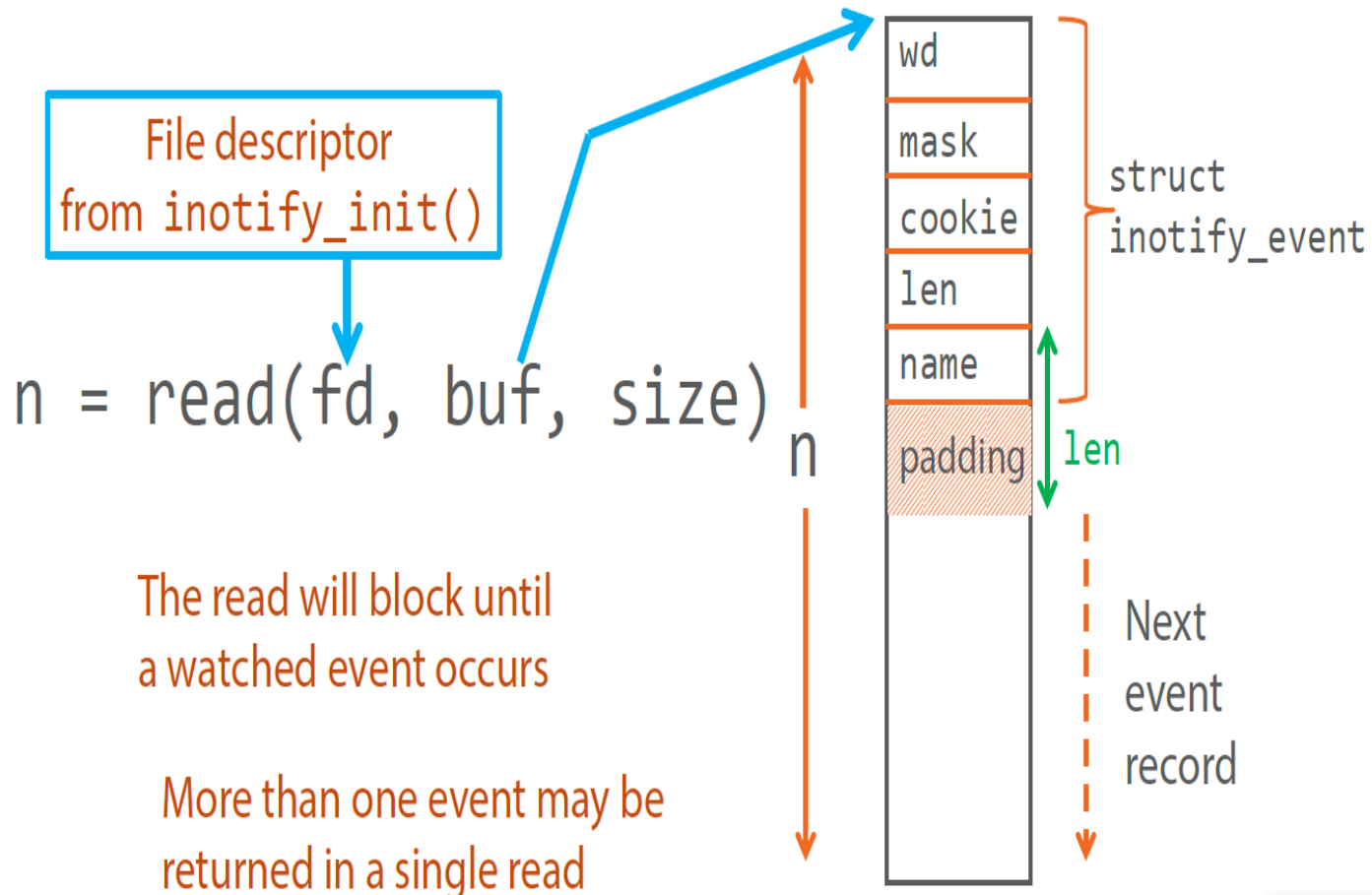


Watches and Events in Detail

Each event is specified by a single-bit constant
-- bitwise OR them together

Bit value	Meaning
IN_ACCESS	File was accessed
IN_ATTRIB	File attributes changed (ownership, permissions etc.)
IN_CREAT	File created inside watched directory
IN_DELETE	File deleted inside watched directory
IN_DELETE_SELF	Watched file deleted
IN_MODIFY	File was modified
IN_MOVE_SELF	File was moved

Reading Events



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