

# UNIT 5

## 1. Pushdown Automata (PDA)

- **Definition:** A PDA is a computational model that extends finite automata with an additional **stack** for memory, allowing it to recognize **context-free languages (CFLs)**.

### Components of PDA:

1. **States:** Finite set of states.
2. **Input Alphabet ( $\Sigma$  \Sigma):** Set of input symbols.
3. **Stack Alphabet ( $\Gamma$  \Gamma):** Set of symbols for the stack.
4. **Transitions:** Defined by a function, where a state, input, and stack top determine the next state, stack operation (push, pop, or no change).
5. **Initial State:** Starting state of the PDA.
6. **Start Stack Symbol:** The initial symbol on the stack.
7. **Final States:** Set of accepting states.

## 2. Representation of PDA

A PDA can be represented in two ways:

1. **Transition Function Representation:**  $\delta(q, a, X) = (p, \gamma)$   $\delta(q, a, X) = (p, \gamma)$ , where:
  - $q, p, \gamma$  = states
  - $a$  = input symbol
  - $X$  = stack top
  - $\gamma$  = new stack contents.
2. **Transition Diagram:** Visual representation with states and labeled transitions showing input, stack operations, and new states.

## 3. Acceptance by PDA

A PDA accepts a string if:

1. **Final State Acceptance:** PDA reaches a final state after processing the input and stack operations.
2. **Empty Stack Acceptance:** PDA empties the stack after processing the input.

## 4. Deterministic and Non-Deterministic PDA

1. **Deterministic PDA (DPDA):**
  - At most one transition for a given state, input, and stack top.
  - Less powerful than NDPDA (not all CFLs are accepted by DPDA).

## 2. Non-Deterministic PDA (NDPDA):

- Multiple transitions possible for a given state, input, and stack top.
- Equivalent in power to context-free languages.

### Comparison:

- **NDPDA** recognizes all CFLs; **DPDA** only recognizes a subset of CFLs.
- Deterministic parsing (LL and LR) typically uses DPDAs.

## 5. Pushdown Automata and Context-Free Languages

- Every CFL can be recognized by a PDA.
- Conversion between CFG and PDA:
  1. From CFG to PDA: For each production  $A \rightarrow \alpha$ , modify the PDA to push or pop stack symbols accordingly.
  2. From PDA to CFG: Define rules based on PDA transitions.

## 6. Closure Properties of CFLs

CFLs are closed under:

- **Union:** If  $L_1$  and  $L_2$  are CFLs,  $L_1 \cup L_2$  is also a CFL.
- **Concatenation:**  $L_1 L_2$  is a CFL.
- **Kleene Star:**  $L^*$  is a CFL.

Not closed under:

- **Intersection**
- **Complementation**

## 7. LL(k) and LR(k) Grammars

### LL(k) Grammars:

- Used for **Top-Down Parsing**.
- Parse input from **Left to Right** and produce **Leftmost Derivation**.
- $k$ : Number of lookahead tokens used to decide the next action.
- Properties:
  - Simpler and faster parsing.
  - Cannot handle left recursion or ambiguous grammars.

## **LR(k) Grammars:**

- Used for **Bottom-Up Parsing**.
- Parse input from **Left to Right** and produce **Rightmost Derivation** in reverse.
- k: Number of lookahead tokens used.
- Properties:
  - Handles larger class of grammars (including left recursion).
  - More complex but efficient for parsing programming languages.

## **8. Parsing**

### **Top-Down Parsing:**

- Starts from the start symbol and tries to derive the input string.
- Includes:
  - **Recursive Descent Parsing**
  - **Predictive Parsing** (LL(1) parsers)

### **Bottom-Up Parsing:**

- Starts with the input string and reduces it to the start symbol.
- Includes:
  - **Shift-Reduce Parsing**
  - **LR Parsing** (Simple LR, Canonical LR, Lookahead LR).

## Important MCQ

### Pushdown Automata (PDA)

1. A PDA is a finite automaton with:

- (A) A queue
- (B) A stack
- (C) A tape
- (D) A counter

**Answer: B**

2. The stack in a PDA is used for:

- (A) Temporary storage of input symbols
- (B) Keeping track of context
- (C) Both A and B
- (D) None of the above

**Answer: C**

3. Which language cannot be accepted by a PDA?

- (A)  $\{a^n b^n \mid n \geq 1\}$
- (B)  $\{ww^R \mid w \in \{a, b\}^*\}$
- (C)  $\{a^n b^m c^n \mid n, m \geq 0\}$
- (D) None of the above

**Answer: B**

4. A PDA accepts a language if:

- (A) It reaches a final state
- (B) Its stack becomes empty
- (C) Both A and B
- (D) Either A or B

**Answer: D**

5. Which of the following languages is context-free?

- (A)  $\{a^n b^n c^n \mid n \geq 1\}$
- (B)  $\{a^n b^m c^{n+m} \mid n, m \geq 1\}$
- (C)  $\{a^i b^j \mid i \neq j\}$
- (D)  $\{a^i b^j \mid i = j\}$

**Answer: D**

6. The transition function of a PDA is of the form:

- (A)  $\delta(q, a) = p$
- (B)  $\delta(q, a, X) = (p, \gamma)$
- (C)  $\delta(q, a, X) = p$
- (D)  $\delta(q, X) = (p, a)$

**Answer: B**

7. PDAs recognize:

- (A) Regular languages
- (B) Context-free languages
- (C) Context-sensitive languages
- (D) All of the above

**Answer: B**

8. What is the key difference between DPDA and NDPDA?

- (A) DPDAs can accept more languages
- (B) NDPDAs can accept more languages
- (C) Both accept the same languages
- (D) None of the above

**Answer: B**

9. The stack in a PDA can hold symbols from:

- (A) Input alphabet
- (B) Stack alphabet
- (C) Both A and B
- (D) None of the above

**Answer: B**

10. A PDA accepting by empty stack:

- (A) Always has final states
- (B) Does not require final states
- (C) Cannot process infinite strings
- (D) Accepts only finite languages

**Answer: B**

### **Context-Free Languages and PDA**

11. Every PDA recognizes:

- (A) A regular language
- (B) A context-free language
- (C) A context-sensitive language
- (D) An undecidable language

**Answer: B**

12. CFGs are equivalent to:

- (A) Finite automata
- (B) PDAs
- (C) Turing machines
- (D) None of these

**Answer: B**

13. The stack is used in PDA to:

- (A) Store current input symbols
- (B) Track context of grammar rules
- (C) Store output symbols
- (D) None of the above

**Answer: B**

14. A DPDA cannot recognize:

- (A) Deterministic CFLs
- (B) Non-deterministic CFLs
- (C) Regular languages
- (D) None of these

**Answer: B**

15. A language is context-free if and only if:

- (A) It can be generated by a CFG
- (B) It can be recognized by a PDA
- (C) Both A and B
- (D) None of these

**Answer: C**

### **Closure Properties of CFLs**

16. CFLs are closed under:

- (A) Union
- (B) Concatenation
- (C) Kleene star
- (D) All of the above

**Answer: D**

17. CFLs are not closed under:

- (A) Intersection
- (B) Complementation
- (C) Both A and B
- (D) None of these

**Answer: C**

18. Intersection of a CFL and a regular language is:

- (A) Regular
- (B) Context-free
- (C) Context-sensitive
- (D) None of these

**Answer: B**

19. Complement of a CFL is always:

- (A) Regular
- (B) CFL
- (C) Non-CFL
- (D) None of these

**Answer: C**

20. CFLs are closed under:

- (A) Homomorphism
- (B) Inverse homomorphism
- (C) Both A and B
- (D) None of the above

**Answer: C**

### **LL(k) and LR(k) Grammars**

21. LL(k) parsing is a type of:

- (A) Top-down parsing
- (B) Bottom-up parsing
- (C) Shift-reduce parsing
- (D) None of these

**Answer: A**

22. LR(k) parsing is:

- (A) Predictive parsing
- (B) Bottom-up parsing
- (C) Top-down parsing
- (D) None of these

**Answer: B**

23. LL(k) parsers cannot handle:

- (A) Left recursion
- (B) Ambiguous grammars
- (C) Both A and B
- (D) None of these

**Answer: C**

24. LR(k) parsers are:

- (A) More powerful than LL(k) parsers
- (B) Faster than LL(k) parsers
- (C) Suitable only for regular languages
- (D) None of these

**Answer: A**

25. The lookahead kkk in LL(k) and LR(k) parsers represents:

- (A) Number of tokens used to decide next move
- (B) Depth of recursion
- (C) Number of production rules
- (D) None of the above

**Answer: A**

### **Parsing Techniques**

26. Top-down parsing starts with:

- (A) Start symbol of the grammar
- (B) Input string
- (C) Both A and B
- (D) None of these

**Answer: A**

27. Bottom-up parsing starts with:

- (A) Start symbol of the grammar
- (B) Input string
- (C) Both A and B
- (D) None of these

**Answer: B**

28. Predictive parsing is a type of:

- (A) Top-down parsing
- (B) Bottom-up parsing
- (C) Shift-reduce parsing
- (D) None of these

**Answer: A**

29. The main limitation of top-down parsers is:

- (A) They cannot handle left recursion
- (B) They cannot handle right recursion
- (C) They are slow
- (D) None of the above

**Answer: A**

30. LR parsers use:

- (A) Shift-reduce parsing technique
- (B) Recursive descent
- (C) Predictive parsing
- (D) None of these

**Answer: A**



## Parsing Techniques

31. In shift-reduce parsing, a shift action means:

- (A) Adding a production rule
- (B) Moving input symbol onto the stack
- (C) Reducing stack contents
- (D) None of these

**Answer: B**

32. A reduce action in shift-reduce parsing involves:

- (A) Moving input symbol onto the stack
- (B) Applying a production rule to replace symbols on the stack
- (C) Removing all stack contents
- (D) None of these

**Answer: B**

33. Which parser uses a parsing table?

- (A) Recursive descent parser
- (B) LL parser
- (C) LR parser
- (D) Both B and C

**Answer: D**

34. Bottom-up parsing can also be called:

- (A) Predictive parsing
- (B) Shift-reduce parsing
- (C) Recursive descent parsing
- (D) None of these

**Answer: B**

35. Which of the following grammars can be parsed by an LR(1) parser but not an LL(1) parser?

- (A) Left-recursive grammars
- (B) Ambiguous grammars
- (C) Right-recursive grammars
- (D) None of these

**Answer: A**

## Pushdown Automata Properties

36. The difference between NDPDA and DPDA lies in:

- (A) The type of stack
- (B) The number of stacks used
- (C) The number of transitions allowed for a single input
- (D) None of these

**Answer: C**

37. How does a PDA accept by final state?

- (A) When the PDA halts
- (B) When PDA reaches a pre-defined accepting state
- (C) When PDA empties the stack
- (D) Both B and C

**Answer: B**

38. A deterministic PDA is:

- (A) More powerful than a non-deterministic PDA
- (B) Equivalent in power to an NDPDA
- (C) Less powerful than an NDPDA
- (D) None of these

**Answer: C**

39. A language accepted by an NDPDA is:

- (A) Always regular
- (B) Always context-free
- (C) Always deterministic
- (D) None of these

**Answer: B**

40. Which of the following statements is true for PDAs?

- (A) They can recognize all context-free languages.
- (B) They can recognize some context-sensitive languages.
- (C) Both A and B.
- (D) None of these.

**Answer: A**

### **LL(k) and LR(k) Parsing Techniques**

41. An LL(1) parser cannot handle:

- (A) Right recursion
- (B) Left recursion
- (C) Ambiguity
- (D) Both B and C

**Answer: D**

42. What is a major advantage of LR parsers?

- (A) Simpler implementation
- (B) Ability to parse a larger class of grammars
- (C) Does not require lookahead
- (D) None of these

**Answer: B**

43. LL parsers derive strings by:

- (A) Rightmost derivation
- (B) Rightmost derivation in reverse
- (C) Leftmost derivation
- (D) None of these

**Answer: C**

44. LR parsers derive strings by:

- (A) Rightmost derivation
- (B) Rightmost derivation in reverse
- (C) Leftmost derivation
- (D) None of these

**Answer: B**

45. The main reason for using lookahead in parsers is to:

- (A) Improve speed
- (B) Resolve ambiguities
- (C) Handle left recursion
- (D) None of these

**Answer: B**

### **Grammar and Parsing Relations**

46. A grammar is said to be ambiguous if:

- (A) It has no derivation for some strings
- (B) It has two different parse trees for the same string
- (C) It cannot be converted to a PDA
- (D) None of the above

**Answer: B**

47. Top-down parsing can fail when:

- (A) The grammar is ambiguous
- (B) The grammar is left-recursive
- (C) The grammar is right-recursive
- (D) Both A and B

**Answer: D**

48. A CFG can always be converted to:

- (A) An ambiguous grammar
- (B) A PDA
- (C) A deterministic grammar
- (D) None of the above

**Answer: B**

49. The stack content of a PDA can represent:

- (A) Leftmost derivations
- (B) Rightmost derivations
- (C) Both A and B
- (D) None of these

**Answer: C**

50. Parsing is used in compilers to:

- (A) Translate source code to machine code
- (B) Generate an intermediate representation
- (C) Analyze syntax structure of the source code
- (D) Both B and C

**Answer: D**