

CS150 Homework 4
due Wednesday, November 17th 5:00 PM

Problem 1.

(a) Construct a CFG for the set of all binary strings of the form $0^i1^j0^k$, where $i+k = j$.

1. $S \rightarrow AB$
2. $A \rightarrow 0A1 \mid \epsilon$
2. $B \rightarrow 1B0 \mid \epsilon$

(b) $0^i1^j0^k$, where $i+k = 2j$.

1. $S \rightarrow AB \mid 0A1B0$
(first prod: i and k are even, second prod: i and k are odd; no other option otherwise $i + k$ is odd)
2. $A \rightarrow 0A1 \mid \epsilon$
2. $B \rightarrow 1B00 \mid \epsilon$

Problem 2.

For the grammar:

$S \rightarrow A1B$

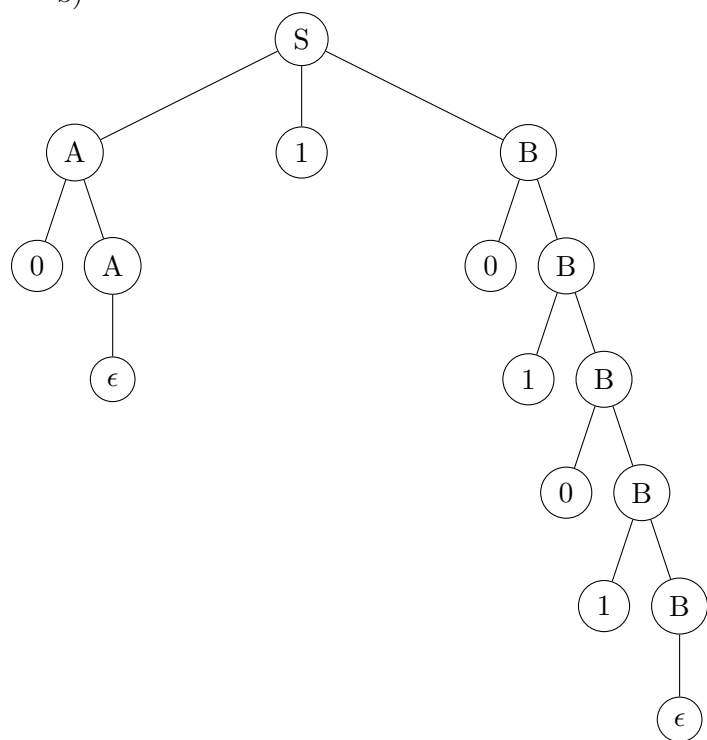
$A \rightarrow 0A \mid \epsilon$

$B \rightarrow 0B \mid 1B \mid \epsilon$

and string 010101, give the a) leftmost derivation and b) parse tree.

a) $S \Rightarrow A1B \Rightarrow 0A1B \Rightarrow 01B \Rightarrow 010B \Rightarrow 0101B \Rightarrow 01010B \Rightarrow 010101B \Rightarrow 010101$

b)



Problem 3.

For this grammar,

$$S \rightarrow A1B$$

$$A \rightarrow 0A \mid \epsilon$$

$$B \rightarrow 0B \mid 1B \mid \epsilon$$

a) Show that this grammar is unambiguous.

Algorithm for reading input string with production:

- 1) Scan the string from left-to-right
- 2) There is only 1 production for expanding S, don't read any input when expanding S.
- 3) If it's required to expand A, use $A \rightarrow 0A$ if the next symbol is a 0, otherwise if it's 1 or at the end, use $A \rightarrow \epsilon$
- 4) If it's required to expand B, use $B \rightarrow 0B$ if the next symbol is a 0, use $B \rightarrow 1B$ if the next symbol is a 1, and use $B \rightarrow \epsilon$ if at the end of input string.

This will produce a unique leftmost derivation for any given input string.

Input string: 001010

Remaining input	Steps of left most derivation
001010	S
001010	A1B
01010	0A1B
1010	00A1B
1010	001B
010	001B
10	0010B
0	00101B
	001010B
	001010

Therefore, this grammar is unambiguous.

b) Find a grammar for the same language that is ambiguous, and demonstrate its ambiguity.

$$S \rightarrow A1B$$

$$A \rightarrow 0A \mid \epsilon$$

$$B \rightarrow 0B \mid B0 \mid 1B \mid \epsilon$$

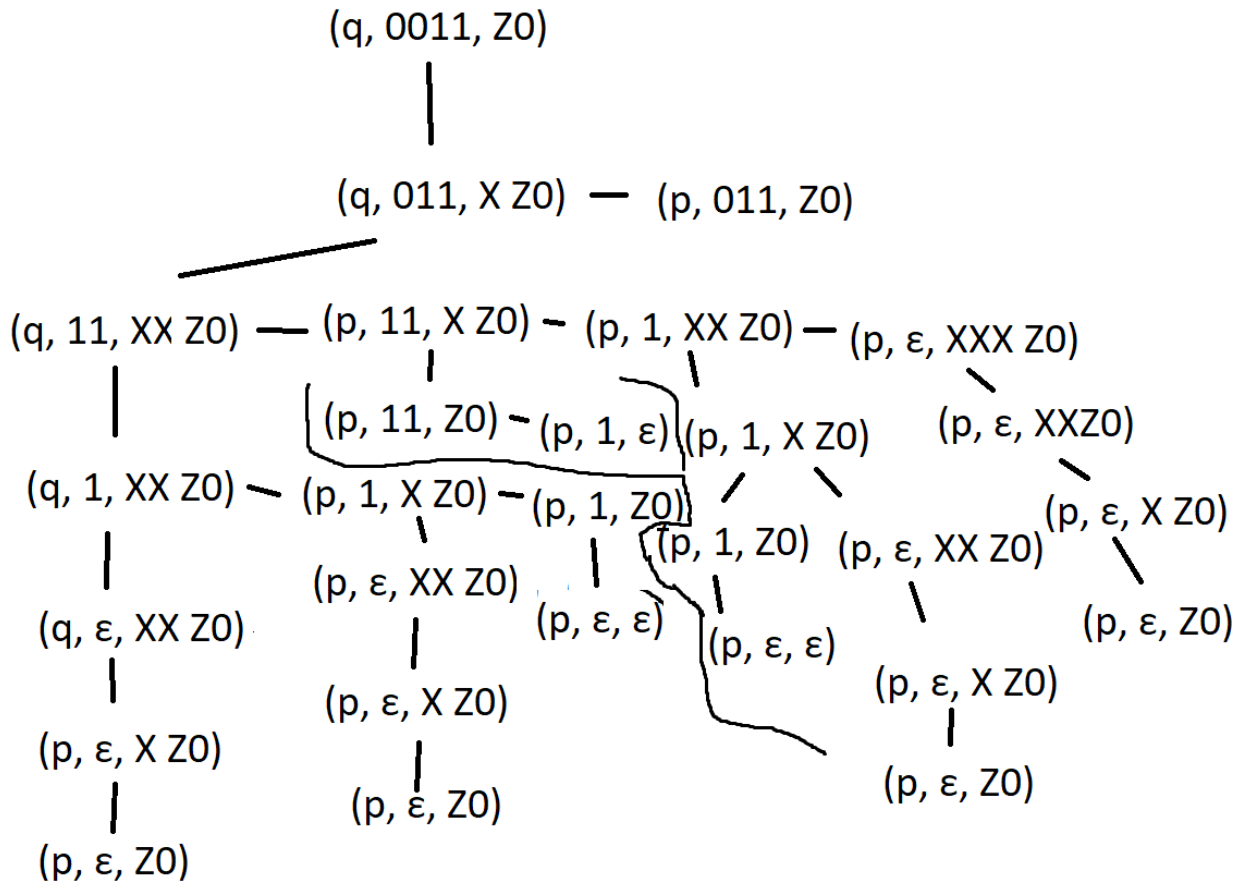
For string $w = 10$, there exist 2 left most derivations as follows:

$$1) S \Rightarrow A1B \Rightarrow 1B \Rightarrow 10B \Rightarrow 10$$

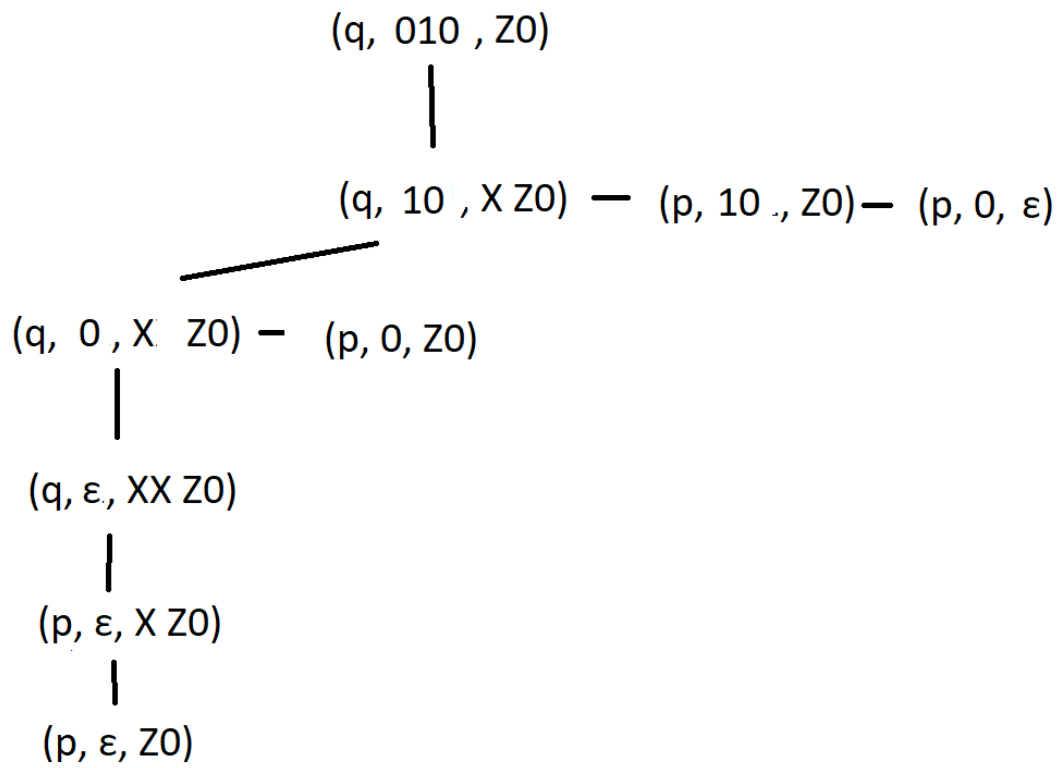
$$2) S \Rightarrow A1B \Rightarrow 1B \Rightarrow 1B0 \Rightarrow 10$$

Thus, this grammar is ambiguous.

b) Show all reachable IDs with input w 0011

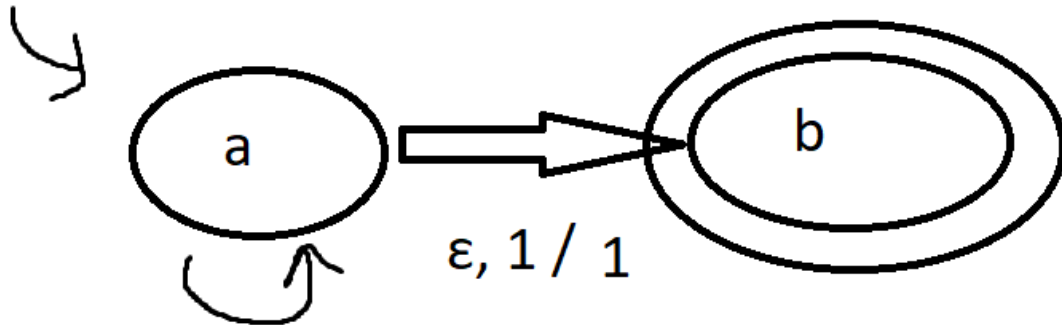


c) Show all reachable IDs with input w 010



Problem 5.

a) PDA for The set of all binary strings with more 1's than 0's. (Transition Diagram)



0, Z0 / 0 Z0

1, Z0 / 1 Z0

0, 0 / 00

1, 1 / 11

0, 1 / ε

1, 0 / ε

b) $0^i 1^j 0^k$, where $i+k = j$.

