## **END EXAMINATIONS - AUGUST**

**B.E:-Computer Science and** Program

Engineering 100 Finite Automata and Formal Languages Max. Marks

Semester

CO1

(06)

**Course Name** Duration 3 Hrs Course Code **CS45** 

## **Instructions to the Candidates:**

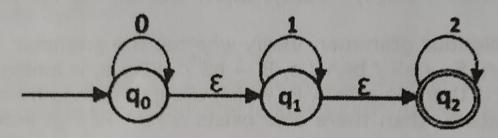
Answer one full question from each unit.

## UNIT - I

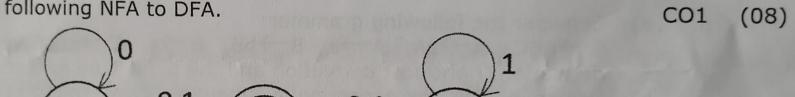
Define DFA. Design a DFA which accepts all strings with a substring 01. (06)CO1 1. a) (06)b)

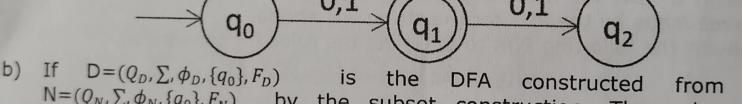
Prove that language L is accepted by some  $\epsilon\text{-NFA}$  if and only if L is CO<sub>1</sub> accepted by some DFA.

(80)CO1 Convert the following  $\varepsilon$ -NFA to DFA. C)



2. Convert the following NFA to DFA.





NFA CO1 (06) $N = (Q_N, \Sigma, \phi_N, \{q_0\}, F_N)$ the subset construction. Then by show  $L_D = L_N$ .

Obtain a DFA to accept

i. L=
$$\{n_a(w) \mod 5=0\}$$
 on  $\Sigma = \{a,b\}$ 

ii. L= $\{n_a \text{ (w)} \mod 3 \neq 0\}$  on  $\Sigma = \{a\}$ 

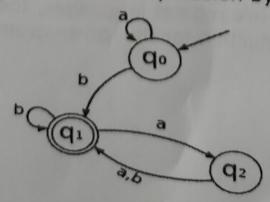
UNIT - II 3.

Obtain the regular expressions to describe the following languages. (i) Strings of a's and b's whose first and last symbols are the same. CO2 (06)(ii)  $L = \{a^nb^n, n > = 1\}$ 

(iii) Strings of 0's and 1's whose lengths are multiples of 3.

Prove that every language defined by a regular expression is also defined by a finite automaton. CO2 (07)

Convert the following into a regular expression by eliminating states. CO2 (07)



- Prove that regular languages are closed under union, complementation (06)4. State the pumping lemma for regular languages. Prove that the set of and difference operations. (05)
  - strings of 0's and 1's of the form ww is not a regular language. CO<sub>2</sub> (09)C)
    - Minimize the DFA given below. Start q0

**UNIT - III** 

- 5. Define PDA. Construct DPDA to accept strings with **CO3** (07)L= $\{x \in \{a, b\}^* \mid n_a(x) = n_b(x)\}$ . Show the moves for the input string abbaba.
  - Define ambiguous grammar. Verify whether the grammar CO3 (05)S  $\rightarrow$  aB / bA, S  $\rightarrow$  aS / bAA / a, B  $\rightarrow$  bS / aBB / b, is ambiguous?
  - Prove that if there is a PDA P<sub>N</sub> which accepts strings from a language L **CO3** (80)by empty stack, then there also exists a PDA PF that accepts L by final state.
- 6. Consider the following grammar: (a)  $S \rightarrow ABC$ ,  $A \rightarrow aA$ ,  $A \rightarrow \epsilon$ ,  $B \rightarrow bB$ ,  $B \rightarrow \epsilon$ ,  $C \rightarrow \epsilon$ . Give the leftmost CO3 (06)derivation, rightmost derivation and the parse tree for the string aabbba.
  - Design a PDA for accepting a language  $\{ww^R | wE(0+1)^*\}$ . Trace the moves made by the PDA for the string w= abbab. CO3 (80)
  - Convert the following PDA to CFG. List the rules for conversion. **CO3**

(06)

$$\delta(q, 1, Z_0) = \{(q, XZ_0)\}$$
 $\delta(q, 1, X) = \{(q, XX)\}$ 
 $\delta(q, 0, X) = \{(p, X)\}$ 
 $\delta(q, \epsilon, X) = \{(q, \epsilon)\}$ 
 $\delta(p, 1, X) = \{(p, \epsilon)\}$ 
 $\delta(p, 0, Z_0) = \{(q, Z_0)\}$ 

**UNIT- IV** a)

7. Obtain the grammar in CNF: S →0A|1B  $A \rightarrow 0AA|1S|1$ CO4 (07)B→1BB|0S|0

- Prove that if L and M are regular languages, then so is L $\cap$ M. Eliminate all unit production for the given grammar: CO4 (07)
  - $A \rightarrow a$ CO4 (06) $B \rightarrow C|b$  $C \rightarrow D$
  - D→E| bC E→dIAb

CO4 (06)

			CO4	(00)
8.	a)	Eliminate all ε production for the given grammar:		100
0.	u	S -> ABC   DD		
		A →BC b		
		$B \rightarrow b \mid \epsilon$		
		$C \rightarrow C \mid \epsilon$ $D \rightarrow d$	CO4	(10)
	h)	For the given grammar:	CO4	(10)
	D)	S->ABC BaB		
		A->aA BaClaaa		
		B->bBbla D		
		C->CAJAC		
		D->E  i) Eliminate E-productions		
		ii) Eliminate unit productions in the resulting grammar.		
		iii) Eliminate any useless symbols in the resulting grammar.		(0.1)
	c)	Define the following:	CO4	(04)
		i. Unit production		
		ii. CNF		100000
		iii. Null-able production		
		iv. Reachable Symbol.		
		UNIT - V		
9.	a)	Write the properties of recursive & recursively enumerable languages.	CO5	(05)
	b)	Obtain a Turing machine to accept the language containing strings of 0's	CO5	
		and 1's ending with 011.		
	c)	Define a Turing Machine. With a neat diagram explain the working of a	CO5	(05)
		Turing Machine.		
10.	a)	Explain in detail about variations of the TM?		
	b)		CO5	
		∑ € { a , b, c }	CO5	(06)
	c)	Define PCP. Verify whether the following lists have a PCP solution.	001	- (06
		(abab) (acabbb) (acb) (ba) (ac)	COS	(06
		(ababaaa), (acabbb), (acab), (ba), (ba), (ab), (acab),		
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